

Computer Architecture and System Programming Laboratory

TA Session 3

EFLAGS

SHIFT

Little Endian

Sections (.bss, .data, .rodata, .text, Heap, Stack)

Stack (push, pop)

C calling convention

gnu debugger – basics

EFLAGS – Flags / Status Register

- **flag** is a single independent bit of (status) information
- each flag has a **two-letter symbol name**

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
O	O	O	O	O	O	O	O	O	O	I D	V I P	V I F	A C	V M	R F	O	N T	I O P L	O F	D F	I F	T F	S F	Z F	O	A F	O	P F	1	C F	

Status Flags

CF = Carry flag

PF = Parity flag

AF = Auxiliary carry flag

ZF = Zero flag

SF = Sign flag

OF = Overflow flag

Control flag

DF = Direction flag

System flags

TF = Trap flag

IF = Interrupt flag

IOPL = I/O privilege level

NT = Nested task

RF = Resume flag

VM = Virtual 8086 mode

AC = Alignment check

VIF = Virtual interrupt flag

VIP = Virtual interrupt pending

ID = ID flag

CF - Carry Flag

CF gets '1' if **result is larger than "capacity" of target operand**

$$11111111_b + 00000001_b = \underbrace{100000000_b}_0 = 0 \rightarrow CF = 1$$

CF gets '1' if **subtract larger number from smaller** (borrow out of "capacity" of target operand)

$$00000000_b - 00000001_b = 11111111_b \rightarrow CF = 1$$

$$\begin{aligned} 00000000_b - 00000001_b &\equiv \\ 100000000_b - 00000001_b &\equiv 11111111_b \end{aligned}$$

```
char x=0b11111111; // x=-1;
unsigned char x=0b11111111; // x=255;
```

unsigned arithmetic $\rightarrow 11111111_b = 255$ (decimal) \rightarrow got wrong answer

signed arithmetic $\rightarrow 11111111_b = -1$ (decimal) \rightarrow ignore CF flag

otherwise CF gets '0'

- **In unsigned arithmetic, watch CF flag to detect errors.**
- In signed arithmetic ignore CF flag.

OF – Overflow Flag

addition of two positive numbers results negative \rightarrow OF gets '1'

$$01000000_b + 01000000_b = 10000000_b \rightarrow \text{OF} = 1$$

addition two negative numbers results positive \rightarrow OF gets '1'

$$10000000_b + 10000000_b = 100000000_b = 0 \rightarrow \text{OF} = 1$$

otherwise OF gets '0'

$$01000000_b + 00010000_b = 01010000_b \rightarrow \text{OF} = 0$$

$$01100000_b + 10010000_b = 11110000_b \rightarrow \text{OF} = 0$$

$$10000000_b + 00010000_b = 10010000_b \rightarrow \text{OF} = 0$$

$$11000000_b + 11000000_b = 10000000_b \rightarrow \text{OF} = 0$$

- **In signed arithmetic, watch OF flag to detect errors.**
- In unsigned arithmetic ignore OF flag.

ZF, SF, PF, and AF Flags

ZF – Zero Flag - set if a result is zero; cleared otherwise

```
mov eax, 0  
add eax, 0 → ZF = 1
```

SF – Sign Flag – set if a result is **negative** (i.e. MSB of the result is 1) ; cleared otherwise

```
00000000b - 00000001b = 11111111b → SF = 1
```

PF – Parity Flag - set if low-order eight bits of result contain even number of 1-bits; cleared otherwise

```
1111111111010000b + 0000000000001000b = 1111111111011000b → 4 bits are '1' → PF = 1)  
1111000011000000b + 0000000000001000b = 1111000011001000b → 3 bits are '1' → PF = 0)
```

AF – Auxiliary Carry Flag (or **Adjust Flag**) is set if there is a carry from low nibble (4 bits) to high nibble or a borrow from a high nibble to low nibble

```
00001111b + 00000001b = 00010000b → AF = 1  
      └─┬─┘      └─┬─┘  
      nibble    nibble
```

```
00010000b - 00000001b = 00001111b → AF = 1  
      └─┬─┘      └─┬─┘  
      nibble    nibble
```

Instructions seen till now – affecting flags

MOV NOT JMP

does not affect flags

AND OR

OF and CF flags are cleared; SF, ZF, and PF flags are set according to the result.

The state of AF flag is undefined.

DEC INC

OF, SF, ZF, AF, and PF flags are set according to the result. CF flag is not affected.

ADD ADC SUB SBB NEG CMP

OF, SF, ZF, AF, PF, and CF flags are set according to the result.

full set of assembly
instructions may
be found [here](#)

ADC - add integers with carry

Example:

adc AX, BX ;(AX = AX+BX+CF)

SBB - subtract with borrow

Example:

sbb AX, BX ;(AX = AX-BX-CF)

Instructions seen till now – affecting flags

MOV NOT JMP

does not affect flags

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ADC - add integers with carry

Example:

adc AX, BX ;(AX = AX+BX+CF)

ADC illustration example:

$$\begin{array}{r} \text{CF}=1 \\ 13 \\ + \quad 9 \\ \hline = 22 \end{array}$$

<instruction> r/m8(16,32) 1/CL/imm8

SHL, SHR – Bitwise Logical Shifts on first operand

- number of bits to shift is given by second operand
- vacated bits are filled with zero
- (last) shifted bit enters to CF flag

shift indeed performs **fast**
division / multiplication by 2

Example:

```
mov al, 10110111b      ; AL = 10110111b
shr al, 1                ; shift right 1 bit → AL = 01011011b, CF = 1
shl al, 4                ; shift left 4 bits → AL = 10110000b, CF = 1
```

SAL, SAR – Bitwise Arithmetic Shift on first operand

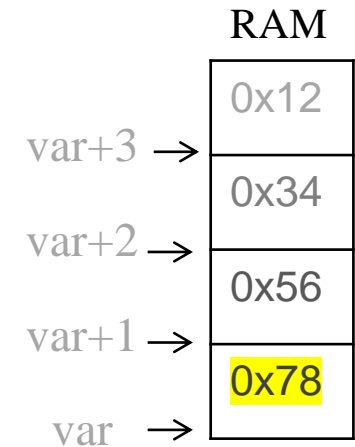
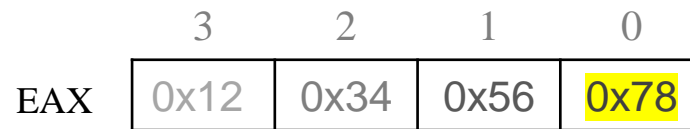
- vacated bits are filled with zero for SAL
- vacated bits are filled with copies of the **original high bit** of the source operand for SAR
- (last) shifted bit enters to CF flag

Example:

```
mov al, 10110111b      ; AL = 10110111b
sar al, 3                ; shift arithmetical right 3 bits → AL = 11110110b, CF = 1
sal al, 2                ; shift (arithmetical) left 2 bits → AL = 11011000b, CF = 1
```


