

## Recursion

Patterns ( triangle , \* )

Bubble sort

Selection sort

\* \* \* \*

\* \* \*

\* \*

\*

triangle ( \* , c )

static void triangle ( int r , int c )

{

if ( r == 0 )

return ;

\*

if ( c < r )

{

cout << "

triangle ( r , c+1 ); }

\* \* \* \*

else

triangle ( r-1 , 0 ); } }

for

A

A A

A A A

( 3 , 1 )

( 3 , 2 ) → ( 2 , 0 )

new  
line

( 2 , 1 ) → ( 1 , 0 ) → ( 0 , 0 )

new  
line

new  
line

# dry run

A A A  
A A A  
A A  
A

looks like

bubble sort



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static void bubble ( int arr[], int r, int c )

{

if ( $r == 0$ )  
return;

$r \Rightarrow$  outer loop

$c \Rightarrow$  inner loop

if ( $c < r$ )

{

swap (arr, c);

bubble (arr, r, c+1);

{

else

} bubble (arr, r-1, 0); }

}

Selection Sort - try yourself.

See Kamal's notes

Merge Sort via recursion

You can watch again the video

Prefer Kamal's notes

Mr. Ch. \*

# Quick Sort

Watch video + Prefer Kunal's notes

for code also

in Competitive Programming & Interview use  
Inbuild sorting algorithms.

Array. sort ( arr ); // this is much more efficient  
internally it uses hybrid  
sorting algorithms  
 $O(n \log n)$  performance  
Faster than the one pivot sort

# Recursion -

subset questions & String questions

- String →
- Q. 1. skip 'a' character
  - Q. 2. skip a string
  - Q. 3. Skip a string if it's not the required string

- str.substring ( ) ;
- str.isEmpty ( ) ;
- str.startsWith ("abc") ;

## Subset introduction

\* Permutations & Combinations

\* subsets → non-adjacent collection

$[3, 5, 9] \rightarrow [3], [3, 5], [3, 9], [3, 5, 9], [5, 9], [5], [9]$   
in sets ordering can not change  $[3, 5]$  is same as  $[5, 3]$

in permutations ordering can be changed

Subset → Arrays

Subsequence → Strings

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str = "abc"

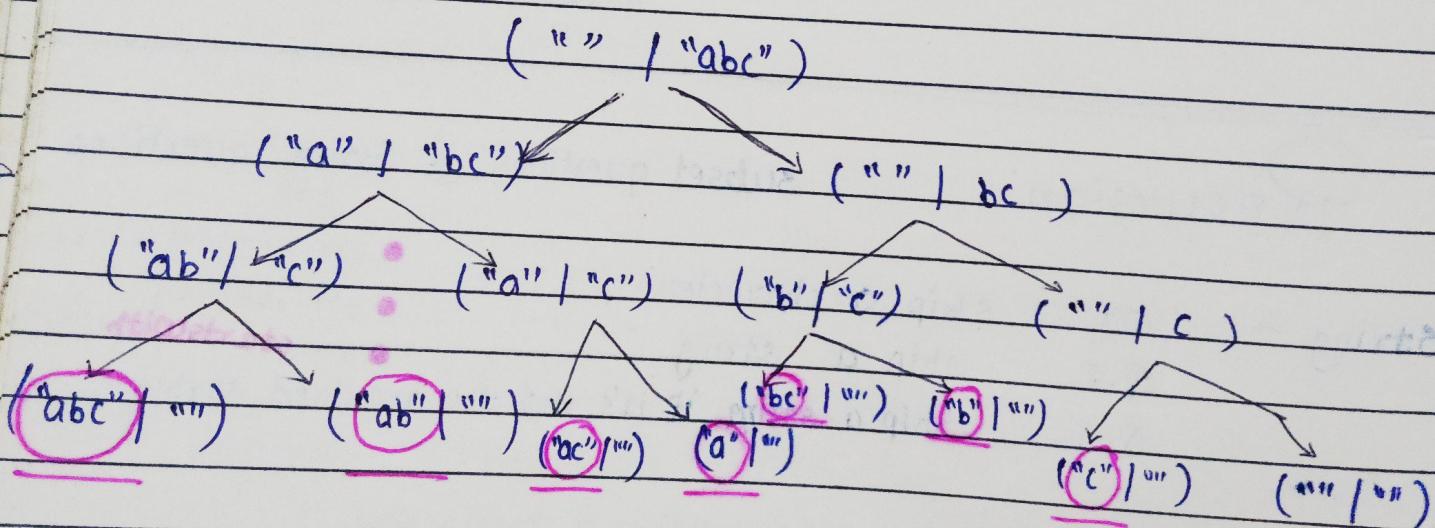
ans = [ "a", "b", "c", "ab", "ac", "bc", "abc" ]

We are trying to create subsets

VI

This pattern of taking some elements & removing some  
is known as Subset pattern

\* (processed, unprocessed)



When in unprocessed → empty, Processed will have answers

Code by yourself for subseq function

# subsequence.java

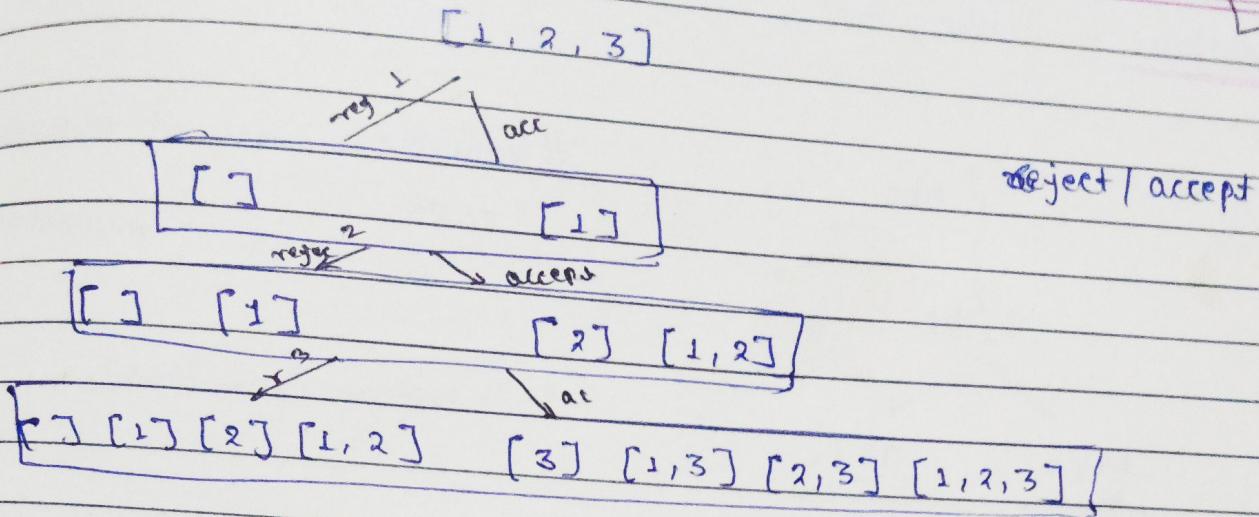
→ instead of printing, return whatever

→ ascii value related q sii, str in A am

# Notes - KUNAL GUPTA

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char ch = 'a';  
sof (ch + 0);  
↓  
97



```
static List<List<Integer>> subset (int[] arr) {
```

```
    List<List<Integer>> outer = new ArrayList<>();
```

```
    outer.add(new ArrayList<>());
    for (int num : arr) {
        int n = outer.size();
        for (int i=0; i < n; i++)
```

```
        List<Integer> internal = new ArrayList<>(outer.get(i));
```

```
        internal.add(num);
```

```
        outer.add(internal);
```

}

return outer;

$O(N * 2^N)$  time complexity

Space  $O(2^N * N)$  total subsets → space taken by each subset

#subset.java

notes after 2 pages →

for duplicates

check → KK victory

(10 minutes)

#SubsetDub.java

str = abc

permutations = [ abc, bac, cab, bca, acb, cba ... ]

$$n! \text{ permutations you can have}$$

$$n=3 \quad 3! \Rightarrow 6$$

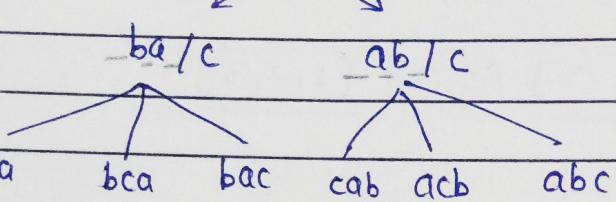
"" / abc

calls.

p+1 (processed+1)

↓  
a/bc

// Variable number of recursive



Program from Kunal

[ ]  
inclusive    exclusive

# PerRec.java

static void permutations ( String p , String up )

{ if ( up.isEmpty() )

{ System.out.println(p);  
return;

}

# PerRecRet.java

char ch = up.charAt(0);

for( int i=0 ; i&lt;=p.length(); i++ )

}

String f = p.substring(0, i)

String s = p.substring(i, p.length());

permutations ( f+ch+s , up.substring(1));

