

一. 解:

(1). 设面试到第 k 名学生.

$$t_k = \frac{k}{n} \left(1 - \frac{(k-1)}{n-1}\right) = \frac{k(n-k)}{n(n-1)}$$

$$q_k = \frac{k}{n} \left(\frac{k-1}{n-1}\right) = \frac{k(k-1)}{n(n-1)}$$

(2). ~~若录取 k 名之后的第一名. $P_1 = \frac{2}{n-k+1} \cdot \frac{1}{n} \cdot \frac{n-m}{m-1} \cdot \frac{k}{n} = \frac{k}{n} \cdot \frac{1}{m-k+1}$~~

若录取 k 名之后综合第一名. $P_1 = \frac{n}{n-k+1} \cdot \frac{1}{n} \cdot \frac{n-m}{n-m} \cdot \frac{k}{n} = \frac{k}{n^2} \cdot \frac{1}{m-k+1}$

若... = 名. $P_2 = \frac{k}{n^2} \cdot \frac{1}{m-k+1} \cdot \frac{1}{m-2}$

$$V_k = \max\{P_1, P_2\}$$

$$\text{则 } \Delta V_k = \frac{2k-2n}{n^2} + \frac{n-1}{n^2} \cdot \frac{1}{m-k} \cdot \frac{1}{m}$$

(3) 由(2)知 k 名后用策略 $\Delta V_k = \frac{2k-2n}{n^2} + \frac{n-1}{n^2} \cdot \frac{1}{m-k} \cdot \frac{1}{m}$.

$$\text{令 } \Delta V_k - \Delta V_{k+1} = \frac{1}{n^2} \cdot \left(2 - \frac{n-1}{k+1}\right) \quad \text{一类二阶导}$$

$\therefore \Delta V_k - \Delta V_{k+1}$ 近似于 0 时 k 的值即为最优解.

$\therefore k = \lfloor \frac{n+1}{2} \rfloor$ 时 ΔV_k 最小.

二. 解.

$$(1) m_1 + m_2 + \dots + m_r = n-1$$

$$\begin{cases} m_i \geq 1 & (1 \leq i \leq r) \end{cases}$$

$$m_1 \leq \left\lfloor \frac{n - \sum_{i=2}^r m_i}{2} \right\rfloor$$



(2) $m_1 + \dots + m_r = 3$

有两种情况: $m_1 + m_2 + m_3 = 1 + 1 + 1 = 3$ $m_1 + m_2 = 2 + 1 = 3$

① 队1水平为 V_1

$$P_1 = \sum_{j=2}^4 \frac{V_1}{V_1 + V_j} \cdot \frac{1}{3} + \frac{1}{2} \left(\frac{V_1}{V_1 + V_j} \right) \left(\frac{V_1}{V_1 + V_k} \right) + \frac{1}{6} \frac{V_1}{V_1 + V_i}$$

j 与 k 不是同一支 i 为任意一支

② $P_2 = \frac{V_1}{V_1 + V_j} \cdot \frac{V_1}{V_1 + V_k}$
 j 与 k 不是同一支

假设其他队伍水平相同. 设其他队伍水平平均值为 V .

$$P_2 - P_1 = \frac{V_1 V (V_1 - V)}{6(V_1 + V)^3} > 0$$

~~若队1水平高则~~

$\therefore P_2$ 大于 P_1

采用第二种赛制获胜概率最大, 第一种最小.

(3) $f_1 = 1 - \frac{2kV_1}{n(V_1 + V)}$

$$f_2 = \left(\frac{j}{n} \cdot \frac{V}{V_1 + V} \cdot \frac{n-j}{n} \right) \left(\frac{2(k-j)}{n-j} \cdot \frac{V}{V_1 + V} + \frac{n+j-2k}{n-j} \right)$$

$$f_1 - f_2 = \frac{j}{n(n-j)} (k-j) \cdot \frac{2V_1}{V_1 + V} \cdot \frac{V - V_1}{V_1 + V} \quad \therefore f_1 > f_2$$

(4) 证明: ~~$n = 2^s$~~ $n \in [2^{s+1} + 1, 2^{s+1} - 1]$. 设 $k = n - 2^s$

由(3)知胜率大, 轮数定少.

\therefore 首次比赛队伍数量不超过 $\lceil \frac{n}{2} \rceil$

$\therefore r \geq s+1$ $P_0 = \left(\frac{V}{V_1 + V} \right)^s \left(1 - \frac{2k}{2^{s+1} + k} \cdot \frac{V_1}{V_1 + V} \right)$

~~其~~只有某次比赛时使得队伍数恰好为 2^i 才可能在 $s+1$ 轮内结束

此时 $P = \left(\frac{V}{V_1 + V} \right)^{s-t} \left(1 - \frac{2k}{2^{s+1} + k} \cdot \frac{V_1}{V_1 + V} \right) \cdot \left(\frac{V}{V_1 + V} \right)^t$

$\therefore P - P_0 > 0$ 得证.