Little Endian - store least significant byte in lowest address

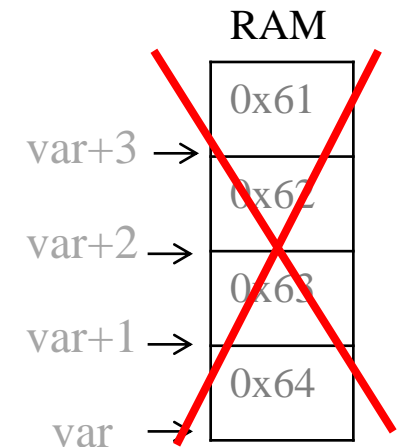
var: dd 0x12345678



var: dd 'abcd'

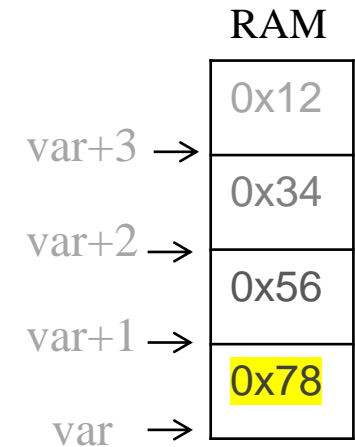
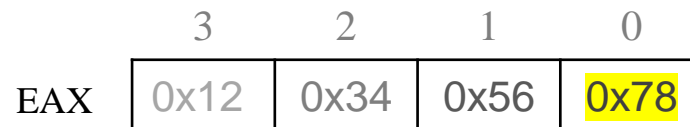


'abcd' = 0x61626364



Little Endian - store least significant byte in lowest address

var: dd 0x12345678

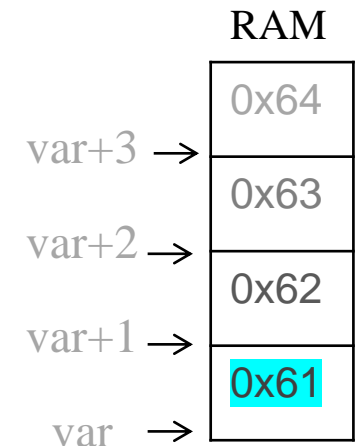
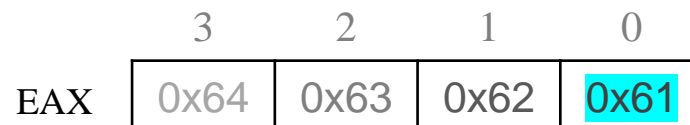


String is translated by NASM into numeric value in reversed order.
Thus, when (numeric value of) string inserted into memory, we get it in source order.

var: dd 'abcd'



var: dd 0x64636261



Sections

Every process consists of sections that are accessible to the process during its execution.

Data

- .bss** - holds uninitialized data
- .data** - holds initialized data
- .rodata** - holds read-only data

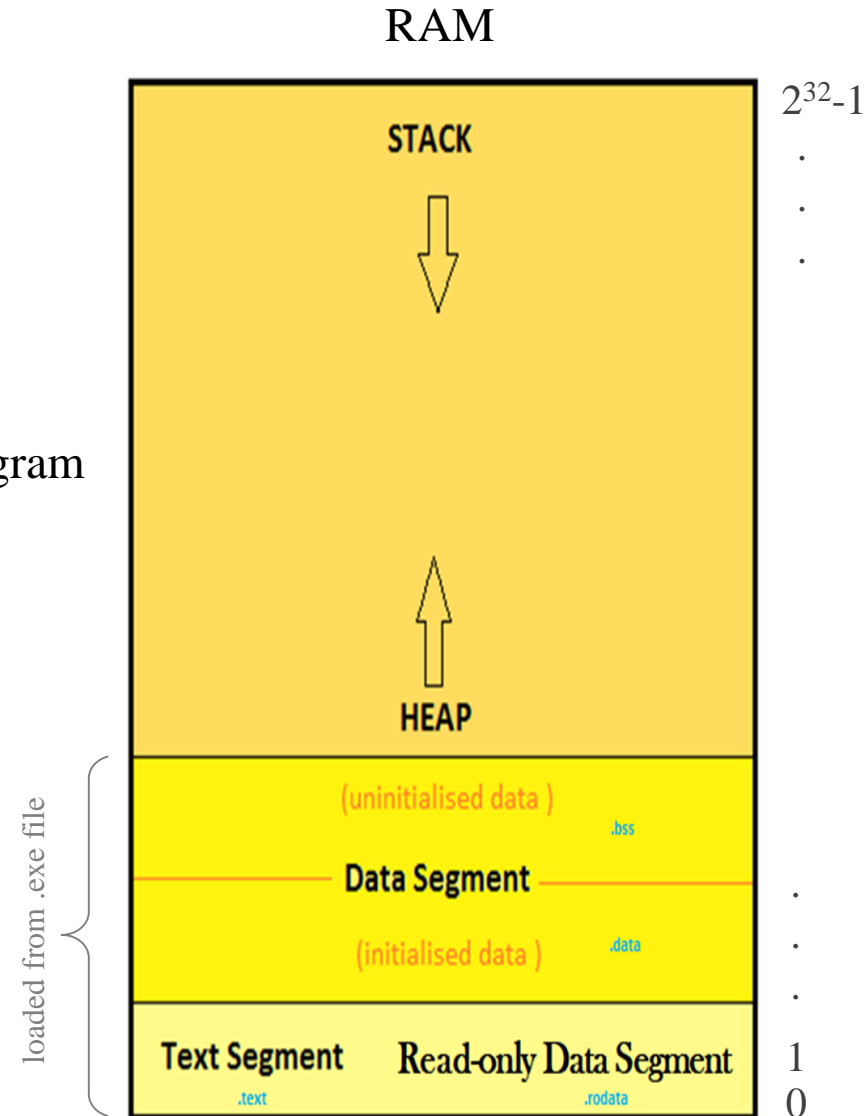
Text

- .text** – holds executable instructions of a program

Stack* – holds functions activation frame

Heap – holds dynamically allocated memory

* Stack pointer is normally initialized to point to the top of stack and proceed downward. In fact, we can point to any location we like and manage any area of memory as stack.



Sections

Every process consists of sections that are accessible to the process during its execution.

