

Concurrency & Concurrent Processes





Outline

- Concurrent Process
- Bernstein Conditions for Concurrency
- Concurrency Representation
 - Begin--end, **and** *par begin* ---*par end*
 - Fork and Join
 - Semaphores

Concurrency

(Concurrent Processes)

Concurrent Processes

- ▶ Two Processes are said to be concurrent if they can be executed simultaneously (**in Parallel**) in a multiprogramming environment.
- ▶ In a **uniprocessor** environment, two processes are said to be concurrent if the **order of execution** doesn't matter, i.e., either of the processes can be executed first.
- ▶ **Example:** Consider a Process p with four statements.

- S1 & S2 are concurrent
- S1 & S3 are sequential (non-concurrent)
- S2 & S3 are also sequential
- S2 & S4 are also sequential (**Transitive concurrency**)
- S3 & S4 are also sequential

Process P1:

```
S1 : a = b + c;  
S2 : d = b - d;  
S3 : e = a * d;  
S4 : f = e / 2;
```

Bernstein Concurrency Conditions

- ▶ Two Processes (statements) are concurrent if the following three conditions are satisfied:
- ▶ Let S1 and S2 are two processes:
 - $R(S1) \cap W(S2) = \emptyset$
 - $W(S1) \cap R(S2) = \emptyset$
 - $W(S1) \cap W(S2) = \emptyset$
- ▶ **R(S1)**: Read the set of processes/statements S1. It is the set of all variables that are read-referenced in the process/statement S1.
 - E.g., $R(S1) = \{b, c\}$
- ▶ **W(S1)**: Write the set of processes/statements S1. It is the set of all those variables on which write operation is performed in the process/statement.
 - E.g., $W(S1) = \{a\}$
- ▶ **Note: Along with the Bernstein condition, transitive dependencies are also to be observed to identify concurrency.**

Representation of Concurrency

► Two Methods:

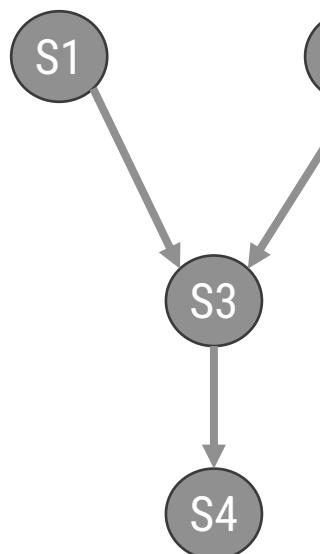
- Precedence Graph
- Using Concurrent Statements (Languages)
 - Begin---end, and *par begin* ---*par end*
 - Fork and Join
 - Semaphores

Representation of Concurrency: Precedence Graph

- *It is a directed graph with a vertex (statements) & edges (between dependent statements).*

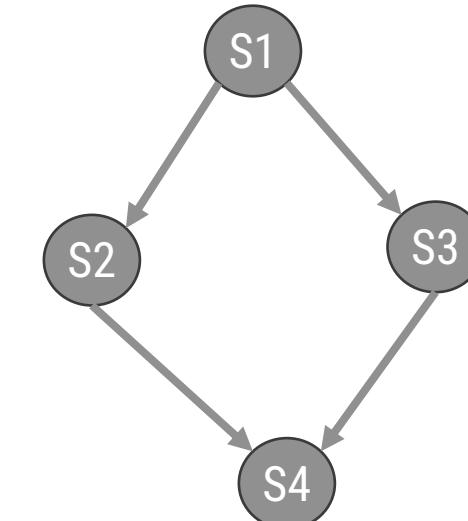
- *Example:*

```
S1 : a = b + c;  
S2 : d = b - d;  
S3 : e = a * d;  
S4 : f = e / 2;
```



- *Example:*

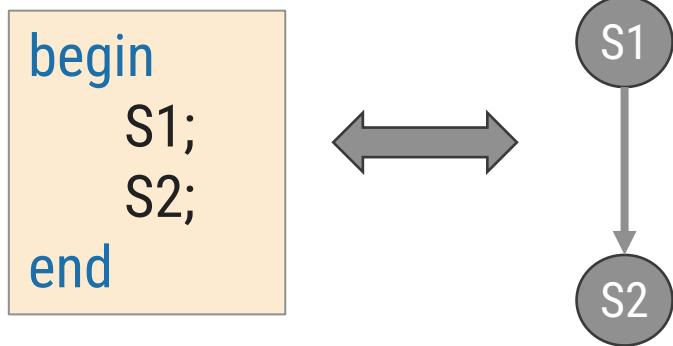
```
S1 : a = b + c;  
S2 : d = a * 2;  
S3 : e = e + a;  
S4 : f = d / e;
```



