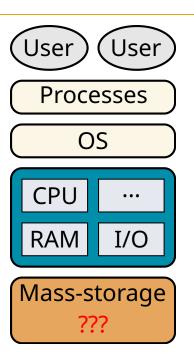
Filesystem Abstraction

Introduction

Long-term data storage requirements

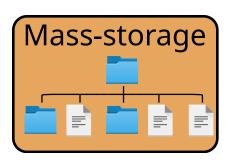
- (Very) large amounts of non-volatile data
- Easy way to find data
- Concurrent access from processes
- Controlled sharing between users
- Performance
- Reliability
 - o Survive power-off, OS crash, process termination



Filesystems (FS)

Abstractions for maximizing the usage of non-volatile storage

- Persistent, named data (files)
- Hierarchical organization (directories, subdirectories)
- Performance
- Access control
- Crash and storage error tolerance



File

- Abstraction that provides persistent, described data
- Logical storage unit

Metadata

- File attributes added and managed by the OS
- Size, creation/modification time, owner, permissions, etc.
- File's content location (index structure)

Data

- What a user puts in the file
- Array of untyped bytes

Metadata

- size: 42 KiB
- created: 1970-01-01

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Data

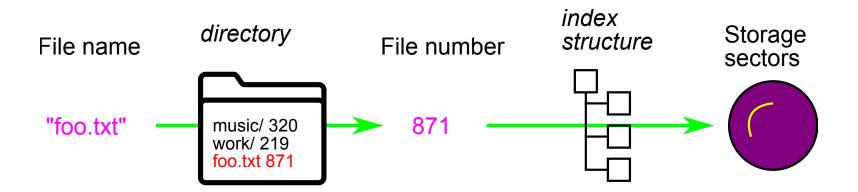
010111011001 101001110001 110111011000

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Directory

Directories provide names for files

- Groups of named files or subdirectories
- Mapping from human-readable file names to file metadata locations



Naming conventions

• Depends on filesystem

Case

- Windows traditionally case-insensitive
- UNIX traditionally case-sensitive

Length

- Before, could be quite limited (11 with MS-FAT)
- Today, usually up to 255 characters

Extension

- Before, separated from filename, and meaningful
- Today, usually part of the filename, and just a hint

```
$ file innocent_text_file.txt
innocent_text_file: ELF 64-bit LSB executable, x86-64,
version 1 (SYSV), dynamically linked

$ file scary_executable_file.x
scary_executable_file.x: ASCII text
```

Digression about executables

How does the OS recognize executable files?

Binary executables

- *Magic number* at very beginning of file, to identify its format
- Example with an ELF executable

```
$ xxd /usr/bin/firefox | head -1
00000000: 7f45 4c46 0201 0100 0000 0000 0000 0000 .ELF.....
```

Scripts

- Shebang as first line of script (#!)
 - Tells OS which *interpreter* should be used to run the script
 - Line ignored by interpreter because it's a comment
- Example with a shell script

```
$ cat test_script.sh
#!/bin/sh
echo Hello
```

```
$ ./test_script.sh
Hello
$ /bin/sh test_script.sh
Hello
```

Path

Absolute

- Path of file from the root directory
- e.g., /home/jporquet/project1/solution

Relative

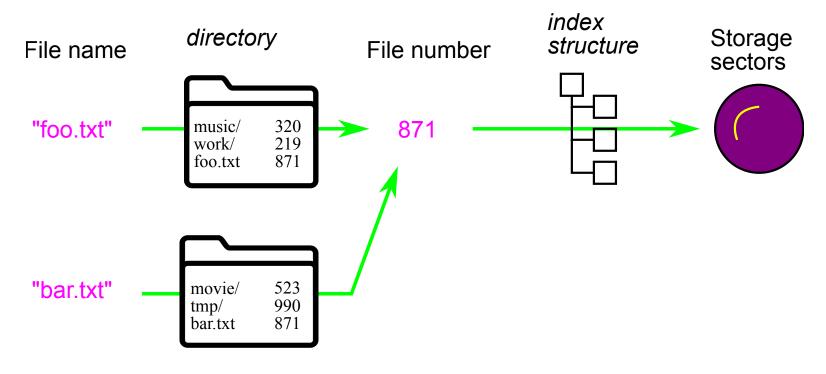
- Path from the current working directory (CWD)
- CWD stored in process' PCB