Data

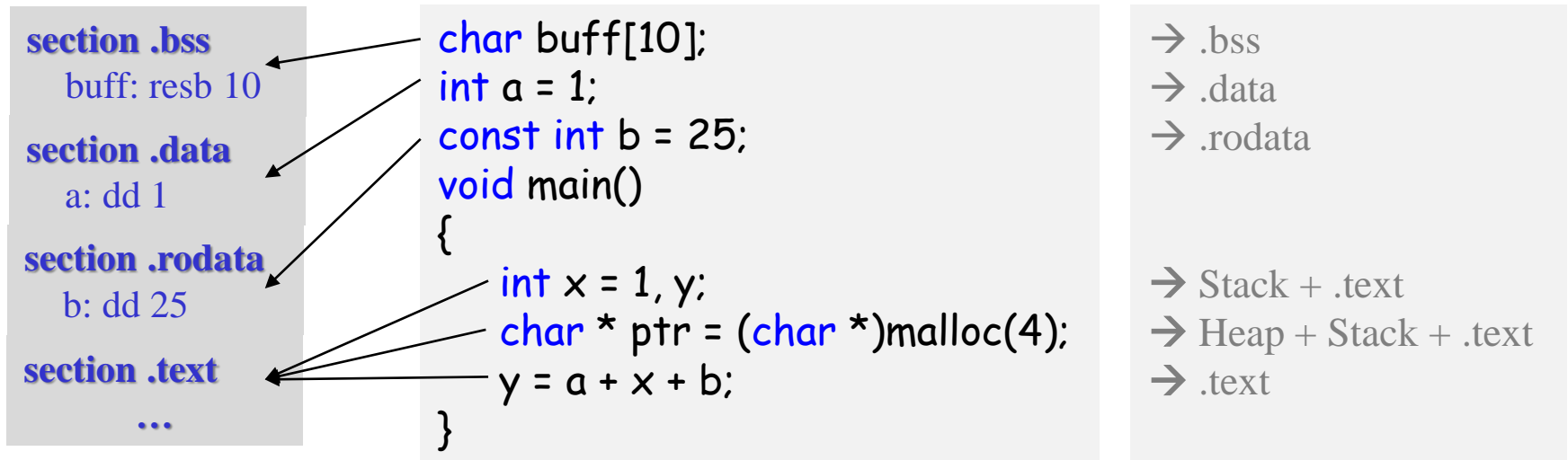
- .bss** - holds uninitialized data
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- .text** – holds executable instructions of a program

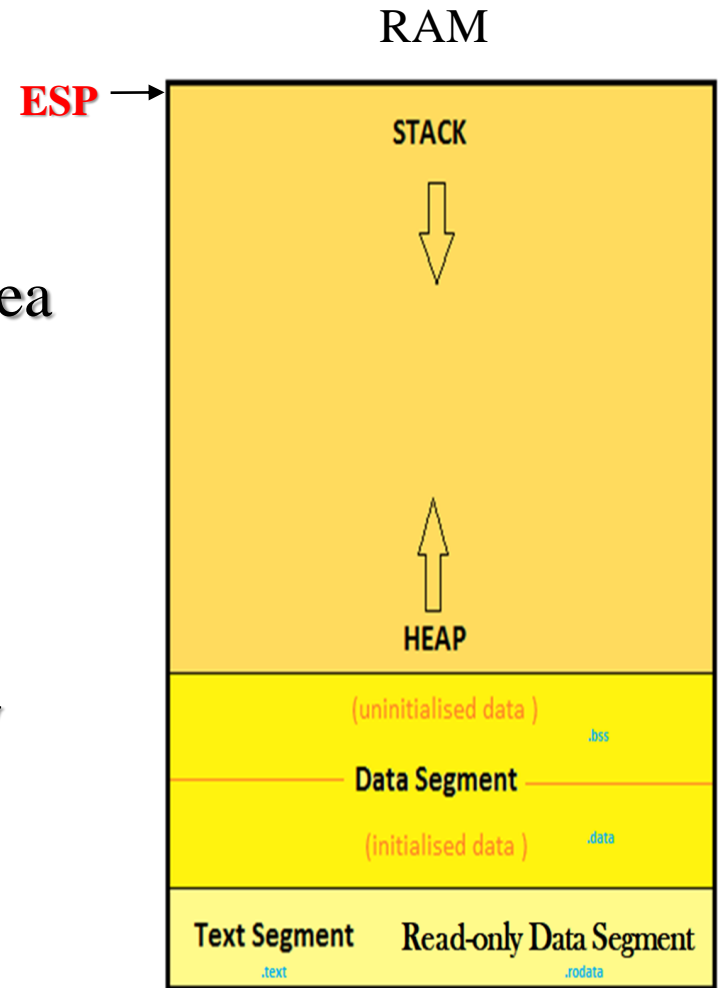
Stack* – holds functions activation frame

Heap – holds dynamically allocated memory



Stack

- Stack is temporary storage memory area
- **ESP** points to the top of Stack
(by default, it is highest RAM address)
- Stack addresses go from high to low



Stack Operations

PUSH - push data on Stack

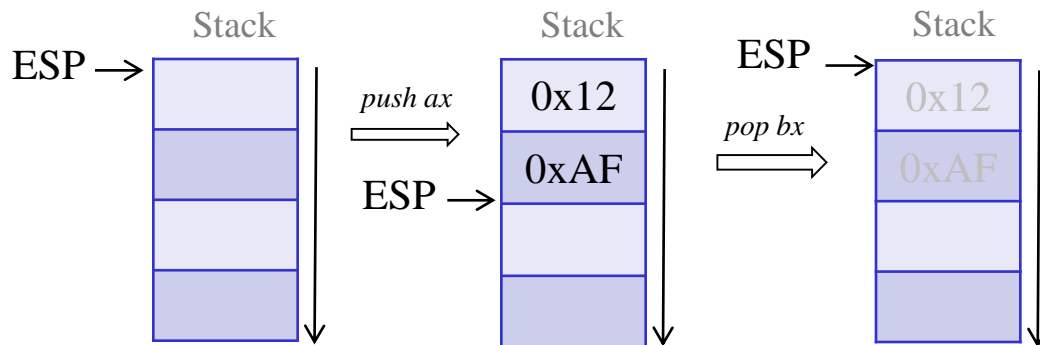
- decrements ESP by 2/4 bytes (according to the operand size)
- stores the operand value at ESP address on Stack (in Little Endian manner)

POP - load a value from Stack

- reads the operand value at ESP address on Stack (in Little Endian manner)
- increment ESP by 2/4 bytes (according to the operand size)

`mov ax, 0x12AF`
`push ax`
`pop bx`

In 32-bit mode, you can push 32 or (with operand-size prefix) 16 bits.



<https://stackoverflow.com/questions/45127993/how-many-bytes-does-the-push-instruction-push-onto-the-stack-when-i-dont-specif>

