

Academic Year: 2020-21.

19CS2109A - Computer Networks + Security

Set-3 Key + Scheme

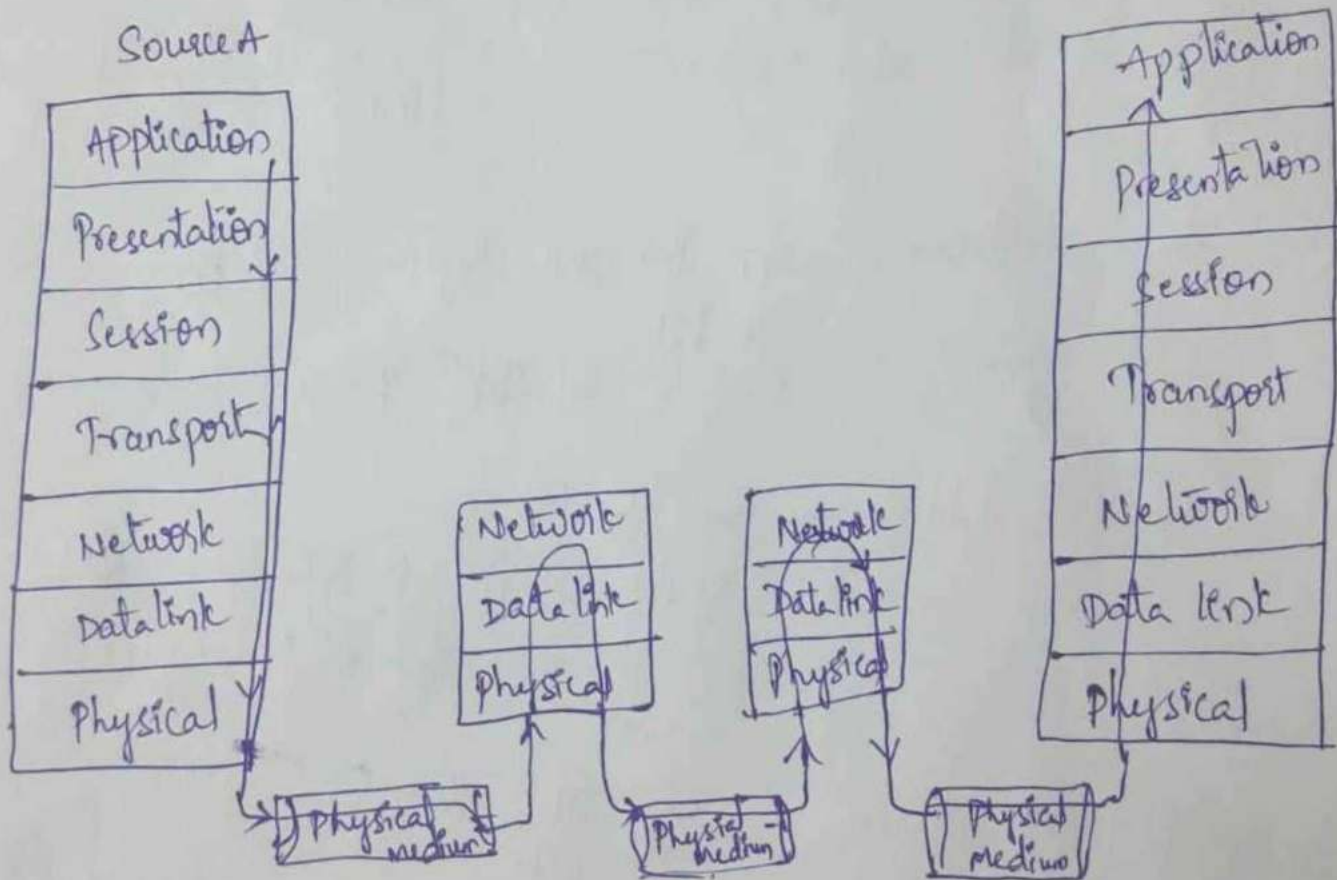
Max Marks: 50m.

- ① Given source = A
Destination = B

Diagram (2.5 M)
Explanation (2 M)

Intermediate nodes = X + Y

How the data flow happens is as follows.



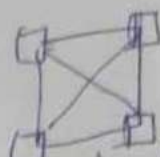
Flow of data through various layers
in various nodes.

