

Backtracking 0/1 knapsack problem.

$$P = [11, 21, 31, 33, 43, \overset{54}{\cancel{53}}, 55, 65] \quad m = 110$$

$$W = [1, 11, 21, 23, 33, 43, 45, 55] \quad n = 8$$

$$P/W = [11, 1.9, 1.47, 1.43, 1.3, 1.23, 1.22, 1.18]$$

Greedy solution -  $x = [1, 1, 1, 1, 1, 21/43, 0, 0]$   $\sum$

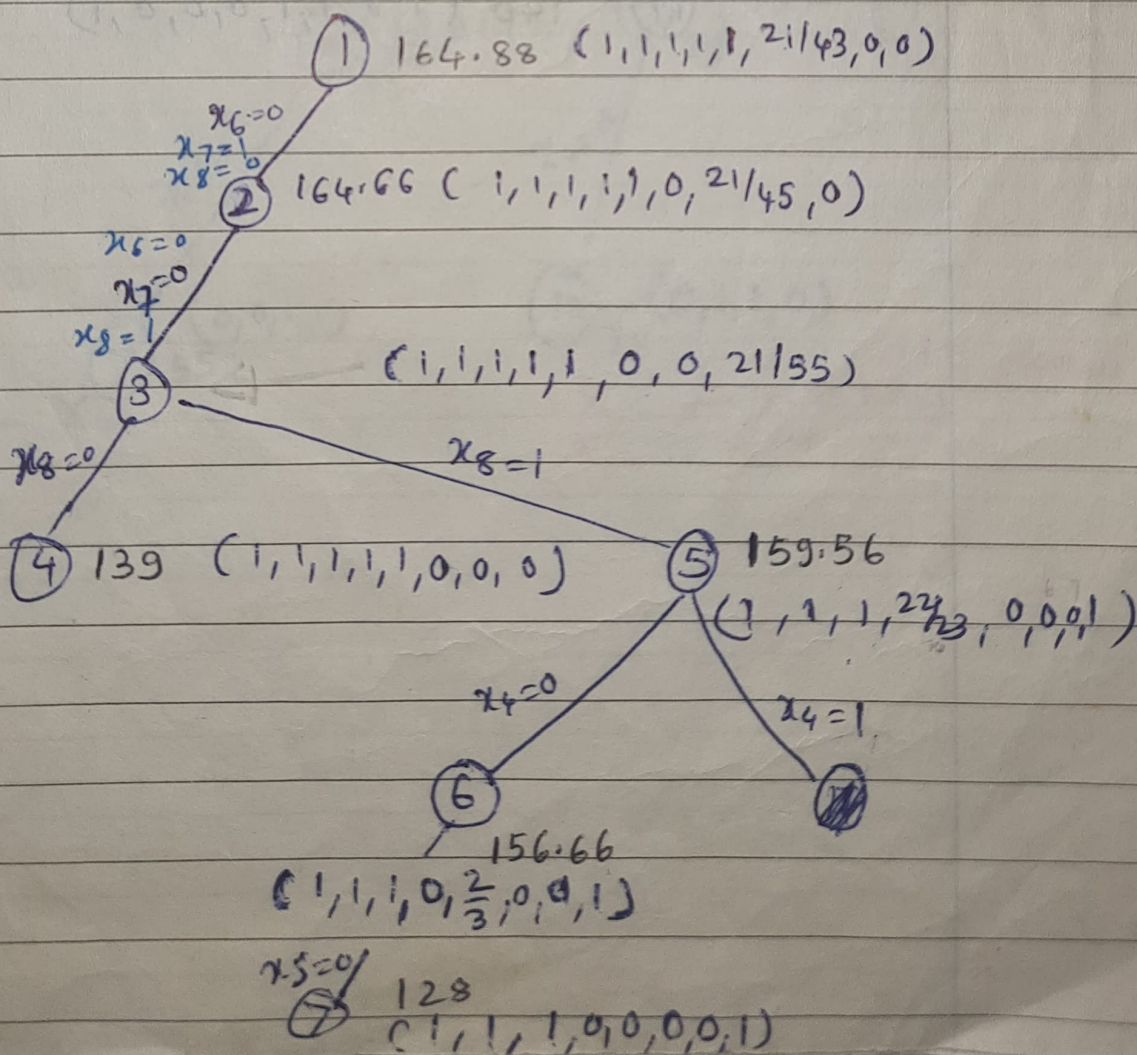
$$\begin{array}{l} 11 + 11 + 21 \\ + 23 + 33 \\ = 89 \\ \text{Rem.} - \\ 21 \end{array}$$

