

## Practise problems 1

### Notation:

$\mathbb{N}$ : Set of natural numbers.

Boldface letters, for example  $\mathbf{x}, \mathbf{x}_i, \mathbf{a}_i, \mathbf{d}, \mathbf{d}_i, \mathbf{b}, \mathbf{y}_i$ , etc, are column vectors, whereas  $x_i$ 's are scalars.

$$Fea(LPP) = \{\mathbf{x} \in \mathbb{R}^n : A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$$

$$= \{\mathbf{x} \in \mathbb{R}^n : \mathbf{a}_i^T \mathbf{x} \leq b_i \text{ for all } i = 1, \dots, m, -\mathbf{e}_j^T \mathbf{x} \leq 0 \text{ for all } j = 1, \dots, n\}.$$

$$\tilde{\mathbf{a}}_i = \mathbf{a}_i \text{ for } i = 1, 2, \dots, m$$

$$= -\mathbf{e}_{i-m} \text{ for } i = m+1, m+2, \dots, m+n.$$

$$\tilde{b}_i = b_i \text{ for } i = 1, 2, \dots, m$$

$$= 0 \text{ for } i = m+1, m+2, \dots, m+n.$$

**Definition:** Two vectors  $\mathbf{a}$  and  $\mathbf{b}$  are said to be orthogonal to each other if  $\mathbf{a}^T \mathbf{b} = 0$ .

1. A furniture company manufactures four models of desks. Each desk is first constructed in the carpentry shop and is next sent to a finishing shop where it is varnished waxed and polished. The number of man hours required in each shop is as shown below.

	Desk 1 (hrs)	Desk 2 (hrs)	Desk 3 (hrs)	Desk 4 (hrs)	Available (hrs)
carpentry shop	4	9	7	10	6000
Finishing shop	1	1	3	40	4000

The profit from the sale of each item is as follows.

	Desk 1	Desk 2	Desk 3	Desk 4
Profit in Rs	12	20	80	40

Assuming that raw materials and supplies are available in adequate supply and all desks produced can be sold, the desk company wants to determine the quantities to make of each type of product which will maximize profit. Formulate this as a linear programming problem (you need not solve the LPP).

2. A furniture manufacturer has 3600 bd.ft of walnut, 4300 bd.ft of Maple, and 6550 bd.ft of Oak in stock. He can produce three types of products using this with input requirements as given below:

	Wood needed ( bd.ft /unit )			Revenue per unit
	Walnut	Maple	Oak	
Table	10	50	100	Rs 1000
Desk	10	30	40	Rs 500
Chair	80	5	0	Rs 100

A table is always sold in combination with 4 chairs; and a desk is always sold in combination with one chair. But a chair can be sold independently. Formulate the problem of finding how many of tables, chairs and desks he is going to produce to maximize his revenue, as a linear programming problem (ignoring the integer requirement of the variables).

3. Formulate the following machine problem as a linear programming problem.

An assembled item consists of two different metal parts. The milling work can be done on different machines: milling machines, turret lathes, or an automatic turret lathes.

The data are available in the following table.

Type of machine	Number of machines	Maximum output per machine per hour	
		First part	Second part
Milling Machines	3	10	20
Turret lathes	3	20	30
Automatic Turret lathes	1	30	80

4. A manufacturer of plastics is planning to blend a new product from four chemical compounds. These compounds are mainly composed of three elements, A, B and C. The composition and unit cost of these chemicals are given below:

Chemical compound	1	2	3	4
Percentage of A	30	20	40	20
Percentage of B	20	60	30	40
Percentage of C	40	15	25	30
Cost in rupees per Kilogram	20	30	20	15

The new product consists of 20 percent of element A, at least 30 percent of element B and at least 20 percent of element C. Due to possible side effects of compounds 1 and 2, they must not exceed 30 percent and 40 percent of the content of the new product. Formulate this as a LPP if the objective is to find the least cost of blending.

5. Consider the following problem (P):

$$\begin{aligned}
 &\text{Max } x_1 - x_2 \\
 &\text{subject to } x_1 + x_2 - x_3 \leq 2 \\
 &\quad 2x_1 + 3x_2 - 5x_3 = 2 \\
 &\quad 3x_1 + 4x_2 - 3x_3 \geq 0 \\
 &\quad x_1 \geq 0, x_2 \geq 0, x_3 \geq 0.
 \end{aligned}$$

Write the above problem in the form,

$$\text{Min } \mathbf{c}^T \mathbf{x}$$

subject to  $A_{4 \times 3} \mathbf{x} \leq \mathbf{b}$ ,  $\mathbf{x} \geq \mathbf{0}$ , and call the new problem as (P').

