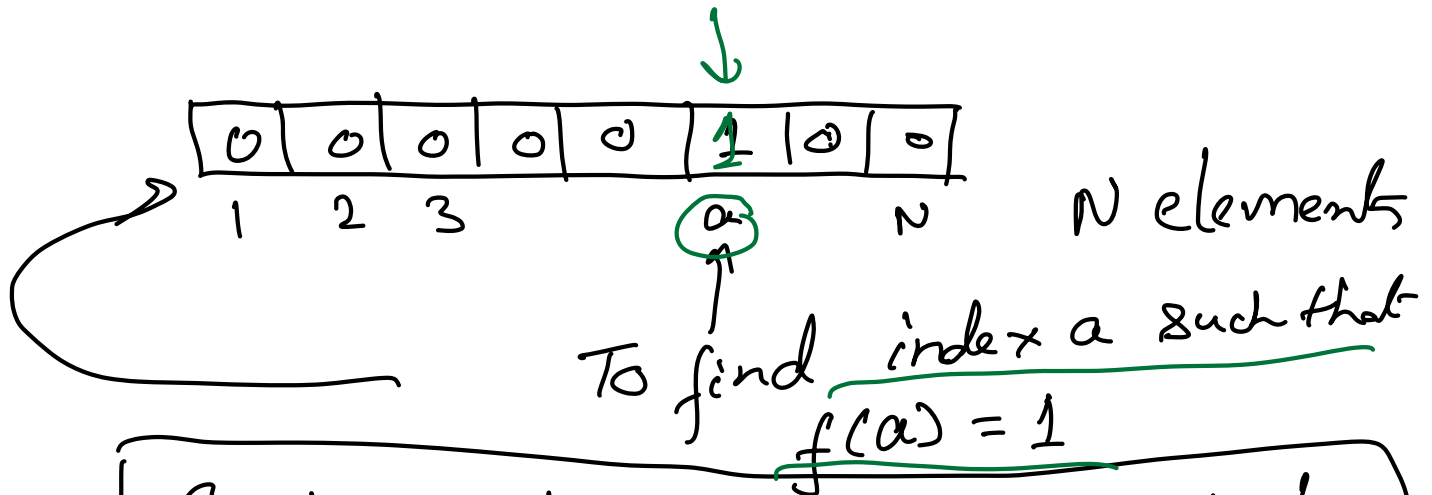


# Grover's Search Algorithm (1996)



Goal is to search for a single element  $a_1, a_2, \dots, a_N$  does there exist an  $x$  such that  $a_x = y$

I want to find " $a$ " such that  
 $f(a) = 1$   $f(x) = 0$  for all other  $x \neq a$

✓ Input:-  $n$  bits ✓ " $N$ " possible  
✓ Output:- a single bit  
0 or 1

bits  $\sqrt{N}$   
 $N = 2$   
Marked entry

$f(x) = 1$  at some  $x = x_0$   
and for other inputs the value of

$f(x) = 0$   
Goal:- To find  $x_0$  for which  $f(x_0) = 1$

→ Classically how much time would it take  $O(N)$

Can we do better?

In a quantum setting, you can do this task  $O(\sqrt{N})$

→ An oracle  $f: \{0,1\}^n \rightarrow \{0,1\}$    
  $N = 2^n$    
 for every input. Find a s.t  $f(x) = 1$