$$P_{\text{profit}} = 11 + 21 + 31 + 33 + 43 + \cancel{53} \times \cancel{21} + \frac{54 \times 21}{43}$$

$$= 142 + 24.71 = 164.88$$

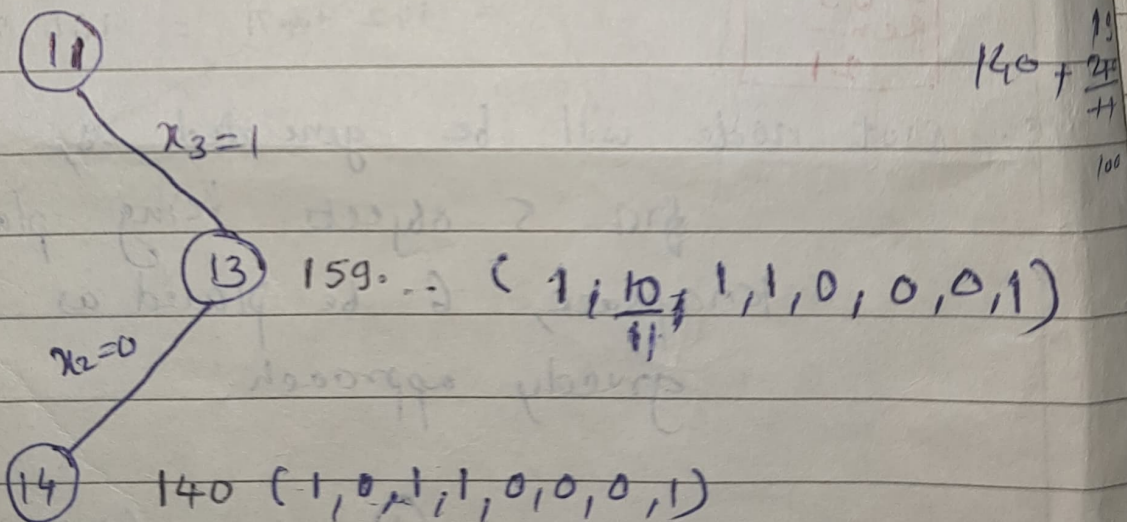
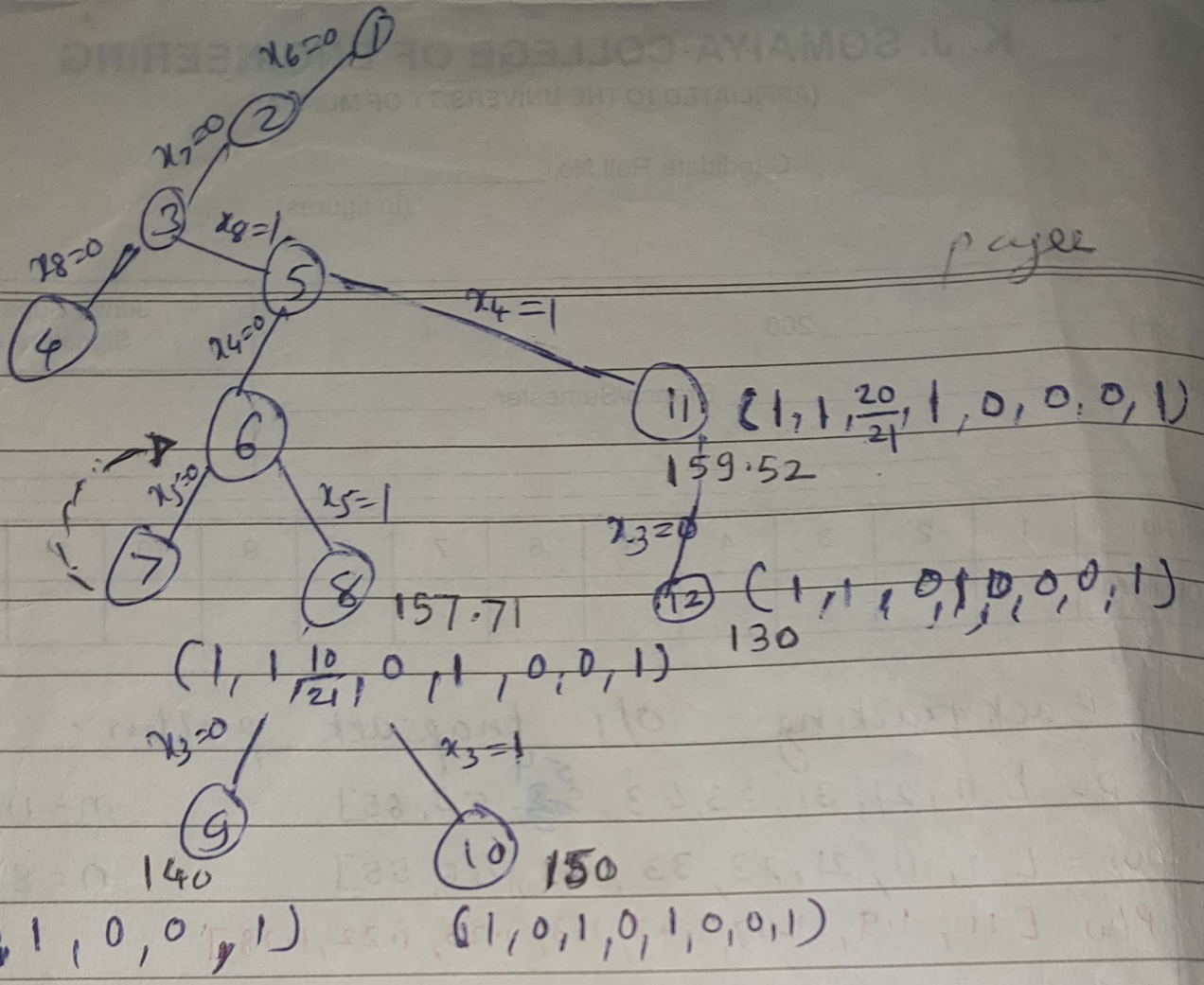
i.e. root node will be generated as,