}

Try code for

returning ArrayList

& count the permutations  
code

or watch again Kunal's

code

returning no. of permutations - } we can return by formula but

using permutations fun. also we can do, without p  
counting

// not getting all the thing  
// Give it a dry letter

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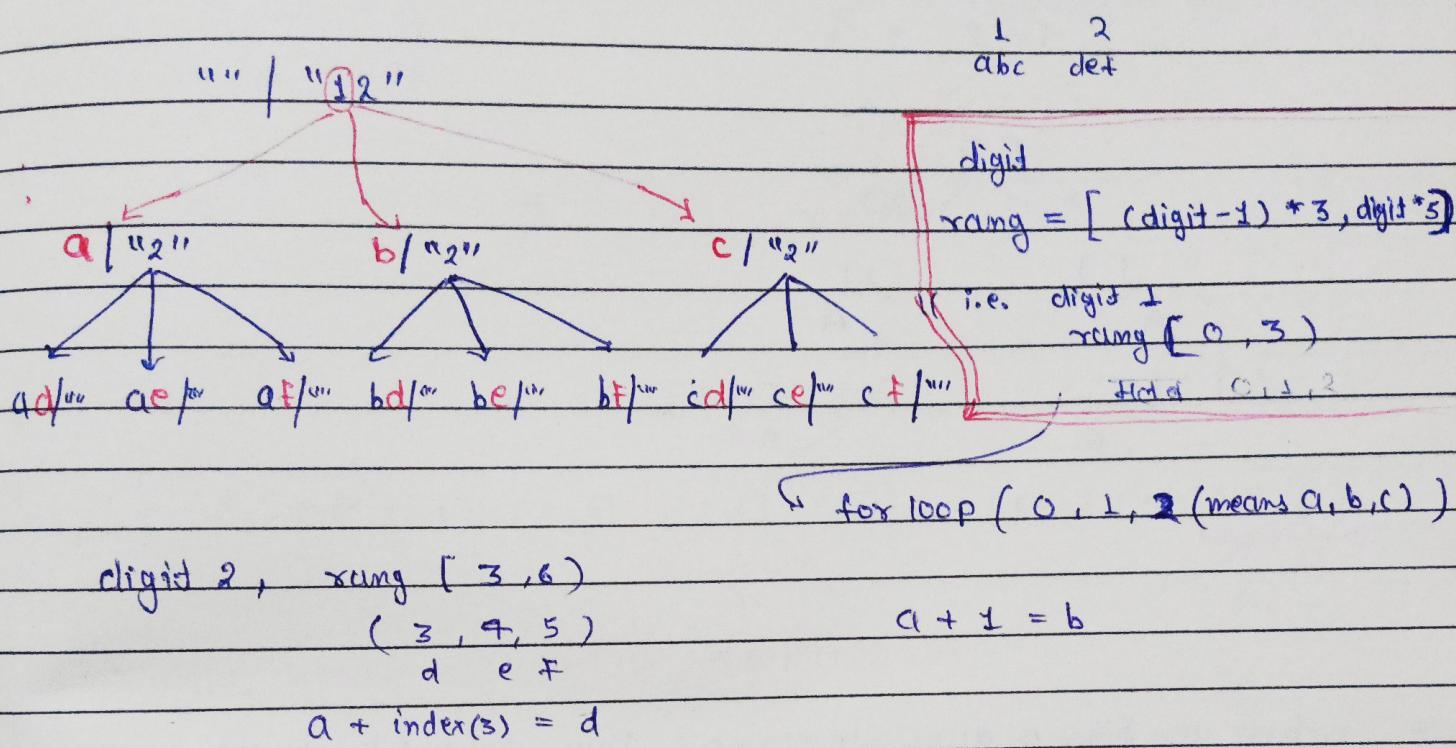
Recursion  
4 videos  
वादा

## Combination questions

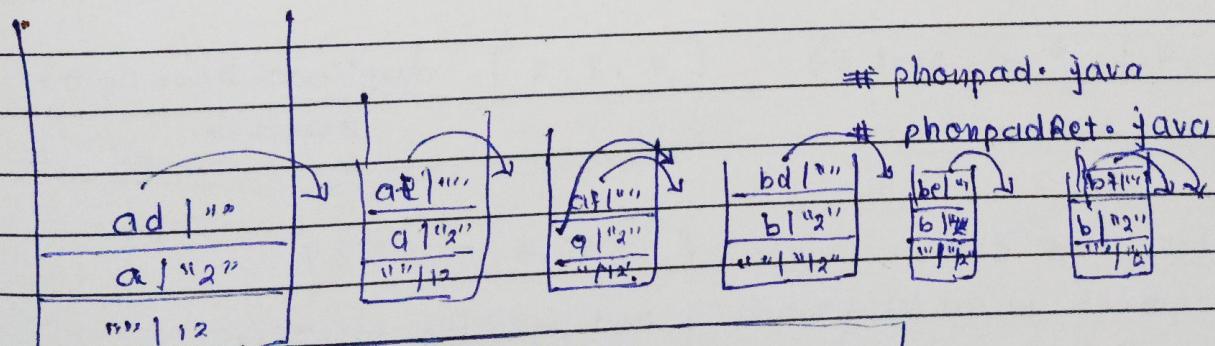
Coming back

13 Jan 2024

### Leet 12. Letter combinations of a phone number



character to add in processed = 'a' + index



permutation using iteration → needs queue

Adv. Sliding window, two pointer → hashmap, heap

Recursion is used everywhere - arrays

- strings, vector, linkedlist, trees, graph

dp

Tack V  
Forward Complexity tip

if constraint  $10^5$   
& we use  $O(n^2)$   
Hence

$(10^5)^2 \Rightarrow (10^{10})$  operation  $\Rightarrow$  TLE

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Subset with duplicate value

[+, 2, 2]

A - accept  
R → reject

A [+]  $\xrightarrow{A^1}$

R [ ]  $\xrightarrow{A^2}$

[+] [+]  $\xrightarrow{R}$  [2] [1, 2]  $\xrightarrow{A^2}$

[+] [+] [2] [1, 2] [2] [1, 2] [2, 2] [1, 2, 2]   
 duplicate set

इसमें add नहीं करो, जैसा होगा ।

\* when you find a duplicate element only add it in the newly created subset of previous step.

