# Common

# Setup

- 1. Terminal: font Monospace 12
- 2. Gedit: Oblivion, font Monospace 12, auto indent, display line numbers, tab 4, highlight matching brackets, highlight current line, F9 (side panel)
- 3. /.bashrc: export CXXFLAGS='-Wall -Wshadow -Wextra -Wconversion -Wnounused-result -Wno-deprecated-declarations -O2 -std=gnu++11 -g -DLOCAL'
- 4. for i in {A..K}; do mkdir \$i; cp main.cpp \$i/\$i.cpp; done

# Template

```
#include <bits/stdc++.h>
using namespace std;
\#define\ pb\ push\_back
#define mp make_pair
#define fst first
#define snd second
\#define \ sz(x) \ (int) \ ((x).size())
#define form(i, n) for (int i = 0; i < (n); ++i)
#define form (i, n) for (int i = (n) - 1; i \ge 0; --i)
#define forab(i, a, b) for (int i = (a); i < (b); ++i)
#define all(c) (c).begin(), (c).end()
using ll = long long;
using vi = vector<int>;
using pii = pair<int, int>;
#define FNAME ""
int main() {
#ifdef LOCAL
  freopen(FNAME".in", "r", stdin);
  freopen(FNAME".out", "w", stdout);
  cin.tie(0);
  ios_base::sync_with_stdio(0);
  return 0;
```

#### 3 **Stress**

```
@echo off
for /L %%i in (1,1,10000000) do (
gen.exe || exit
main.exe || exit
stupid.exe || exit
fc .out 2.out || exit
echo Test %%i OK
```

### Java

```
import java.io.BufferedReader;
import java.io.FileNotFoundException;
import java.io.FileReader;
import java.io.IOException;
import java.io.InputStreamReader;
import java.io.PrintWriter;
import java.util.*;
public class Main {
 FastScanner in;
 PrintWriter out;
 void solve() {
   int a = in.nextInt();
    int b = in.nextInt();
    out.print(a + b);
  void run() {
```

```
trv {
      in = new FastScanner("input.txt");
      out = new PrintWriter("output.txt");
      solve();
      out.flush():
      out.close();
    \ \} \ \mathtt{catch} \ (\mathtt{FileNotFoundException} \ \mathtt{e}) \ \{
      e.printStackTrace();
      System.exit(1);
  }
  class FastScanner {
    BufferedReader br:
    StringTokenizer st;
    public FastScanner() {
      br = new BufferedReader(new InputStreamReader(System.in));
    public FastScanner(String s) {
       br = new BufferedReader(new FileReader(s));
      } catch (FileNotFoundException e) {
        e.printStackTrace();
    }
    String nextToken() {
      while (st == null || !st.hasMoreElements()) {
          st = new StringTokenizer(br.readLine());
        } catch (IOException e) {
          e.printStackTrace();
      }
      return st.nextToken();
    }
    int nextInt() {
      return Integer.parseInt(nextToken());
    7
    long nextLong() {
      return Long.parseLong(nextToken());
    double nextDouble() {
      return Double.parseDouble(nextToken());
    char nextChar() {
      try {
       return (char) (br.read());
      } catch (IOException e) {
        e.printStackTrace();
      return 0;
    String nextLine() {
      try {
       return br.readLine();
      } catch (IOException e) {
        e.printStackTrace();
      return "";
  7
  public static void main(String[] args) {
   new Main().run();
}
```

