

## Wasserstein

The Wasserstein distance is defined as:

$$D_w \left[ \Pr(x) \| q(x) \right] = \min_{\mathbf{P}} \left[ \sum_{i,j} P_{ij} \cdot |i - j| \right], \tag{15.10}$$

subject to the constraints that:

$$\sum_{j} P_{ij} = \Pr(x = i) \quad \text{(initial distribution of } \Pr(x)\text{)},$$

$$\sum_{i} P_{ij} = q(x = j) \quad \text{(initial distribution of } q(x)\text{)},$$

$$P_{ij} \geq 0 \quad \text{(non-negative masses)}.$$
(15.11)

In other words, the Wasserstein distance is the solution to a constrained minimization problem that maps the mass of one distribution to the other. This is inconvenient as we must solve this minimization problem over the elements  $P_{ij}$  every time we want to compute the distance.

To rewrite the problem in matrix form, let:

$$\mathbf{p} = \text{vec}(\mathbf{P}), \quad \mathbf{c} = \text{vec}(\mathbf{C}),$$

where **p** and **c** are vectorized forms of **P** and **C**. Define the constraint matrix **A** and vector **b** such that:

$$Ap = b$$

where:

$$\mathbf{b} = \begin{bmatrix} \Pr(x=1) \\ \Pr(x=2) \\ \vdots \\ q(x=1) \\ q(x=2) \\ \vdots \end{bmatrix} \quad \mathbf{p} = \begin{bmatrix} P_{11} \\ P_{12} \\ \vdots \\ P_{1K} \\ P_{21} \\ P_{22} \\ \vdots \\ P_{2K} \\ \vdots \\ P_{K1} \\ P_{K2} \\ \vdots \\ P_{KK} \end{bmatrix} \quad \text{vec}(\mathbf{C}) = \begin{bmatrix} C_{11} \\ C_{12} \\ \vdots \\ C_{1K} \\ C_{21} \\ C_{22} \\ \vdots \\ C_{2K} \\ \vdots \\ C_{K1} \\ C_{K2} \\ \vdots \\ C_{KK} \end{bmatrix}.$$

The matrix **A** encodes the row and column constraints. It has dimensions  $(2K) \times K^2$ , where the first K rows correspond to the row constraints, and the next K rows correspond to the column constraints.

• Row Marginal Constraints  $(\sum_{j} P_{ij} = \Pr(x = i))$ :

$$\mathbf{A}_{\text{rows}} = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 0 & 0 & 0 & \cdots & 1 \end{bmatrix}.$$

 $\bullet$  Column Marginal Constraints  $(\sum_i P_{ij} = q(x=j))$  :

$$\mathbf{A}_{\text{columns}} = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 0 & \cdots & 1 \\ 0 & 1 & 0 & 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \end{bmatrix}.$$

Combine these into:

$$\mathbf{A} = egin{bmatrix} \mathbf{A}_{ ext{rows}} \ \mathbf{A}_{ ext{columns}} \end{bmatrix}.$$

The primal problem becomes:

$$\min_{\mathbf{p}} \mathbf{c}^{\top} \mathbf{p},$$

subject to 
$$\mathbf{Ap} = \mathbf{b}$$
,  $\mathbf{p} \ge 0$ .

## Lagrangian Formulation

Introduce the Lagrange multipliers f for the equality constraint Ap = b. The Lagrangian is:

$$\min_{\mathbf{p}} \max_{\mathbf{f}} L(\mathbf{p}, \mathbf{f}) = \mathbf{c}^{\top} \mathbf{p} + \mathbf{f}^{\top} (\mathbf{b} - \mathbf{A} \mathbf{p}).$$

Simplify the Lagrangian:

$$\min_{\mathbf{p}} \max_{\mathbf{f}} L(\mathbf{p}, \mathbf{f}) = \mathbf{f}^{\top} \mathbf{b} + \mathbf{p}^{\top} (\mathbf{c} - \mathbf{A}^{\top} \mathbf{f}).$$

To ensure L is finite as we are minimizing w.r.t p, we require:

$$\mathbf{c} - \mathbf{A}^{\top} \mathbf{f} \ge 0.$$

## **Dual Problem**

Minimizing  $L(\mathbf{p}, \mathbf{f})$  with respect to  $\mathbf{p}$  gives:

$$L(\mathbf{p}, \mathbf{f}) = \mathbf{f}^{\top} \mathbf{b}$$
, subject to  $\mathbf{c} - \mathbf{A}^{\top} \mathbf{f} \ge 0$ .

Thus:

Primal Form	Dual Form
minimize $\mathbf{c}^T \mathbf{p}$ , such that $\mathbf{A}\mathbf{p} = \mathbf{b}$ , $\mathbf{p} \ge 0$	$\begin{vmatrix} \text{maximize } \mathbf{b}^T \mathbf{f}, \\ \text{such that } \mathbf{A}^T \mathbf{f} \leq \mathbf{c} \end{vmatrix}$

## Interpreting the Dual Problem

Returning to the original Wasserstein problem, the dual variables  $\mathbf{f}$  correspond to potentials  $f_i$  and  $f_j$  for the bins i and j. The dual constraint  $\mathbf{A}^{\top}\mathbf{f} \leq \mathbf{c}$  implies:

$$f_i - f_j \le C_{ij}, \quad \forall i, j.$$

The dual objective becomes:

$$D_w \left[ \Pr(x) \| q(x) \right] = \max_{\mathbf{f}} \left[ \sum_{i} \Pr(x = i) f_i - \sum_{j} q(x = j) f_j \right]$$

where  $f_i$  and  $f_j$  satisfy the constraints