$$x \rightarrow \boxed{U_f} \rightarrow f(x)$$

$$\boxed{U_f |x, y\rangle = |x, y \oplus f(x)\rangle}$$

$$U_f |x, 0\rangle = |x, f(x)\rangle$$

$$U_f |x, 1\rangle = |x, 1 \oplus f(x)\rangle = |x, \overline{f(x)}\rangle$$

ancilla qubit

$$|y\rangle = \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

$$\boxed{U_f |x, -\rangle = (-1)^{f(x)} |x, -\rangle} \quad \text{--- (1)}$$

Start with  $N = 2^n$

$$|\psi_0\rangle = \frac{1}{\sqrt{N}} \sum_{x \in \{0,1\}^n} |x\rangle$$

Assume that  $f(a) = 1$   
 $f(x) = 0 \quad \forall x \neq a$

$$|\psi_0\rangle = \frac{1}{\sqrt{N}} \sum_{x \in \{0,1\}^n} |x\rangle$$

$$= \frac{1}{\sqrt{N}} \left[ |a\rangle + \sum_{x \neq a} |x\rangle \right]$$

$$= \frac{1}{\sqrt{N}} |a\rangle + \frac{\sqrt{N-1}}{\sqrt{N}} \left[ \frac{\sum_{x \neq a} |x\rangle}{\sqrt{N-1}} \right]$$

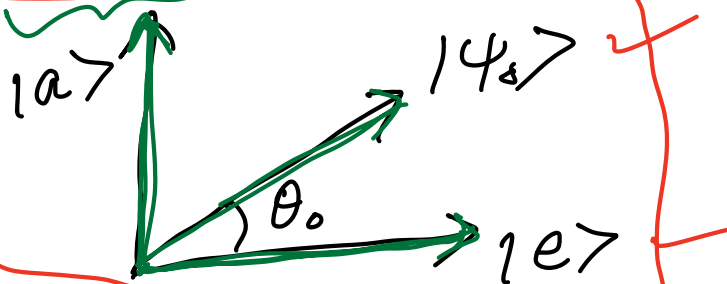
$$|\psi_0\rangle$$

$$= \frac{1}{\sqrt{N}} |a\rangle + \sqrt{1 - \frac{1}{N}} |e\rangle \quad \text{--- (1)}$$

$$\cos \theta_0 = |\langle e | \psi_0 \rangle|$$

$$= \sqrt{1 - \frac{1}{N}}$$

$$\sin \theta_0 = \frac{1}{\sqrt{N}}$$



$$|\psi_0\rangle = \sin \theta_0 |a\rangle + \cos \theta_0 |e\rangle \quad \text{--- (1)}$$

If  $N$  is very large,  $\theta_0$  is very small

Oracle

$$U_f |x\rangle \rightarrow = (-1)^{f(x)} |x\rangle \rightarrow \quad - (*)$$

$$f(x)=0 \\ \forall x \neq a \\ f(a)=1$$

$$U_f |a\rangle \rightarrow = (-1)^{f(a)} |a\rangle \rightarrow$$

$$U_f |a\rangle \rightarrow = -|a\rangle \rightarrow \quad - (2)$$

$$U_f |e\rangle \rightarrow = U_f \left| \frac{\sum_{x \neq a} |x\rangle}{\sqrt{N-1}} \right\rangle \rightarrow$$

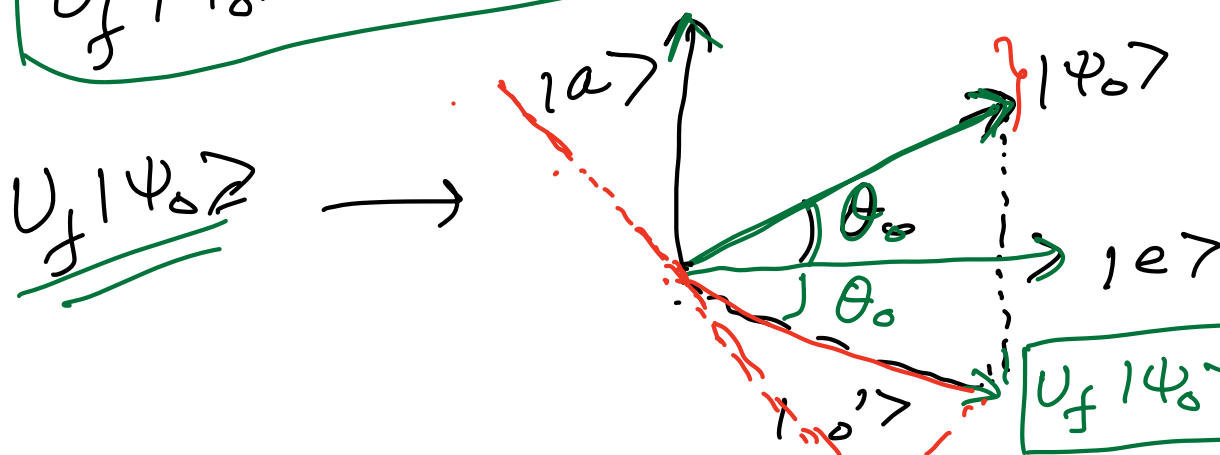
$$= (-1)^{f(a)} \left| \frac{\sum_{x \neq a} |x\rangle}{\sqrt{N-1}} \right\rangle \rightarrow$$

$$U_f |e\rangle \rightarrow = |e\rangle \rightarrow \quad - (3)$$

$$U_f |\psi_0\rangle \rightarrow = U_f \{ \sin \theta_0 |a\rangle + \cos \theta_0 |e\rangle \} \rightarrow$$

$$= \sin \theta_0 \{ -|a\rangle \} + \cos \theta_0 |e\rangle \rightarrow$$

$$U_f |\psi_0\rangle = -\sin \theta_0 |a\rangle + \cos \theta_0 |e\rangle$$



$$U_f |\psi_0\rangle = R_{|e\rangle} |\psi_0\rangle$$

Q: Can we achieve rotation closer to  $|a\rangle$  through reflection?

→ Can achieve this from 2 consecutive reflections!

$$|\psi_0\rangle = \sin\theta_0 |a\rangle + \cos\theta_0 |e\rangle$$

$$U_f |\psi_0\rangle = -\sin\theta_0 |a\rangle + \cos\theta_0 |e\rangle$$

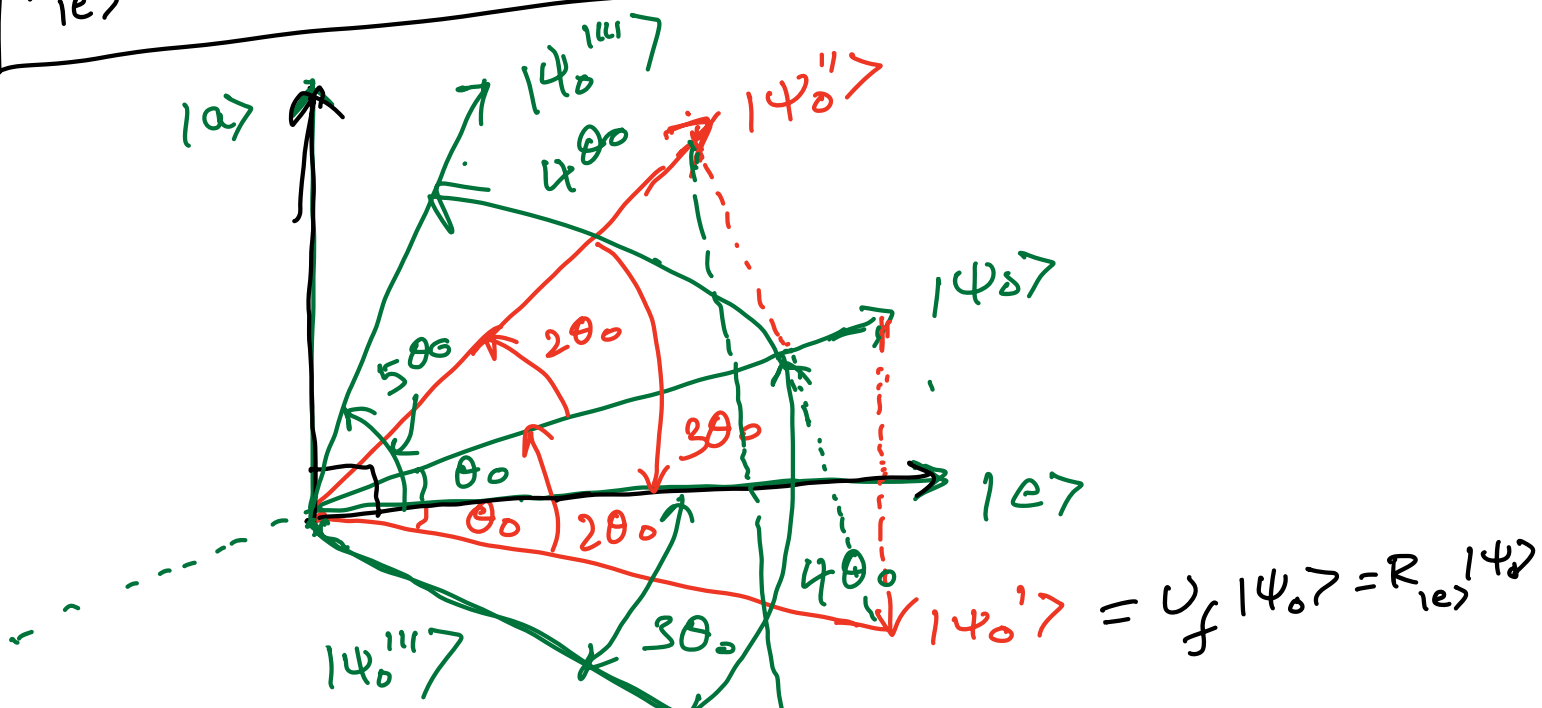
$$R_{|e\rangle} |\psi_0\rangle = (\underbrace{I - 2|a\rangle\langle a|}_{\text{reflection}}) |\psi_0\rangle$$

$$= (|\psi_0\rangle - 2|a\rangle\langle a|\psi_0\rangle)$$

$$= (\sin\theta_0 |a\rangle + \cos\theta_0 |e\rangle$$

$$- 2|a\rangle [\sin\theta_0])$$

$$R_{|e\rangle} |\psi_0\rangle = -\sin\theta_0 |a\rangle + \cos\theta_0 |e\rangle \rightarrow |\psi_0'\rangle$$



$$|\psi_0''\rangle = R_{|\psi_0\rangle} |\psi_0'\rangle \Rightarrow \sin 3\theta_0 |a\rangle + \cos 3\theta_0 |e\rangle$$

$$|\psi_0''\rangle = R_{|\psi_0\rangle} R_{|e\rangle} |\psi_0\rangle = \underbrace{\sin 3\theta_0 |a\rangle + \cos 3\theta_0 |e\rangle}$$

$$G = R_{|\psi_0\rangle} R_{|e\rangle}$$

$$|\psi_0''' \rangle = R_{|e\rangle} [R_{|\psi_0\rangle} R_{|e\rangle} |\psi_0\rangle]$$

$$= R_{|e\rangle} |\psi_0''\rangle$$

$$= -\sin 3\theta_0 |a\rangle + \cos 3\theta_0 |e\rangle$$

$$R_{|\psi_0\rangle} |\psi_0''' \rangle = \sin 5\theta_0 |a\rangle + \cos 5\theta_0 |e\rangle$$

$G \rightarrow$  Grover operator

$$\underbrace{R_{|\psi_0\rangle}}_{\text{marked}} R_{|e\rangle} = R_{|\psi_0\rangle} \cup f$$

$$G^k |\psi_0\rangle = \sin(\underbrace{(2k+1)\theta_0}) |a\rangle + \underbrace{\cos((2k+1)\theta_0)} |e\rangle$$

$$(2k+1)\theta_0 \approx \frac{\pi}{2}$$

$$(2k+1)\theta_0 = \frac{\pi}{2} \Rightarrow$$

$$\begin{aligned} 2k\theta_0 + \theta_0 \\ = \frac{\pi}{2} \end{aligned}$$

$$\Rightarrow 2k\theta_0 = \frac{\pi}{2} - \theta_0$$

$$\Rightarrow k = \underbrace{\frac{\pi}{4\theta_0}} - \frac{1}{2}$$

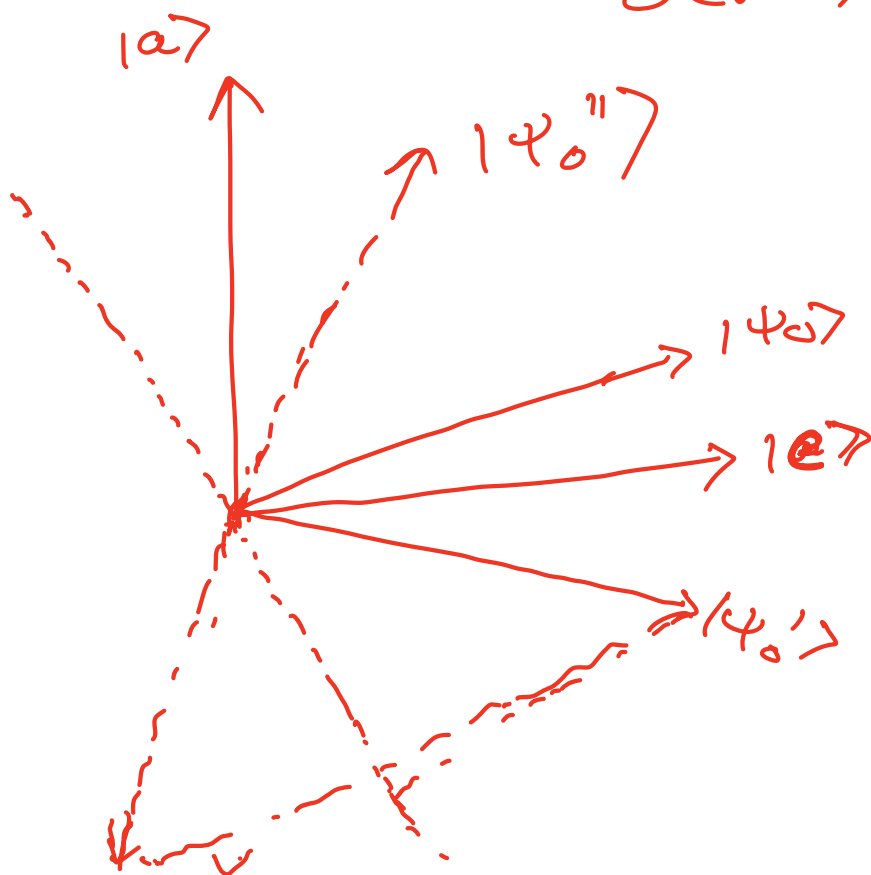
$$k \approx \frac{\sqrt{N}\pi}{4}$$

$$O(\sqrt{N})$$

$$\sin\theta_0 = \frac{1}{\sqrt{N}}$$

$$\sin\theta_0 \approx \theta_0$$

$$\theta_0 \approx \frac{1}{\sqrt{N}}$$



$$R_{|\psi_0\rangle} = -(I - 2|\psi_0\rangle\langle\psi_0|)$$

$$R_{|\psi_0\rangle} = 2|\psi_0\rangle\langle\psi_0| - I$$

$$R_{|\psi_0\rangle} = H^{\otimes n} (2|0^{\otimes n}\rangle\langle 0^{\otimes n}| - I) H^{\otimes n}$$

$$H^{\otimes n} |0^{\otimes n}\rangle = |\psi_0\rangle \quad (H^{\otimes n})^2 = I$$

$$|0^{\otimes n}\rangle = (H^{\otimes n})^\dagger |\psi_0\rangle$$

$$= H^{\otimes n} |\psi_0\rangle$$

$$R_{|\psi_0\rangle} = \left[ 2 \underbrace{H^{\otimes n} |0^{\otimes n}\rangle\langle 0^{\otimes n}|}_{|\psi_0\rangle\langle\psi_0|} - H^{\otimes n} \right] H^{\otimes n}$$

$$= \left[ 2 |\psi_0\rangle\langle\psi_0| - \underbrace{H^{\otimes n}}_{I} \right] \underbrace{H^{\otimes n}}_{I}$$

$$R_{|\psi_0\rangle} = 2 |\psi_0\rangle\langle\psi_0| - I$$

Amplitude amplification

$|\psi_0\rangle$  which is prepared by

oracle  $U_{\psi_0}$  i.e.  $U_{\psi_0} |0^{\otimes n}\rangle = |\psi_0\rangle$

$$|\psi_0\rangle = \sqrt{p_0} |\psi_{\text{good}}\rangle + \sqrt{1-p_0} |\psi_{\text{bad}}\rangle$$



Grover's algorithm kind  
 of approach can be used  
 to amplify the coefficient  
 in front of  $|\psi_{\text{good}}\rangle$  i.e. one  
 can devise a quantum circuit  
 similar to Grover's algorithm  
 to construct states which have  
 large overlap  $|\psi_{\text{good}}\rangle$

$$G = R_{\psi_0} R_{\text{good}}$$

## Quantum Phase Estimation

### Phase estimation

$U \rightarrow$  unitary operator

$|\psi\rangle \rightarrow$  eigenvector of  $U$   
 i.e.

$$U|\psi\rangle = e^{i\theta} |\psi\rangle$$

$$\theta = 2\pi\phi, \quad \phi \in [0, 1)$$

$$U|\psi\rangle = e^{i2\pi\phi}|\psi\rangle$$

$\swarrow$  Given       $\downarrow$  Given       $\downarrow$  Estimate  $\phi$ ?

$$Qx = \lambda x$$

$$x^\dagger Q x = \lambda \frac{x^\dagger x}{x^\dagger x}$$

$$\lambda = \frac{x^\dagger Q x}{x^\dagger x}$$

## Quantum Phase Estimation

Algo  $\rightarrow$

### (I) Hadamard Test

$\Downarrow$

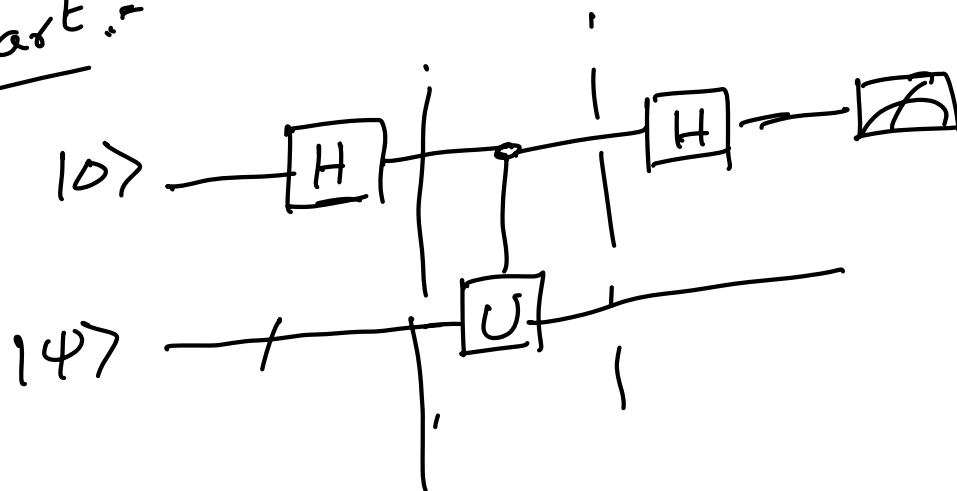
Compute expectation value  
of unitary operator with respect  
to a state  $\langle\psi|U|\psi\rangle$   
 $\hookrightarrow$  complex number

Since  $U$  is unitary this quantity

$\langle \psi | U | \psi \rangle$  is a complex number

and one needs to measure real and imaginary part of  $\langle \psi | U | \psi \rangle$  separately!

Real part:-



$$\text{Re } \langle \psi | U | \psi \rangle$$

$$|0\rangle \otimes |\psi\rangle \xrightarrow{H \otimes I} \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \otimes |\psi\rangle$$

$\downarrow U$

$$\frac{1}{\sqrt{2}} (|0\rangle \otimes |\psi\rangle + |1\rangle \otimes U|\psi\rangle)$$

$\downarrow H \otimes I$

$$\frac{1}{\sqrt{2}} (|+\rangle \otimes |\psi\rangle + |-\rangle \otimes U|\psi\rangle)$$

$$= \frac{1}{\sqrt{2}} \left( \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \otimes |\psi\rangle + \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \otimes U|\psi\rangle \right)$$

$$= \frac{1}{2} \left[ |0\rangle \otimes (|\psi\rangle + U|\psi\rangle) + \frac{i}{2} \left[ |1\rangle \otimes (|\psi\rangle - U|\psi\rangle) \right] \right]$$

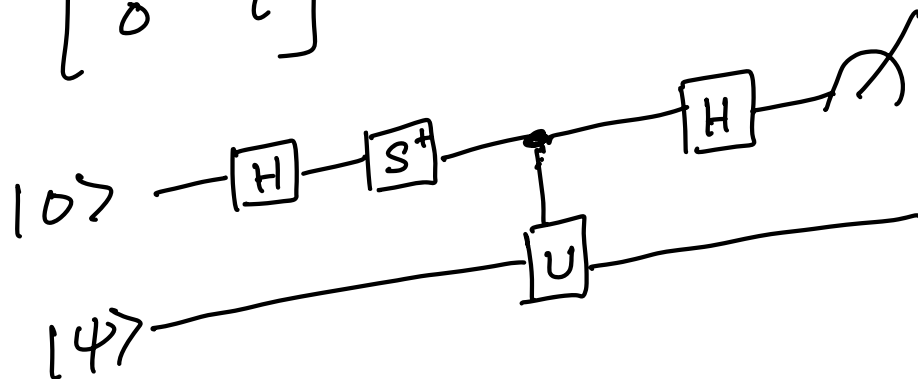
Probability of measuring qubit 0  
to be in state  $|0\rangle$

$$p(0) = \frac{1}{2} (1 + \text{Re} \langle \psi | U | \psi \rangle)$$

$$|\psi\rangle + U|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

To define the imaginary part

$$S = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \quad \text{phase Gate}$$



$$\frac{1}{2} [ |0\rangle \otimes (|\psi\rangle - iU|\psi\rangle) ] + \frac{1}{2} [ |1\rangle \otimes (|\psi\rangle + iU|\psi\rangle) ]$$

Probability of measuring qubit 0  
(1<sup>st</sup> qubit)

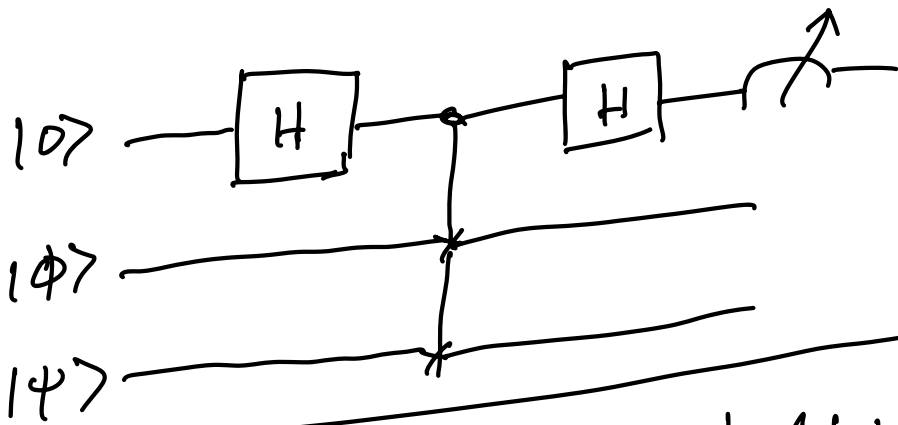
to be in state  $|0\rangle$

$$p(0) = \frac{1}{2} (1 + \text{Im}(\langle\psi|U|\psi\rangle))$$

Combining the results from  
two circuits we obtain the  
estimate to  $\langle\psi|U|\psi\rangle$

Overlap estimate:- (Application of Hadamard test)

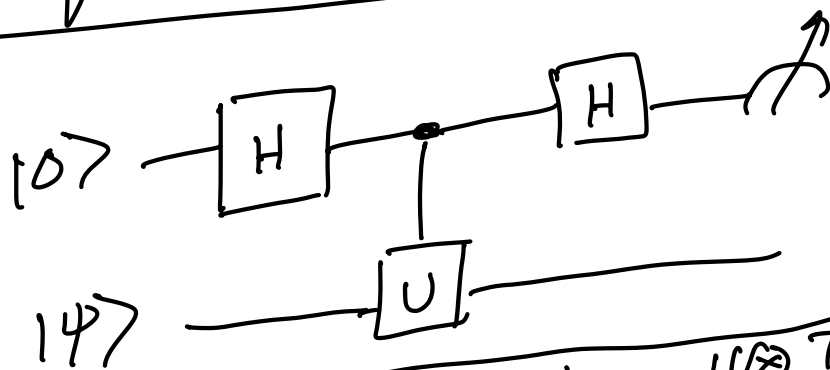
Swap test  $\rightarrow$  estimate overlap of  
two quantum state  
 $|\langle\phi|\psi\rangle|$



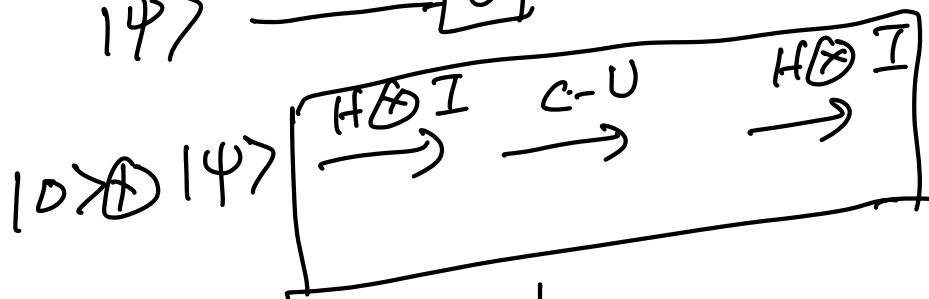
$$|p(0)| = \frac{1}{2} (1 + |\langle \phi | \psi \rangle|^2)$$

1<sup>st</sup> Qubit to be in state  $|0\rangle$

Single qubit phase estimation



$|\psi\rangle$  is an eigen vector of  $U$



$$\frac{1}{2} (|0\rangle \otimes (|\psi\rangle + U|\psi\rangle) + |1\rangle \otimes (|\psi\rangle - U|\psi\rangle))$$

Probability of measuring  $1^{\text{st}}$  qubit  
to be state  $|1\rangle$

$$p(1) = \frac{1}{2} (1 - \text{Re} \langle \psi | U | \psi \rangle)$$

$$= \frac{1}{2} (1 - \cos(2\pi\phi))$$

$$\phi = \pm \frac{\cos^{-1}(1 - 2p(1))}{2\pi}$$

$p(1)$  is close to 0 or 1  
or somewhere in between

The number of samples needed  
is  $O(1/\epsilon^2)$  where  $\epsilon$  is

the precision to determine  
" $\phi$ "

QPE  $\rightarrow$  Can we do better than  
 $O(1/\epsilon^2)$  sampling to estimate  
 $\phi$ ?