\* [2, 1, 2]  $\Rightarrow$  to sort करी [1, 2, 2] duplicates have to be together.

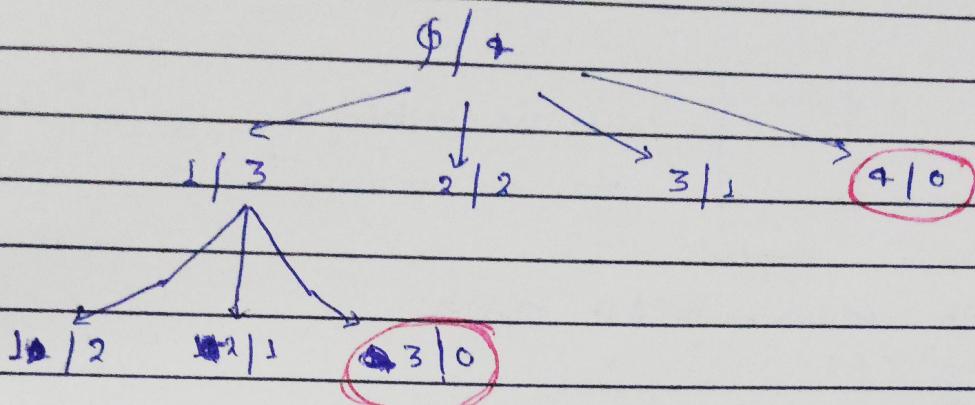
you have some data

[1, 2, 3, 4, 5, 6]

in the answers you are taking something and ignoring something  
This follows processed and unprocessed

Ans [1, 2, 3, 4, 5, 6] , possible combinations for making

Ans [4, 31, 211, 1111, 22, ... ]



```
static ArrayList<String> diceRet (String p , int target )
```

```
}
```

```
if (target == 0)
```

```
{ ArrayList<String> list = new ArrayList<>();
```

```
list.add(p);
```

```
return list;
```

```
}
```

```
ArrayList<String> list = new ArrayList<>();
```

```
for (int i=1 ; i <= 6 && i <= target ; i++ )
```

```
{ list.addAll(diceRet( p+i , target-i ));
```

```
}
```

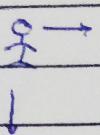
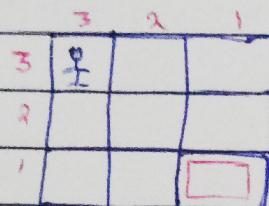
```
return list;
```

```
}
```