Check that (P) and (P') have the same optimal solution set.

If 20 is the optimal value of (P), then what is the optimal value of (P')?

6. Consider the following problem (P):

$$\begin{aligned}
 &\text{Max } x_1 + 4 \\
 &\text{subject to } x_1 + x_2 \leq 8 \\
 &\quad x_1 + x_2 \geq 2.
 \end{aligned}$$

Write the above problem in the form,

$$\text{Min } \mathbf{c}^T \mathbf{x}$$

subject to  $A_{2 \times 2} \mathbf{x} \leq \mathbf{b}$ ,  $\mathbf{x} \geq \mathbf{0}$ , and call the new problem as (P'), such that (P) and (P') have the same solution set.

Does the feasible region of (P) (that is the collection of all elements of  $\mathbb{R}^2$  which satisfies the two constraints) have any extreme point? Are the defining hyperplanes of (P) LI?

7. Let (P) be a LPP with  $(m+n)$  constraints, and let  $\mathbf{x}_0 + \alpha \mathbf{d}$  lie in both  $H_1$  and  $H_2$ , for all  $\alpha \geq 0$ , where  $\mathbf{d} \in \mathbb{R}^n$  and  $H_1, H_2$  are two defining hyperplanes of  $\text{Fea}(P)$ . Does it imply that  $\mathbf{x}_0 + \alpha \mathbf{d}$  also lies in  $H_3$  for all  $\alpha \geq 0$ , where  $H_1, H_2, H_3$  are LD?
8. If  $\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3, \mathbf{x}_4 \in \mathbb{R}^2$ , such that no two of them lie on the same line passing through the origin, then sketch the regions given by:
- (a)  $W_0 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda \mathbf{x}_2, \lambda \in \mathbb{R}\}$ .
  - (b)  $W_1 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda \mathbf{x}_2, \lambda \geq 0\}$ .
  - (c)  $W_2 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda \mathbf{x}_2, 0 \leq \lambda \leq 1\}$ .
  - (d)  $S = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \mathbf{x}_1 + \lambda \mathbf{x}_2, \lambda \in \mathbb{R}\}$ .
  - (e)  $S_{-2} = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \mathbf{x}_1 + \lambda \mathbf{x}_2, \lambda \geq 0\}$ .
  - (f)  $S_{-1} = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \mathbf{x}_1 + \lambda \mathbf{x}_2, 0 \leq \lambda \leq 1\}$ .
  - (g)  $S_0 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2, \lambda_1, \lambda_2 \in \mathbb{R}\}$ .
  - (h)  $S_1 = \{\mathbf{x} : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2, \lambda_1, \lambda_2 \geq 0, \lambda_1 + \lambda_2 = 1\}$ .
  - (i)  $S_2 = \{\mathbf{x} : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2, \lambda_1, \lambda_2 \geq 0, \lambda_1 + \lambda_2 \leq 1\}$ .
  - (j)  $S_3 = \{\mathbf{x} : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2, \lambda_1, \lambda_2 \geq 0\}$ .
  - (k)  $S_4 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \frac{1}{2} \mathbf{x}_3, \lambda_1, \lambda_2 \geq 0, \lambda_1 + \lambda_2 = \frac{1}{2}\}$ .
  - (l)  $S_5 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \frac{1}{2} \mathbf{x}_3, \lambda_1, \lambda_2 \geq 0\}$ .
  - (m)  $S_6 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \frac{1}{2} \mathbf{x}_3, \lambda_1, \lambda_2 \in \mathbb{R}\}$ .
  - (n)  $S_7 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \lambda_3 \mathbf{x}_3, \lambda_1, \lambda_2, \lambda_3 \in \mathbb{R}\}$ .
  - (o)  $S_8 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \lambda_3 \mathbf{x}_3, \lambda_1, \lambda_2, \lambda_3 \geq 0\}$ .
  - (p)  $T_{-1} = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \lambda_3 \mathbf{x}_3, \lambda_1, \lambda_2, \lambda_3 \geq 0, \lambda_1 + \lambda_2 + \lambda_3 \leq 1\}$ .
  - (q)  $T = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \lambda_3 \mathbf{x}_3, \lambda_1, \lambda_2, \lambda_3 \geq 0, \lambda_1 + \lambda_2 + \lambda_3 = 1\}$ .
  - (r)  $T_1 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \alpha(\lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \lambda_3 \mathbf{x}_3) + (1 - \alpha) \mathbf{x}_4, 0 \leq \alpha \leq 1, \lambda_1, \lambda_2, \lambda_3 \geq 0, \lambda_1 + \lambda_2 + \lambda_3 = 1\}$ .
  - (s)  $T_2 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \lambda_3 \mathbf{x}_3 + \lambda_4 \mathbf{x}_4, \lambda_1, \lambda_2, \lambda_3, \lambda_4 \geq 0, \lambda_1 + \lambda_2 + \lambda_3 + \lambda_4 = 1\}$ .
  - (t)  $T_3 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \alpha(\lambda_1 \mathbf{x}_1 + (1 - \lambda_1) \mathbf{x}_2) + (1 - \alpha)(\lambda_2 \mathbf{x}_3 + (1 - \lambda_2) \mathbf{x}_4), 0 \leq \alpha \leq 1, \lambda_1, \lambda_2 \geq 0\}$ .
  - (u)  $T_4 = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \alpha(\lambda_1 \mathbf{x}_1 + (1 - \lambda_1) \mathbf{x}_2) + (1 - \alpha)(\lambda_2 \mathbf{x}_3 + (1 - \lambda_2) \mathbf{x}_4), 0 \leq \alpha \leq 1, 0 \leq \lambda_1, \lambda_2 \leq 1\}$ .