# 2 Big numbers

# 5 Big Int

```
constexpr int BASE = 1000000000;
constexpr int BASE_DIGITS = 9;
struct BigInt {
 // value == 0 is represented by empty z
 vi z; // digits
  // sign == 1/-1 <==> value >=/< 0
 int sign;
   BigInt(): sign(1) {}
 BigInt(ll v) { *this = v: }
   {\tt BigInt\&\ operator=(11\ v)\ \{}
    sign = v < 0 ? -1 : 1; v *= sign;
    z.clear(); for (; v > 0; v = v / BASE) z.pb((int) (v \%
    → BASE)):
    return *this;
 }
   BigInt& operator+=(const BigInt& other) {
    if (sign == other.sign) {
     for (int i = 0, carry = 0; i < sz(other.z) || carry; ++i) {</pre>
       if (i == sz(z)) z.pb(0);
       z[i] += carry + (i < sz(other.z) ? other.z[i] : 0);
        carry = z[i] >= BASE;
       if (carry) z[i] -= BASE;
    } else if (other != 0 /* prevent infinite loop */) {
      *this -= -other;
    }
   return *this;
 }
    friend BigInt operator+(BigInt a, const BigInt% b) { return a
    \hookrightarrow += b; }
    BigInt& operator = (const BigInt& other) {
    if (sign == other.sign) {
     if ((sign == 1 && *this >= other) || (sign == -1 && *this
      for (int i = 0, carry = 0; i < sz(other.z) || carry; ++i)
         z[i] -= carry + (i < sz(other.z) ? other.z[i] : 0);</pre>
          carry = z[i] < 0;
          if (carry)
           z[i] += BASE;
       }
       trim():
      } else {
        *this = other - *this:
        this->sign = -this->sign;
     }
   } else
      *this += -other;
   return *this;
    friend BigInt operator-(BigInt a, const BigInt% b) { return a
     \rightarrow -= b; }
    BigInt& operator*=(int v) {
    if (v < 0) sign = -sign, v = -v;
    for (int i = 0, carry = 0; i < sz(z) || carry; ++i) {
     if (i == sz(z))
       z.pb(0);
      11 cur = (11) z[i] * v + carry;
      carry = (int) (cur / BASE);
     z[i] = (int) (cur \% BASE);
    }
   trim();
    BigInt operator*(int v) const { return BigInt(*this) *= v; }
    friend pair < BigInt, BigInt > divmod(const BigInt & a1, const
    → BigInt& b1) {
    int norm = BASE / (b1.z.back() + 1);
   BigInt a = a1.abs() * norm;
    BigInt b = b1.abs() * norm;
   BigInt q, r;
    q.z.resize(sz(a.z));
    fornr (i, sz(a.z)) {
     r *= BASE, r += a.z[i];
      int s1 = sz(b.z) < sz(r.z) ? r.z[sz(b.z)] : 0;
```

```
int s2 = sz(b.z) - 1 < sz(r.z) ? r.z[sz(b.z) - 1] : 0;
    int d = (int) (((11) s1 * BASE + s2) / b.z.back());
    r -= b * d;
    while (r < 0) r += b, --d;
    q.z[i] = d;
  q.sign = a1.sign * b1.sign, r.sign = a1.sign;
  q.trim(), r.trim();
  return {q, r / norm};
  BigInt operator/(const BigInt& v) const { return divmod(*this,
  \rightarrow v).fst; }
  BigInt operator%(const BigInt& v) const { return divmod(*this,
  → v).snd: }
  BigInt& operator/=(int v) {
  if (v < 0) sign = -sign, v = -v;
  int rem = 0;
  formr (i, sz(z)) {
   11 \text{ cur} = z[i] + \text{rem} * (11) BASE;
    z[i] = (int) (cur / v);
    rem = (int) (cur % v);
  trim();
  return *this;
}
  BigInt operator/(int v) const { return BigInt(*this) /= v; }
  int operator%(int v) const {
  if (v < 0) v = -v;
  int m = 0;
  formr (i, sz(z))
    m = (int) ((z[i] + m * (11) BASE) % v);
  return m * sign;
  BigInt\& operator*=(const BigInt\& v) { return *this = *this *}
  → v; }
BigInt& operator/=(const BigInt& v) { return *this = *this / v;
  bool operator<(const BigInt& v) const {</pre>
  if (sign != v.sign) return sign < v.sign;</pre>
  if (sz(z) != sz(v.z)) return sz(z) * sign < sz(v.z) * v.sign;
  formr (i, sz(z))
    if (z[i] != v.z[i])
      return z[i] * sign < v.z[i] * sign;</pre>
  return false;
  bool operator>(const BigInt& v) const { return v < *this; }</pre>
bool operator<=(const BigInt& v) const { return !(v < *this); }</pre>
bool operator>=(const BigInt& v) const { return !(*this < v); }</pre>
  bool operator==(const BigInt& v) const { return !(*this < v)</pre>
  \hookrightarrow && !(v < *this); }
  bool operator!=(const BigInt \& v) const { return *this < v || v
  \hookrightarrow < *this; }
  void trim() {
  while (!z.empty() && z.back() == 0) z.pop_back();
  if (z.empty()) sign = 1;
}
bool isZero() const { return z.empty(); }
friend BigInt operator-(BigInt v) {
  if (!v.z.empty()) v.sign = -v.sign;
  return v;
BigInt abs() const {
  return sign == 1 ? *this : -*this;
void read(const string& s) {
  sign = 1, z.clear();
  int pos = 0;
  while (pos < s.size() && (s[pos] == '-' || s[pos] == '+')) {
    if (s[pos] == '-') sign = -sign;
    ++pos;
  }
  for (int i = sz(s) - 1; i >= pos; i -= BASE_DIGITS) {
    int x = 0;
    forab (j, max(pos, i - BASE_DIGITS + 1), i)
      x = x * 10 + s[j] - '0';
    z.pb(x);
  }
  trim();
friend ostream &operator << (ostream & stream, const BigInt & v) {
```

```
if (v.sign == -1)
     stream << '-';
    stream << (v.z.empty() ? 0 : v.z.back());
    fornr (i, sz(v.z) - 1)
      stream << setw(BASE_DIGITS) << setfill('0') << v.z[i];</pre>
    return stream;
  }
  static vi convertBase(const vi& a, int oldDigits, int
  \hookrightarrow newDigits) {
   vector<ll> p(max(oldDigits, newDigits) + 1);
    p[0] = 1;
    for (int i = 1; i < sz(p); i++)
     p[i] = p[i - 1] * 10;
    vi res:
    11 cur = 0;
    int curDigits = 0;
    for (int v : a) {
     cur += v * p[curDigits];
     curDigits += oldDigits;
     while (curDigits >= newDigits) {
       res.pb(int(cur % p[newDigits]));
        cur /= p[newDigits];
        curDigits -= newDigits;
     }
    }
    res.pb((int) cur);
    while (!res.empty() && res.back() == 0) res.pop_back();
    return res;
  7
  static vll karatsubaMultiply(const vll& a, const vll& b) {
    int n = sz(a);
    vll res(n + n):
    if (n <= 32) \{
     for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++)
         res[i + j] += a[i] * b[j];
      return res;
    int k = n >> 1;
    vll a1(a.begin(), a.begin() + k), a2(a.begin() + k, a.end());
    vll b1(b.begin(), b.begin() + k), b2(b.begin() + k, b.end());
    vll a1b1 = karatsubaMultiply(a1, b1);
    vll a2b2 = karatsubaMultiply(a2, b2);
    forn (i, k) a2[i] += a1[i];
    forn (i, k) b2[i] += b1[i];
    vll r = karatsubaMultiply(a2, b2);
    forn (i, sz(a1b1)) r[i] -= a1b1[i];
    forn (i, sz(a2b2)) r[i] -= a2b2[i];
    forn (i, sz(r)) res[i + k] += r[i];
    forn (i, sz(a1b1)) res[i] += a1b1[i];
    forn (i, sz(a2b2)) res[i + n] += a2b2[i];
    return res;
  BigInt operator*(const BigInt& v) const {
    vi a6 = convertBase(this->z, BASE_DIGITS, 6);
    vi b6 = convertBase(v.z, BASE_DIGITS, 6);
    vll a(all(a6)), b(all(b6));
    while (sz(a) < sz(b)) a.pb(0);
    while (sz(b) < sz(a)) b.pb(0);
    while (sz(a) & (sz(a) - 1)) a.pb(0), b.pb(0);
    vll c = karatsubaMultiply(a, b);
    BigInt res;
    res.sign = sign * v.sign;
    int carry = 0;
    forn (i, sz(c)) {
     ll cur = c[i] + carry;
     res.z.push_back((int) (cur % 1000000));
     carry = (int) (cur / 1000000);
   res.z = convertBase(res.z, 6, BASE_DIGITS);
   res.trim();
   return res:
};
```

