

Dynamic Programming - III

1

Tutorial Exercise#9

- Suppose the job of a firm is to manage the construction of billboards on the Surat-Dumas Gauravpath that runs east-west for M kms. The possible sites for the billboards are given by numbers $x_1, x_2, x_3, x_4, x_5, \dots, x_n$ each in the interval $[0..M]$. x_i 's are indicating the position of the billboards along the Gauravpath in kms, measured from its western end. If it is decided to place a billboard at location x_i , then the firm receives a revenue of $r_i > 0$.

Regulations of the SMC required that no two of the billboards be **within less than or equal to t kms** of each other. Thus, as part of the optimization problem, the firm has to decide where to place the billboards at a subset of the sites $x_1, x_2, x_3, x_4, x_5, \dots, x_n$ so as to maximize the total revenue, subject to this constraint.

Give an algorithm that takes in an instance of this problem as input and returns the maximum total revenue that can be obtained from any valid subset of sites. Give its running time also.

2

2

Tutorial Exercise#9...

- Illustration1:
Suppose $M = 20$, $n = 4$, $t=5$
Distances $[x_1, x_2, x_3, x_4] = [6, 7, 12, 14]$
revenue $[r_1, r_2, r_3, r_4] = [5, 6, 5, 1]$
Then, what would be the optimal solution ?
- Illustration2:
Input : $M = 20$, separation distance $t=5$ kms
Distances $x[] = \{6, 7, 12, 13, 14\}$
revenue $[] = \{5, 6, 5, 3, 1\}$ Then, Output: 10 ?
- Illustration3:
Input : $M = 15$, separation distance $t=2$ kms
Distances $x[] = \{6, 9, 12, 14\}$
revenue $[] = \{5, 6, 3, 7\}$ Then, Output : 18

3

Tutorial Exercise#10

- Suppose you own two stores, A and B. On each day you can be either at A or B. If you are currently at store A (or B) then moving to store B the next day (or A) will cost C amount of money. For each day i , $i = 1, \dots, n$, we are also given the profits $P^A(i)$ and $P^B(i)$ that you will make if you are store A or B on day i respectively. Give a schedule which tells where you should be on each day so that the overall money earned (profit minus the cost of moving between the stores) is maximized.
–Approach:
 - Define two arrays TA[] and TB[]. TA[i] gives the most profitable schedule for days i, \dots, n given that we start at store A on day i . Define T B [i] similarly.
 - Write the recurrences... ..
- $TA[i] = PA(i) + \max(TA[i + 1], TB[i + 1] - C)$.

4

Tutorial Exercise#11

- You are going on a long trip. You start on the road at mile post 0. Along the way there are n hotels, at mile posts $a_1 < a_2 < \dots < a_n$, where each a_i is measured from the starting point. The only places you are allowed to stop are at these hotels, but you can choose which of the hotels you stop at. You must stop at the final hotel (at distance a_n), which is your destination.
- You'd ideally like to travel 200 miles a day, but this may not be possible (depending on the spacing of the hotels). If you travel x miles during a day, the *penalty* for that day is $(200 - x)^2$. You want to plan your trip so as to **minimize the total penalty**—that is, the sum, over all travel days, of the daily penalties.
- Give an efficient algorithm that determines the optimal sequence of hotels at which to stop.
- Solution:
 - To get $\text{OPT}(i)$, we consider all possible hotels j we can stay at the night before reaching hotel i . For each of these possibilities, the minimum penalty to reach i is the sum of:
 - the minimum penalty $\text{OPT}(j)$ to reach j ,
 - and the cost $(200 - (a_i - a_j))^2$ of a one-day trip from j to i .

5

5

Optimal Binary Search Trees

6