# Backtracking + Maze Problems - Theory + code + Tips

KK

why backtracking, how to know backtracking will be used

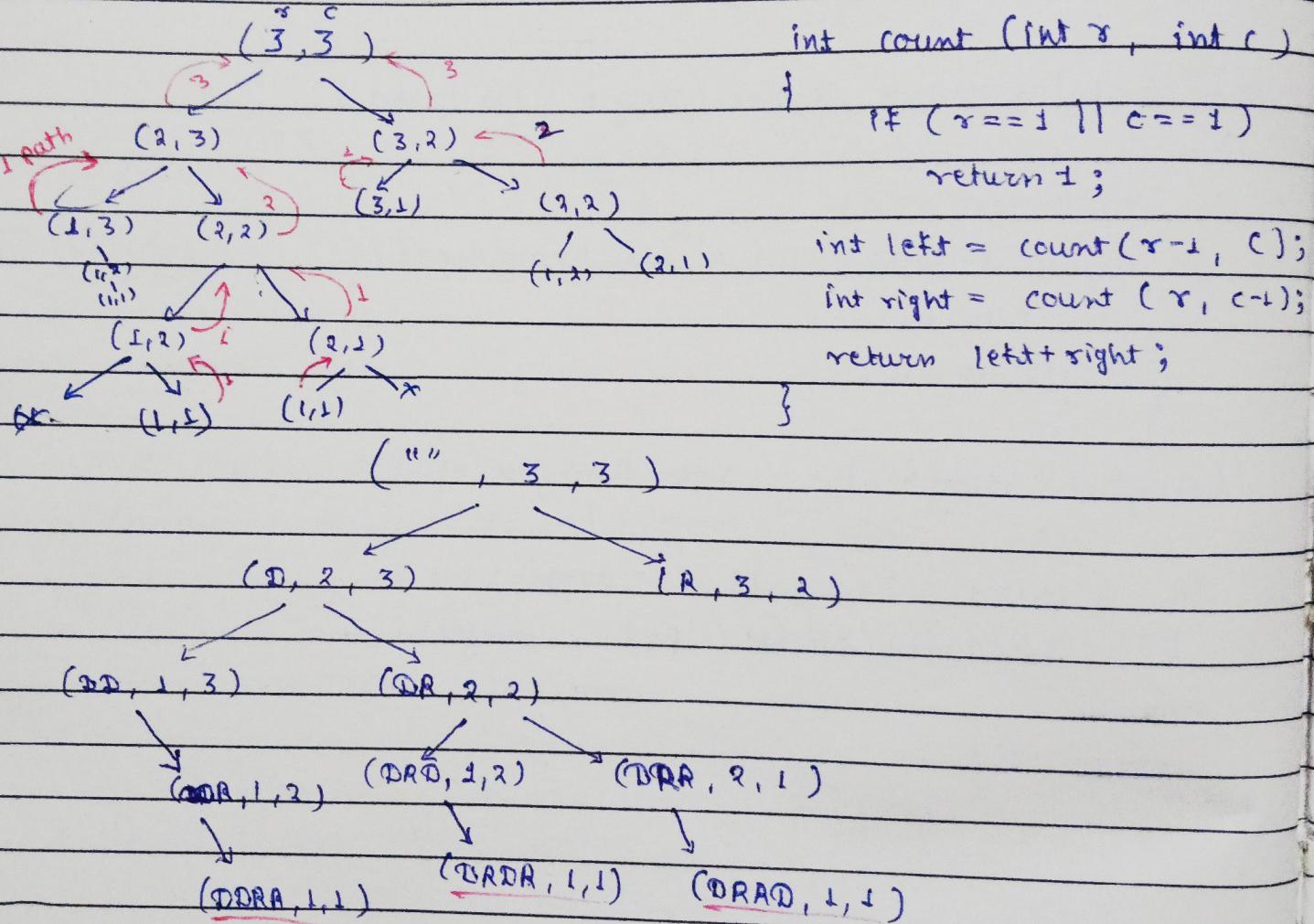


what we are doing?

How we are doing?

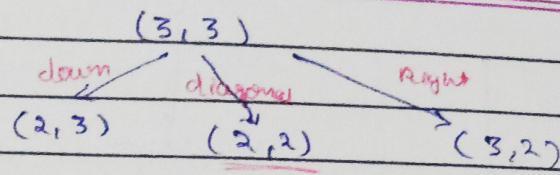
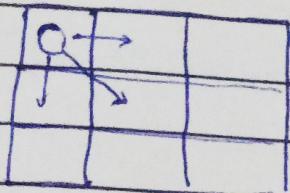
Ans → RADD TDRA RDAD RDAD CADR

O/P → no. of paths (6)



```

void printp (String p, int r, int c)
{
    if (r == 1 && c == 1) { s.o.p (p); return; }
    if (r > 1) { printp (p + 'D', r-1, c); }
    if (c > 1) { printp (p + 'R', r, c-1); }
}
    
```

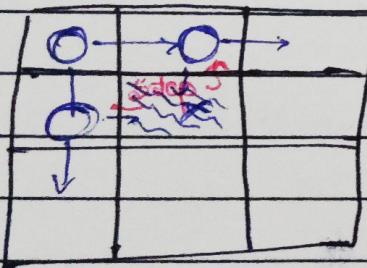


```

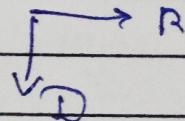
    if ( x == 1 && c == 1 ) {
        condition = true;
        return;
    }
    if ( x > 1 && c > 1 )
        list.add( pathRetDiagonal( p + 'D', x - 1, c - 1 ) );
    if ( x > 1 )
        list.add( pathRetDiagonal( p + 'V', x - 1, c ) );
    if ( c > 1 )
        list.add( pathRetDiagonal( p + 'H', x, c - 1 ) );
}

```

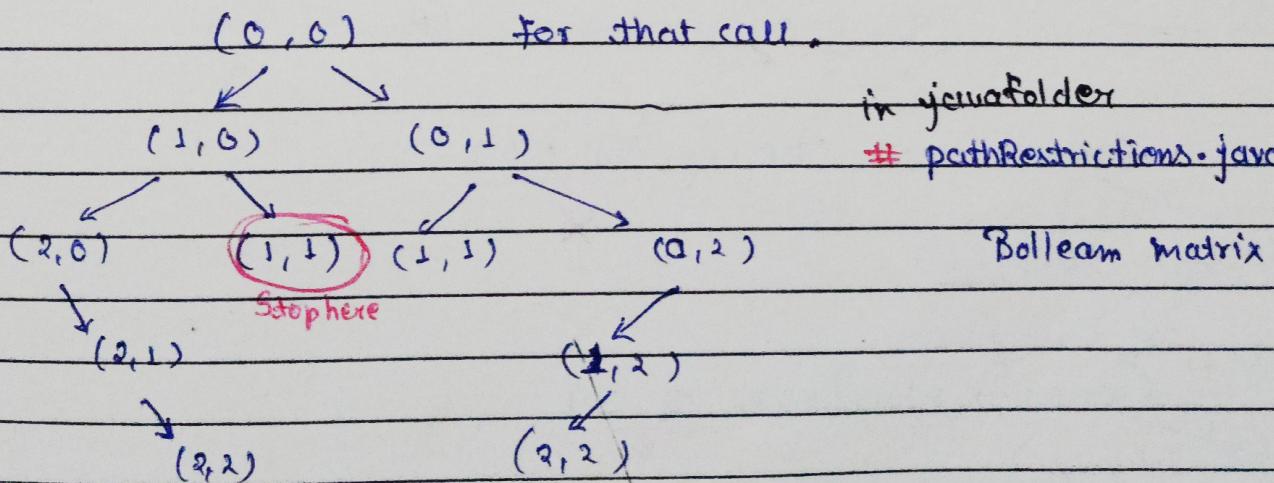
## Maze with Opticals



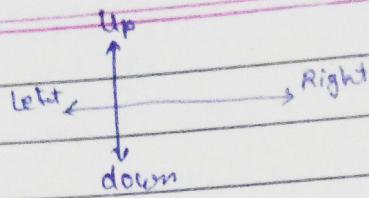
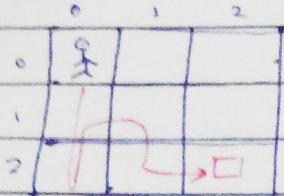
False-river Boolean matrix



