

# Matrix Multi dimensional array

Hello world

Shrinath  
Date  
Page No.

~~Top Spiral~~

Multidimensional array means  $\rightarrow$  Array of Array

Day 1

Array is system to row wise feed data & value

1<sup>st</sup> row    2<sup>nd</sup> row    3<sup>rd</sup> row

given a matrix  $m \times n$

int rows = arr.length;  $\Rightarrow$  length of array

int columns = rows[0].length;  $\Rightarrow$  length of row

H.W. description gfg. question

Day 2

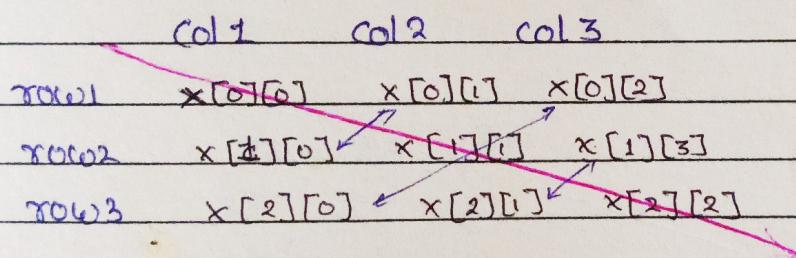
gfg. Sum of upper and lower triangles

Day 3

## Transpose of Matrix gfg

- Transpose of square matrix of size  $N \times N$
- $N \rightarrow$  number of Row and Column
- changing rows to columns and columns to rows

$$\begin{array}{ccc|c}
 1 & 1 & 1 & R \rightarrow C \\
 2 & 2 & 2 & \xrightarrow{C \rightarrow R} \\
 3 & 3 & 3 & 1 \quad 2 \quad 3 \\
 & & & 1 \quad 2 \quad 3
 \end{array}$$



swap until diagonal

```
for ( i=0; i <= n; i++ )
```

```
    for ( j=0; j < i; j++ )
```

```
{ int temp = arr[i][j];
```

```
    arr[i][j] = arr[j][i];
```

```
    arr[j][i] = temp;
```

```
}
```

$O(n)$

Day + Rotate by  $90^\circ$  degrees +

without extra space

Given a square matrix of size  $N \times N$ .

$N \rightarrow$  Number of Row and column

Rotate it by  $90^\circ$  degrees in anti-clockwise

$$\begin{array}{ccc}
 1 & 2 & 3 \\
 + & 5 & 6 \\
 7 & 8 & 9
 \end{array}
 \xrightarrow[90^\circ]{\quad}
 \begin{array}{ccc}
 3 & 6 & 9 \\
 2 & 5 & 8 \\
 1 & 4 & 7
 \end{array}$$

input

O/P

$$\begin{array}{ccc}
 11 & 12 & 13 \\
 21 & 22 & 23 \\
 31 & 32 & 33
 \end{array}$$

$$\begin{array}{ccc}
 13 & 23 & 33 \\
 12 & 22 & 32 \\
 11 & 21 & 31
 \end{array}$$

$$\begin{array}{ccc}
 11 & 21 & 31 \\
 22 & 22 & 32 \\
 13 & 23 & 33
 \end{array}$$

transpose

~~Transpose~~ | changing upper &  
lower most

$$\begin{array}{ccc}
 13 & 23 & 33 \\
 12 & 22 & 32 \\
 11 & 21 & 31
 \end{array}$$

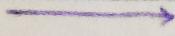
Answer

(n)

9 x 4

$$\begin{matrix} 1 & 4 & 7 & 7 \\ 2 & 5 & 8 & 8 \\ 3 & 6 & 9 & 9 \\ 3 & 6 & 9 & 9 \end{matrix}$$

transpose



$$\begin{matrix} 1 & 2 & 3 & 3 \\ 9 & 5 & 6 & 6 \\ 7 & 8 & 9 & 9 \\ 7 & 8 & 9 & 9 \end{matrix}$$



$$\begin{matrix} 7 & 8 & 9 & 9 \\ 7 & 8 & 9 & 9 \\ 9 & 5 & 6 & 6 \\ 1 & 2 & 6 & 6 \end{matrix}$$

$$\begin{matrix} 2 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{matrix}$$

$$\begin{matrix} 1 & 4 & 7 \\ 3 & 5 & 8 \\ 3 & 6 & 9 \end{matrix}$$

GFG

Q.