Stack Operations

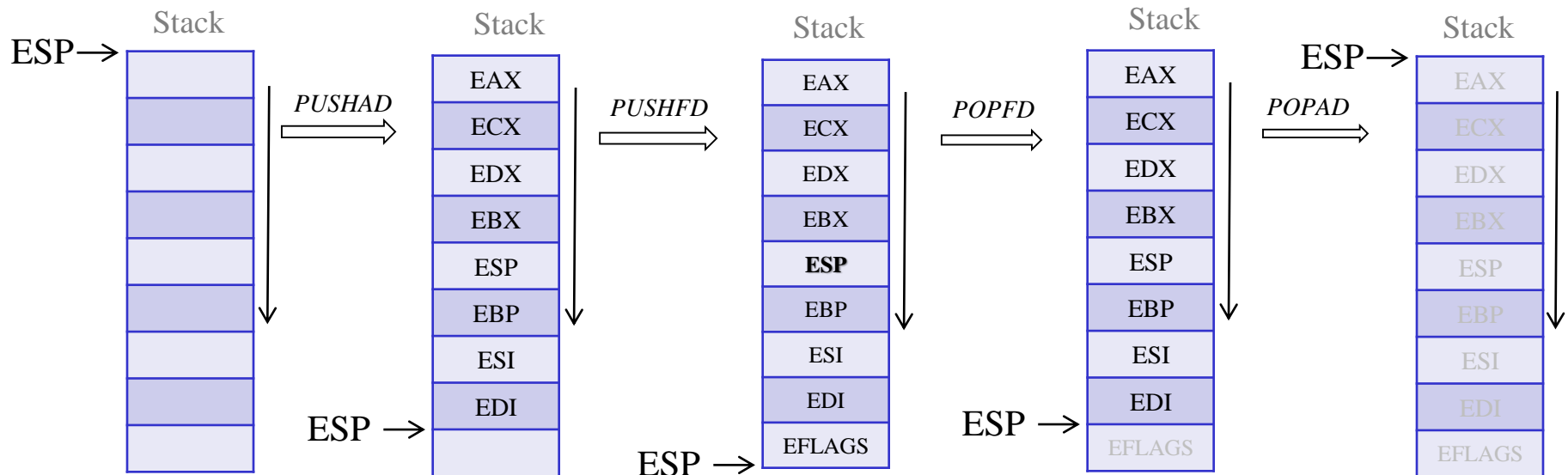
PUSHAD (push all double) - pushes values of EAX, ECX, EDX, EBX, **ESP**, EBP, ESI, and EDI onto Stack

original ESP value before PUSHAD

PUSHFD (push flags double) - push value of EFLAGS onto Stack

POPAD (pop all double) - pop a dword from Stack into each one of (successively) EDI, ESI, EBP, nothing (placeholder for ESP), EBX, EDX, ECX, and EAX

POPFD (pop flags double) - pop a dword from Stack into EFLAGS



C Calling Convention

```
int Func(int x, int y) {  
    int sum;  
    sum = x + y;  
    return sum;  
}  
int result;  
void main() {  
    result = Func(1, 2);  
}
```


C Calling Convention

section .bss

result: resd 1

section .text

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int Func(int x, int y) {  
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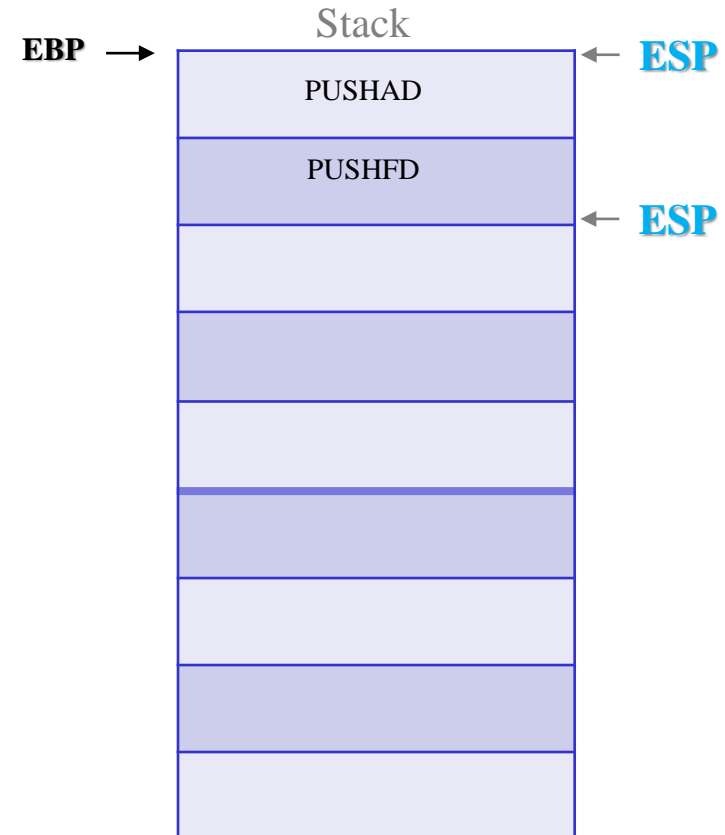
main: ; caller code

pushad ; backup registers
pushfd ; backup EFLAGS

```
int Func(int x, int y) {  
    int sum;  
    sum = x + y;  
    return sum;  
}  
  
int result;  
void main() {  
    result = Func(1, 2);  
}
```

According to C calling convention, the called function is allowed to mess up values of EAX, ECX and EDX registers. EBX, ESI, and EDI registers' values should be preserved and restored by the called function.

But in fact it is not 100% sure that all C code obeys this (may be compiler implementation dependent), so it is a good idea to backup and restore all registers.



C Calling Convention

section .bss

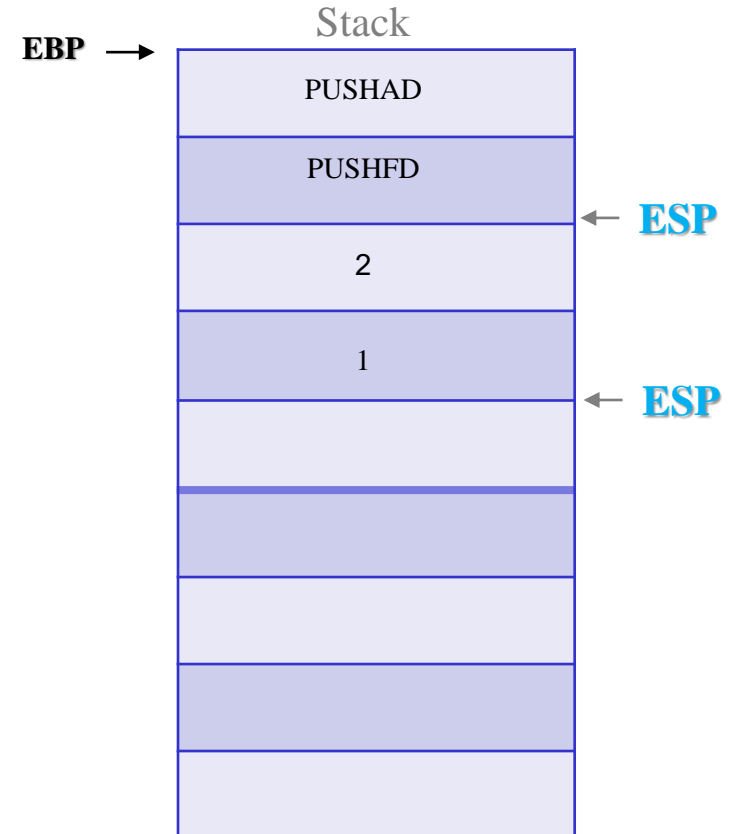
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main: ; caller code

pushad ; backup registers
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push dword 2 ; push second argument
push dword 1 ; push first argument