Special entries

- 2 special entries in each UNIX directory
 - .: current directory
 - . .: parent directory

Hard link

Link from name to metadata location



Hard links and cycles

• No difference between an "original" file name and a hard-link to it

```
$ touch test
$ ls -li test
21892004 -rw-r--r-- 1 joel joel 0 2017-03-02 17:15 test

$ ln test test2
$ ls -li test*
21892004 -rw-r--r-- 2 joel joel 0 2017-03-02 17:15 test
21892004 -rw-r--r-- 2 joel joel 0 2017-03-02 17:15 test
```

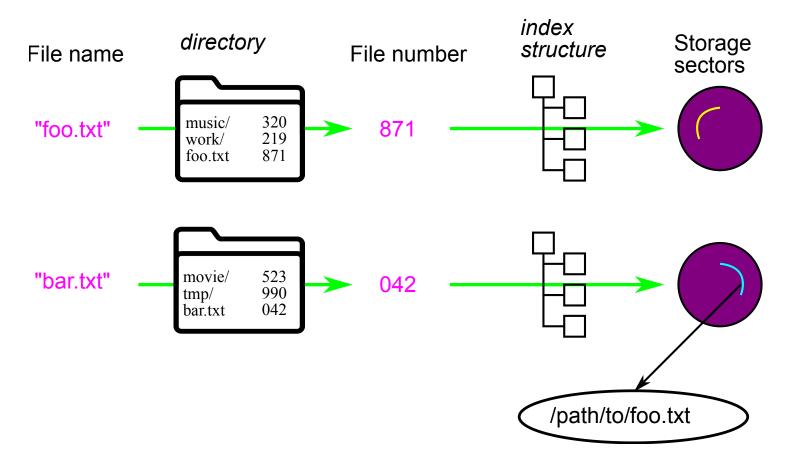
- Creation of a directory loop would make any file tree walker error-prone (e.g. du or fsck)
 - o Unless keeping track of all the inodes that are being traversed
- Hard-links to directory are forbidden

```
$ mkdir mydir
$ ln mydir mydirlink
ln: mydir: hard link not allowed for directory
```

Solution: special type of files dedicated for links.

Soft link (aka symbolic link)

Link from name to alternate name



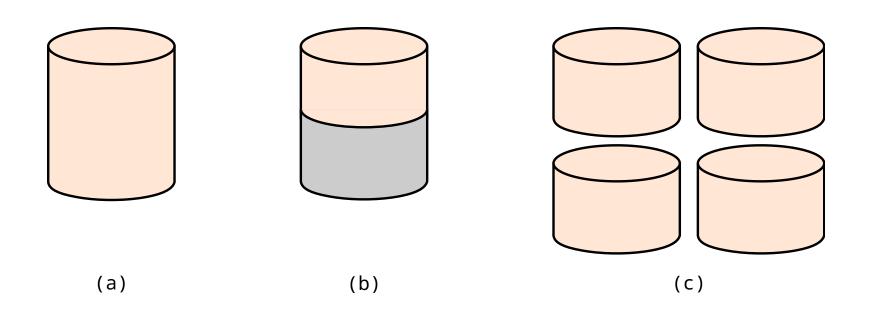
- Symlinks can be well-identified
- File tree walkers can safely ignore them
- POSIX limit: #define _POSIX_SYMLOOP_MAX 8

Volume

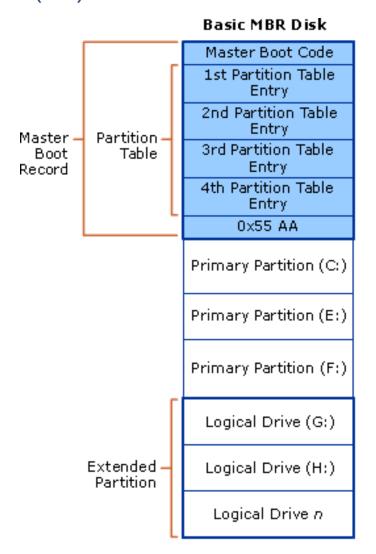
A volume is a collection of physical storage resources that form a logical storage device containing a file-system.