② List the number of cables required for connecting a set of nodes when the nodes follow

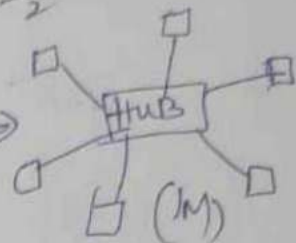
(a) Mesh - Diagram + explanation $n=4$ (1.5M)

$$\frac{n(n-1)}{2} \text{ where } n = \text{nodes}$$

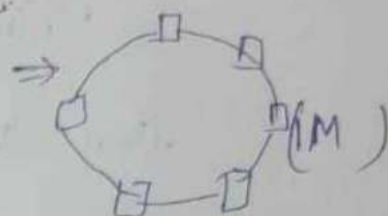
$$\therefore \frac{4(4-1)}{2} = 6$$



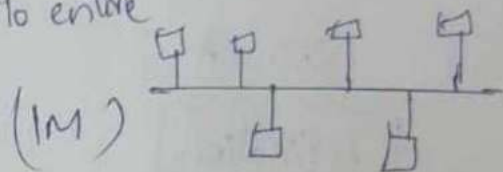
(b) Star - Single hub at center
Diagram with explanation \Rightarrow



(c) Ring - 2 nodes connecting cable



(d) Bus - Backbone failure leads to entire n/w failure



Port Address - In transport layer (2M)

16 bit

Used to identify process in a node

Logical Address - In n/w layer

32 bit if IPv4 is followed (2M)

Used to identify node uniquely in Internet

Physical Address - In DLL layer

48 bit

Identify node within the network (2M)

Data Link Layer Responsibility \Rightarrow Framing, Error control, Flow control

Network layer Responsibility \Rightarrow Routing, congestion, IP addressing

⑦ Simple Parity

→ Used for error detection alone.

— No. of parity bits are less in number
(2M)

2D parity / 2 dimensional Parity

— Used for Error detection + also for Single bit error correction

— No. of parity bits are more in number than simple parity.
(2M)

Single Bit Error

Is when only one bit in the data unit has changed.

(2M)

0 0 0 0 1 0 1 0

Received

0 changed to 1

0 0 0 0 0 0 1 0

sent

Burst Errors

— Means that two or more consecutive bits in the data unit have changed.

(2M)

⑤ Go Back N + selective Repeat

- ① In Go Back N protocol if the sent frame are found Suspected then all the frames are retransmitted from the lost pkt.
- ② Less Complex
- ③ Ack type is Cumulative
- ④ Order of pkts is ensured
- ⑤ More retransmissions
- ① In selective Repeat, only those frames are retransmitted which are found Suspected.
- ② More complex
- ③ Ack type is individual
- ④ No order ^{for} of packet arriving.
- ⑤ Reduces retransmissions unnecessarily.

$$(5 \times 2 = 10M)$$

— In stop & wait protocol sender sends the packet & waits for the ACK of the packet. Once the ACK reaches the sender, it transmits the next packet in row. If the ACK is not received, it retransmits the previous pkt again. Maximum no. of sequence numbers are '2' (0 or 1). $\Rightarrow (2.5M)$

To speed up pkts transmission rate pipelined transmission are adopted in Go back N + selective Repeat.

⑥ (i) Given code word is 221

Convert to binary 0010 0010 0001.

After stuffing 01111110 0010 0010 0001 01111110

* Each character can be converted into 3 bits/8 bits instead of 4 bits as above.

* 221 can be converted (two hundred + twenty one) into binary + then perform bit stuffing (4M)

(ii) $D_p = 0110111111001111011111110000$.

Flag = 01111110

Stuffed data

(4M)

01111110 011011111101100111110011110111110000 01111110

(iii) Flag - DLE

Data - IDLE

Stuffed data DLE IDLE DLE DLE (4.5)

Disadv: Wastage of Bandwidth by more unnecessary stuffed data transmission.

⑦ Given 4 bits for representing sequence number

⇒ 1.5M

Possible no. of sequence numbers are $2^4 = 16$.

Sequence numbers are 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16.

— Sequence numbers need Explanation → 3M.

③ pipelined transmission.

-(4.5M)

③

- Frames will be transmitted one by one
- Before receiving the acknowledgement of previous frame, next frame can be transmitted.

Eg:- Go-Back N protocol } explanation
Selective Repeat protocol }

⑦ In Go back N, if we don't receive Ack of a pkt whole window (lost pkt) is sent again.

Assume window size = 3.

Initially, window have 1, 2, 3. As Ack of '1' receives window slides to 4. 4 is transmitted. Now when 4th packet ack received 7th pkt is sent & when 5th packet ack is received 8th packet is sent.

Now Ack of '6' is not received so window of '6' i.e. 6, 7, 8 packets are retransmitted. Now 6th pkt from there is 9, so 9, 10 will be retransmitted.

∴ These are the serial transmissions of packets:

1, 2, 3, 4, 5, 6, 7, 8, 6, 7, 8, 9, 10, 9, 10

Total no. of transmissions = 15.

Go back N Explanation (5M)

(10) Pure Aloha + Slotted Aloha Explanation (2M.)
Derivation — (6M)

Pure Aloha: K — station's number
 N — New frames generation meantime
 $0 < N < 1$

If $N > 1 \Rightarrow$ overloading the channel.

Then $G \Rightarrow$ Mean for frames (new) + Retransmission frames

$N=0$ then $G \approx N$ S-throughput

$N > 1$ then $G > N$.

