

Problem - 0

Implement the fibonacci series and debug the code.

```
#include <stdio.h>
int fib(int n)
{
    if (n == 0)
        return 0;
    else if (n == 1)
        return 1;
    else
        return fib(n-1) + fib(n-2);
}
```

```
int main ()
{
    int result;
    result = fib(5);
    return 0;
}
```

If we analyse the above algorithm and pass the 'n' value as 5, the following are the recursive calls and function call stack.

fib(5) ; n=5 ; fib(5) calls fib(4) and fib(3)

fib(4) ; n=4 ; fib(4) calls fib(3) and fib(2)

fib(3) ; n=3 ; fib(3) calls fib(2) and fib(1)

fib(2) ; n=2 ; fib(2) calls fib(1) and fib(0)

fib(1) ; n=1 ; fib(1) returns 1

fib(0) ; n=0 ; fib(0) returns 0

fib(0) returns 0 fib(1) returns 1

$$\text{fib}(2) ; n=2 \Rightarrow 1 + \text{fib}(0)$$

$$\text{fib}(0) ; n=0 \Rightarrow \text{returns } 0$$

$$\text{fib}(2) \text{ returns } 1 + 0 = 1 \text{ \# the sum of } \text{fib}(1) \text{ \& } \text{fib}(0)$$

$$\text{fib}(3) ; n=3 \Rightarrow \text{fib}(2) + \text{fib}(1)$$

$$\Rightarrow 1 + \text{fib}(1)$$

$$\text{fib}(1) \text{ returns } 1$$

$$\text{fib}(3) \text{ returns } 1 + 1 = 2 \text{ \# the sum of } \text{fib}(2) \text{ \& } \text{fib}(1)$$

$$\text{fib}(4) ; n=4 \Rightarrow 2 + \text{fib}(2)$$

$$\text{fib}(2) ; n=2 \Rightarrow \text{fib}(1) + \text{fib}(0) = 1$$

$$\text{fib}(4) \text{ returns } 2 + 1 = 3 \text{ \# the sum of } \text{fib}(3) \text{ \& } \text{fib}(2)$$

$$\text{fib}(5) ; n=5 \Rightarrow \text{fib}(4) + \text{fib}(3)$$

$$= 3 + \text{fib}(3)$$

$$\text{fib}(3) ; n=3 \Rightarrow \text{returns } 2$$

$$\text{fib}(5) ; n=5 \Rightarrow 3 + 2 = 5 \text{ \# the sum of } \text{fib}(4) \text{ \& } \text{fib}(3)$$

\therefore Hence, $\text{fib}(5) = 5$ by using the above recursive calls.

Problem - 1 : Merge Arrays :

a) code uploaded in Github

b) Time Complexity : for the function mergeKarrays
Since there is 1 for loop, it iterates over each array in the arrays list and merges them using merge function. So, it runs 'n' times.

So, time complexity is $O(n)$

Since, we are iterating over 'k' arrays and merging them one by one. The overall time complexity will be given as :

$$T(n) = \Theta(k * n)$$

c) ways to improve implementation.

• The current algorithm involves repeatedly appending elements to the list. This could increase the need of memory reallocation. So, instead of this, I could have created the space for the merged array.

• If the input arrays are larger, appending or concatenating them would be critical. So, we can implement a recursive sorting algorithm or consider optimising the input reading process.

Problem - 2 : Remove - duplicates

a) Code has been uploaded in Github.

b) Time complexity :

The for loop runs from 0 to $n-1$ where 'n' is length of input array. In each iteration, the loop checks the current element different from next.

So, here time complexity is $\Theta(n)$

The 'if' runs in constant time for each iteration.

ie : $\Theta(1)$

The dominant factors in time complexity is 'n' and the small factors are negligible.

So, time complexity = $\Theta(n)$

c) Ways to improve implementation :

I could have used more efficient and simpler algorithm. Though it has time complexity $\Theta(n)$. We can achieve linear time complexity by comparing adjacent elements and copying the different elements to the result array by avoiding duplicates.