# 6 FFT

```
int rev[MAX_N];
//typedef complex<dbl> Num;
struct Num {
  dbl x, y;
  Num() {}
  Num(dbl _x, dbl _y): x(_x), y(_y) {}
  inline dbl real() const { return x; }
  inline dbl imag() const { return y; }
  inline Num operator+(const Num &B) const { return Num(x + B.x, y
  \rightarrow + B.y); }
  inline Num operator-(const Num &B) const { return Num(x - B.x, y
  \hookrightarrow - B.v); }
  inline Num operator*(dbl k) const { return Num(x * k, y * k); }
  inline Num operator*(const Num &B) const { return Num(x * B.x -
  \hookrightarrow y * B.y, x * B.y + y * B.x); }
  inline void operator+=(const Num &B) { x += B.x, y += B.y; }
 inline void operator/=(dbl k) { x /= k, y /= k; }
 inline void operator*=(const Num &B) { *this = *this * B; }
};
Num rt[MAX_N];
inline Num sqr(const Num &x) { return x * x; }
inline Num conj(const Num &x) { return Num(x.real(), -x.imag());
inline int getN(int n) {
 int k = 1;
  while(k < n)
   k <<= 1;
  return k:
}
void fft(Num *a, int n) {
  assert(rev[1]); // don't forget to init
  int q = MAX_N / n;
  forn (i, n)
    if(i < rev[i] / q)
     swap(a[i], a[rev[i] / q]);
  for (int k = 1; k < n; k <<= 1)
    for (int i = 0; i < n; i += 2 * k)
     forn (j, k) {
        const Num z = a[i + j + k] * rt[j + k];
        a[i + j + k] = a[i + j] - z;
        a[i + j] += z;
void fftInv(Num *a, int n) {
 fft(a, n);
  reverse(a + 1, a + n);
  forn (i, n)
    a[i] /= n;
void doubleFft(Num *a, Num *fa, Num *fb, int n) { // only if you
\hookrightarrow need it
 fft(a, n);
  const int n1 = n - 1;
 forn (i, n) {
    const Num &z0 = a[i], &z1 = a[(n - i) & n1];
    fa[i] = Num(z0.real() + z1.real(), z0.imag() - z1.imag()) *
    fb[i] = Num(z0.imag() + z1.imag(), z1.real() - z0.real()) *
 }
}
Num tmp[MAX_N];
template<class T>
void mult(T *a, T *b, T *r, int n) { // n = 2^n k
  forn (i, n)
    tmp[i] = Num((dbl) a[i], (dbl) b[i]);
  fft(tmp, n);
  const int n1 = n - 1;
  const Num c = Num(0, -0.25 / n);
  fornr (i, n / 2 + 1) {
```

```
const int j = (n - i) \& n1;
    const Num z0 = sqr(tmp[i]), z1 = sqr(tmp[j]);
   tmp[i] = (z1 - conj(z0)) * c;
    tmp[j] = (z0 - conj(z1)) * c;
 fft(tmp, n);
 forn (i, n)
   r[i] = (T) round(tmp[i].real());
void init() { // don't forget to init
 forn(i, MAX N)
    rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (LOG - 1));
 rt[1] = Num(1, 0);
 for (int k = 1, p = 2; k < LOG; k++, p *= 2) {
   const Num x(cos(PI / p), sin(PI / p));
    forab (i, p / 2, p)
     rt[2 * i] = rt[i], rt[2 * i + 1] = rt[i] * x;
 }
```

# 7 FFT by mod and FFT with digits up to $10^6$

```
Num ta[MAX_N], tb[MAX_N], tf[MAX_N], tg[MAX_N];
const. int. HALF = 15:
void mult(int *a, int *b, int *r, int n, int mod) {
  int tw = (1 << HALF) - 1;</pre>
  forn (i, n) {
    int x = int(a[i] % mod);
    ta[i] = Num(x \& tw, x >> HALF);
  forn (i, n) {
   int x = int(b[i] % mod);
    tb[i] = Num(x \& tw, x >> HALF);
  fft(ta, n), fft(tb, n);
  forn (i, n) {
    int j = (n - i) & (n - 1);
    Num a1 = (ta[i] + conj(ta[j])) * Num(0.5, 0);
    Num a2 = (ta[i] - conj(ta[j])) * Num(0, -0.5);
    Num b1 = (tb[i] + conj(tb[j])) * Num(0.5 / n, 0);
    Num b2 = (tb[i] - conj(tb[j])) * Num(0, -0.5 / n);
    tf[j] = a1 * b1 + a2 * b2 * Num(0, 1);
    tg[j] = a1 * b2 + a2 * b1;
  fft(tf, n), fft(tg, n);
  forn (i, n) {
    11 aa = 11(tf[i].x + 0.5);
    11 bb = 11(tg[i].x + 0.5);
    11 cc = 11(tf[i].y + 0.5);
   r[i] = int((aa + ((bb \% mod) << HALF) + ((cc \% mod) << (2 *)
    → HALF))) % mod);
 }
int tc[MAX_N], td[MAX_N];
const int MOD1 = 1.5e9, MOD2 = MOD1 + 1;
void multLL(int *a, int *b, ll *r, int n){
 mult(a, b, tc, n, MOD1), mult(a, b, td, n, MOD2);
    r[i] = tc[i] + (td[i] - tc[i] + (11)MOD2) * MOD1 % MOD2 *
    \hookrightarrow MOD1;
```