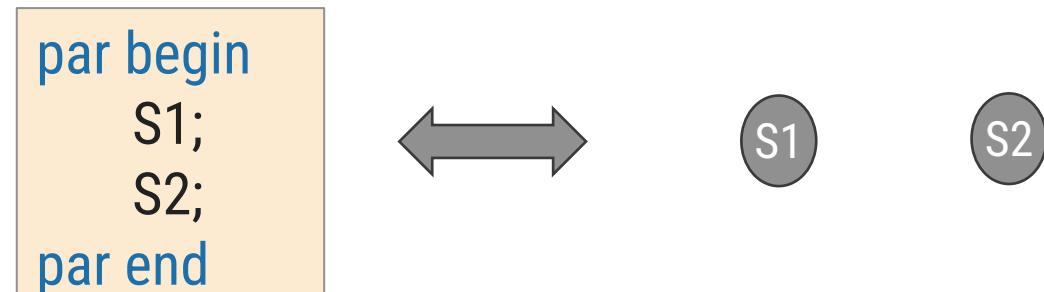
Representation of Concurrency: Using Concurrent Statements

- ▶ **1. *begin---end* and *par begin----par end***

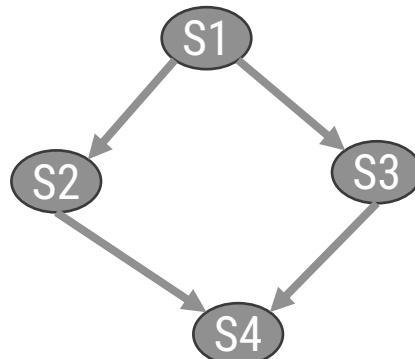
Sequential



Concurrent

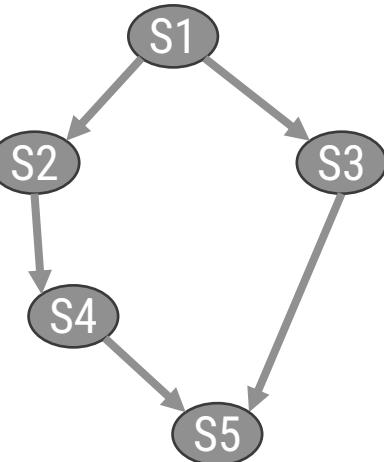
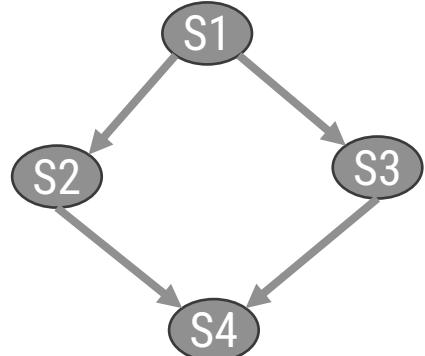


- ▶ **Example: Write concurrent statements for the following graph.**



Representation of Concurrency: Using Concurrent Statements

► Ques. Write concurrent statements for the following graph.



► Solution:

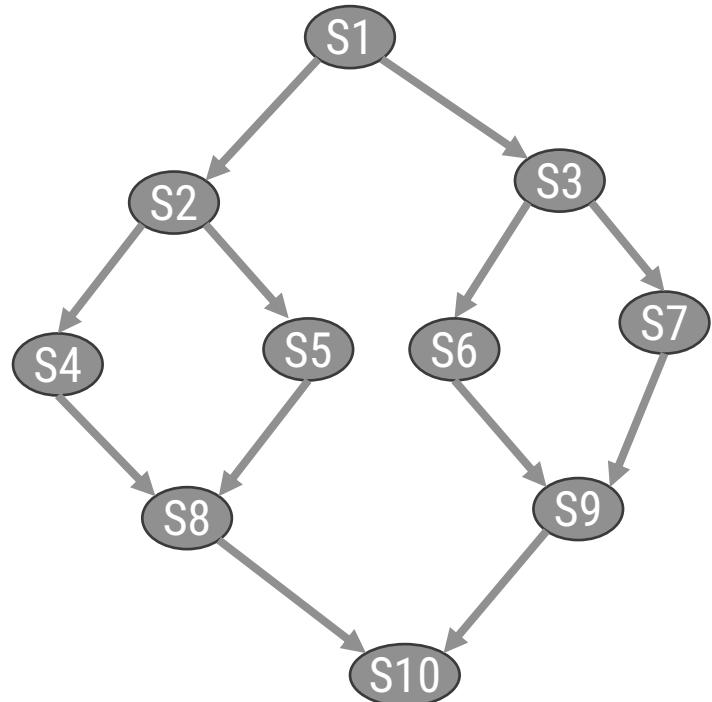
```
begin  
    S1;  
    par begin  
        S2;  
        S3;  
    par end  
    S4;  
end
```

```
begin  
    S1;  
    par begin  
        S3;  
        begin  
            S2;  
            S4;  
        end  
    par end  
    S5;  
end
```

```
begin  
    S1;  
    par begin  
        begin  
            S2;  
            S4;  
        End  
        S3;  
    par end  
    S5;  
end
```

Representation of Concurrency: Using Concurrent Statements

► Ques. Write concurrent statements for the following graph.



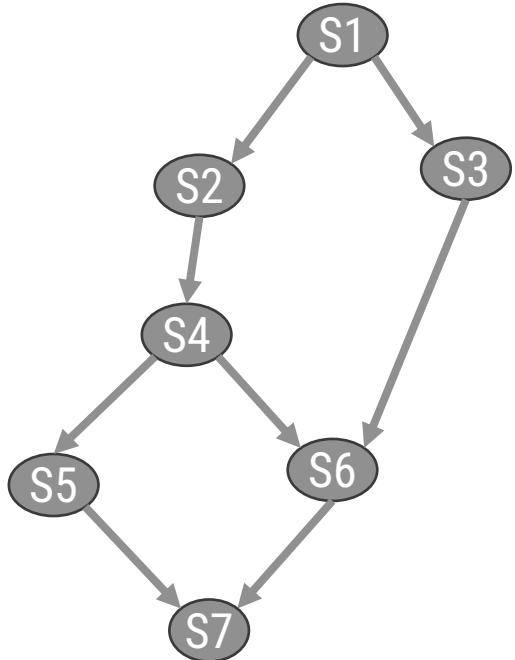
Solution:
begin
S1;
par begin
begin
S2;
par begin
S4;
S5;
par end
S8;
end

begin
S3;
par begin
S6;
S7;
par end
S9;
end
par end
S10;
end

► **Solution:**

Representation of Concurrency: Using Concurrent Statements

► Ques. Write concurrent statements for the following graph.



► Solution:

→ Not Possible with only **begin–end**, and **par begin – par end**, method.

Representation of Concurrency: Using Concurrent Statements

► 2. Fork & Join Construct:

► Fork:

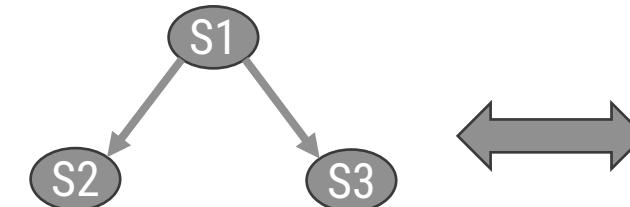
→ E.g.,

Syntax for using Fork:

Fork: L1

L1:

Concurrent



S1;
Fork: L1
S2;
Goto Next
L1: S3;
Next:

- **L1 (Level):** Used to jump. When we create a level (L1), **two concurrent processes** are created.
- First one below the **Fork**, & second one is start from level **L1**.

