

# Combinatorial Optimization Final Report

Nicoló Ruggeri

## 1 Introduction and First Notations

In this report we analyze the performances of two different methods for solving a particular instance of the TSP problem. Specifically, we solve the problem presented below with a linear programming network flow formulation, as requested in homework 1, and with a genetic algorithm presented in section 3.

In our case we formulate the problem on a complete graph  $G = (V, E)$  where the adjacency matrix  $A \in \mathbb{R}^{|V|}$  of  $G$  is symmetric and dense. In our case we consider a weighted graph, so that  $a_{i,j} \in \mathbb{R}_{\geq 0}$  is the cost of moving from point  $j$  to  $i$ . We will explain how these distances are generated in section 2.

The particular properties of the matrix reflect two assumptions that we make on our model, namely:

- $A$  symmetric. This means that the cost of going from  $j$  to  $i$  is the same as going the other way around. This makes sense for the underlying practical problem at hand, i.e. the one of a drill moving from position  $i$  to  $j$
- $A$  dense (and in the particular case of our experiments, with almost all elements strictly positive, see sections 2 and 4). With a slight abuse of terminology, here we mean that all entries of  $A$  are available (or, equivalently for the problem, lesser than infinite). This reflects the fact that  $ij \in E \ \forall i, j$ , i.e. the drill can move between any two arbitrary points

In the following section we will refer to points equivalently as nodes, and to the connections between them as edges (from the graph formulation).

## 2 Data Generation

### Distance Matrix

We frame the problem as that of moving between points in the unitary square  $[0, 1]^2$ . Notice that this doesn't represent a restriction, since it can be seen as a rescaling of a problem of any size. Given that we have a finite number of points, we can number them from 1 to  $n$ , so that the vertex set of the graph is  $V = \{1, \dots, n\}$ . Now we consider these two dimensional sampled vectors

$x_i \in [0, 1]^2$ , that represent the position of every point in the drilling board. If we want to generate the cost matrix  $A$  from these points, a natural interpretation could be to assume a movement time proportional to the distance between the points. In fact in the experiments section 4 we consider a distance matrix given by

$$a_{ij} := \|x_i - x_j\|_2 \quad (1)$$

(CHECK IF THE TWO NORM IS SQUARED OR NOOOOOOOOOOOOOOOOOOT)

The choice of the 2-norm is arbitrary, and could be replaced by any euclidean norm, which could in principle affect the final optimal solution.

Notice that this approach is easily generalized in many ways. For example to exclude some edges from the final solution we just need to impose  $a_{ij} = +\infty$ . To differentiate the cost between points we can substitute the formulation above with  $a_{ij} = f(x_i, x_j)$  for some function  $f$  (that doesn't even need to be symmetric in the arguments, as the genetic algorithm from section ).

## Experimental Data

Following the process above, we just need to generate any number  $n$  of points  $x_i \in [0, 1]^2$ , and then build the cost matrix that we will input to the solvers. In our case we decided to generate these points at random in the following way:

- a number  $n_1$  of points is generated according to a mixture of equiprobable two-dimensional gaussians with given variance  $\sigma^2$  and means  $\mu_1, \dots, \mu_k$ . Since the sampled points could end up being outside the allowed space, we clip all the points' coordinates in  $[0, 1]$ .
- the remaining  $n - n_1$  points are equally distributed on a regular grid on  $[0, 1]^2$ . This allows to have some clusters, given by the gaussians, and some other points scattered on the square.

CHECK VERY WELL THIS FORMULATION IF ON THE CODE IS THE SAMEEEEEEEEEEEEEEEEEEEEEEEEEEEEEEEEE Notice that this procedure allows, with small probability, overlapping points. This is not a problem for our model, as the solvers need to be able to recognize this peculiarity and find a solution where such points are visited sequentially, since this movement has cost 0 according to equation (1).

## 3 The Genetic Algorithm

After the implementation of the OPL model for homework 1, we decided to implement a heuristic for the solution of the problem described above. Specifically

## 4 Experiments