first 5 objects being placed in knapsack, 6 be placed as  $21/43$  by greedy approach.





page

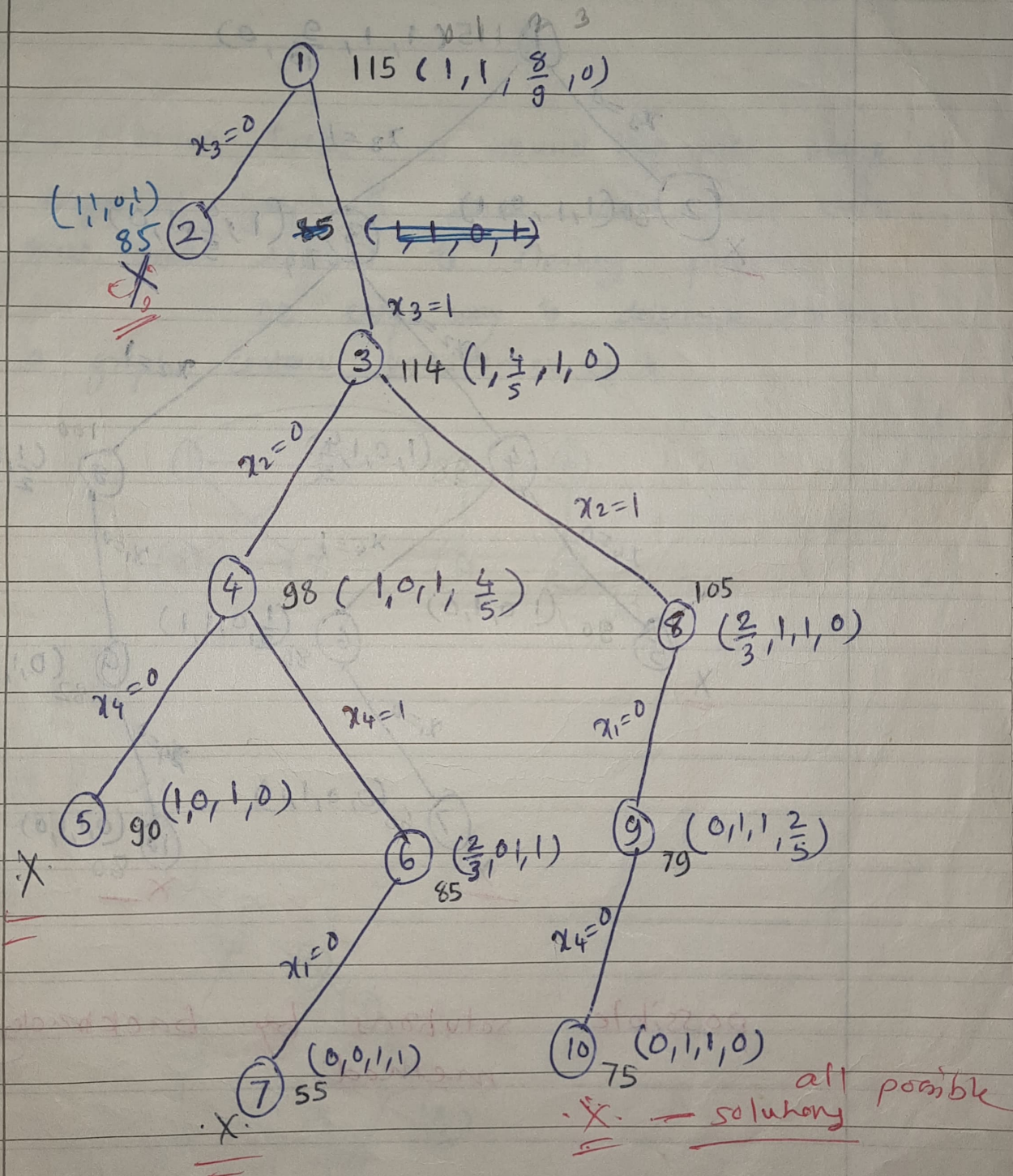




$$m = 16, n = 4, P = [45, 30, 45, 10]$$

$$P_w = [15, 6, 5, 2]$$

$$w = [3, 5, 9, 5]$$