$S = G * P_0$ * P_0 is the probability of frames transmitted without collision.

$$P_r(K) = \frac{G^K \cdot e^{-G}}{K!} \Rightarrow P_0 = \frac{G^0 \cdot e^{-G}}{0!} \Rightarrow P_0 = \frac{1 \cdot e^{-G}}{1!}$$

$P_0 = e^{-G}$ — substitute in $S = G P_0 \Rightarrow S = G \cdot e^{-G}$.

Vulnerable time — $2T_r$ $\therefore G$ becomes $2G$.

$$S = G \cdot e^{-2G} \Rightarrow S = \frac{1}{2} \cdot e^{-1} \quad [G=0.5] \text{ for pure Aloha}$$

$$\Rightarrow S = \frac{1}{2} \cdot e^{-1} \Rightarrow \frac{1}{2e} \Rightarrow S = 0.183 \Rightarrow 18\%$$

Slotted Aloha

Vulnerable time is T_r $\therefore S = G \cdot e^{-G}$.

$G=1$ for Slotted Aloha

$$S = 1 \cdot e^{-1} \Rightarrow S = \frac{1}{e} \Rightarrow 0.378 \Rightarrow 38\%$$

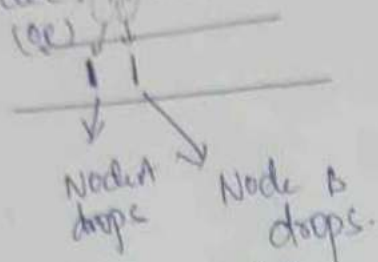
\therefore slotted Aloha gives more throughput than Pure Aloha.

(11) (a)

Node A: 01110

Node B: 10011

Node C: 11111



\therefore winning station is node-C
 \therefore channel is conquered by 'C'.

Explanation - 3M

Example problem - 3.5M.

Disadvantage - 6M.

Disadv of Bit map over binary countdown -

- Floats contention slots + asks all the stations to reserve the slots, allocated the channel as per the reservations.

Eg:- If a station received data from its up layer after completing the contention slots. That station has and floats contention slots again.

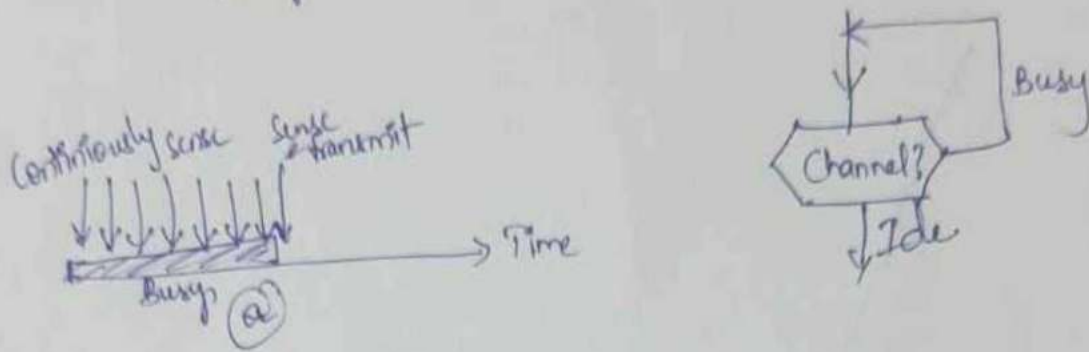
\therefore waiting time may be more for few stations sometimes.

- In Binary countdown a station waits only until any one other station completes its transmission waiting time may be decreased.

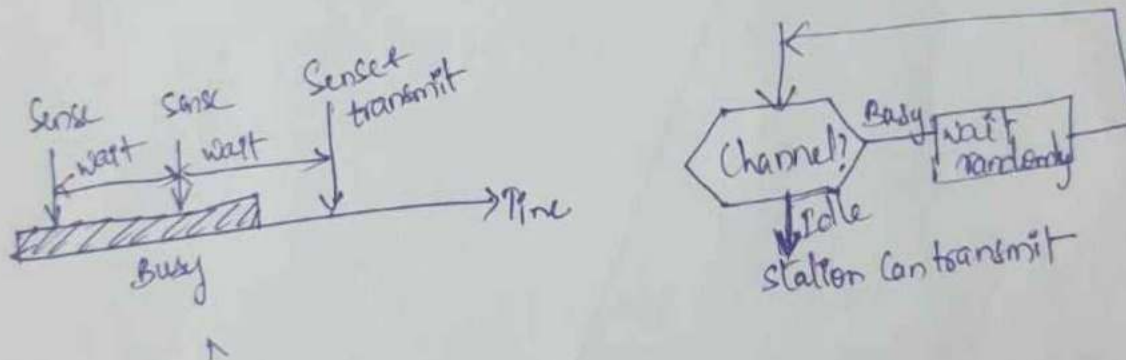
(12) Persistent methods — 1 persistent
P Persistent
N-persistent.

1-Persistent Explanation + flowchart - 3/8/11 (4M)

Diagram with explanation



Non-Persistent Explanation + flowchart - (4M)



p-Persistent Explanation + flowchart - (4.5M)

