

At the beginning I want to say that files in the first exercise are generated in a way:

<first element><space character><second element><space char>...<nth element><space><slash at the end of the row>

4 files with different graph representations: -adjacency matrix;

-incidence matrix;

-edge list;

-incidence list;

In the second exercise I am using adjacency matrix representation.

Final conclusions:

Edge list – the slowest one – searching through edge list is a linear search of complexity $O(E)$ but it gets $E*[el1,el2]$ sublists in memory.

Advantages: easy to implement; take small amount of RAM to storage

Disadvantages: takes a great deal of time to find an edge in it

Adjacency matrix – a symmetric matrix which takes $O(V^2)$ RAM memory and give result immediately $O(1)$ complexity

Advantages: fast searching; easy to implement; it takes a lot of RAM but it is not that bad

Disadvantages: I think that memory allocation can be a minus when dealing with big matrices with little number of edges

Adjacency/Incidence List – list in which every sublist contains vertices connected to the vertice which number is a sublist's index in an Incidence list

Advantages: fast searching elements in constant $O(1)$ time(in the worst case $O(d)$ where d is the degree of vertex i in sublist because that's the of i 'th sublist. .

Disadvantages: needed a good imagination to implement it properly. Memory storage in the worst case is $O(2E)$

Incidence Matrix – matrix in which every row consists of two 1's and a bunch of zeros which indicates which vertices are connected with each other.

Advantages: easy to implement; slightly better than edge list $O(E)$ search complexity

Disadvantages: very big memory allocated $O(E*V)$ it takes long to generate a file with it.

For the second exercise I have chosen the adjacency matrix for graph representation because of it's $O(1)$ complexity of searching edge in it and because it is easy to implement such a matrix.

Topological sorting in my program uses Breadth First Search traverse through a graph with time complexity $O(E+V)$ I have created a Directed Acyclic Graph using property of adjacency matrices such that if you randomly fill a matrix below the main diagonal(with saturation of 0.3) You will always get an adjacency matrix representing Directed Acyclic Graph. Then I have measured the time of sorting this graph in the topological order and put it on a graph with $\log(n)$ and $n\log(n)$ line graph for comparison purposes.

Here are the graphs for the first and the second task:

