

Number Theory for Information Security



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Contents



- Quick Review of Modular Arithmetic
- Congruences, Exponentiation
- Review of Groups, Rings, Fields
- Galois Fields
- Euler's Totient Function
- Euler's Phi Function
- Fermat's Little Theorem
- Euler's Theorem
- Generator, Order of a group

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One-to-one & onto functions



- def: one-to-one:
 - A function is 1-1, if each element in the codomain Y is the image of **at most** one element in the domain X .
- def: onto:
 - A function is onto, if each element in the codomain Y is the image of **at least** one element in the domain X . A function $f: X \rightarrow Y$ is onto, if $\text{Im}(f) = Y$.

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Tutorial#1



- Consider a function F whose domain-range are $f: \{a,b,c,d,e,f...z\} \rightarrow \{0,1,2,3,4,5.....25\}$ with the definition as follows:

$$f(i^{\text{th}} \text{ letter of alphabet}) = i-1$$

Analyze whether this function is one-to-one and onto or not ?

- Consider a function g whose domain-range are $g: \{\text{binary bit strings of length 4}\} \rightarrow \{\text{binary bit strings of length 3}\}$ with the definition as follows:

$$g(b_1b_2b_3b_4) = b_1b_2b_4$$

Analyze whether this function is one-to-one and onto or not ?

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Bijection of a function



- def: If a function $f: X \rightarrow Y$ is 1-1 and $\text{Im}(f) = Y$, then f is called a **bijection**.
- Obser#1: If $f: X \rightarrow Y$ is 1-1 then $f: X \rightarrow \text{Im}(f)$ is a bijection
 - i.e. if $f: X \rightarrow Y$ is 1-1 and X and Y are finite sets of the same size. Why the latter ?

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Inverse of a function



- def: If f is a bijection from X to Y ,
 - then **there exists a bijection g from Y to X** also i.e.
 - for each $y \in Y$, $g(y) = x$, where $x \in X$ and $f(x) = y$.
 - Then, the function g so obtained from f is called **the inverse function** of f i.e $g = f^{-1}$

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Bijection & inverse of a function (contd)



- Let $X = \{a,b,c,d,e\}$ and $Y = \{1,2,3,4,5\}$ and let f be defined such that
 $f(a) = 5, f(b) = 3, f(c)=4, f(d)=1, f(e)=2$, then
 - f is one-to-one
 - Since $\text{Im}(f) = \{1,2,3,4,5\} = Y$, f is onto and it is bijection
 - The inverse function of f can be formed by defining a g such that.....
 - If f is a bijection, so is f^{-1}
- Bijections are the heart of the cryptography.....Why ?

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Bijection & inverse of a function (contd)



- In cryptography,
 - bijections are used to as a tool for encryption and the
 - inverse are used for decryption
 - Why bijections are required for encryption ?

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Tutorial#3: Bijection Functions



- Are the DES or the AES – the symmetric key cryptography algorithms bijections ?
- Is the RSA function a bijection ?

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Ciphers and the property of Determinism



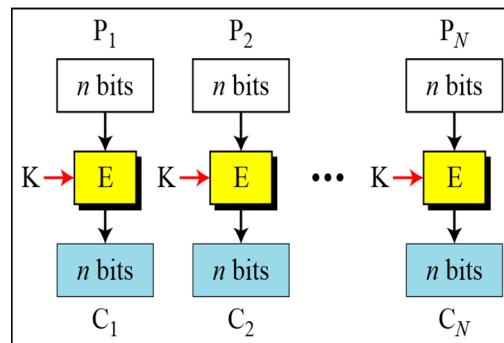
- AES, DES, RC5.....are these ciphers deterministic or probabilistic ?
- Determinism and Semantic Security
- How to introduce probabilistic nature in cipher implementation ?

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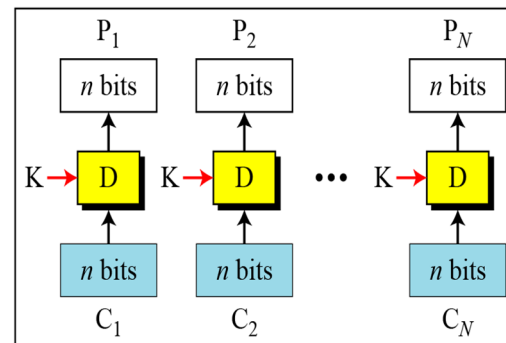
ECB Block Cipher Mode



E: Encryption D: Decryption
 P_i : Plaintext block i C_i : Ciphertext block i
 K: Secret key



Encryption



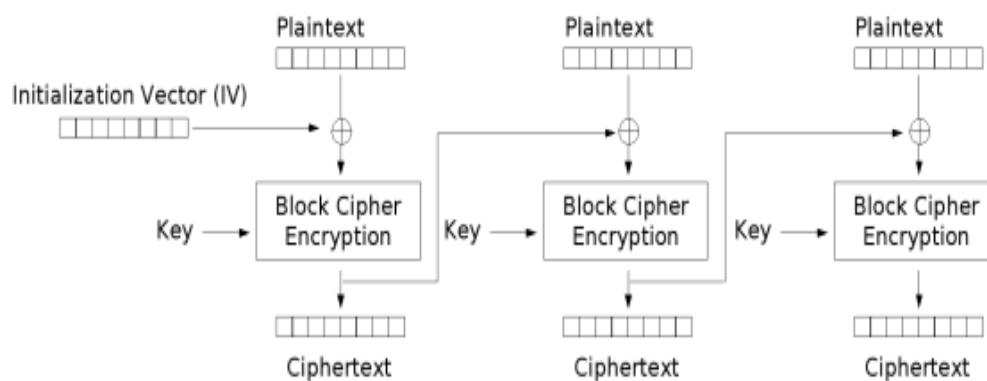
Decryption

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CBC Encryption & Use of an IV [wiki]



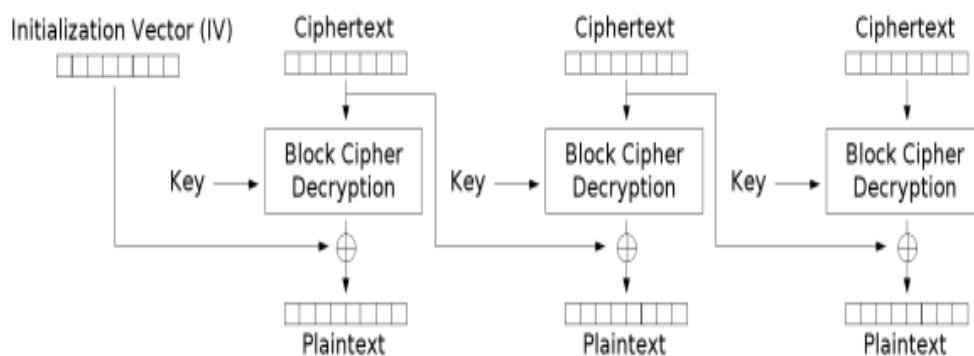
Cipher Block Chaining (CBC) mode encryption

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CBC decryption & Use of an IV [wiki]



Cipher Block Chaining (CBC) mode decryption

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One way functions



- def: A function $f: X \rightarrow Y$ is called a one way function
 - if $f(x)$ is *easy* to compute for all $x \in X$, but
 - for “essentially all” elements of $y \in \text{Im}(f)$, it is **computationally infeasible** to find any $x \in X$, such that $f(x)=y$.

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One way function (contd)



- Illustration: Let $X = \{1, 2, 3, \dots, 16\}$ and let $f(x) = r_x$ for all $x \in X$, where r_x is the remainder when 3^x is divided by 17. What is then $f(x)$?
 - Is it feasible to compute $f(x)$ from x ?
- | | | | | | | | | | | | | | | | |
|---|---|----|----|---|----|----|----|----|----|----|----|----|----|----|----|
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| 3 | 9 | 10 | 13 | 5 | 15 | 11 | 16 | 14 | 8 | 7 | 4 | 12 | 2 | 6 | 1 |
- Is it feasible to compute x from $f(x)$?

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The Trapdoor oneway functions



- Illustration: Let
 - primes $p = 48611$ and $q = 53993$, number $n = pq = 2624653723$ and let $X = \{1, 2, 3, 4, \dots, n-1\}$.
 - let a function $f_x = r_x$ be defined for each $x \in X$, where r_x is the remainder when x^3 is divided by n .
 - e.g. $f(248991) = 1981394214$ as
 - $248991^3 = 5881949859 * n + 1981394214$
 - IS it easy to compute the value of $f(x)$ given x ?
- Finding the reverse ...i.e. is it easy to compute x given $f(x)$?
 - Computation of modular cuberoot with modulus n
 - if the factors of n are unknown and large then it is a difficult problem.
- Such functions are the trapdoor oneway functions.....

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The Trapdoor oneway functions (contd)



- def: a function $f: X \rightarrow Y$ is a trapdoor one way function, it is one-way function
 - with the additional property that given some extra trapdoor information it becomes feasible to find for any given $y \in \text{Im}(f)$, an $x \in X$, such that $f(x) = y$.
- In the example above, knowing p and q (each five digits long), it is easy to invert the function.
- What should be the length of digits in p and q to make it infeasible?
 - at least 100 digits
 - well-known integer factorization problem.
- The existence of such functions is difficult to rigorously prove, mathematically.

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Modular Arithmetic



- Any integer a can be expressed as $a = qn + r$; $0 \leq r < n$; $q = \lfloor a/n \rfloor$
 - e.g. in modulo 7 arithmetic, $11 = 1 \times 7 + 4$ i.e. $r = 4$ and
 - $-11 = ?$
 - $-11 = -2 \times 7 + 3$ to yield $r = 3$.
- def: **modulo operator** “ $a \bmod n$ ” is defined as the remainder b
 - when a is divided by n , b is called the **residue** of $a \bmod n$
 - usually choose **the smallest positive remainder** as residue $0 \leq b < n$
 - the process is known as **modulo reduction**

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Congruent modulo n



- Finite fields have become increasingly important in cryptography.
- Two integers a and b are said to be **congruent modulo n** if $(a \bmod n) = (b \bmod n)$
 - i.e. when divided by n , a & b have the same remainder
 - i.e. when n divides $b-a$
 - e.g. $(100 \bmod 11) = (34 \bmod 11)$
 - denoted as $100 \equiv 34 \bmod 11$
 - Is $-12 \equiv -5 \bmod 7$ true ?
 - Is $2 \equiv 9 \bmod 7$ true ?
 - Is $73 \equiv 4 \bmod 23$ true ?
 - Is $21 \equiv -9 \bmod 10$ true?

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Tutorial #4:



- State whether true or false:
 - Is $13 \equiv 523 \bmod 17$?
 - Is $-15 \equiv 6 \bmod 7$ true ?
 - Is $-14 \equiv 1 \bmod 3$ true ?
 - Is $73 \equiv 4 \bmod 23$ true ?
 - Is $21 \equiv -9 \bmod 10$ true?
 - Is $82 \equiv 1 \bmod 9$ true ?
 - Is $-82 \equiv 1 \bmod 9$ true ?
 - Is $63 \equiv 8 \bmod 11$ true ?
 - Is $-63 \equiv 3 \bmod 11$ true ?
 - Is $121 \equiv 1 \bmod 15$ true
 - Is $-119 \equiv 1 \bmod 15$ true ?

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Congruence



- Congruence modulo n is an **equivalence relation** on the integers.
- What is an **equivalence relation** ?
- Congruence modulo n is an equivalence relation on the integers.....given the three properties are true. Which ones ?
 - any integer is congruent to itself modulo n (reflexivity). **How ?**
 - $a \equiv b \pmod{n}$ implies that $b \equiv a \pmod{n}$ (symmetry). **How?**
 - $a \equiv b \pmod{n}$ and $b \equiv c \pmod{n}$ implies that $a \equiv c \pmod{n}$ (transitivity). **How ?**

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Z_n - The integers modulo n



- def: The set of integers modulo n i.e. Z_n is the set of (equivalence classes of) integers $\{0, 1, 2, \dots, n-1\}$.
- All the operations in Z_n viz.
 - multiplication, addition and subtraction are performed modulo n .
 - e.g. $Z_{25} = \{0, 1, 2, 3, \dots, 24\}$. Then,
 - $6 + 14 = ?$ in Z_{25} $14 + 14 = ?$ in Z_{25}
 - $15 + 35 = ?$ in Z_{25} $20 + 32 = ?$ in Z_{25}
 - e.g. $Z_{49} = \{0, 1, 2, 3, \dots, 48\}$. Then
 - $21 + 23 = ?$ in Z_{49}
 - $35 + 35 = ?$ in Z_{49}

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The additive inverse



- The additive inverse of a number a in modular arithmetic is the integer y such that $x + y = 0 \pmod n$.
- e.g. addition arithmetic modulo 8 is as shown in the table.
- What are the AIs of 1, 2, 3, 5 in modulo 8?

+	0	1	2	3	4	5	6	7
0	0	1	2	3	4	5	6	7
1	1	2	3	4	5	6	7	0
2	2	3	4	5	6	7	0	1
3	3	4	5	6	7	0	1	2
4	4	5	6	7	0	1	2	3
5	5	6	7	0	1	2	3	4
6	6	7	0	1	2	3	4	5
7	7	0	1	2	3	4	5	6

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The multiplicative inverse



- The multiplicative inverse of a number a is a number b such that $a * b = 1 \pmod n$.
 - if exists, it is unique
- e.g. the table shows the multiplication modulo 7
- unlike additive inverse, the multiplicative inverse of a number may not exist e.g.
 - what are the MIs of 2,3,4?
 - what are the MIs of 4 in modulo 8?

×	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

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Abstract Algebra



- Finite fields
 - are of increasing importance in cryptography
 - AES, Elliptic Curve, IDEA, Public Key
 - concern with operations on “numbers” where
 - what *constitutes* a “number” and the *type* of operations varies considerably
 - start with concepts of groups, rings, fields from abstract algebra

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Group



- A **Group** is a set of elements or “numbers” with some operation $*$ such that
 - closure: whose result is also in the set
 - associative law: $(a.b).c = a.(b.c)$
 - has identity $e: e.a = a.e = a$
 - has inverses $a^{-1}: a.a^{-1} = e$
- Semigroup, Monoid, **Group**in that order
- A group
 - if is commutative $a.b = b.a$ then it forms an **abelian group**
- Finite group, order of a finite group
- Infinite group

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Group...



- def: A group $(G, *)$ consists of a set G with a binary operation $*$ on G satisfying the following three axioms:
 - the group operation is **associative** i.e. $a*(b*c) = (a*b)*c$ for all $a, b, c \in G$.
 - there is an element $1 \in G$, called the **identity element**, such that $a * 1 = 1 * a = a$ for all $a \in G$
 - for each $a \in G$ there exists an element $a^{-1} \in G$, called the **inverse** of a such that $a * a^{-1} = a^{-1} * a = 1$
- for a group G , if $a * b = b * a$ for all $a, b \in G$, then the group G is **abelian or commutative**.

Group – illustrations



- The set of integers Z_n with the operation of addition modulo n forms a group of order n . Identity element = ? Inverse of a = ?
 - Is it an abelian group, too?
- The set of real numbers under multiplication is an abelian group.
- Is the set of integers Z_n with the operation of multiplication modulo n , a group of order n ?
- Is the set of integers Z_n with the operation of multiplication modulo n , a monoid ?

The multiplicative inverse



- e.g. the table shows the multiplication modulo 7
- unlike additive inverse, the multiplicative inverse of a number may not exist e.g.
 - what are the MIs of 2,3,4 ?
- what are the MIs of 4 in modulo 8 ?

×	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

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Group – illustration



- Ex: Let set $G_{XOR} = \{EVEN, ODD\}$ and a binary operation \oplus be defined as

\oplus	EVEN	ODD
EVEN	EVEN	ODD
ODD	ODD	EVEN

- Is it a closed under operation \oplus ?
- Does it exhibit associativity ?
- What is the identity element?
 - EVEN
- Does every element have an inverse?
 - What are the inverses of ODD and EVEN elements?

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Ring



- A set of “numbers”
 - with two operations (addition and multiplication) denoted as $(R, +, \times)$ and
 - which forms an abelian group with addition operation (identity 0)
 - multiplication operation
 - has closure
 - is associative i.e. $a \times (b \times c) = (a \times b) \times c$ for all $a, b, c \in R$
 - distributive over addition i.e. $a \times (b + c) = a \times b + a \times c$
- i.e. a ring is a set in which we can do addition, subtraction and multiplication without leaving the set.
- e.g. the set of integers \mathbb{Z} with $+$ supported is a ring
- e.g. is the set of integers \mathbb{Z} with \times supported a ring?

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Invertible element and Field



- An element a of a ring R is called a unit or an invertible element
 - if there is an element $b \in R$ such that $a \times b = 1$.
- A **FIELD** is a set in which we can do addition, subtraction, multiplication, and division without leaving the set.
 - Division is defined with the following rule: $a/b = a(b^{-1})$. We denote a Field as $\{F, +, \cdot\}$

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Field...



- def: A field is a commutative ring in which all the non-zero elements have multiplicative inverses.
 - e.g. the set of integers under the $+$ & \times operations is not a field. Why?
- Are the sets (of) rational numbers, real numbers, complex numbers a field?
- Z_n is a field iff n is a prime number.
- These have hierarchy with more axioms/laws
 - group \rightarrow ring \rightarrow field

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Galois Fields



- Infinite fields are of not much interest. But, finite fields play a key role in cryptography.
- The number of elements in a finite field
 - i.e. the order of a finite field must be a power of a prime p^n , $n \geq 1$
 - the finite field of the order of p^n are known as **Galois fields**
- denoted $GF(p^n)$
- in particular often use the fields
 - $GF(p)$
 - $GF(2^n)$

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Galois Fields $GF(p)$



- $GF(p)$ is the set of integers $Z_p = \{0, 1, \dots, p-1\}$ with arithmetic operations modulo prime p
- these form a finite field - since each element has multiplicative inverse
- hence arithmetic is “well-behaved” and
 - can do addition, subtraction, multiplication, and division without leaving the field $GF(p)$

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Addition modulo 7

+	0	1	2	3	4	5	6	7
0	0	1	2	3	4	5	6	7
1	1	2	3	4	5	6	7	0
2	2	3	4	5	6	7	0	1
3	3	4	5	6	7	0	1	2
4	4	5	6	7	0	1	2	3
5	5	6	7	0	1	2	3	4
6	6	7	0	1	2	3	4	5
7	7	0	1	2	3	4	5	6

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GF(7) Multiplication Example



×	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

How to find inverses when the numbers involved are very large ???

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Finding Inverses – Extended Euclidean algorithm



EXTENDED_EUCLID(m, b)

1. $(A1, A2, A3) = (1, 0, m);$
 $(B1, B2, B3) = (0, 1, b)$
2. **if** $B3 = 0$
return $A3 = \text{gcd}(m, b);$ no inverse
3. **if** $B3 = 1$
return $B3 = \text{gcd}(m, b); B2 = b^{-1} \bmod m$
4. $Q = A3 \text{ div } B3$
5. $(T1, T2, T3) = (A1 - Q B1, A2 - Q B2, A3 - Q B3)$
6. $(A1, A2, A3) = (B1, B2, B3)$
7. $(B1, B2, B3) = (T1, T2, T3)$
8. **goto** 2

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Inverse of 17 in GF(29)



i.e. calling Extended_Euclid(29, 17)

Q	A1	A2	A3	B1	B2	B3
—	1	0	29	0	1	17
1	0	1	17	1	-1	12
1	1	-1	12	-1	2	5
2	-1	2	5	3	-5	2
2	3	-5	2	-7	12	1

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Inverse of 17 in GF(29)



i.e. calling Extended_Euclid(29, 17)

Q	A1	A2	A3	B1	B2	B3
—	1	0	29	0	1	17
1	0	1	17	1	-1	12
1	1	-1	12	-1	2	5
2	-1	2	5	3	-5	2
2	3	-5	2	-7	12	1

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Inverse of 37 in GF(49)



i.e. calling Extended_Euclid(49, 37)

Q	A1	A2	A3	B1	B2	B3
—	1	0	49	0	1	37
1	0	1	37	1	-1	12
3	0	1	12	-3	4	1

• Hence $37^{-1} \equiv 4 \pmod{49}$ OR $4 = 37^{-1} \pmod{49}$

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Inverse of 550 in GF(1759)



i.e. calling Extended_Euclid(1759, 550)

Q	A1	A2	A3	B1	B2	B3
—	1	0	1759	0	1	550
3	0	1	550	1	-3	109
5	1	-3	109	-5	16	5
21	-5	16	5	106	-339	4
1	106	-339	4	-111	355	1

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Inverse of 49 in GF(37)



i.e. calling Extended_Euclid(37, 49)

Q	A1	A2	A3	B1	B2	B3
—	1	0	37	0	1	49
0	0	1	49	1	0	37
1	1	0	37	-1	1	12
3	-1	1	12	4	-3	1

- Hence $49^{-1} \equiv (-3) \pmod{37}$
- But, $-3 \pmod{37} \equiv 34 \pmod{37}$. Hence,
- $34 = 37^{-1} \pmod{49}$

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Tutorial #11



- Find the inverse of the following elements in the GF as indicated::
 1. Inverse of 8 GF(19)
 2. Inverse of 17 in GF(29)
 3. Inverse of 13 in GF(29)
 4. Inverse of 49 in GF(37)
 5. Inverse of 351 in GF(771)
 6. Inverse of 17 in GF(331)

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The Euler Totient function



- def: For $n \geq 1$, let $\phi(n)$ denote the number of integers in the interval $[1, n]$ which **are relatively prime** to n .
- The function ϕ is called the Euler Totient function.
- Note that, when we are doing arithmetic modulo n
 - **the complete set of residues** is : $0, \dots, n-1$, whereas,
 - **the reduced set of residues** is those numbers (residues) which are **relatively prime** to n
 - eg for $n=10$,
 - the complete set of residues is $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$
 - the reduced set of residues is $\{1, 3, 7, 9\}$

The Euler Totient function - Properties



- So, the number of elements in **reduced set of residues** is called the **Euler Totient Function $\phi(n)$**
- Properties:
 1. If p is prime then $\phi(p) = p - 1$.
 2. The function ϕ is multiplicative i.e. if $\gcd(m, n) = 1$, then $\phi(mn) = \phi(m) \cdot \phi(n)$.
 3. If $n = p_1^{e_1} p_2^{e_2} p_3^{e_3} \dots p_k^{e_k}$ is the prime factorization of n , then,

$$\phi(n) = n \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \left(1 - \frac{1}{p_3}\right) \dots \dots \dots \left(1 - \frac{1}{p_k}\right)$$

Euler Totient Illustration



- $\phi(1), \phi(2), \phi(3), \phi(4), \phi(6), \phi(7), \phi(14), \phi(23), \phi(15)$
 - $\phi(1) = 0$ - given by $p = p-1$
 - $\phi(2) = |\{1\}|$ - given by $p = p-1$
 - $\phi(3) = |\{1,2\}|$ - given by $p = p-1$
 - $\phi(4) = |\{1,3\}|$ - 4 is not a prime
 - $\phi(6) = |\{1,5\}|$ - 6 is not a prime
 - $\phi(6) = \phi(3) * \phi(2) = 2 * 1 = 2$
 - $\phi(7) = |\{1, 2, 3, 4, 5, 6\}|$ - given by $p = p-1$
 - $\phi(14) = |\{1,3,5,9,11,13\}|$ - 14 is not a prime
 - $\phi(14) = \phi(7) * \phi(2) = 6 * 1 = 6$
 - $\phi(23) = |\{1,2,3,\dots,22\}|$ - 23 is prime
 - $\phi(15) = ?$
 - $4 * 2 = 8$

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Euler's Totient Function



- If $n = p_1^{e_1} p_2^{e_2} p_3^{e_3} \dots p_k^{e_k}$ is the prime factorization of n , then,

$$\phi(n) = n \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \left(1 - \frac{1}{p_3}\right) \dots \left(1 - \frac{1}{p_k}\right)$$
- e.g. $616 = 2^3 * 7 * 11$
- Therefore,
 - $\phi(616) = 616 * (1 - 1/2) * (1 - 1/7) * (1 - 1/11)$

$$= 616 * 1/2 * 6/7 * 10/11$$

$$= 240.$$

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Tutorial #12



■ Find the Euler's Totient Function of the following:

- $\Phi(273)$
- $\Phi(393)$
- $\Phi(495)$
- $\Phi(289)$
- $\Phi(169)$
- $\Phi(274)$
- $\Phi(472)$
- $\Phi(65)$
- $\Phi(127)$
- $\Phi(133)$
- $\Phi(201)$
- $\Phi(333)$

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Applications of finding $\Phi(n)$ - RSA



- Each user generates a public/private key pair by the following process
 - select two large distinct primes at random - p, q .
 - compute their system modulus $n=p \cdot q$
 - compute $\phi(n)$ How ?
 - select at random the encryption key e
 - where $1 < e < \phi(n)$, $\gcd(e, \phi(n)) = 1$
 - solve the following equation to find decryption key d
 - $e \cdot d = 1 \pmod{\phi(n)}$ and $0 \leq d \leq n$. How ?
 - publish the public encryption key: $PU = \{e, n\}$
 - keep secret private decryption key: $PR = \{d, n\}$
- It is critically important that the factors p & q of the modulus n are kept secret

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Multiplicative Group



- The multiplicative group of Z_n is denoted by Z_n^*
- def: defined as $Z_n^* = \{a \in Z_n \mid \gcd(a, n) = 1\}$
 - If n is prime, then $Z_n^* = \{a \mid 1 \leq a \leq n-1\}$
 - If $a \in Z_n^*$ and $b \in Z_n^*$, then $a.b \in Z_n^*$.
- Let $n = 21$. Then, $Z_{21}^* = \{1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, 20\}$

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Multiplicative Group



- The order of a multiplicative group Z_n^* - denoted $|Z_n^*|$ is defined as
 - $|Z_n^*|$ i.e. the number of elements in Z_n^* .
- Recollect that if n is prime, then $Z_n^* = \{a \mid 1 \leq a \leq n-1\}$
- Illustration:
 - Let $n = 21$. Then, $Z_{21}^* = \{1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, 20\}$
 - Now, $\phi(21) =$
 - $\phi(7) \cdot \phi(3) = 6 \cdot 2 = 12 = |Z_{21}^*|$

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Euler's theorem



- Let $n \geq 2$ be an integer. Then if $a \in \mathbb{Z}_n^*$,
 $a^{\phi(n)} \equiv 1 \pmod{n}$
- e.g.
 - $a=3; n=10; \phi(10)=4;$
 hence $3^4 = 81 \equiv 1 \pmod{10}$
 - $a=2; n=11; \phi(11)=10;$
 hence $2^{10} = 1024 \equiv 1 \pmod{11}$
- If n is a product of distinct primes,
 - and if $r \equiv s \pmod{\phi(n)}$, then $a^r \equiv a^s \pmod{n}$
 - i.e. when working with modulo such as n , exponents can be reduced modulo $\phi(n)$

What about $a=7$
 i.e. $7^4 \pmod{10}$?
 And $a=5$?

