Boyce Codd normal form (BCNF)

- o BCNF is the advance version of 3NF. It is stricter than 3NF.
- \circ A table is in BCNF if every functional dependency $X \to Y$, X is the super key of the table.
- o For BCNF, the table should be in 3NF, and for every FD, LHS is super key.

Example: Let's assume there is a company where employees work in more than one department.

EMPLOYEE table:

EMP_ID	EMP_COUNTRY	EMP_DEPT	DEPT_TYPE	EMP_DEPT_NO
264	India	Designing	D394	283
264	India	Testing	D394	300
364	UK	Stores	D283	232
364	UK	Developing	D283	549

In the above table Functional dependencies are as follows:

- 1. EMP ID \rightarrow EMP COUNTRY
- 2. $EMP_DEPT \rightarrow \{DEPT_TYPE, EMP_DEPT_NO\}$

Candidate key: {EMP-ID, EMP-DEPT}

The table is not in BCNF because neither EMP DEPT nor EMP ID alone are keys.

To convert the given table into BCNF, we decompose it into three tables:

EMP COUNTRY table:

EMP_ID	EMP_COUNTRY
264	India
264	India

EMP_DEPT table:

EMP_DEPT	DEPT_TYPE	EMP_DEPT_NO
Designing	D394	283
Testing	D394	300
Stores	D283	232
Developing	D283	549

EMP_DEPT_MAPPING table:

EMP_ID	EMP_DEPT
D394	283
D394	300
D283	232
D283	549

Functional dependencies:

- 1. $EMP_ID \rightarrow EMP_COUNTRY$
- 2. $EMP_DEPT \rightarrow \{DEPT_TYPE, EMP_DEPT_NO\}$

Candidate keys:

For the first table: EMP_ID

For the second table: EMP_DEPT

For the third table: {EMP_ID, EMP_DEPT}

Now, this is in BCNF because left side part of both the functional dependencies is a

key.