A volume can be:

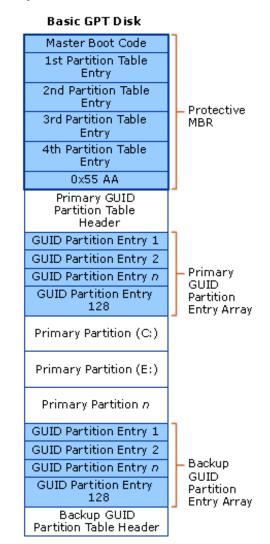
- (a) A whole disk
- (b) A partition on a disk
- (c) Multiple disks



Disk partitions MBR (old)

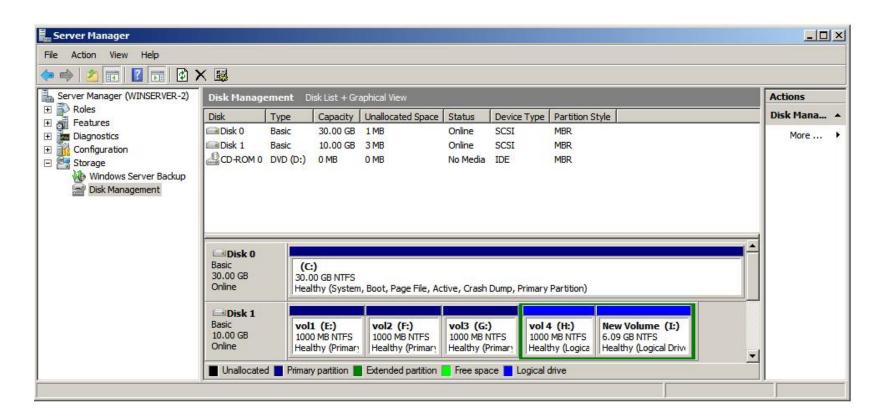


GPT (new)

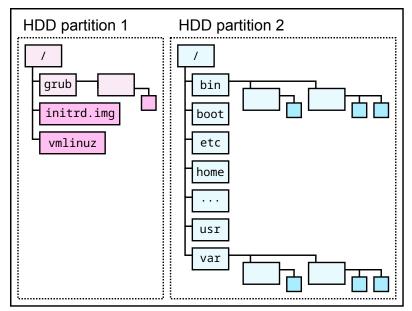


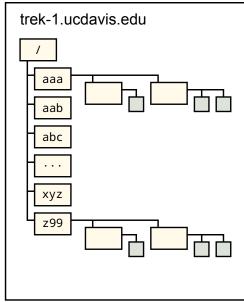
Drive letter assignment (Windows)

- Each volume holds a fully independent tree, in its own namespace
- Letter assignment typical order:
 - A:, B:: first and second floppy disk drives
 - C:: first active primary partition of the first physical hard drive
 - Then, first active primary partition of subsequent physical hard drives, etc.



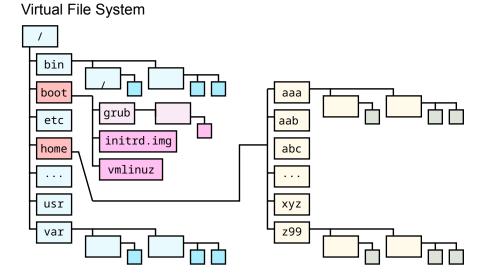
Mounting (UNIX)