Representation of Concurrency: Using Concurrent Statements

► Fork & join Construct:

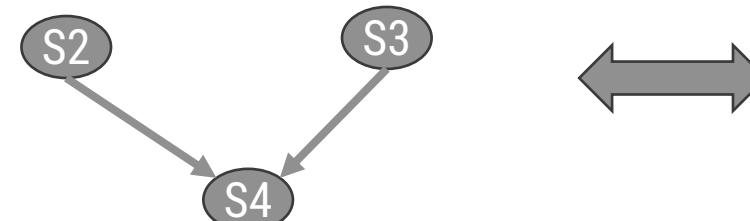
► Join:

→ E.g.,

Syntax for using Join:

initialize counter

Join Counter



C =2;
Fork: L1
S2;
Goto L2

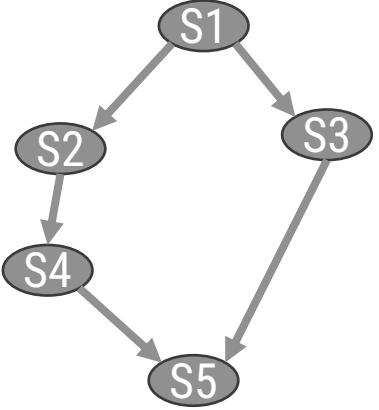
L1: S2;

L2: join C
S4;

- On a visit to the Join statement, the counter variable is decremented by one every time.
- If the counter variable becomes zero after decrement, then the join statement allows the control to execute the next below mentioned statements.
- If the counter is not zero, block the control.

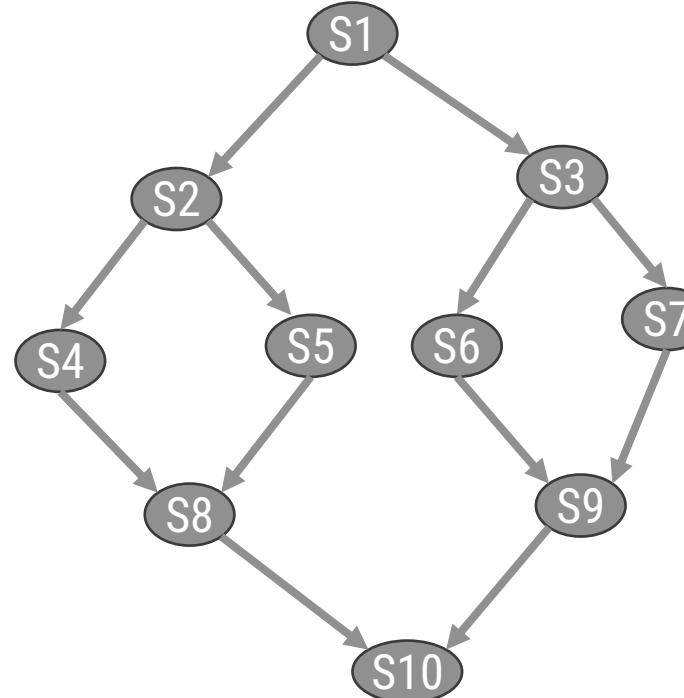
Representation of Concurrency: Using Concurrent Statements

► Ques. Write concurrent statements (using Fork and Join) for the following graph.



► Solution:

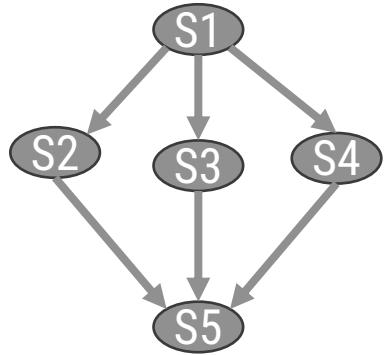
```
C=2;  
S1;  
Fork: L1  
S3;  
Goto L2  
L1: S2;  
S4;  
L2: join C  
S5;
```



Solution: ?

Representation of Concurrency: Using Concurrent Statements

► Ques. Write concurrent statements (using Fork and Join) for the following graph.

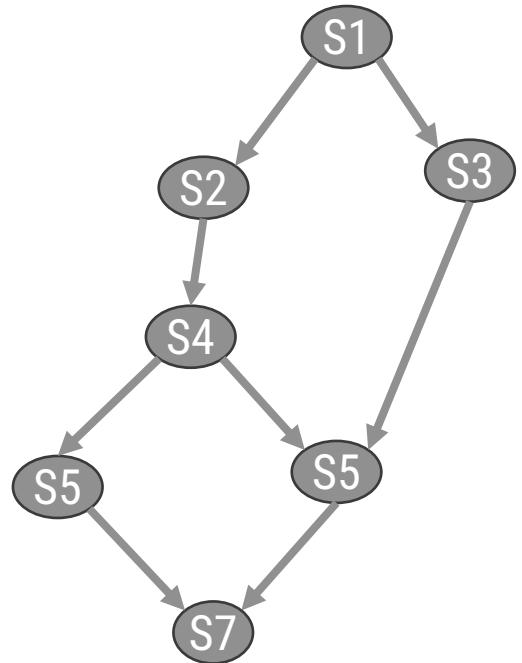


► Solution:

```
C=3;  
S1;  
Fork: L1  
S2;  
Goto L3  
  
L1: Fork: L2;  
S3;  
Goto L3  
  
L2: S4;  
  
L3: join C  
S5;
```

Representation of Concurrency: Using Concurrent Statements

► Ques. Write concurrent statements (*using Fork and Join*) for the following graph.



► Solution:

Representation of Concurrency: Concurrent Statements (Semaphore)

► 3. Semaphore:

- Semaphores are a special type of variable that **cannot be used (modified)** in arithmetic & logical operations. These variables can only be altered with predefined operations called **wait(P)** and **signal (V)**.

► Types of Semaphore

- Boolean (value- 0 or 1)
- Counting

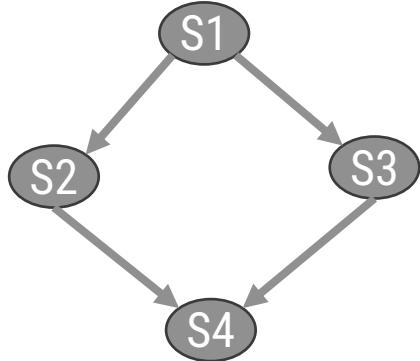
► Example:

```
a : Semaphore  
----  
----  
---- wait(a)  or P(a)  
----  
----signal(a)  or V(a)  
----  
----
```

- **Wait(a):** It decrements the value of a and allows control to go to the next statement if a is 0. Otherwise, wait or block the control.
- **Signal(a):** No wait, it only increments the value by 1, passes the control to go to the next statement.
- **Semaphore can be initialized** using built-in functions as
 - **a, b: counting/Boolean semaphores**
 - **Initialize(a, b =0)**

Representation of Concurrency: Concurrent Statements (Semaphore)

► Ques. Write concurrent statements using Semaphores for the following graph.



► Solution:

Solution:

a, b, c, d: Boolean semaphores

Initialize (a, b c, d, =0)

par begin

begin S1; V(a); V(b); end

begin P(a); S2; V(c); end

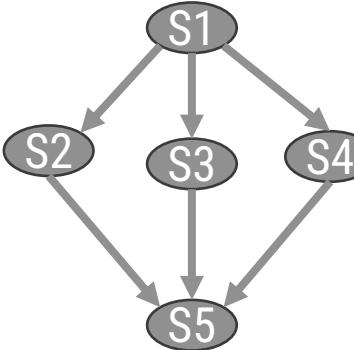
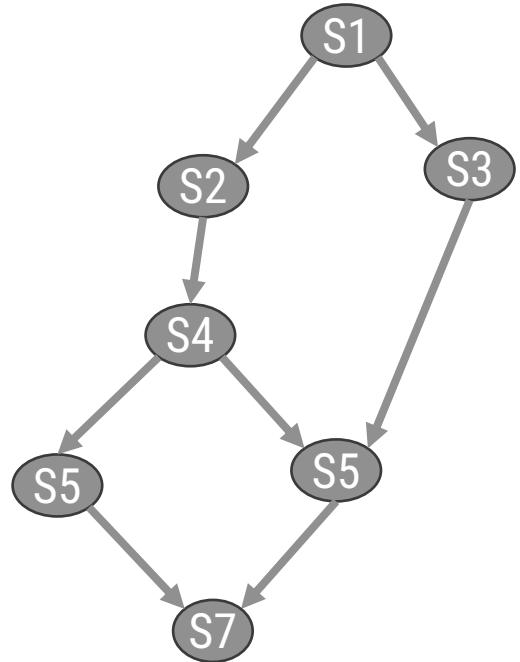
begin P(b); S3; V(d); end

begin P(c); P(d); S4; end

par end

Representation of Concurrency: Concurrent Statements (Semaphore)

► Ques. Write concurrent statements using Semaphores for the following graphs.



► Solution:

***Thank
You***