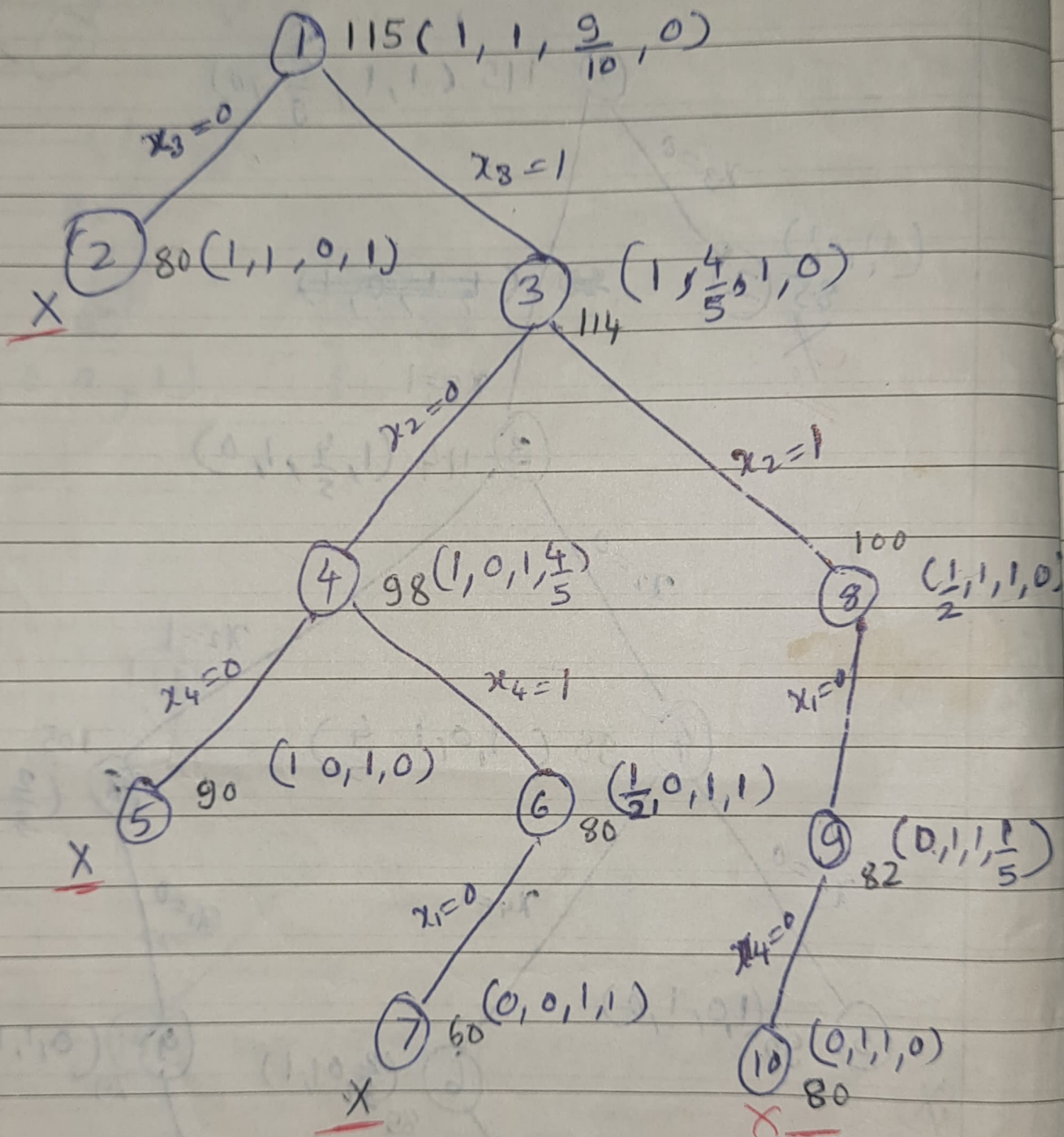


$$n=4, m=16$$

$$P = [40, 30, 50, 10]$$

$$w = [20, 6, 5, 2]$$

$$w = [2, 5, 10, 5]$$



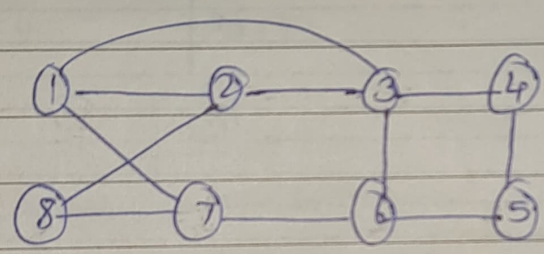
possible solutions by backtracking method.

m  
2) knapsack backtracking.  