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main: ; caller code

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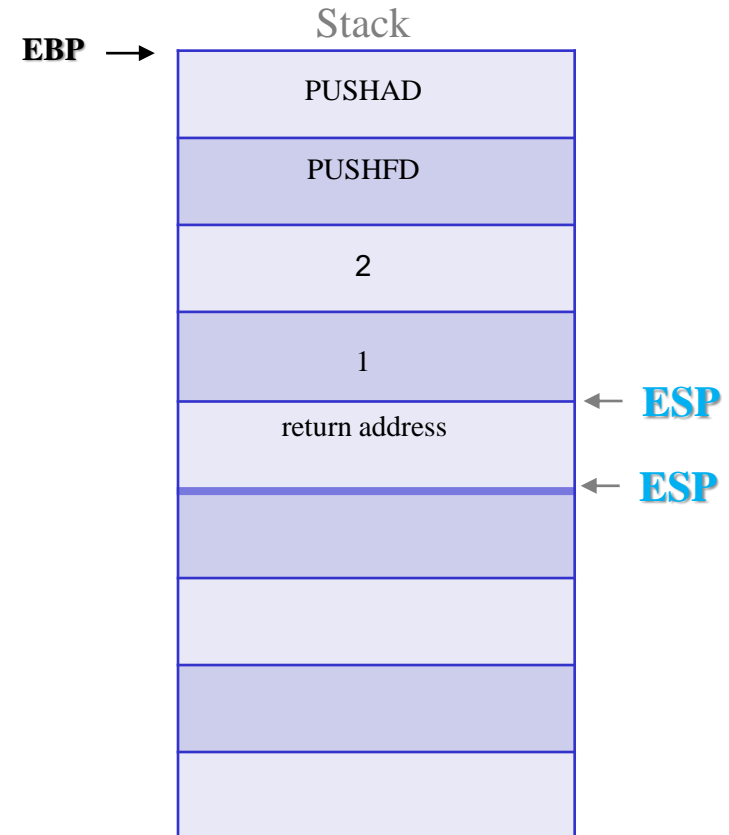
pushfd ; backup EFLAGS