Are the sets  $T_1, T_2, T_3$  equal?

9. Given  $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k \in \mathbb{R}^n$  and  $\lambda_1, \lambda_2, \dots, \lambda_k \geq 0$  with  $\lambda_1 + \lambda_2 + \dots + \lambda_k = 1$ ,  $\lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2 + \dots + \lambda_k \mathbf{x}_k$  is called a **convex combination** of  $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k$ . Show that if  $S \subseteq \mathbb{R}^n$  is a convex set and  $\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3 \in S$ , then any **convex combination** of  $\mathbf{x}_1, \mathbf{x}_2, \mathbf{x}_3$  belongs to  $S$ . In general show that if  $S \subseteq \mathbb{R}^n$  is a convex set, and  $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k \in S$ , ( $k \in \mathbb{N}$ ), then any convex combination of  $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_k$  belongs to  $S$ . (**Hint:** Use induction.)

10. (a) Give an example of a convex set say  $S \subset \mathbb{R}^2$  such that  $\mathbf{x}_1, \mathbf{x}_2 \in S$ , but  $T = \{\mathbf{x} \in \mathbb{R}^2 : \mathbf{x} = \lambda_1 \mathbf{x}_1 + \lambda_2 \mathbf{x}_2, 0 \leq \lambda_1, \lambda_2 \leq 1, \lambda_1 + \lambda_2 \leq 1\}$  is not a subset of  $S$ .  
 (b) Let  $S \subset \mathbb{R}^2$  be a convex set such that  $\mathbf{0} \in S$ , and let  $\mathbf{x}_1, \mathbf{x}_2$  and  $T$  be as defined in part(a), then is  $T = S$ ?
11. If possible give an example of a LPP in three variables such that  $Fea(LPP) \subset \mathbb{R}^3$  has exactly 3 corner points (extreme points).
12. If possible give an example of a LPP in three variables such that  $Fea(LPP) \subset \mathbb{R}^3$  has exactly 5 corner points (extreme points).
13. If possible give an example of a LPP in three variables and three constraints (other than the nonnegativity constraints) such that  $Fea(LPP) \subset \mathbb{R}^3$  has exactly 8 corner points (extreme points).
14. If possible find a nonempty polyhedral subset  $S$  of  $\mathbb{R}^3$ , such that  $S$  has three **LI** defining hyperplanes (not necessarily the nonnegativity constraints) but does not have an extreme point.
15. Consider the following problem:

$$\begin{aligned}
 &\text{Max} \quad -x_1 + 2x_2 \\
 &\text{subject to} \quad x_1 + x_2 \geq 2 \\
 &\quad \quad \quad -2x_1 + x_2 \geq 2 \\
 &\quad \quad \quad x_2 \geq 3 \\
 &\quad \quad \quad x_1 \geq 0, x_2 \geq 0.
 \end{aligned}$$

