

Operating Systems-2

CS 3510 Spring 2019

Programming Assignment 2:

Implementing TAS, CAS and Bounded Waiting CAS

Mutual Exclusion Algorithms

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Task

The goal of this assignment is to implement TAS, CAS and Bounded Waiting with CAS mutual exclusion (ME) algorithms studied in the class. Implement these algorithms in C++.

Approach and Implementation

1. To achieve our above mentioned goal make a *testCS* function which takes *thread_index* as a parameter. We will pass this function along with the thread index to the each thread and calculate the average waiting time and max waiting time for all three mutual exclusion algorithm.

```
vector<thread> t; // Array of n threads
for(int i=0;i<n;i++) t.push_back(thread(testCS,i));
for(int i=0;i<n;i++) t[i].join();
```

2. Functions and data-types of the chrono library (and other libraries) like
 - a. **std::chrono::system_clock, std::chrono::system_clock::now()**
 - b. **std::chrono::time_point**
 - c. **struct tm - Time structure**
 - d. **struct tm *localtime(const time_t *timer)**were used to calculate the average waiting time and max waiting time.
3. **int usleep(useconds_t usec)** function was used to suspend the execution of the thread for microsecond intervals
4. **template <class RealType = double> class exponential_distribution:** This is a random number distribution that produces floating-point values according to an exponential distribution, which is described by the following probability density function:
$$p(x|\lambda) = \lambda e^{-\lambda x} , x > 0$$
5. We make two exponential distributions and pass the value of $1/\lambda_1$, $1/\lambda_2$ in the constructor. Later this can be used to obtain random numbers *t1* and *t2* with values that are exponentially distributed with an average of λ_1 , λ_2 seconds.

```
distribution1 = new exponential_distribution<double>(1/lt1);
distribution2 = new exponential_distribution<double>(1/lt2);
```

6. The implementation of the critical section varies in *testCS* function with the change in mutual exclusion algorithm.

```

void testCS(int id){          // test function
    chrono::time_point<chrono::system_clock> start_time,end_time;
    // Variables for calculating waiting time
    for(int i=0;i<k;i++){
        // Entry Section Starts
        string reqEnterTime=getSysTime();
        start_time = chrono::system_clock::now();

        /*
            Mutual Exclusion algorithm implementation
            --Only this part changes
        */

        // Critical Section Starts
        end_time = chrono::system_clock::now();
        fout<<i<<"th CS Requested at "<<reqEnterTime;
        fout<<" by thread "<< id <<endl;
        tmp=(end_time - start_time);
        waiting_time +=tmp;
        // Calculating waiting time
        max_waiting_time = max(max_waiting_time,tmp);
        // Updating max time taken
        fout<<i<<"th CS Entered at  "<<getSysTime();
        fout<<" by thread "<< id <<endl;
        usleep((*distribution1)(generator)*scale);
        fout<<i<<"th CS Exited at  "<<getSysTime();
        fout<<" by thread "<< id <<endl;
        flag.clear();
        // Critical Section ends

        // Reminder Section starts
        usleep((*distribution2)(generator)*scale);
        // Reminder Section ends
    }
}

```

7. Implementation of TAS

```

while(atomic_flag_test_and_set(&flag));

```

8. Implementation of CAS

```
while(!atomic_compare_exchange_strong(&flag,&flag1,true)) flag1=false;
```

9. Implementation of CAS-Bounded waiting

Here the **waiting[i]==true** represents that thread is waiting to enter the critical section.

A thread can enter the critical section if **waiting[i]==false or key==false**.

The later part of the code ensures that no thread starves for a long time.

```
while(waiting[id] && key){//Compare and Swap(Bounded waiting)implementation
    if(atomic_compare_exchange_strong(&flag,&flag1,true)) key=false;
    else flag1=false;
}
// Critical Section Starts
end_time = chrono::system_clock::now();
waiting[id]=false;
usleep((*distribution1)(generator)*scale);
int j=(id+1)%n;
while(j!=i && !waiting[j]) j=(j+1)%n;
if(j==i)flag =false;
else waiting[j]=false;
```