# 3 Data Structures

# 8 Centroid Decomposition

```
vi g[MAX_N];
int d[MAX_N], par[MAX_N], centroid;
//d par -
int find(int v, int p, int total) {
```

```
int size = 1, ok = 1;
  for (int to : g[v])
   if (d[to] == -1 \&\& to != p) {
      int s = find(to, v, total);
      if (s > total / 2) ok = 0;
      size += s;
    }
  if (ok && size > total / 2) centroid = v;
  return size;
void calcInComponent(int v, int p, int level) {
  // do something
 for (int to : g[v])
   if (d[to] == -1 && to != p)
      calcInComponent(to, v, level);
//fill(d, d + n, -1)
//decompose(0, -1, 0)
void decompose(int root, int parent, int level) {
  find(root, -1, find(root, -1, INF));
  int c = centroid;
  par[c] = parent, d[c] = level;
  \verb| calcInComponent(centroid, -1, level); \\
  for (int to : g[c])
    if (d[to] == -1)
      decompose(to, c, level + 1);
}
```

# 9 Convex Hull Trick

```
struct Line {
  int k. b:
  Line() {}
  Line(int _k, int _b): k(_k), b(_b) {}
  ll get(int x) { return b + k * 111 * x; }
  bool operator<(const Line &1) const { return k < 1.k; } //</pre>
};
                    (a,b)
                             (a,c)
inline bool check(Line a, Line b, Line c) {
  return (a.b - b.b) * 111 * (c.k - a.k) < (a.b - c.b) * 111 *
  \hookrightarrow (b.k - a.k);
struct Convex {
  vector<Line> st:
  inline void add(Line 1) {
    while (sz(st) \ge 2 \&\& ! check(st[sz(st) - 2], st[sz(st) - 1],
      st.pop_back();
    st.pb(1);
  int get(int x) {
    int 1 = 0, r = sz(st);
    while (r - 1 > 1) {
      int m = (1 + r) / 2; //
      if (st[m - 1].get(x) < st[m].get(x))
       1 = m;
      else
        r = m;
    }
    return 1;
  Convex() {}
  Convex(vector<Line> &lines) {
    st.clear():
    for(Line &1 : lines)
      add(1);
  Convex(Line line) { st.pb(line); }
  Convex(const Convex &a, const Convex &b) {
    vector<Line> lines;
    lines.resize(sz(a.st) + sz(b.st));
    merge(all(a.st), all(b.st), lines.begin());
    st.clear();
    for(Line &1 : lines)
      add(1):
```

forn (j, 2)

toPush[path][2 \* v + j] = toPush[path][v];

```
};
                                                                             t[path][v] = toPush[path][v];
                                                                           toPush[path][v] = -1;
                                                                        }
10
    DSU
                                                                      }
int pr[MAX_N];
                                                                      int getST(int path, int v, int vl, int vr, int ind) {
                                                                        pushST(path, v, vl, vr);
int get(int v) {
                                                                        if (vl == vr - 1)
 return v == pr[v] ? v : pr[v] = get(pr[v]);
                                                                         return t[path][v];
                                                                        int vm = (vl + vr) / 2;
                                                                        if (ind >= vm)
bool unite(int v, int u) {
                                                                          return getST(path, 2 * v + 1, vm, vr, ind);
  v = get(v), u = get(u);
                                                                        return getST(path, 2 * v, v1, vm, ind);
  if (v == u) return 0;
  pr[u] = v;
  return 1;
                                                                      void setST(int path, int v, int vl, int vr, int l, int r, int val)
                                                                       ← {
                                                                        if (vl >= l && vr <= r) {
void init(int n) {
                                                                          toPush[path][v] = val;
 forn (i, n) pr[i] = i;
                                                                          pushST(path, v, v1, vr);
                                                                          return;
    Fenwick Tree
                                                                        pushST(path, v, v1, vr);
                                                                        if (vl >= r || l >= vr)
int t[MAX_N];
                                                                          return;
                                                                        int vm = (vl + vr) / 2;
int get(int ind) {
                                                                        setST(path, 2 * v, vl, vm, l, r, val);
  int res = 0;
                                                                        setST(path, 2 * v + 1, vm, vr, l, r, val);
  for (; ind >= 0; ind &= (ind + 1), ind--)
                                                                        t[path][v] = min(t[path][2 * v], t[path][2 * v + 1]);
   res += t[ind];
                                                                      }
  return res:
                                                                      bool isUpper(int v, int u) {
                                                                        return tin[v] <= tin[u] && tout[v] >= tout[u];
void add(int ind, int n, int val) {
  for (; ind < n; ind |= (ind + 1))
    t[ind] += val;
                                                                       int getHLD(int v) {
                                                                        return getST(comp[v], 1, 0, sz(t[comp[v]]) / 2, num[v]);
int sum(int 1, int r) { // [l, r)
 return get(r - 1) - get(l - 1);
                                                                      int setHLD(int v, int u, int val) {
                                                                        int ans = 0, w = 0;
                                                                        forn (i, 2) {
                                                                          while (!isUpper(w = top[comp[v]], u))
  setST(comp[v], 1, 0, sz(t[comp[v]]) / 2, 0, num[v] + 1,
12 Hash Table
using H = 11;
                                                                             \hookrightarrow val), v = pr[w];
const int HT_SIZE = 1<<20, HT_AND = HT_SIZE - 1, HT_SIZE_ADD =</pre>
                                                                          swap(v, u);

→ HT_SIZE / 100;

H ht[HT_SIZE + HT_SIZE_ADD];
                                                                        setST(comp[v], 1, 0, sz(t[comp[v]]) / 2, min(num[v], num[u]),
int data[HT_SIZE + HT_SIZE_ADD];
                                                                        \rightarrow max(num[v], num[u]) + 1, val);
                                                                        return ans;
int get(const H &hash){
 int k = ((11) hash) & HT_AND;
  while (ht[k] && ht[k] != hash) ++k;
                                                                      void dfs(int v, int p) {
  return k;
                                                                        tin[v] = curTime++;
                                                                        size[v] = 1:
                                                                        pr[v] = p;
void insert(const H &hash, int x){
                                                                        for (int u : g[v])
  int k = get(hash);
                                                                          if (u != p) {
  if (!ht[k]) ht[k] = hash, data[k] = x;
                                                                            dfs(u, v);
                                                                            size[v] += size[u];
                                                                          7
bool count(const H &hash, int x){
                                                                        tout[v] = curTime++;
  int k = get(hash);
  return ht[k] != 0;
                                                                      void build(int v) {
                                                                        if (v == 0 \mid \mid size[v] * 2 < size[pr[v]])
                                                                           top[curPath] = v, comp[v] = curPath, num[v] = 0, curPath++;
     Heavy Light Decomposition
                                                                          comp[v] = comp[pr[v]], num[v] = num[pr[v]] + 1;
                                                                        lst[comp[v]].pb(v);
int size[MAX_N], comp[MAX_N], num[MAX_N], top[MAX_N], pr[MAX_N],
                                                                        for (int u : g[v])

    tin[MAX_N], tout[MAX_N];

                                                                           if (u != pr[v])
vi t[MAX_N], toPush[MAX_N], lst[MAX_N];
                                                                            build(u):
int curPath = 0, curTime = 0;
void pushST(int path, int v, int vl, int vr) {
                                                                      void initHLD() {
  if (toPush[path][v] != -1) {
                                                                        dfs(0, 0);
    if (vl != vr - 1)
                                                                        build(0);
```

```
forn (i, curPath) {
   int curSize = 1;
   while (curSize < sz(lst[i]))
      curSize *= 2;
   t[i].resize(curSize * 2);
   toPush[i] = vi(curSize * 2, -1);
   //initialize t[i]
}</pre>
```