Mountpoints:

hda2 => / hda1 => /boot trek-1.ucdavis.edu => /home

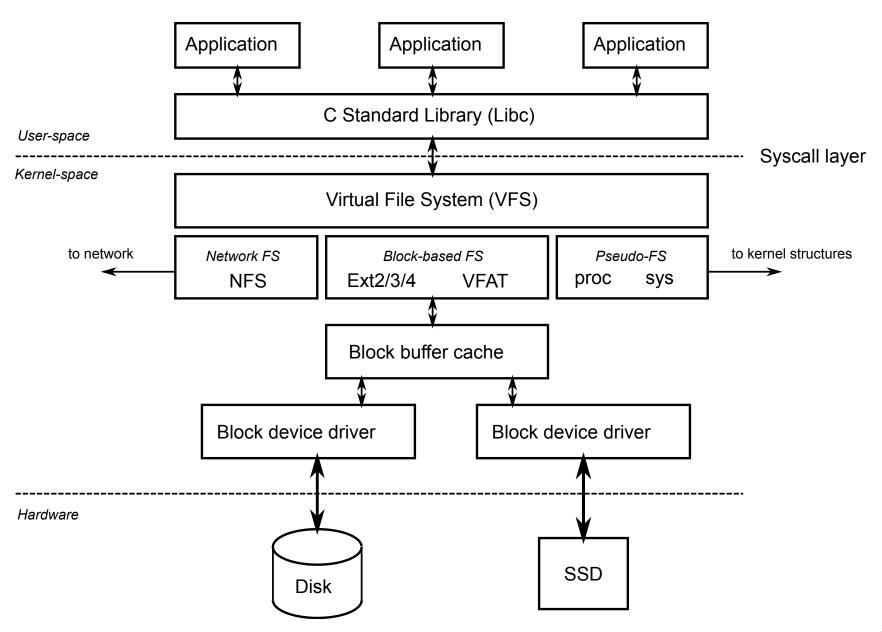


Mounting (UNIX)

Mounting multiple volumes arbitrarily in a single logical hierarchy

```
$ cat /etc/fstab
/dev/sda1
                                       ext3
/dev/sda2
                    /boot
                                       ext3
/dev/sdb1
                   /home
                                       ext3
/dev/cdrom
                 /mnt/cdrom
                                       iso9660
                                  vfat
/dev/sdc1
              /mnt/usbkey
10.0.0.1:/shared /srv/shared
                                       nfs
# the following entry would not make much sense (duplicate) but is possible
#/dev/sdb1
                    /mnt/home
                                            ext3
```

Filesystem software layers



Filesystem typical API

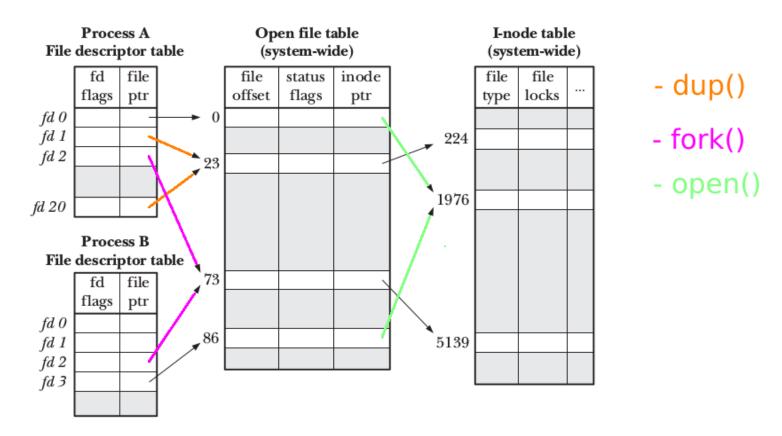
Create and delete files

- create(filename, mode)
 - o create new (empty) file, including metadata and name in directory
- link(existing_filename, new_filename)
 - o create new name for same underlying file as existing filename
- unlink(filename)
 - remove name for a file from its directory

Filesystem typical API

Open and close

- open(filename)
 - prepare to access specified file and return file descriptor
- close(filename):
 - release resources associated with open file



Filesystem typical API

File access

- read(), write(): sequential reading/writing
- seek(): change file's current position for random access

```
fd = open(...)
lseek(fd, 42, SEEK_CUR)
/* Write in open file at offset 42 */
char c = 'a';
write(fd, &c, 1);
close(fd);
```

- mmap(): create mapping between file's content and memory
- munmap(): destroy mapping

```
fd = open(...)
char *address = mmap(0, len, PROT_WRITE, MAP_SHARED, fd, 0)
address[42] = 'a';
munmap(address, len);
close(fd);
```