push dword 2 ; push second argument

push dword 1 ; push first argument

call Func ; call the function → **push return address**
; into Stack and jump to function code

```
int Func(int x, int y) {  
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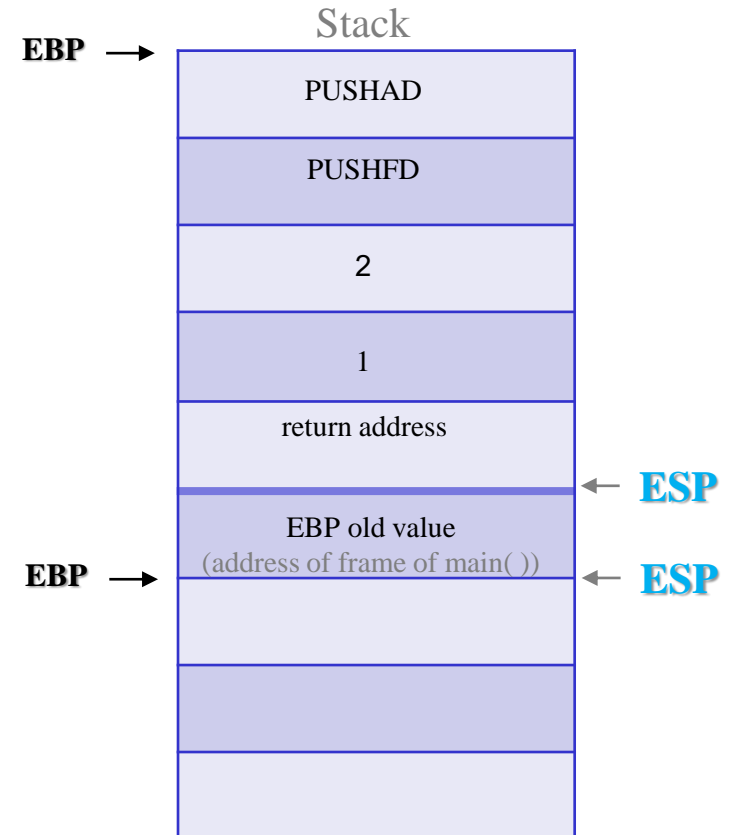
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Func: ; callee code

push ebp ; backup EBP
mov ebp, esp ; reset EBP to the current ESP

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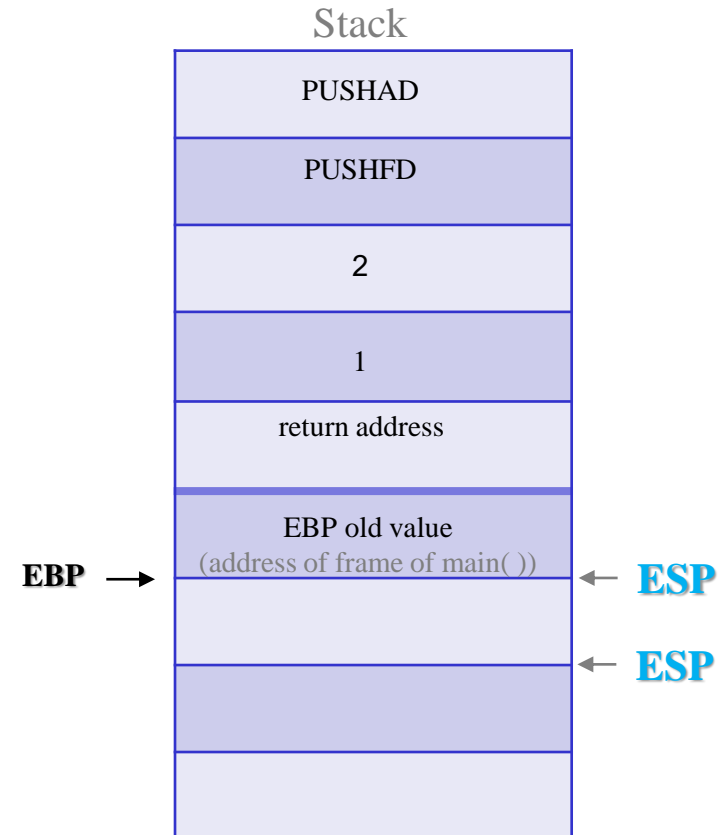
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sub esp, 4 ; allocate space for local variable sum

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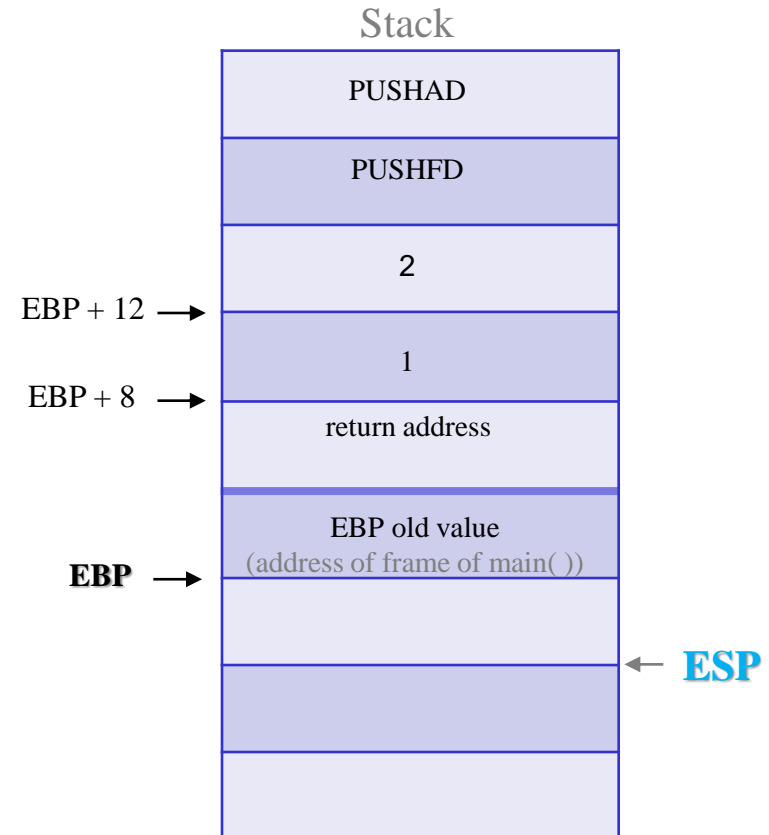
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mov ebx, [ebp+8] ; get first argument
mov ecx, [ebp+12] ; get second argument

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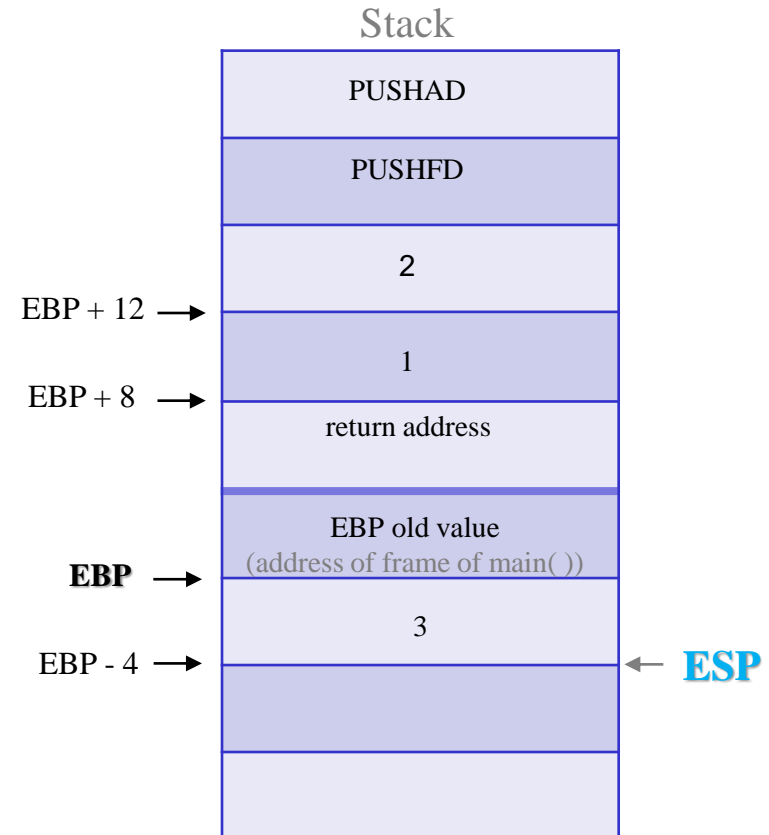
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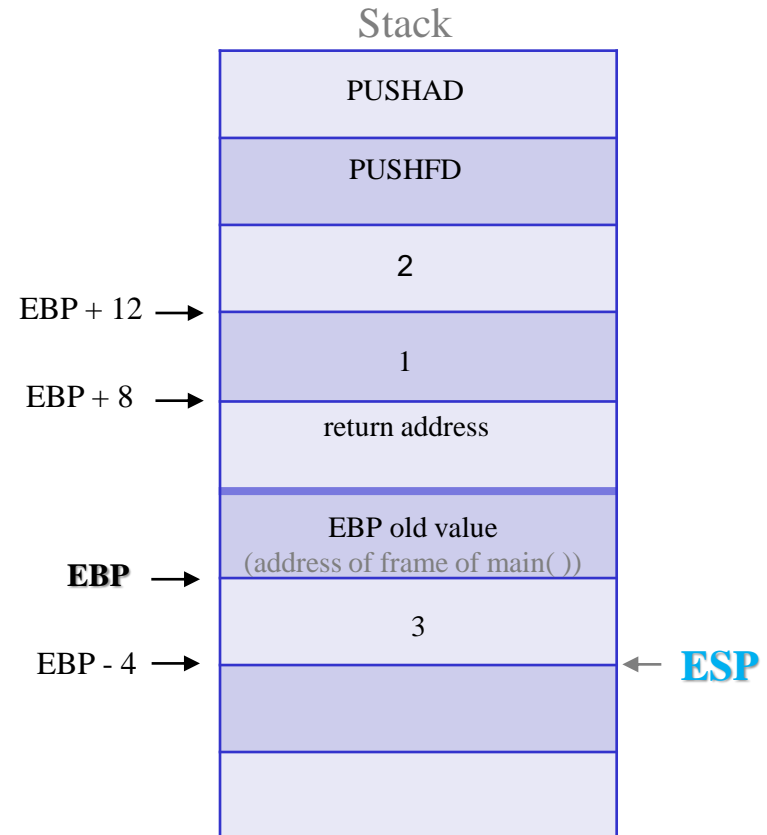
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mov eax, [ebp-4] ; put return value into EAX

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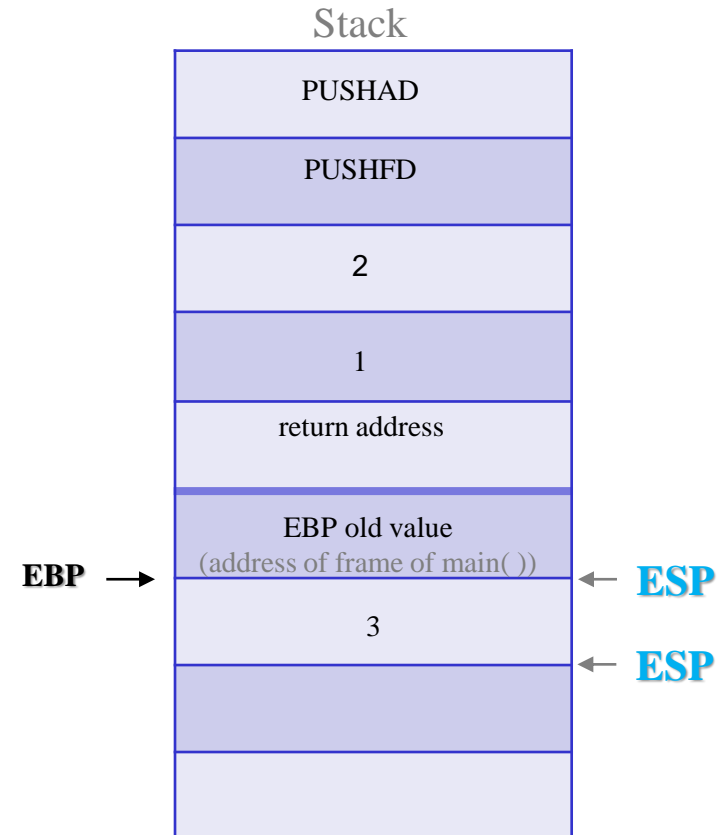
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mov eax, [ebp-4] ; put return value into EAX
mov esp, ebp ; “free” function activation frame

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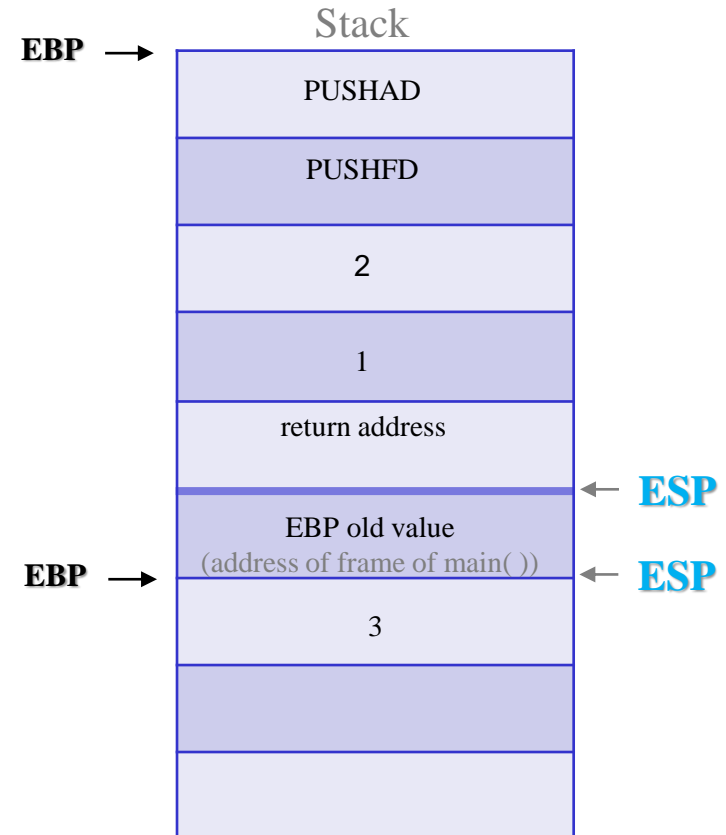
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mov esp, ebp ; “free” function activation frame
pop ebp ; restore activation frame of main()

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int Func(int x, int y) {  
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void main() {  
