@ Ambiguous Grammar!

L) An Ambiguous grammar is a grammar with either more than one leftmost derivation tree or more than one right-most derivation trees.

Eliminating Ambiguity of a Context-Free Grammar;

let us take an example!

Girammar!

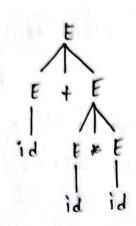
$$E \rightarrow E + E$$

$$E \rightarrow E \times E$$

$$E \rightarrow id$$

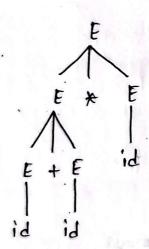
Input string: id+id *id

The leftmost derivation can be done in two ways!



Also, E > E*E

- → E+E*E
 - → id+ExE
 - > id+idxid



For the given input string, we get two left most derivation trees. We need to eliminate the ambiguity in the grammar.

L) Ambiguity from all grammar cannot be eliminated.

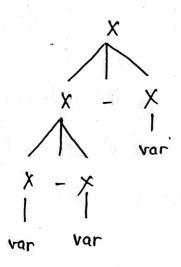
L) We can use Precedence and Associativity to remove the ambiguity from some grammar.

let us take an example!

Grammar!

$$X \rightarrow X - X$$

A string a-b-c has two liftmost doivations

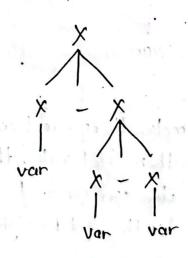


$$\lambda \rightarrow \lambda - \lambda$$

$$\rightarrow \lambda - \lambda - \lambda$$

$$\rightarrow var - var - var$$

$$\begin{array}{c} X \rightarrow X - X \\ \rightarrow var - X \\ \rightarrow var - X - X \\ \rightarrow var - var - X \\ \rightarrow var - var - var \\ \rightarrow a - b - c \end{array}$$



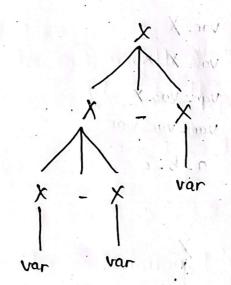
for example, if we take the values a=2, b=3 and c=4 a-b-c=2-3-4=-5

In the first dorivation tree, expression will be evaluated as

$$(a-b)-c=(2-9)-4=-5$$

In the second derivation tree, expression will be evaluated as

Observe that both parce trees aren't giving the same value. They have differed meaning. Therefore, the first derivation tree is the correct parce tree for grammar.



In the above example, there are two same operators in the expression. According to the mathematical rules, the expression must be evaluated based on the associativity of the operator used. The operator (-) which gives left to right associativity. So, for the left to right associative operators, the passe tree has to be left associative.

let us convert the grammar into unambiguous grammar:

 $X \rightarrow X_{-}x$ $X \rightarrow varlconst$

we need to make the grammar left-recursive we need to place a random non-terminal in place of the right non-terminal.

X-) X-PIP P > var/cont

Now for the string, a-b-c

 $X \rightarrow X - P$ $\Rightarrow X - X - P$ $\Rightarrow P - P - var$ $\Rightarrow var - var - var$

for the Gramman!

E → E * E E → E * E

- Ly this grammar will give two left derivation trees for the string, ideid id
- Ly we can't use associativity har as there are two different operators + and x. Hence we need to use "howedonce".
- is In the string id tidzid, the order of evaluation must be id+(idzid) as 'z' has more precedence than 't'. The operator at with the highest priority must be evaluated first. Hence, the operators with high priority are to be arranged in the lower levels of the passe tree.

E>F+PIP
P>P*919
Q>id