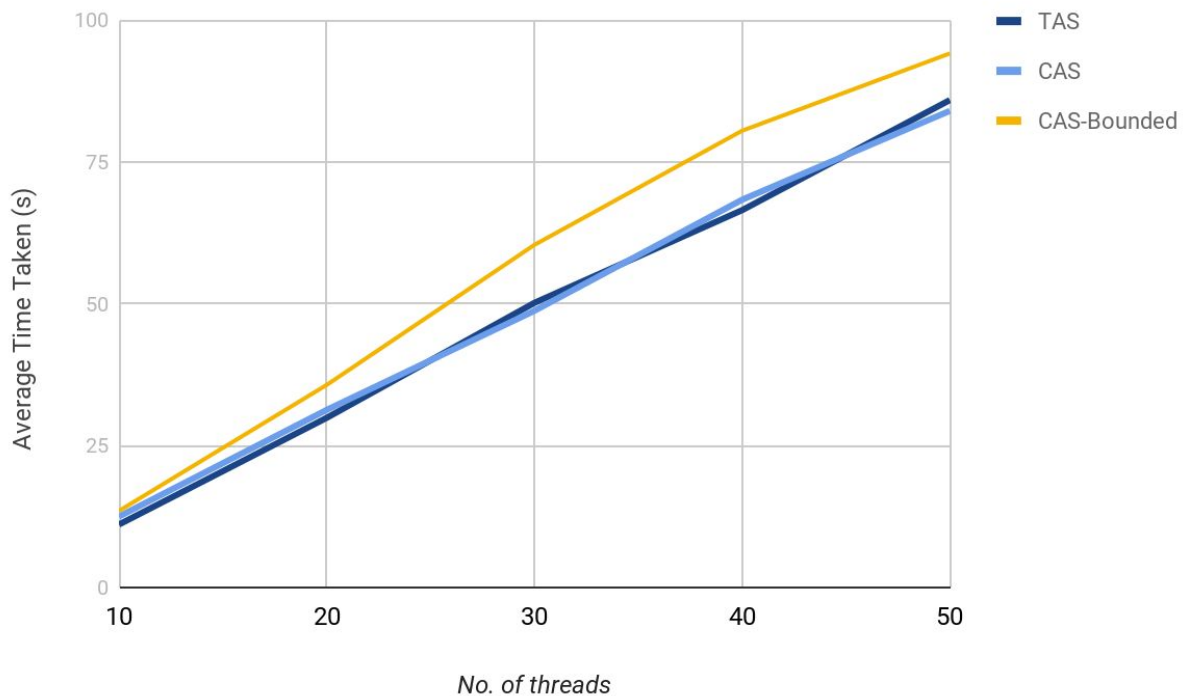
Input Parameters

N: No. of threads vary from 10 to 50 in steps of 10.

K = 10: No. of CS request by each thread

$\lambda_1=2$, $\lambda_2=2$

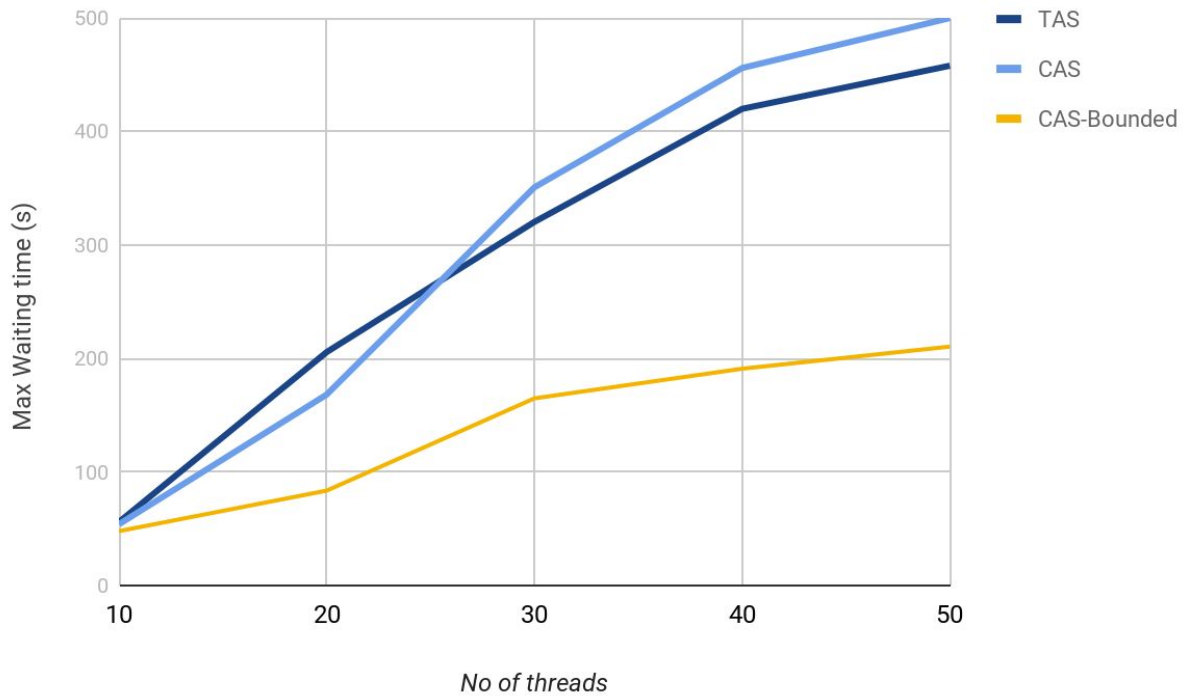
Graph 1 : Avg Waiting Time vs No. of Threads



Output analysis of Graph 1

1. Average waiting time taken by TAS and CAS algorithm is almost the same.
2. Average waiting time taken by Bounded-CAS is greater than both TAS and CAS as it ensures that no thread starve.

Graph 2 : Worst Waiting Time vs No of Threads



Output analysis of Graph 2

1. Worst waiting time taken by TAS and CAS algorithm is almost the same.
2. Worst waiting time taken by Bounded-CAS is much less than both TAS and CAS as it ensures that no thread starves.
3. The difference between the worst waiting times of Bounded-CAS and (TAS,CAS) with increases as the no of threads increase.