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Order of elements of an MG

- Let $a \in \mathbb{Z}_n^*$. Then, the **order of a** , denoted by **ord(a)**,
 - is the **least** positive integer t such that $a^t \equiv 1 \pmod{n}$
 - e.g. consider again $\mathbb{Z}_{21}^* = \{1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, 20\}$
 - $\phi(21)=12 = |\mathbb{Z}_{21}^*|$.
 - Now the orders of various elements in \mathbb{Z}_{21}^* are:

a	1	2	4	5	8	10	11	13	16	17	19	20
Ord(a)	1	6	3	6	2	6	6	2	3	6	6	2

- $\text{Ord}(a) = \text{mod}(\text{power}(a, Ai), 21)$ in Excel sheet

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Generator, Cyclic group

- Let $\alpha \in Z_n^*$.
 - if the order of α is $\phi(n)$, then α is said to be a generator or a primitive element of Z_n^* .
 - Are there any generators in the group Z_{21}^* ?

a	1	2	3	4	5	6	7	8	9	10
Ord(a)	1	6	–	3	6	–	–	2	–	6
a	11	12	13	14	15	16	17	18	19	20
Ord(a)	6	–	2	–	–	3	6	–	6	2

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Generator, Cyclic group

- IF Z_n^* has a generator, then Z_n^* is said to be a cyclic group.
 - In the above example, Z_{21}^* is not a cyclic group, since no generator is equal to $\phi(n)$ i.e. 12.

a	1	2	4	5	8	10	11	13	16	17	19	20
Ord(a)	1	6	3	6	2	6	6	2	3	6	6	2

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Generator, Cyclic group (contd)

□ Consider now a group Z_{25}^*

□ $Z_{25}^* = \{1, 2, 3, 4, 6, 7, 8, 9, 11, 12, 13, 14, 16, 17, 18, 19, 21, 22, 23, 24\}$

□ i.e. $\Phi(25) = |Z_{25}^*| = 20$

□ Now the orders of various elements in Z_{25}^* are:

Use the formula $\text{Ord}(a) = \text{mod}(\text{power}(a, \text{Ai}), 25)$ in Excelsheet												
a	1	2	3	4	6	7	8	9	10	11	12	13
Ord(a)	1	20	20	10	5	5	20	10	–	5	?	?
a	14	15	16	17	18	19	21	23	24			
Ord(a)	?	?	?	?	?	?						

□ Thus, Z_{25}^* is indeed a cyclic group because 2,3,8,... are the generators of the group.

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Generator, Cyclic group (contd)..

Members of Z_{25}^* are {1,2,3,4,6,7,8,9,11,12,13,14,16,17,18,19,21,22,23,24} and $\Phi(25)=20$

Z_{25}^* with a=2	Z_{25}^* with a=3	Z_{25}^* with a=4	Z_{25}^* with a=6	Z_{25}^* with a=7	Z_{25}^* with a=8	Z_{25}^* with a=9	Z_{25}^* with a=12	Z_{25}^* with a=13	Z_{25}^* with a=14	Z_{25}^* with a=16	Z_{25}^* with a=17	Z_{25}^* with a=18	Z_{25}^* with a=19	Z_{25}^* with a=21	Z_{25}^* with a=22	Z_{25}^* with a=23	Z_{25}^* with a=24		
1	2	3	4	6	7	8	9	11	12	13	14	16	17	18	19	21	22	23	24
2	4	9	16	11	24	14	6	21	19	19	21	6	14	24	11	16	9	4	1
3	8	2	14	16	18	12	4	6	3	22	19	21	13	7	9	11	23	17	24
4	16	6	6	21	1	21	11	16	11	16	11	21	1	21	6	6	16	1	
5	7	18	24	1	7	18	24	1	7	18	24	1	7	18	24	1	7	18	24
6	14	4	21	6	24	19	16	11	9	9	11	16	19	24	6	21	4	14	1
7	3	12	9	11	18	2	19	21	8	17	4	6	23	7	14	16	13	22	24
8	6	11	11	16	1	16	21	6	21	21	6	21	16	1	16	11	11	6	1
9	12	8	19	21	7	3	14	16	2	23	9	11	22	18	4	6	17	13	24
10	24	24	1	1	24	24	1	1	24	24	1	1	24	24	1	1	24		
11	23	22	4	6	18	17	9	11	13	12	14	16	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
12	21	16	16	11	1	11	6	21	6	6	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
13	17	23	14	16	7	13	4	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
14	9	19	6	21	24	4	11	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
15	18	7	24	1	18	7	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
16	11	21	21	6	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
17	22	13	9	11	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
18	19	14	11	#NUM!	#NUM!	9	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
19	13	17	18	#NUM!	#NUM!	22	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
20	1	1	1	#NUM!	#NUM!	1	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!	#NUM!
21	2	3	4			8													
22	4	9	16																
23	8	2	#NUM!																