- (a) Check that the above problem does not have an optimal solution.
- (b) Check that if  $\mathbf{x}_0 \in Fea(LPP)$  then  $\mathbf{x}_0 + \alpha[0, 1]^T \in Fea(LPP)$  for all  $\alpha > 0$ , and  $\mathbf{c}^T \mathbf{x}_0 < \mathbf{c}^T (\mathbf{x}_0 + \alpha[0, 1]^T)$  and for all  $\alpha > 0$ .
- (c) Can you give one more vector other than the vector  $[0, 1]^T$ , say  $\mathbf{d}_1$  which satisfies the same conditions as  $[0, 1]^T$  (as given in part (b)) at  $\mathbf{x}_0$ ?
- (d) Give an example of a set of two linearly dependent hyperplanes say  $H_1, H_2$  defining the feasible region and an  $\mathbf{x}_0$  such that  $\mathbf{x}_0 \in H_1$  but  $\mathbf{x}_0$  does not belong to  $H_2$ .
- (e) Find a point in  $Fea(LPP)$  which does not lie in any of the hyperplanes defining the feasible region.
- (f) Find a point in  $Fea(LPP)$  which lies in exactly one of the hyperplanes defining the feasible region.
- (g) If possible give an example of an  $\mathbf{x}_1 \in \mathbb{R}^2$  such that  $\mathbf{x}_1$  lies in exactly one linearly independent hyperplanes defining  $Fea(LPP)$  but does not belong to  $Fea(LPP)$ .
- (h) If possible give an example of an  $\mathbf{x}_1 \in \mathbb{R}^2$  such that  $\mathbf{x}_1$  lies in two linearly independent hyperplanes defining  $Fea(LPP)$  but is not a corner point of  $Fea(LPP)$ .

- (i) By changing the objective function **only** in the above LPP, give a LPP which has a unique optimal solution. Give the optimal solution and the optimal value in this case.
- (j) By changing the objective function **only** in the above LPP, give a LPP which has infinitely many optimal solutions. How many optimal solutions are also extreme points of Fea(LPP)? Give the optimal solution and the optimal value in this case.
- (k) By changing the second constraint **only** in the above LPP, give a LPP which always has an optimal solution for any objective function. Then find the optimal solution and the optimal value for the given objective function.
16. If the feasible region for a linear programming problem (P) given below as  

$$\text{Min } \mathbf{c}^T \mathbf{x}$$
subject to  $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$ ,  
is unbounded, then is it always possible to choose a  $\mathbf{c}$  (feasible region remaining same as given in (P)) such that the LPP does not have an optimal solution? Also is it possible to choose a  $\mathbf{c}$  (feasible region remaining same as (P)) such that the LPP has an optimal solution?
17. If  $\text{Min } \mathbf{c}^T \mathbf{x}_{2 \times 1}$   
subject to  $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$ ,  
and  $\text{Max } \mathbf{c}^T \mathbf{x}_{2 \times 1}$   
subject to  $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$ ,  
both have optimal solution/s then can the common feasible region of the above problems be unbounded?
18. If  $\text{Min } \mathbf{c}^T \mathbf{x}_{2 \times 1}$   
subject to  $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$ ,  
and  $\text{Max } (\mathbf{c}^T + [1, 2])\mathbf{x}_{2 \times 1}$   
subject to  $A\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$ ,  
both have optimal solution/s then can the common feasible region of the above problems be unbounded?
19. Let the linear programming problem (LPP)  

$$\text{Min } \mathbf{c}^T \mathbf{x}$$
subject to  $A_{3 \times 2} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}$ ,  
has  $\mathbf{x}_0$  as an optimal solution.
- (a) If  $\mathbf{b}' = \mathbf{b} + [1, 2, 0]^T$  ( $[1, 2, 0]^T$  is a column vector with components 1, 2 and 0) then will the above LPP with  $\mathbf{b}$  replaced by  $\mathbf{b}'$  again have an optimal solution?
- (b) If  $\mathbf{b}' = \mathbf{b} + [-1, 0, 0]^T$  ( $[-1, 0, 0]^T$  is a column vector with components -1, 0 and 0) and the above LPP with  $\mathbf{b}$  replaced by  $\mathbf{b}'$  (everything else remaining same) again has  $\mathbf{x}_0$  as a feasible solution then will  $\mathbf{x}_0$  also be an optimal solution of the changed LPP?
20. Consider the following problem:

$$\text{Max } x_1 + x_2$$

subject to  $x_1 + x_2 - x_3 \leq 2$

$$2x_1 + 3x_2 - 2x_3 \geq 2$$

$$3x_1 + 4x_2 - 3x_3 \geq 0$$

$$x_1 \geq 0, x_2 \geq 0, x_3 \geq 0.$$

- (a) Check that the hyper planes corresponding to the first three constraints are linearly dependent.
- (b) Check that  $\tilde{x} = [1, 1, 1]^T$  is a feasible solution of the LPP.
- (c) Does  $\tilde{x} = [1, 1, 1]^T$  lies in any of the hyperplanes defining  $Fea(LPP)$ ?
- (d) Find a  $\mathbf{d} \neq \mathbf{0}$  such that  $\tilde{x} + \alpha\mathbf{d}, \tilde{x} - \alpha\mathbf{d} \in Fea(LPP)$ , for all  $\alpha > 0$ , sufficiently small, where  $\tilde{x}$  is as given in part(b).
- (e) Is this  $\mathbf{d}$  necessarily a direction of  $Fea(LPP)$ ?
- (f) Find a  $\mathbf{d}_0 \neq \mathbf{0}$  such that  $\mathbf{y}_0 = \tilde{x} - \alpha_0\mathbf{d}_0$ , where  $\alpha_0 = \max\{\alpha > 0 : \tilde{x} - \alpha\mathbf{d}_0 \in Fea(LPP)\}$ , lies in exactly one LI hyperplane defining  $Fea(LPP)$ .
- (g) If possible find a  $\mathbf{d}_1 \neq \mathbf{0}$  such that  $\mathbf{y}_1 = \tilde{x} - \alpha_1\mathbf{d}_1$ , where  $\alpha_1 = \max\{\alpha > 0 : \tilde{x} - \alpha\mathbf{d}_1 \in Fea(LPP)\}$ , lies in exactly two LI hyperplane defining  $Fea(LPP)$ .
- (h) If possible find a  $\mathbf{d}_2 \neq \mathbf{0}$  such that  $\mathbf{y}_2 = \tilde{x} - \alpha_2\mathbf{d}_2$ , where  $\alpha_2 = \max\{\alpha > 0 : \tilde{x} - \alpha\mathbf{d}_2 \in Fea(LPP)\}$ , lies in exactly three LI hyperplane defining  $Fea(LPP)$ .
- (i) Find an extreme (or corner point) point of  $Fea(LPP)$ .
- (j) How many extreme points ( corner points) does this feasible region have?
- (k) Does this problem have an optimal solution?

21. **The Transportation Problem.** Let there be  $m$  supply stations  $S_1, \dots, S_m$  for a particular product and  $n$  destination stations  $D_1, D_2, \dots, D_n$  where the product is to be transported. Let  $c_{ij}$  be the cost of unit amount of the product from  $S_i$  to  $D_j$ . Let  $s_i$  be the available amount of the product at  $S_i$  and let  $d_j$  be the demand at  $D_j$ . The problem is to find  $x_{ij}$ ,  $i = 1, 2, \dots, m$ ,  $j = 1, 2, \dots, n$ , where  $x_{ij}$  is the amount of the product to be transported from  $S_i$  to  $D_j$  such that the cost of transportation is minimum. The problem is given by

$$\text{Min } \sum_{i,j} c_{ij}x_{ij}$$

$$\text{subject to } \sum_{j=1}^n x_{ij} \leq s_i, i = 1, 2, \dots, m$$

$$\sum_{i=1}^m x_{ij} \geq d_j, j = 1, 2, \dots, n, x_{ij} \geq 0 \text{ for } i = 1, 2, \dots, m, j = 1, 2, \dots, n.$$

Consider a transportation problem with 2 supply stations and 5 destinations such the maximum capacity of the two supply stations  $S_1$  and  $S_2$  are given by 30 units and 40 units respectively. The demands at each of the 5 destinations are 10 units each.