Given a matrix of  $R \times C$ ,  
 Traverse the matrix in spiral form.

input  $R = 4, C = 4$

O/P

$$\left[ \begin{array}{cccc} 1 & 2 & 3 & 4 \\ 5 & 6 & 7 & 8 \\ 9 & 10 & 11 & 12 \\ 13 & 14 & 15 & 16 \end{array} \right] \Rightarrow 12348121615141395671110$$

~~$i=0, j=0$~~        ~~$i++, j$~~

$$\begin{matrix} 11 & 12 & 13 & 14 \\ 21 & 22 & 23 & 24 \\ 31 & 32 & 33 & 34 \\ 41 & 42 & 43 & 44 \end{matrix}$$

$i = j++$ , when $j < n$ until $j = n$	$j = i++$ until $i < n$	$i = j--$ until $j \geq 0$
-------------------------------------------	-------------------------	----------------------------

$\boxed{i = j++}$        $i, j++$

Left	Right			
↓	↓			
Top →	1	2	3	4
	5	6	7	8
	9	10	11	12
Bottom →	13	14	15	16

Top = 0

bottom = R - 1

left = 0

right = C - 1

1110

left → right      1 2 3 + top++ top = 1

top → bottom      8 12 16      right-- right = 2

right → left      15 14 13      bottom-- bottom = 2

bottom → top      13 9 5      left++ left = 1

while

( continue loop until ( top &lt;= bottom &amp;&amp; left &lt;= right ) )

```

for( int i = left ; i <= right ; i++ )
    ans.add( matrix[ top ][ i ] );
    top++;
}

```

```

for( i = top ; i <= bottom ; i++ )
    matrix[ i ][ right ];
    right--;
}
if ( top <= bottom )

```

```

for( i = right ; i >= left ; i-- )
    matrix[ bottom ][ i ];
    bottom++;
}
if ( left <= right )

```

```

for( i = bottom ; i >= top ; i-- )
    matrix[ left ][ i ];
    left++;
}

```

	left	right
Top →	↓	↓
1		
5		
9		
bottom ↘	13	

Top = 0

bottom = 3

left = 0

right = 0

while ( top &lt;= bottom &amp;&amp; left &lt;= right )

{

for( left → right ) add matrix[Top][i];

i

Top ++

~~5 9 13~~ for ( Top → bottom ) add matrix[i][right]

right -- ;

top &lt;= bottom      0 &lt;= 3

for

6 12 18

36 10 22

34 28 29

Medium

set matrix zero done ↗  
done ↗ Find first & last position of element in sorted array

1Q. Jump game - did not understand 40.

## Binary Search in 2D Array

### Searching in matrices

linear search      target = 91

	nums		
	0	1	2
0	18	9	12
1	36	-7	91
2	44	33	16

```

for (x=0; x<n; x++)
{
    for (c=0; c<n; c++)
        if ( nums[x][c] == target )
            return {x, c};
}
return {-1, -1};

```

Answer [1, 2]

$N \times N = n^2$  = time complexity  $\Rightarrow O(N^2)$

Q. The matrix is sorted in a row wise & column wise manner

	0	1	2	3	
0	10	20	30	40	
1	17	25	35	45	target = 37
2	28	28	37	49	
3	33	34	38	50	

valid for  $m \times n$  also      case 1      if ( nums[r][c] == target ) return;

case 2      if ( element > target )  
                c--;

case 3      if ( element < target )  
                r++;

```

static int[] search ( int[][] matrix, int target )
{
    int r = 0;
    int c = matrix.length - 1;

    while ( r < matrix.length && c >= 0 )
    {
        if ( matrix[r][c] == target )
            return new int[] { r, c };

        if ( matrix[r][c] < target )
            r++;
        else
            c--;
    }
}

```

### f. Search in a sorted matrix

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	16

1<sup>st</sup> approach linear search

2<sup>nd</sup> approach convert it into a 1d array and then apply binary search

But we are not going to do that

	0	1	2	3
0	1	2	3	4
1	5	6	7	8
2	9	10	11	12
3	13	14	15	16

Tack middle column & perform binary search on it.

let target = 2

mid = target       $\therefore \text{mid} = 6, \text{tar} = 2$

1	2	3	4
5	6	7	8
9	10	11	12
13	14	15	16

$\rightarrow \text{mid} > \text{target}$

so all these element  $>$  target

so we can exclude the rows

case 1 if element = target //ans

case 2 if element  $>$  target  $\rightarrow$  ignore rows after it

case 3 if element  $<$  target  $\rightarrow$  ignore above rows

$\log(N)$

in the end 2 rows are remaining.

1	2	3	4
5	6	7	8

① check whether the mid column you are at contain the answer  
[2, 6]

② consider the four parts

$\log(m)$

$O(\log(N) + \log(m))$