Snapshot of Z_{25}^* computation from the Excel sheet

Snapshot of Z_{25}^* computation from the Excel sheet

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Generator, Cyclic group (contd)

- Consider now a multiplicative group Z_{13}^*
 - $Z_{13}^* = \{1, 2, 3, 4, 6, 7, 8, 9, 11, 12\}$
 - i.e. $\Phi(13) = |Z_{13}^*| = 12$
 - Compute the orders of various elements in Z_{13}^* :

α	0	1	2	3	4	5	6	7	8	9	10	11
$\alpha^i \bmod 13$	1	6	12	3	7	4	12	12	4	3	6	12

- Thus,
 - $\alpha = 2, 6, 7, 11$ are the generators of the group.
 - Note the case of $5^t \bmod 13$ with $t=4,12$.

Generators.....



- How many Generators can be there of a group if Z_n^* is a cyclic group ?
 - if Z_n^* is cyclic, then the number of generators is $\Phi(\Phi(n))$.
 - e.g. Z_{21}^* is not cyclic – doesn't have a generator because n does not satisfy any of the conditions above in first
- Are $Z_{11}^*, Z_7^*, Z_{13}^*, Z_{17}^*, Z_{19}^*$ cyclic ?
- Is Z_{30}^* cyclic ? $\Phi(30)$ is $\Phi(6) * \Phi(5) = 2 * 4 = 8$.

How to test for a given number to be a Generator



- Consider a MG Z_p^* , where p is a prime.
- Then, it is easy to test whether a given element is its generator or not. How ?
 - As p is a prime, $\Phi(p) = p-1$, and
 - the number of generators in it is $\Phi(p-1)$,
 - now, if $p_1, p_2, p_3, \dots, p_k$ are the distinct prime factors of $p-1$, then,
 - g is a generator of Z_p^* if and only if

$$g^{(p-1)/p_i} \not\equiv 1 \pmod{p} \text{ for all } p_i, 1 \leq i \leq k$$

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How to test for a given number to be a



- e.g. consider Z_{13}^* . Check whether 7 is a generator or not.
- Now,
 - $\Phi(13) = p-1 = 12$, and
 - the number of generators in it is $\Phi(p-1) = \Phi(12) = 4$.
 - Also, the distinct prime factors of $p-1$ i.e. 12 are 2,3. Hence, $p_1=2$, $p_2=3$.
 - Then,
 - $g^{(p-1)/p_1} = 7^{12/2} = 7^6 \pmod{13} = 12 \pmod{13} \not\equiv 1 \pmod{13}$, and
 - $g^{(p-1)/p_2} = 7^{12/3} = 7^4 \pmod{13} = 9 \pmod{13} \not\equiv 1 \pmod{13}$
- Hence, 7 is indeed a generator of Z_{13}^*

$$g^{(p-1)/p_i} \not\equiv 1 \pmod{p} \text{ for all } p_i, 1 \leq i \leq k$$

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How to test for a given number to be a



- e.g. consider Z_{13}^* . Now, check whether 8 is a generator or not.
- Now,
 - $\Phi(13) = p-1 = 12$, and
 - the number of generators in it is $\Phi(p-1) = \Phi(12) = 4$.
 - Also, the distinct prime factors of $p-1$ i.e. 12 are 2, 3. Hence, $p_1=2$, $p_2=3$.
 - Then,
 - $g^{(p-1)/p_1} = 8^{12/2} = 8^6 \bmod 13 = 12 \bmod 13 \neq 1 \bmod 13$, and
 - $g^{(p-1)/p_2} = 8^{12/3} = 8^4 \bmod 13 = 1 \bmod 13$
- Hence, **8 is NOT a generator of Z_{13}^***

$$g^{(p-1)/p_i} \neq 1 \bmod p \text{ for all } p_i, 1 \leq i \leq k$$

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Thank You !!!