- (a) Obtain a feasible solution for this problem. Without trying to calculate the optimal solution, justify that this problem will always have an optimal solution.
- (b) By changing **only** the demands at the five supply stations of the above problem, find a transportation problem with no feasible solution.
- (c) Hence find a necessary and sufficient condition for the transportation problem with  $m$  supply stations and  $n$  destination stations to have an optimal solution.

22. Consider the following problem (LPP):

$$\begin{aligned}
 &\text{Max} \quad -x_1 + 2x_2 + x_3 \\
 &\text{subject to} \quad x_1 + 2x_2 \geq 1 \\
 &\quad \quad \quad -2x_1 + x_2 + x_3 \leq 5 \\
 &\quad \quad \quad -x_1 + x_2 \leq 2 \\
 &\quad \quad \quad x_1 \geq 0, x_2 \geq 0, x_3 \geq 0.
 \end{aligned}$$

- (a) Find all the extreme points of  $\text{Fea}(\text{LPP})$ .
- (b) Find the set of all extreme directions of  $\text{Fea}(\text{LPP})$ .
- (c) Does the above problem has an optimal solution?
- (d) If the above problem is changed to a minimization problem ( everything else remaining same ) then will the new problem have an optimal solution?
- (e) Let  $\mathbf{d}_1$  be an extreme direction and let  $H_1, H_2$  be a set of two LI hyperplanes (defining  $\text{Fea}(\text{LPP})$ ) the normals of which are orthogonal to  $\mathbf{d}_1$ . Then if possible find an extreme point which lies on both  $H_1$  and  $H_2$ .
- (f) Let  $\mathbf{x}_1$  be an extreme point of  $\text{Fea}(\text{LPP})$  (you can choose any one), then if possible find another extreme point  $\mathbf{x}_2$ , such that  $\mathbf{x}_1$  and  $\mathbf{x}_2$  lie on a common set of  $(n-1)$  ( here  $n = 3$  ) LI hyperplanes defining  $\text{Fea}(\text{LPP})$  (such extreme points are called adjacent extreme points).
- (g) Let  $\mathbf{x}_1$  be an extreme point of  $\text{Fea}(\text{LPP})$  (you can choose any one), then if possible find an  $\mathbf{x}_2 \in \text{Fea}(\text{LPP})$ , such that  $\mathbf{x}_2$  is not an extreme point and  $\mathbf{x}_1$  and  $\mathbf{x}_2$  lie on a common set of  $(n-1)$  ( here  $n = 3$  ) LI hyperplanes defining  $\text{Fea}(\text{LPP})$ .
- (h) Let  $\mathbf{x} \in \text{Fea}(\text{LPP})$  be such that it lies on **only** one LI hyperplane (defining  $\text{Fea}(\text{LPP})$ ) given by  $-x_1 + x_2 = 2$ , then give  $\mathbf{d}_1, \mathbf{d}_2$  LI and an  $\alpha > 0$  such that each of  $\mathbf{x} + \alpha\mathbf{d}_1, \mathbf{x} - \alpha\mathbf{d}_1, \mathbf{x} + \alpha\mathbf{d}_2, \mathbf{x} - \alpha\mathbf{d}_2$  belongs to  $\text{Fea}(\text{LPP})$  (that is starting from  $\mathbf{x}$  you can move along the positive and negative direction of these two LI vectors for some distance and still be inside  $\text{Fea}(\text{LPP})$ ).
- (i) Let  $\mathbf{x}_0 \in \text{Fea}(\text{LPP})$  be such that it lies on **only** two LI hyperplanes (defining  $\text{Fea}(\text{LPP})$ ) given by  $-x_1 + x_2 = 2$  and  $x_1 = 0$ , then give  $\mathbf{d} \neq \mathbf{0}$  and an  $\alpha > 0$  such that each of  $\mathbf{x}_0 + \alpha\mathbf{d}, \mathbf{x}_0 - \alpha\mathbf{d}$  belongs to  $\text{Fea}(\text{LPP})$ . Can you find  $\mathbf{d}_1, \mathbf{d}_2$  LI and an  $\alpha > 0$  such that each of  $\mathbf{x}_0 + \alpha\mathbf{d}_1, \mathbf{x}_0 - \alpha\mathbf{d}_1, \mathbf{x}_0 + \alpha\mathbf{d}_2, \mathbf{x}_0 - \alpha\mathbf{d}_2$  belongs to  $\text{Fea}(\text{LPP})$  ?
- (j) Given  $\mathbf{x} = [1, 1, 0]^t$ , check that  $\mathbf{x} \in \text{Fea}(\text{LPP})$  and it lies on only one hyperplane (defining  $\text{Fea}(\text{LPP})$ ) given by  $x_3 = 0$ .
  - i. Check that  $\mathbf{d} = [1, 5, 0]^t$  is orthogonal to the normal of the hyperplane  $x_3 = 0$ . Find  $\gamma$  such that  $\mathbf{y}_1 = \mathbf{x} + \gamma\mathbf{d} \in \text{Fea}(\text{LPP})$  but  $\mathbf{x} + \beta\mathbf{d}$  does not belong to  $\text{Fea}(\text{LPP})$  for  $\beta > \gamma$ . Write the point  $\mathbf{y}_1$ . Check that this  $\mathbf{y}_1$  lies on two LI hyperplanes (defining  $\text{Fea}(\text{P})$ ) one of which is  $x_3 = 0$ .

