

PROBLEM 247. GOLOMB'S SELF-DESCRIBING SEQUENCE.

The *Golomb's self-describing sequence*  $\{G(n)\}$  is the only nondecreasing sequence of natural numbers such that  $n$  appears exactly  $G(n)$  times in the sequence. The values of  $G(n)$  for the first few  $n$  are

$n$	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
$G(n)$	1	2	2	3	3	4	4	4	5	5	5	6	6	6	6

You are given that  $G(10^3) = 86$ ,  $G(10^6) = 6137$ . You are also given that  $\sum_{n=1}^{999} G(n^3) = 153506976$ .

Find  $\sum_{n=1}^{999999} G(n^3)$ .

SOLUTION

Since there is no explicit formula for  $G(n)$  and apparently no special property for  $G(n^3)$ , this problem is about storing and retrieving the terms in  $\{G(n)\}$  efficiently.

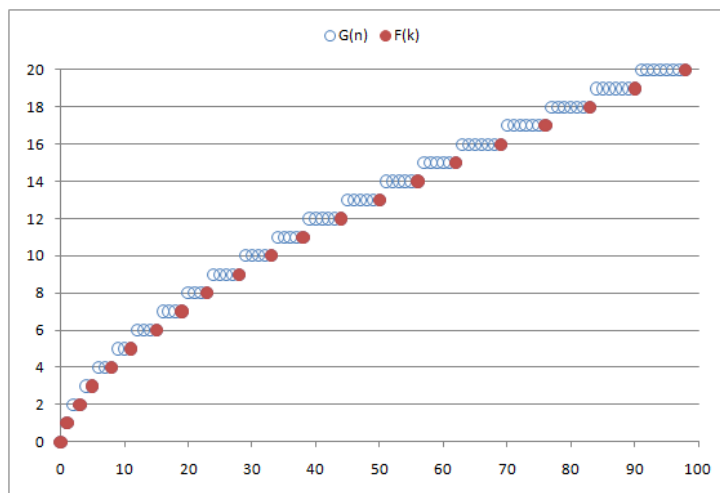
First, note that  $\{G(n)\}$  can be stored in compact form by omitting duplicate terms and keeping only the largest  $n$  for a given value. Let  $F$  denote the generalized inverse function of  $G$ :  $F(k)$  is the location of the last appearance of  $k$  in  $\{G(n)\}$ . Formally,

$$F(k) = \max\{n \mid G(n) = k\}.$$

The first 10 values of  $F(k)$  are

$k$	1	2	3	4	5	6	7	8	9	10
$F(k)$	1	3	5	8	11	15	19	23	28	33

Below is a plot of the first 98 values of  $G(n)$  and first 20 values of  $F(k)$ :



It is easy to find the recurrence relation

$$F(k) = F(k-1) + G(k), \quad (1)$$

with the starting value  $F(0) = 0$ . To understand this, note that  $F(k-1)$  is the location of the last appearance of  $(k-1)$  in  $\{G(n)\}$ , and  $k$  appears after  $(k-1)$  for  $G(k)$  times. Therefore the location of the last appearance of  $k$  is  $F(k-1) + G(k)$ .

Once we have obtained the values of  $\{F(k)\}$ , we can back out  $G(n)$  as follows. Find the unique  $k$  such that  $F(k-1) < n \leq F(k)$ , then  $G(n) = k$ .

Now we move one step further, and claim that  $F(k)$  is also the largest value that appears  $k$  times in  $\{G(n)\}$ . This is because all  $\{G(n)\}$  where  $F(k-1) < n \leq F(k)$  are equal to  $k$ , so every  $n$  where  $F(k-1) < n \leq F(k)$  appears  $k$  times in the sequence. But  $F(k)$  is, by definition, the largest such  $n$ . So it is the largest value that appears  $k$  times.

Next, let  $F(F(k))$  be the location of the last appearance of  $F(k)$  in  $\{G(n)\}$ . The first few values of  $F(F(k))$  are

$k$	1	2	3	4	5	6	7	8	9	10
$F(k)$	1	3	5	8	11	15	19	23	28	33
$F(F(k))$	1	5	11	23	38	62	90	122	167	217

We find the following recurrence relation for  $F(F(k))$ :

$$F(F(k)) = F(F(k-1)) + kG(k), \quad (2)$$

with the starting value  $F(F(0)) = 0$ . To understand this, note that each  $n$  where  $F(k-1) < n \leq F(k)$  appears  $k$  times in the sequence, and there are  $G(k)$  values that appear  $k$  times.

Once we have obtained the sequence  $\{F(k), F(F(k))\}$ , we can back out  $G(n)$  as follows. Given  $n$ , find the unique  $k$  such that  $F(F(k-1)) < n \leq F(F(k))$ . By construction, every  $n$  where  $F(k-1) < n \leq F(k)$  appears  $k$  times, and  $G(F(F(k))) = F(k)$ . It follows that

$$G(n) = F(k) - \left\lfloor \frac{F(F(k)) - n}{k} \right\rfloor.$$

The algorithm is outlined as follows. We use equation (1) to build and store the sequence  $\{F(k)\}$ , which is then used to find out  $G(n)$  for small  $n$ . Then we use equation (2) to compute the sequence  $\{F(F(k))\}$ . As we compute, we back out the values of  $G(n^3)$ . So there is no need to store  $\{F(F(k))\}$ .

To estimate the complexity of the algorithm, we run a regression on the first one million values of  $F(k)$  (with 1000 samples at  $k = 1000, 2000, \dots$ ) and get the following empirical relation:

$$F(k) \approx \varphi^{1-\varphi} k^\varphi,$$

where  $\varphi = (1 + \sqrt{5})/2 \approx 1.618$  is the golden ratio. A few useful results follow:

$$\begin{aligned} F^{-1}(n) &\approx \varphi^{2-\varphi} n^{\varphi-1}. \\ F^{-1}(F^{-1}(n)) &\approx \varphi^{\varphi-1} n^{2-\varphi}. \\ F^{-1}(F^{-1}(F^{-1}(n))) &\approx \varphi^{4-2\varphi} n^{2\varphi-3}. \end{aligned}$$

In this problem, we are asked to find  $F(F(k)) \geq n$  for  $n = 10^{18}$ . This corresponds to  $k \approx F^{-1}(F^{-1}(n)) = 10105311$ . To use equation (2) to compute  $F(F(k))$ , we need to know  $G(k)$ , which can in turn be backed out using equation (1). This requires a storage of  $F^{-1}(k) \approx 25638$ .

#### COMPLEXITY

Time complexity:  $\mathcal{O}(1.35 \times n^{1.15})$

Space complexity:  $\mathcal{O}(1.44 \times n^{0.71})$

#### ANSWER

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