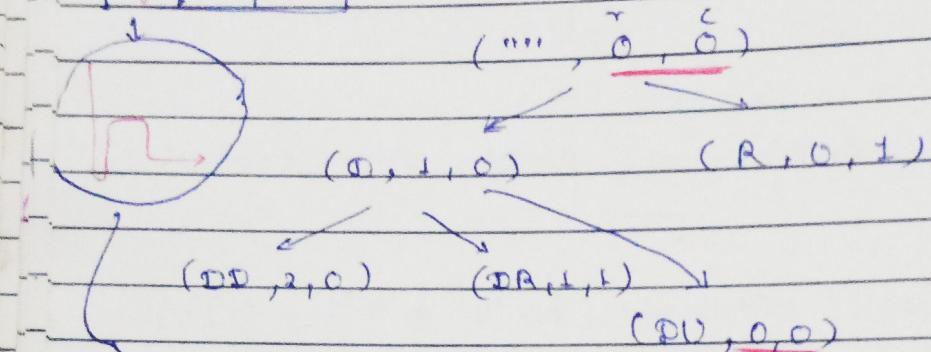
Note - check, when you land on a new cell, check whether is river or not  
 If you land on river, stop recursion for that call.



## Problem



all directions are allowed



pathAll

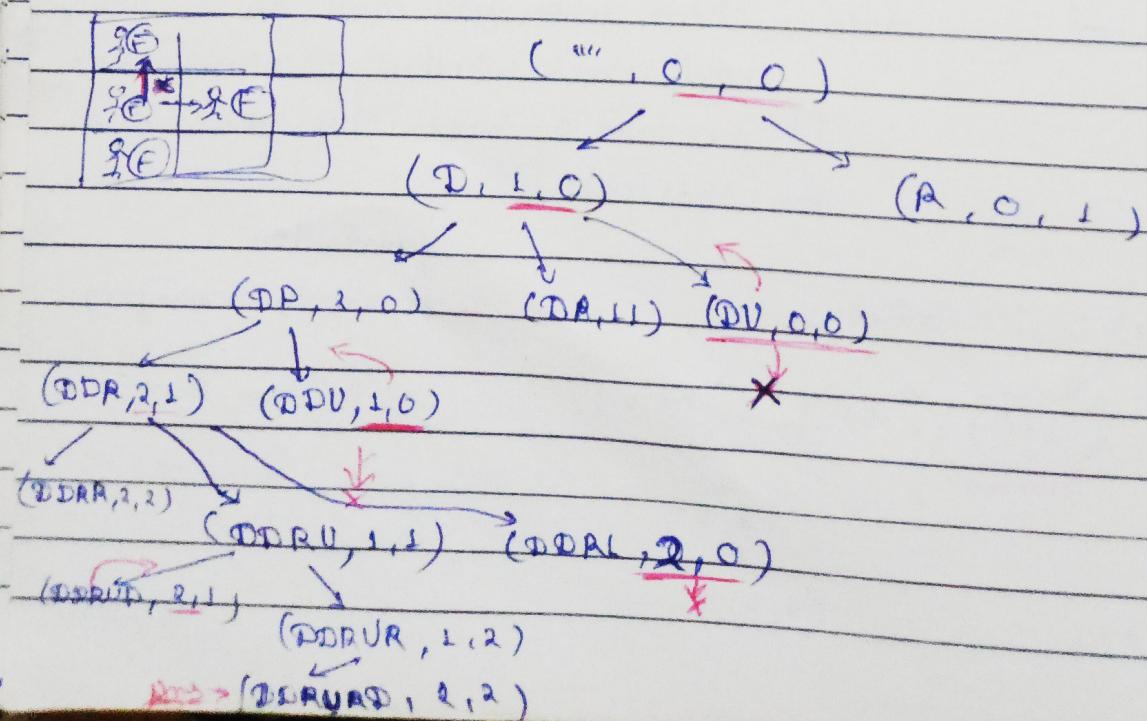
can not be an answer.

- Do not move back the same path. This is known as backtracking problem.
- it's coming back to the same point where it was started, here it's stucking in endless loop.

## How to solve

All cells that are visited mark those as false.

so that it does not go there



marking false == I have that cell in my current path.

So, when that path is over. ex: you are in another recursion call, those cells should not be false.

\* while you are moving back, you restore the maze at it was.

When do we go back  $\Rightarrow$  happens when function is ~~executed~~ returned.

When you come out of the recursive function  $\rightarrow$  you are now in above recursion call. Hence re-mark the cell as true.