p = [25 45 12 7 6 10 5]  
w = [5 11 3 2 2 7 4]

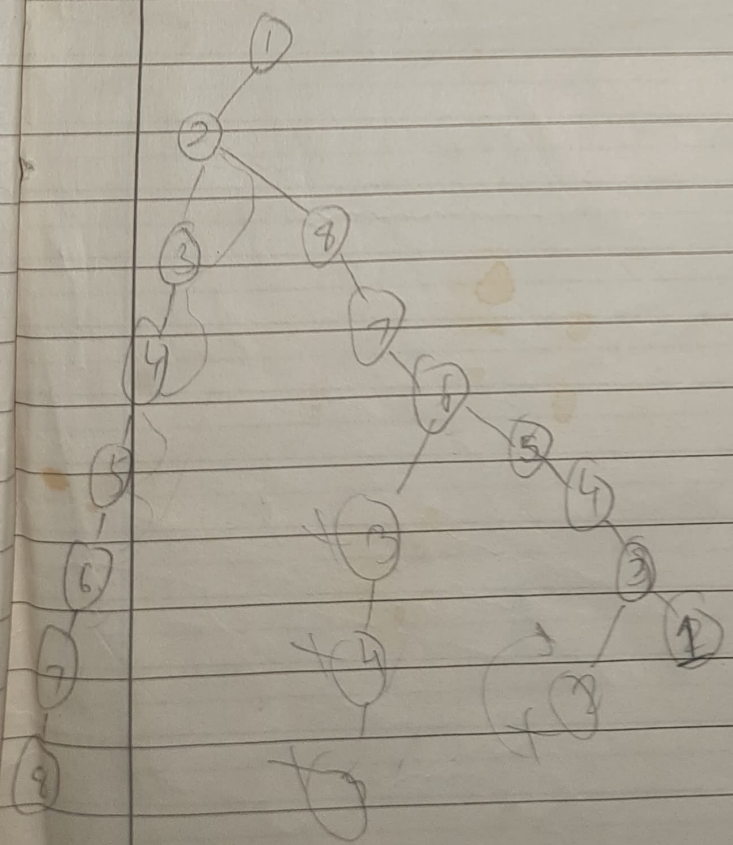
n = 6  
m = 15

Hamiltonian cycles - a round trip path along 'n' edges of graph G that visits every vertex once and returns to starting position.