- ii. Now find a  $\mathbf{d}_1$ , such that by starting from  $\mathbf{y}_1$  and moving along  $\mathbf{d}_1$  gives an extreme point which lies on each of the two defining hyperplanes on which  $\mathbf{y}_1$  lies.
  - i. Check that  $\mathbf{d} = [1, 2, 0]^t$  is orthogonal to the normal of the hyperplane  $x_3 = 0$ . Find  $\gamma$  such that  $\mathbf{y}_1 = \mathbf{x} + \gamma\mathbf{d} \in \text{Fea}(LPP)$  but  $\mathbf{x} + \beta\mathbf{d}$  does not belong to  $\text{Fea}(LPP)$  for  $\beta > \gamma$ . Write the point  $\mathbf{y}_1$ . Check that this  $\mathbf{y}_1$  lies on two LI hyperplanes (defining  $\text{Fea}(P)$ ) one of which is  $x_3 = 0$ .
  - ii. Now find a  $\mathbf{d}_1$ , such that by starting from  $\mathbf{y}_1$  and moving along  $\mathbf{d}_1$  gives an extreme point which lies on each of the two defining hyperplanes on which  $\mathbf{y}_1$  lies.
- (k) Check that  $[1, 5, 0]^t, [1, 2, 0]^t$  are LI. Can you give 5 more vectors such that if you start from  $\mathbf{x}$  (where  $\mathbf{x}$  is as defined in the previous part) and move along the positive direction of these vectors then you will be able to apply the same process as that in the previous part, to get an extreme point?
- (l) By changing **only** the third constraint (such that the first three constraints are distinct constraints) in the above problem (LPP) if possible give another problem LPP' and  $\text{Fea}(LPP')$ , such that  $\text{Fea}(LPP')$  has an extreme point which lies on 4 hyperplanes (defining  $\text{Fea}(LPP')$ ).
- (m) By changing **only** the second constraint (such that the first three constraints are distinct constraints) in the above problem (LPP) if possible give another problem LPP' and  $\text{Fea}(LPP')$ , such that  $\text{Fea}(LPP')$  has an element which lies on 3 hyperplanes (defining  $\text{Fea}(LPP')$ ) and is not an extreme point of  $\text{Fea}(LPP')$ .

23. Let  $\text{Fea}(LPP)$  of a LPP be given by  $\text{Fea}(LPP) = \{\mathbf{x} \in \mathbb{R}^n : A_{m \times n}\mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}\}$ .

- (a) Let  $\mathbf{d}$  be an extreme direction of  $\text{Fea}(LPP)$  (that is  $\mathbf{d}$  is a direction and is orthogonal to the normals of  $(n - 1)$  LI hyperplanes defining the feasible region  $\text{Fea}(LPP)$ ) then does there always exist an extreme point  $\mathbf{x}$  of  $\text{Fea}(LPP)$  such that  $\mathbf{x}$  lies on the same  $(n - 1)$  LI hyperplanes, the normals to which are orthogonal to  $\mathbf{d}$ ?