→ This is known as Backtracking

When the function is return change the changes that you made previously

Backtracking # allpath.java

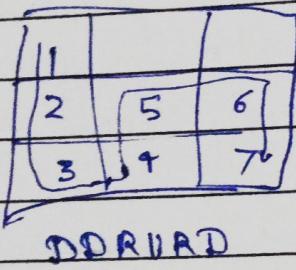
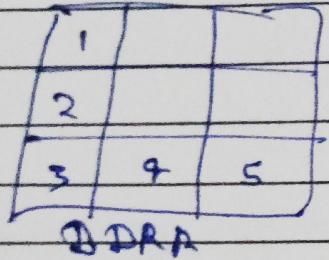
# = ~~callpathPath.java~~

Backtracking -

you are making some changes while going in the below recursion calls, So when you go outside of those recursion calls the changes that made by those recursion calls should also not be available

thought process - make a change, reverse the change, when that work is done

Q. give the O/P of previous problem, as print the path in the matrix formate. matrix and the path



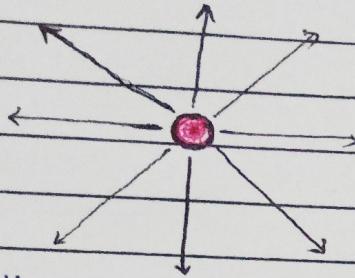
- \* Take a step variable
- \* update the path array
- \* print it in base condition
- \* Backtrack

pathprint  
# pathAllMatPrint.java  
in your folder

 queen (Q)

N-queens, N-knights - Sudoku Solver

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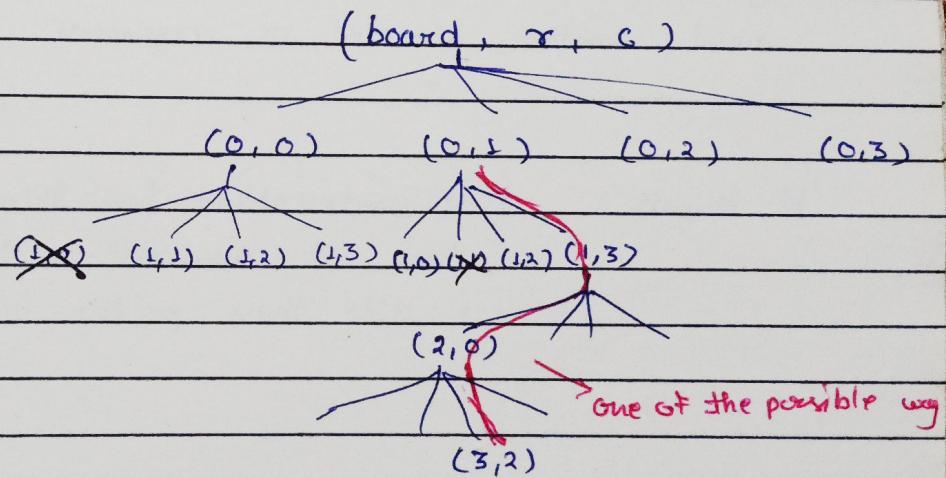
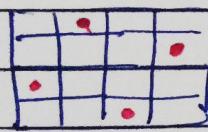
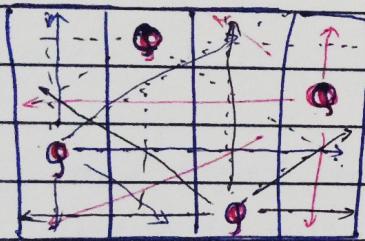
directions - in which queen can attack

$n = 4$ , mean the question says you are given a  $n \times n$  board and you have to place  $n$  queens on it.



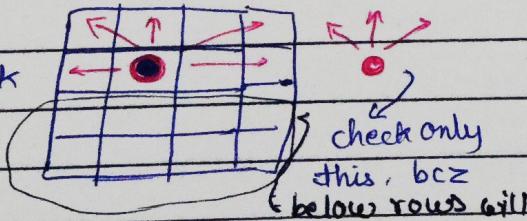
Find the total number of ways to place the queens.

→ Place queens so that no two are eliminating one another.

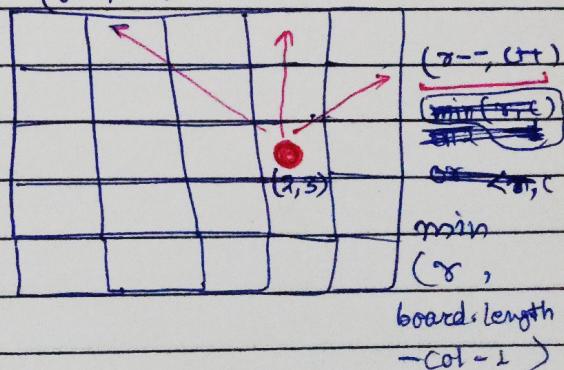


How to check whether the particular condition is safe or not?