# 14 Next Greater in Segment Tree

# 15 Sparse Table

```
int st[MAX_N] [MAX_LOG];
int lg[MAX_N];

int get(int 1, int r) { // [l, r)
    int curLog = lg[r - 1];
    return min(st[1][curLog], st[r - (1 << curLog)][curLog]);
}

void initSparseTable(int *a, int n) {
    lg[1] = 0;
    forab (i, 2, n + 1) lg[i] = lg[i / 2] + 1;
    forn (i, n) st[i][0] = a[i];
    forn (j, lg[n])
        forn (i, n - (1 << (j + 1)) + 1)
            st[i][j + 1] = min(st[i][j], st[i + (1 << j)][j]);
}</pre>
```

### 16 Fenwick Tree 2D

 $\hookrightarrow$  get(x\_1 - 1, y\_1 - 1);

```
11 a[4][MAX_N][MAX_N];
int n, m;
inline int f(int x) { return x & ~(x - 1); }
inline void add(int k, int x, int y, ll val) {
  for (; x \le n; x += f(x))
   for (int j = y; j \le m; j += f(j))
      a[k][x][j] += val;
inline ll get(int k, int x, int y) {
 11 s = 0;
 for (; x > 0; x -= f(x))
   for (int j = y; j > 0; j -= f(j))
     s += a[k][x][j];
 return s;
inline ll get(int x, int y) {
 return ll(x + 1) * (y + 1) * get(0, x, y) - (y + 1) * get(1, x, y)
      -(x + 1) * get(2, x, y) + get(3, x, y);
inline void add(int x, int y, ll val) {
 add(0, x, y, val);
 add(1, x, y, val * x);
 add(2, x, y, val * y);
 add(3, x, y, val * x * y);
inline ll get(int x_1, int y_1, int x_2, int y_2) {
  return get(x_2, y_2) - get(x_1 - 1, y_2) - get(x_2, y_1 - 1) +
```

```
// Adds val to corresponding rectangle
inline void add(int x_1, int y_1, int x_2, int y_2, ll val) {
  add(x_1, y_1, val);
  if (y_2 < m) add(x_1, y_2 + 1, -val);
  if (x_2 < n) add(x_2 + 1, y_1, -val);
  if (x_2 < n && y_2 < m) add(x_2 + 1, y_2 + 1, val);
}</pre>
```

# 17 Segment Tree 2D

```
int tSize = (1 << 10);</pre>
struct Node1D {
    Node1D *1, *r;
    ll val, need;
    Node1D(): l(nullptr), r(nullptr), val(0), need(0) {}
    inline void norm() {
        if(!1) 1 = new Node1D();
         if(!r) r = new Node1D();
     11 get(int q1, int qr, int v1 = 0, int vr = tSize) {
        if(vl >= qr || ql >= vr)
             return 0;
         if(ql <= vl && vr <= qr)
            return val;
         int a = max(vl, ql), b = min(vr, qr), vm = (vl + vr) / 2;
         norm():
         return l->get(ql, qr, vl, vm) + r->get(ql, qr, vm, vr) + need
         \rightarrow * 11(b - a);
     }
     void add(int ql, int qr, int x, int vl = 0, int vr = tSize) {
         if (ql >= vr || vl >= qr)
            return;
         if (ql <= vl && vr <= qr){
             need += x;
             val += x * 11(vr - v1);
             return;
        int vm = (v1 + vr) / 2;
         norm();
         1->add(q1, qr, x, v1, vm), r->add(q1, qr, x, vm, vr);
         val = 1->val + r->val + need * (vr - vl);
    }
};
struct Node2D {
    Node2D *1. *r:
    Node1D *val, *need;
    Node2D(): 1(nullptr), r(nullptr), val(new Node1D()), need(new
     \rightarrow Node1D()) {}
    inline void norm() {
        if(!1) 1 = new Node2D();
        if(!r) r = new Node2D();
     ll get(int ql0, int qr0, int ql1, int qr1, int vl = 0, int vr =
     if(vl >= qr0 || ql0 >= vr)
            return 0:
         if(q10 <= v1 && vr <= qr0)
            return val->get(ql1, qr1);
         int a = max(v1, q10), b = min(vr, qr0), vm = (v1 + vr) / 2;
         norm():
         return 1->get(q10, qr0, q11, qr1, v1, vm) + r->get(q10, qr0,
         \rightarrow ql1, qr1, vm, vr) + need->get(ql1, qr1) * ll(b - a);
     void add(int q10, int qr0, int q11, int qr1, int x, int v1 = 0,

    int vr = tSize) {
        if (ql0 >= vr || vl >= qr0)
            return:
         if (ql0 <= vl && vr <= qr0){
             need->add(ql1, qr1, x);
              val->add(ql1, qr1, x * ll(vr - vl));
             return;
         }
         int a = max(q10, v1), b = min(qr0, vr), vm = (v1 + vr) / 2;
         norm();
         1-> add(q10, \ qr0, \ q11, \ qr1, \ x, \ vl, \ vm), \ r-> add(q10, \ qr0, \ q11, \ qr1, \ r-> add(q10, \ qr0, \ q11, \ qr1, \ 
         \hookrightarrow qr1, x, vm, vr);
```

```
val->add(ql1, qr1, x * ll(b - a));
};
```