- no easy way to determine whether a graph contains hamiltonian cycle



hamiltonian cycle.  
1, 2, 3, 4, 5, 6, 7, 8, 1  
~~+ 7 6 5 4 3 2~~  
1 3 4 5 6 7 8 2 1





0/1 Knapsack problem using Branch & Bound

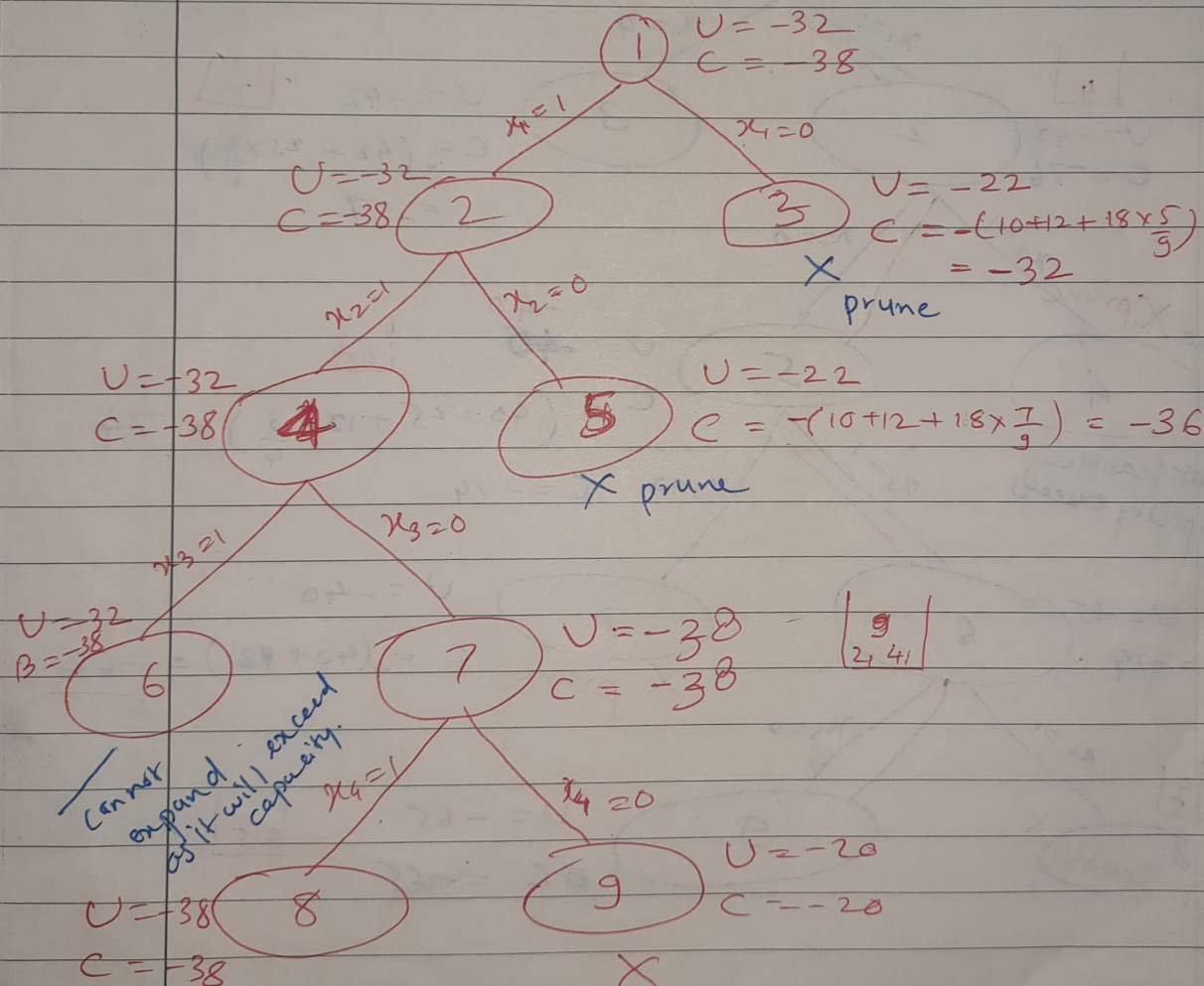
$n=4$ , Profit:  $\{10, 10, 12, 18\}$   
 $M=15$  Weight:  $\{2, 4, 6, 9\}$

Initial upper bound  $U = 32$   $\boxed{2, 4, 6}$

(By adding first 3 elements in bag)

Cost function:  $-10 + 10 + 12 + 18 \times 3/9$   
 $C = 38$

— Negate the values —



Hence soln is  $\{1, 1, 0, 1\}$

If entire tree is explored complexity is  $O(2^n)$   
 best case possible if only one path explored,  $O(n)$

$n=4$   
 $M=10$

Profit  $\{ 40, 42, 25, 12 \}$   
 Weight  $\{ 4, 7, 5, 3 \}$   
 $\frac{P}{W} : \{ 10, 6, 5, 4 \}$

Initial upper bound  $V: 40$

Cost:  $40 + 42 \times \frac{6}{7}$