**Hint:** WLOG let  $\tilde{\mathbf{a}}_1, \dots, \tilde{\mathbf{a}}_{n-1}$  be the normals to the  $(n - 1)$  LI hyperplanes to which  $\mathbf{d}$  is orthogonal. Let  $\mathbf{x}$  be such that  $\tilde{\mathbf{a}}_i\mathbf{x} = \tilde{b}_i$  for  $i = 1, 2, \dots, (n - 1)$  and in addition  $\tilde{\mathbf{a}}_k\mathbf{x} = \tilde{b}_k$ , for some  $(n - 1) < k \leq (m + n)$ . If this  $\mathbf{x} \in \text{Fea}(LPP)$  then done, if not then atleast one of  $\tilde{\mathbf{a}}_i\mathbf{x} > \tilde{b}_i$  for  $i \neq k, i > (n - 1)$ . But since  $\mathbf{d}$  is a direction so  $\tilde{\mathbf{a}}_i\mathbf{d} \leq 0$  for all  $i = 1, \dots, (m + n)$ . Hence you can choose  $\alpha > 0$  suitably large such that  $\tilde{\mathbf{a}}_i(\mathbf{x} + \alpha\mathbf{d}) = \tilde{b}_i$  for all  $i = 1, 2, \dots, (n - 1)$  and  $\tilde{\mathbf{a}}_i(\mathbf{x} + \alpha\mathbf{d}) \leq \tilde{b}_i$  for all  $i > (n - 1)$ . Let  $\gamma = \min\{\alpha > 0 : \mathbf{x} + \alpha\mathbf{d} \in \text{Fea}(LPP)\}$ . Check whether  $\mathbf{x} + \gamma\mathbf{d}$  is the required extreme point.

- (b) Let  $\mathbf{x}$  be an extreme point of  $\text{Fea}(LPP)$ , which has atleast two extreme points. Then does there always another extreme point  $\mathbf{y}$  of  $\text{Fea}(LPP)$  such that  $\mathbf{x}$  and  $\mathbf{y}$  are adjacent extreme points (that is both  $\mathbf{x}$  and  $\mathbf{y}$  lie on a common set of  $(n - 1)$  LI hyperplanes defining  $\text{Fea}(LPP)$ )?

**Hint:** Starting from  $\mathbf{x}$  move along the positive and negative direction of  $\mathbf{d} \neq \mathbf{0}$ ,



where  $\mathbf{d}$  is orthogonal to  $(n - 1)$  LI hyperplanes on which  $\mathbf{x}$  lies, always remaining inside the feasible region. You cannot move indefinitely in both directions (because of nonnegativity constraints). Consider the point beyond which if you move, you will fall out of  $Fea(LPP)$ .

24. Consider the problem (P) given by :

$$\text{Max } \mathbf{c}^T \mathbf{x}$$

$$\text{subject to } A_{5 \times 2} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq \mathbf{0}.$$

- (a) Check whether the above problem (P) with  $\mathbf{c} = [-1, 0]^T$  and  $\mathbf{b} = [0, 3, 1, 1, 0]^T$  has an optimal solution.
- (b) If the above problem is changed to a minimization problem, then find a  $\mathbf{c} \neq \mathbf{0}$  and a  $\mathbf{b} \neq \mathbf{0}$  such that the changed problem has an optimal solution (no matter what the matrix  $A$  is).
- (c) Given  $\mathbf{b} = [-1, -2, 0, 0, 0]^T$  suggest entries of  $\bar{\mathbf{a}}_1 \neq \mathbf{0}$  the first column of  $A$ , such that the feasible region of the above problem (P) is nonempty (no matter what the other columns of  $A$  are).
- (d) If  $\bar{\mathbf{a}}_1$ , the first column of  $A$  is given as  $[-21, -1, 0, -1, -1]^T$ , then is  $Fea(P)$  nonempty? If no justify. If yes, then suggest an extreme direction of  $Fea(P)$ .
- (e) Suggest entries of  $\bar{\mathbf{a}}_1$ , the first column of  $A$  such that the feasible region of the above problem (P) is unbounded (no matter what the other columns of  $A$  are).
- (f) If the first row of  $A$  is given by  $[1, 2]$  and if  $Fea(P) \neq \phi$ , then will the above problem always (no matter what the other rows of  $A$  are) have an optimal solution ?

All the parts in the above question are independent.