# 4 Dynamic Programming

### **18 LIS**

```
int longestIncreasingSubsequence(vi a) {
  int n = sz(a);
  vi d(n + 1, INF);
  d[0] = -INF;
  forn (i, n)
    *upper_bound(all(d), a[i]) = a[i];
  fornr (i, n + 1) if (d[i] != INF) return i;
  return 0;
}
```

### 19 DP tree

```
int dp[MAX_N] [MAX_N], a[MAX_N];
vi g[MAX_N];
int dfs(int v, int n) {
  forn (i, n + 1)
    dp[v][i] = -INF;
  dp[v][1] = a[v];
  int curSz = 1;
  for (int to : g[v]) {
    int toSz = dfs(to, n);
    for (int i = curSz; i >= 1; i--)
        fornr (j, toSz + 1)
            dp[v][i + j] = max(dp[v][i + j], dp[v][i] + dp[to][j]);
    curSz += toSz;
  }
  return curSz;
}
```

### 20 Masks tricks

```
int dp[(1 << MAX_MASK)][MAX_MASK];

void calcDP(int n) {
   forn(mask, 1 << n) {
      dp[mask][n] = 1;
   fornr(i, n) {
       dp[mask][i] = dp[mask][i + 1];
      if ((1 << i) & mask)
            dp[mask][i] += dp[mask ^ (1 << i)][i + 1];
    }
}</pre>
```

# 5 Flows

### 21 Utilities

```
22 Ford-Fulkerson
int used[MAX_N], pr[MAX_N];
int curTime = 1;
int dfs(int v, int can, int toPush, int t) {
 if (v == t) return can;
  used[v] = curTime;
 for (int edge : g[v]) {
    auto &e = edges[edge];
   if (used[e.u] != curTime && e.c - e.f >= toPush) {
      int flow = dfs(e.u, min(can, e.c - e.f), toPush, t);
      if (flow > 0) {
        addFlow(edge, flow), pr[e.u] = edge;
        return flow;
     }
   }
 }
 return 0;
int fordFulkerson(int s, int t) {
 int ansFlow = 0, flow = 0;
  // Without scaling
  while ((flow = dfs(s, INF, 1, t)) > 0)
   ansFlow += flow, curTime++;
  // With scaling
 fornr (i, INF_LOG)
   for (curTime++; (flow = dfs(s, INF, (1 \ll i), t)) > 0;
    ansFlow += flow;
  return ansFlow;
23 Dinic
int pr[MAX_N], d[MAX_N], q[MAX_N], first[MAX_N];
int dfs(int v, int can, int toPush, int t) {
 if (v == t) return can;
  int sum = 0;
  for (; first[v] < (int) g[v].size(); first[v]++) {</pre>
   auto &e = edges[g[v][first[v]]];
    if (d[e.u] != d[v] + 1 || e.c - e.f < toPush) continue;
   int flow = dfs(e.u, min(can, e.c - e.f), toPush, t);
    addFlow(g[v][first[v]], flow);
   can -= flow, sum += flow;
    if (!can)
   return sum:
  return sum;
bool bfs(int n, int s, int t, int curPush) {
  forn (i, n) d[i] = INF, first[i] = 0;
  int head = 0, tail = 0;
  q[tail++] = s;
  d[s] = 0;
  while (tail - head > 0) {
   int v = q[head++];
   for (int edge : g[v]) {
     auto &e = edges[edge];
      if (d[e.u] > d[v] + 1 \&\& e.c - e.f >= curPush)
        d[e.u] = d[v] + 1, q[tail++] = e.u;
  }
  return d[t] != INF;
int dinic(int n, int s, int t) {
 int ansFlow = 0;
  // Without scaling
  while (bfs(n, s, t, 1))
   ansFlow += dfs(s, INF, 1, t);
  // With scaling
 fornr (j, INF_LOG)
    while (bfs(n, s, t, 1 \ll j))
     ansFlow += dfs(s, INF, 1 \ll j, t);
```

return ansFlow;

}

# 24 Hungarian

```
const int INF = 1e9;
int a[MAX_N][MAX_N];
// min = sum of a[pa[i],i]
// you may optimize speed by about 15%, just change all vectors to
  static arrays
vi Hungarian(int n) {
  vi pa(n + 1, -1), row(n + 1, 0), col(n + 1, 0), la(n + 1);
  forn (k, n) {
    vi u(n + 1, 0), d(n + 1, INF);
    pa[n] = k;
    int 1 = n, x;
    while ((x = pa[1]) != -1) {
     u[1] = 1:
      int minn = INF, tmp, 10 = 1;
      forn (j, n)
        if (!u[j]) {
          if ((tmp = a[x][j] + row[x] + col[j]) < d[j])
            d[j] = tmp, la[j] = 10;
          if (d[j] < minn)</pre>
            minn = d[j], 1 = j;
        }
      forn (j, n + 1)
        if (u[j])
         col[j] += minn, row[pa[j]] -= minn;
          d[j] -= minn;
    while (1 != n)
      pa[1] = pa[la[1]], 1 = la[1];
 return pa;
```