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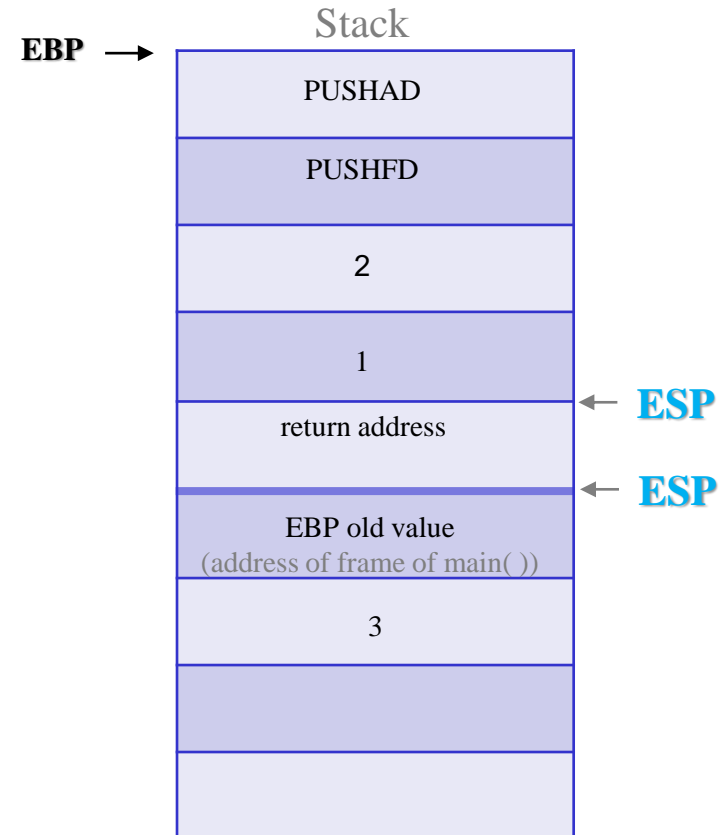
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RET ; return from the function

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int Func(int x, int y) {  
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main: ; caller code

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 pushfd ; backup EFLAGS
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 ; **into Stack and jump to function code**

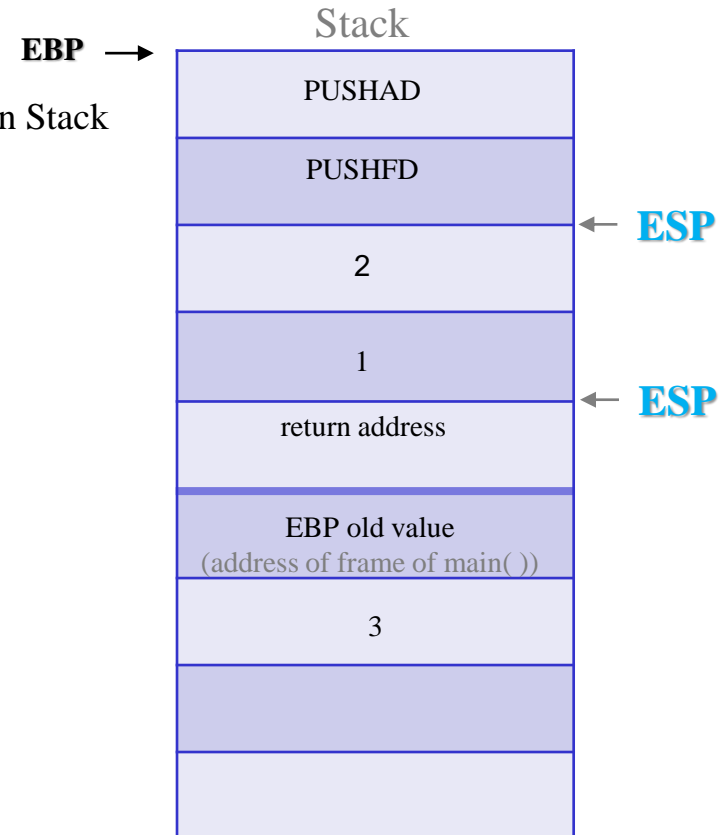
mov [result], eax ; retrieve return value from EAX
 add esp, 8 ; "free" space allocated for function arguments in Stack

Func: ; callee code

push ebp ; backup EBP
 mov ebp, esp ; set EBP to Func activation frame
 sub esp, 4 ; allocate space for local variable sum
 mov ebx, [ebp+8] ; get first argument
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 add ebx, ecx ; calculate x+y
 mov [ebp-4], ebx ; set sum to be x+y
 mov eax, [ebp-4] ; put return value into EAX
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call Func ; call the function → **push return address**
; **into Stack and jump to function code**

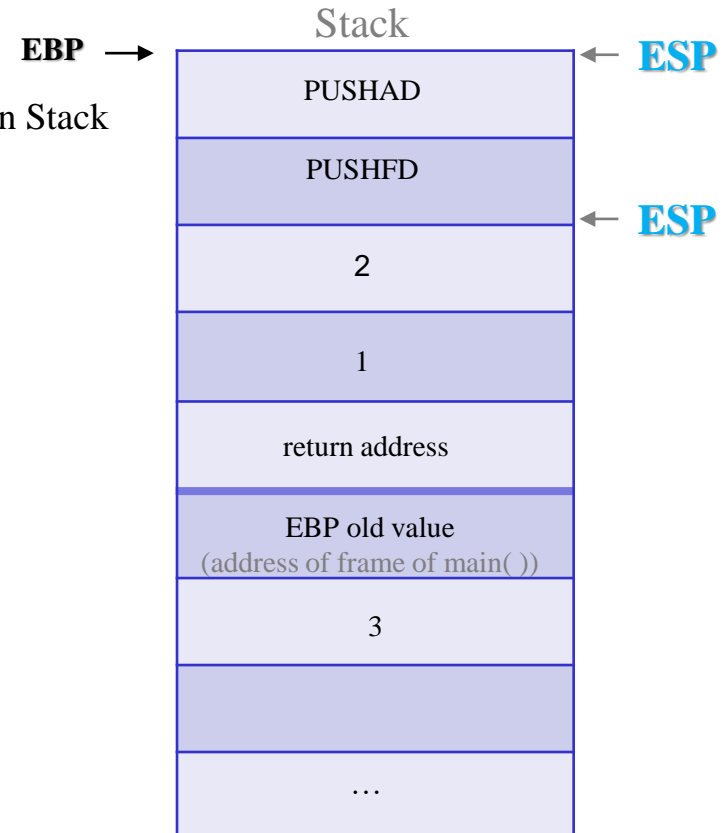
mov [result], eax ; retrieve return value from EAX
add esp, 8 ; "free" space allocated for function arguments in Stack
popfd ; restore flags register
popad ; restore registers

...

Func: ; callee code

push ebp ; backup EBP
mov ebp, esp ; set EBP to Func activation frame
sub esp, 4 ; allocate space for local variable sum
mov ebx, [ebp+8] ; get first argument
mov ecx, [ebp+12] ; get second argument
add ebx, ecx ; calculate x+y
mov [ebp-4], ebx ; set sum to be x+y
mov eax, [ebp-4] ; put return value into EAX
mov esp, ebp ; "free" function activation frame
pop ebp ; restore activation frame of main()
RET ; return from the function