To check



$O(n^2 \times C)$   
 $\min(r, c) \text{ at } 1, 2$   
 $(r--, c--)$



Code  
= NQueens.java in java folder

Knights  $\Rightarrow$    $\Rightarrow$  houses

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Time Complexity:

Recurrence Relation:

$$T(N) = N * T(N-1) + O(N^2)$$

Akraha Bazzi formula

$$O(N^3 + N!) = O(N!)$$

Time complexity lecture

Tip  $\Rightarrow$  you can eliminate for loops with conditions, but then you need another variable in argument

K-Knights, first understand how thing work, knight works  
then if available on lee, give it a try  
if not else then prefer vidic.

feetcode 79 world search

Leetcode 37

## Sudoku Solver

9x9 का board होता है, 😊, ये ज्ञानी पता करें question  
+ पता solve होता है 😊

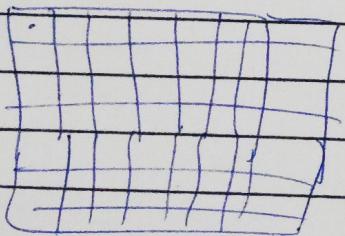
⇒ Sudoku puzzle is 9x9 board, where you have to put the numbers from 1 to 9 with some rules.

1. each of digits 1-9 must occur exactly once in each row.
2. each of digit 1-9 must occur exactly once in each column.
3. ——— once in each of 9, 3x3 sub-boxes of the grid

1 - 9  
must once occur

same

same



we will firstly see what are empty indexes,

and we will try to check from 1-9 which even will be suitable we will put

letter am if we realize, (how we will realize if my last grid end of column will remain empty and I am unable to put the prob number which I have)

put something wrong, and I need to change it, for this I will do backtrace and whatever choices I have made I will check them and I try to put that number there

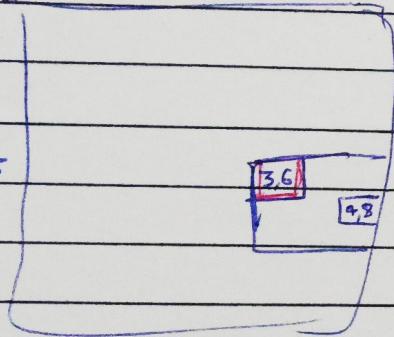
Note - When a choice can effect future answers  
use Backtracking

You are given a incomplete sudoko

isSafe function  $\Rightarrow$

rows, columns  $\rightarrow$   $\square$

How will you find the  $3 \times 3$  box, mean the indexes  
where we are currently at, where does they belong.



$(3, 6)$  where it's box is starting

$$\text{So row} \Rightarrow \text{row} - \text{row} \% 3 \Rightarrow \frac{3}{3} \text{ bcz } 3 - 3 \% 3 = 3$$

$$\text{Col} \Rightarrow \text{col} - \text{col} \% 3$$

$$6 - 6 \% 3 = \underline{\underline{6}}$$

Helpful  
there  
box start  
from  $(3, 6)$

$(7, 8)$  where it's box is starting

$$\text{row} = \text{row} - \text{row} \% 3 \Rightarrow 7 - 7 \% 3 = 3$$

$$\text{col} = \text{col} - \text{col} \% 3 \Rightarrow 8 - 8 \% 3 = 6 \rightarrow (3, 6)$$

We need it bcz we are ~~not~~ passing only board  
not the small box <sup>in our arguments.</sup> indexes.

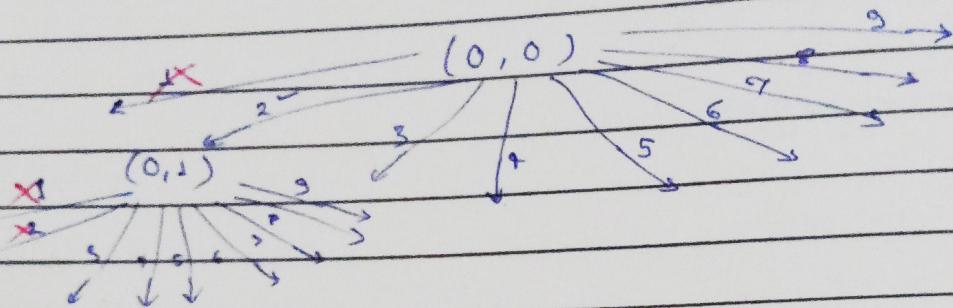
static boolean isSafe ( int[][] board, int row, int col, int num )

{     // three thing we will check  $\rightarrow$   $\square$   $\blacksquare$  }

static boolean solve ( int[][] board ) {

Shradha Thapar, Better explanation than Kunal Trivedi for  
Sudoku

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public boolean helper (char[][] board, int row, int col)

{ if (row == board.length) { return true; }

int nrow = 0;

int ncol = 0;

if (col != board.length - 1)

{ nrow = row; ncol = col + 1; }

else

{ nrow = row + 1; ncol = 0; }

if (board[row][col] == '-') { if (helper (board, nrow, ncol)) return true; }

else {

for (int i=1; i<=9; i++)

{ if (isSafe (board, nrow, ncol, i))

{ board[nrow][ncol] = (char)(i+'0');

if (helper (board, nrow, ncol))

return true;

else board[nrow][ncol] = '-';

} } } } isSafe function → write by yourself

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