# 25 Min Cost Max Flow

forn (i, n) pot[i] += d[i];

```
const int MAX_M = 1e4;
int pr[MAX_N], in[MAX_N], q[MAX_N * MAX_M], used[MAX_N],

    d[MAX_N], pot[MAX_N];

vi g[MAX_N];
struct Edge {
 int v, u, c, f, w;
  Edge() {}
  \label{eq:condition} \mbox{Edge(int $\_$v, int $\_$u, int $\_$c, int $\_$w): $v(\_$v), $u(\_$u), $c(\_$c),}
  \hookrightarrow f(0), w(_w) {}
vector<Edge> edges;
inline void addFlow(int e, int flow) {
  edges[e].f += flow, edges[e ^ 1].f -= flow;
inline void addEdge(int v, int u, int c, int w) {
  g[v].pb(sz(edges)), edges.pb(Edge(v, u, c, w));
  g[u].pb(sz(edges)), edges.pb(Edge(u, v, 0, -w));
int dijkstra(int n, int s, int t) {
  forn (i, n) used[i] = 0, d[i] = INF;
  d[s] = 0;
  while (1) {
    int v = -1;
    forn (i, n)
      if (!used[i] && (v == -1 \mid \mid d[v] > d[i]))
    if (v == -1 \mid \mid d[v] == INF) break;
    used[v] = 1;
    for (int edge : g[v]) {
      auto &e = edges[edge];
      int w = e.w + pot[v] - pot[e.u];
      if (e.c > e.f && d[e.u] > d[v] + w)
        d[e.u] = d[v] + w, pr[e.u] = edge;
  }
  if (d[t] == INF) return d[t];
```

```
return pot[t];
}
int fordBellman(int n, int s, int t) {
 forn (i, n) d[i] = INF;
 int head = 0, tail = 0;
 d[s] = 0, q[tail++] = s, in[s] = 1;
 while (tail - head > 0) {
   int v = q[head++];
   in[v] = 0;
   for (int edge : g[v]) {
     auto &e = edges[edge];
      if (e.c > e.f && d[e.u] > d[v] + e.w) {
       d[e.u] = d[v] + e.w;
       pr[e.u] = edge;
       if (!in[e.u])
          in[e.u] = 1, q[tail++] = e.u;
     }
   }
 }
 return d[t];
int minCostMaxFlow(int n, int s, int t) {
 int ansFlow = 0, ansCost = 0, dist;
 while ((dist = dijkstra(n, s, t)) != INF) {
   int curFlow = INF;
   for (int cur = t; cur != s; cur = edges[pr[cur]].v)
     curFlow = min(curFlow, edges[pr[cur]].c -

    edges[pr[cur]].f);

   for (int cur = t; cur != s; cur = edges[pr[cur]].v)
     addFlow(pr[cur], curFlow);
   ansFlow += curFlow;
   ansCost += curFlow * dist;
 return ansCost:
```

# 6 Games

# 26 Retrograde Analysis

```
int win[MAX_N], lose[MAX_N], outDeg[MAX_N];
vi rg[MAX_N];
void retro(int n) {
  queue<int> q;
  forn (i, n)
   if (!outDeg[i])
     lose[i] = 1, q.push(i);
  while (!q.empty()) {
   int v = q.front();
   q.pop();
   for (int to : rg[v])
     if (lose[v]) {
        if (!win[to])
          win[to] = 1, q.push(to);
     } else {
        outDeg[to]--;
        if (!outDeg[to])
          lose[to] = 1, q.push(to);
  }
}
```

# 7 Geometry

# 27 ClosestPoints (SweepLine)

```
11 d2 = 8e18, d = (11) sqrt(d2) + 1;
Pnt p[N];
inline 11 sqr(int x){
 return (11)x * x;
inline void relax(const Pnt &a, const Pnt &b){
 11 tmp = sqr(a.x - b.x) + sqr(a.y - b.y);
 if (tmp < d2)
    d2 = tmp, d = (11)(sqrt(d2) + 1 - 1e-9); // round up
inline bool xless(const Pnt &a, const Pnt &b){
  return a.x < b.x;
int main() {
 int n;
  scanf("%d", &n);
  forn(i, n)
   scanf("%d%d", &p[i].x, &p[i].y), p[i].i = i;
  sort(p, p + n, xless);
  set <Pnt> s:
  int 1 = 0;
 forn(r, n){
   set<Pnt>::iterator it_r = s.lower_bound(p[r]), it_l = it_r;
    for (; it_r != s.end() && it_r->y - p[r].y < d; ++it_r)
     relax(*it_r, p[r]);
    while (it_l != s.begin() && p[r].y - (--it_l)->y < d)
     relax(*it_l, p[r]);
    s.insert(p[r]);
    while (1 <= r \&\& p[r].x - p[1].x >= d)
      s.erase(p[1++]);
 printf("%.9f\n", sqrt(d2));
 return 0;
```

### 28 ConvexHull

```
typedef vector<Pnt> vpnt;
inline bool byAngle(const Pnt &a, const Pnt &b) {
 dbl x = a \% b:
 return eq(x, 0) ? a.len2() < b.len2() : x < 0;
vpnt convexHull(vpnt p) {
 int n = sz(p);
  assert(n > 0);
  swap(p[0], *min_element(all(p)));
  forab(i, 1, n)
 p[i] = p[i] - p[0];
  sort(p.begin() + 1, p.end(), byAngle);
                              (1) (2)
 (1):
  int k = p.size() - 1;
  while(k > 0 \; \&\& \; eq((p[k-1] - p.back()) \; \% \; p.back(), \; 0))
  reverse(pi.begin() + k, pi.end());*/
  int rn = 0;
  vpnt r(n);
  r[rn++] = p[0];
  forab(i, 1, n){
   Pnt q = p[i] + p[0];
    while(rn >= 2 && geq((r[rn - 1] - r[rn - 2]) % (q - r[rn -
    \leftrightarrow 2]), 0)) // (2) ge
      --rn;
   r[rn++] = q;
 r.resize(rn);
 return r;
```