```
int Func(int x, int y) {  
    int sum;  
    sum = x + y;  
    return sum;  
}  
  
int result;  
void main() {  
    result = Func(1, 2);  
}
```



C Calling Convention

section .bss

result: resd 1

section .text

main: ; caller code

pushad ; backup registers
 pushfd ; backup EFLAGS
 push dword 2 ; push second argument
 push dword 1 ; push first argument
 call Func ; call the function → push return address
 ; into Stack and jump to function code

mov [result], eax ; retrieve return value from EAX
 add esp, 8 ; "free" space allocated for function arguments in Stack
 popfd ; restore flags register
 popad ; restore registers

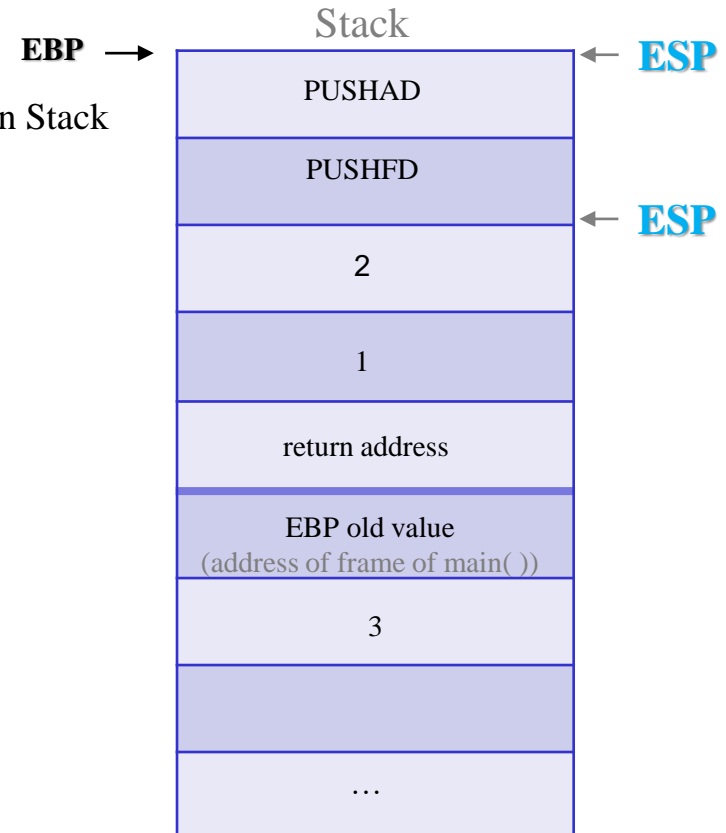
...

Func: ; callee code

push ebp ; backup EBP
 mov ebp, esp ; set EBP to Func activation frame
 sub esp, 4 ; allocate space for local variable sum
 mov ebx, [ebp+8] ; get first argument
 mov ecx, [ebp+12] ; get second argument
 add ebx, ecx ; calculate x+y
 mov [ebp-4], ebx ; set sum to be x+y
 mov eax, [ebp-4] ; put return value into EAX
 mov esp, ebp ; "free" function activation frame
 pop ebp ; restore activation frame of main()
 RET ; return from the function

```
int Func(int x, int y) {
    int sum;
    sum = x + y;
    return sum;
}

int result;
void main() {
    result = Func(1, 2);
}
```



gdb debugger – very basic usage

- ❑ run Gdb from the console by typing:

gdb executableFileName

- ❑ add breaking points by typing:

break label

- ❑ start debugging by typing:

run parameters (argv)

(gdb) set disassembly-flavor intel — change presentation of assembly-language instructions from the default Motorola conventions, that are used by gcc, to the Intel conventions that are used by nasm, that is, from opcode source, dest to opcode dest, src

(gdb) layout asm — this will display the assembly language

(gdb) layout regs – this will display registers

- **s/si** – one step forward
- **c** – continue to run the code until the next break point.
- **q** – quit gdb
- **p/x \$eax** – prints the value in eax
- **x \$esp** – prints esp value (address) and value (dword) that is stored in this address. It is possible to use label instead of esp.
Type **x** again will print the next dword in memory.