# 29 GeometryBase

```
const dbl EPS = 1e-9;
const int PREC = 20:
inline bool eq(dbl a, dbl b) { return abs(a-b)<=EPS; }</pre>
inline bool gr(dbl a, dbl b) { return a>b+EPS; }
inline bool geq(dbl a, dbl b) { return a>=b-EPS; }
inline bool ls(dbl a, dbl b) { return a < b - EPS; }</pre>
inline bool leq(dbl a, dbl b) { return a<=b+EPS; }</pre>
struct Pnt {
    dbl x,y;
    Pnt(): x(0), y(0) {}
    Pnt(dbl xx, dbl yy): x(xx), y(yy) {}
    inline Pnt operator +(const Pnt &p) const { return Pnt(x +
    \rightarrow p.x, y + p.y); }
    inline Pnt operator -(const Pnt &p) const { return Pnt(x -
     \rightarrow p.x, y - p.y); }
    inline dbl operator *(const Pnt &p) const { return x * p.x + y
    \rightarrow * p.y; } // ll
    inline dbl operator %(const Pnt &p) const { return x * p.y - y
    \hookrightarrow * p.x; } // ll
    inline Pnt operator *(dbl k) const { return Pnt(x * k, y * k);
    inline Pnt operator /(dbl k) const { return Pnt(x / k, y / k);
     → }
    inline Pnt operator -() const { return Pnt(-x, -y); }
    inline void operator +=(const Pnt &p) { x += p.x, y += p.y; }
    inline void operator -=(const Pnt &p) { x -= p.x, y -= p.y; }
    inline void operator *=(dbl k) { x*=k, y*=k; }
    inline bool operator ==(const Pnt &p) const { return
    \ \hookrightarrow \ \ abs(x-p.x) <= EPS \ \&\& \ abs(y-p.y) <= EPS; \ \}
    inline bool operator !=(const Pnt &p) const { return
    \ \hookrightarrow \ \ abs(x-p.x)>EPS \ |\ | \ abs(y-p.y)>EPS; \ \}
    inline bool operator <(const Pnt &p) const { return
    \rightarrow abs(x-p.x)<=EPS ? y<p.y-EPS : x<p.x; }
    inline dbl angle() const { return atan2(y, x); } // \mathit{ld}
    inline dbl len2() const { return x*x+y*y; } // ll
    inline dbl len() const { return sqrt(x*x+y*y); } // ll, ld
    inline Pnt getNorm() const {
        auto 1 = len();
        return Pnt(x/1, y/1);
    }
    inline void normalize() {
        auto 1 = len();
        x/=1, y/=1;
    inline Pnt getRot90() const { //counter-clockwise
        return Pnt(-y, x);
    inline Pnt getRot(dbl a) const { // ld
        dbl si = sin(a), co = cos(a);
        return Pnt(x*co - y*si, x*si + y*co);
    inline void read() {
        int xx, yy;
    cin >> xx >> yy;
        x = xx, y = yy;
    }
    inline void write() const{
         cout << fixed << (double)x << " " << (double)y << 'n';
}:
struct Line{
    dbl a, b, c;
    Line(): a(0), b(0), c(0) {}
    // normalizes
    Line(dbl aa, dbl bb, dbl cc) {
      dbl norm = sqrt(aa * aa + bb * bb);
      aa /= norm, bb /= norm, cc /= norm;
      a = aa, b = bb, c = cc;
    }
```

```
SPb HSE (Labutin, Podguzov, Bogomolov)
    Line(const Pnt &A, const Pnt &p){ // it normalizes (a,b),
    \rightarrow important in d(), normalToP()
       Pnt n = (p-A).getRot90().getNorm();
       a = n.x, b = n.y, c = -(a * A.x + b * A.y);
    inline dbl d(const Pnt &p) const { return a*p.x + b*p.y + c; }
    inline Pnt no() const {return Pnt(a, b);}
    inline Pnt normalToP(const Pnt &p) const { return Pnt(a,b) *
    \rightarrow (a*p.x + b*p.y + c); }
    inline void write() const{
     cout << fixed << (double)a << " " << (double)b << " " <<
      GeometryInterTangent
inline dbl sqr(dbl x) { return x * x; }
void buildTangent(Pnt p1, dbl r1, Pnt p2, dbl r2, Line &1) { //
   Pnt p = p2 - p1;
    1.c = r1;
```

```
\hookrightarrow r1, r2 = radius with sign
    dbl c2 = p.len2(), c1 = sqrt(c2 - sqr(r2));
    1.a = (-p.x * (r1 - r2) + p.y * c1) / c2;
    1.b = (-p.y * (r1 - r2) - p.x * c1) / c2;
    1.c -= 1.no() * p1;
    assert(eq(l.d(p1), r1));
    assert(eq(1.d(p2), r2));
struct Circle {
    Pnt p;
    dbl r:
vector<Pnt> v; // to store intersection
// Intersection of two lines
int line_line(const Line &1, const Line &m){
    dbl z = m.a * 1.b - 1.a * m.b;
  dbl x = m.c * l.b - l.c * m.b;
  dbl y = m.c * 1.a - 1.c * m.a;
    if(fabs(z) > EPS){
        v.pb(Pnt(-x/z, y/z));
        return 1;
    }else if(fabs(x) > EPS || fabs(y) > EPS)
       return 0; // parallel lines
    else
        return 2; // same lines
// Intersection of Circle and line
int circle_line(const Circle &c, const Line &l){
    dbl d = 1.d(c.p);
    if(fabs(d) > c.r + EPS)
        return 0;
    if(fabs(fabs(d) / c.r - 1) < EPS) {
        v.pb(c.p - 1.no() * d);
        return 1;
    } else {
        dbl s = sqrt(fabs(sqr(c.r) - sqr(d)));
        v.pb(c.p - 1.no() * d + 1.no().getRot90() * s);
        v.pb(c.p - 1.no() * d - 1.no().getRot90() * s);
        return 2;
}
// Intersection of two circles, 3 = inf
int circle_circle(const Circle &a, const Circle &b) {
  if (a.p == b.p \&\& eq(a.r, b.r))
    return 3;
  Pnt diff = b.p - a.p;
  dbl dist = diff.len();
  if (ls(a.r + dist, b.r) || ls(b.r + dist, a.r))
    return 0:
```

```
Line line(diff.x * 2, diff.y * 2, a.p.len2() - b.p.len2() +
    \hookrightarrow sqr(b.r) - sqr(a.r));
    return circle_line(a, line);
// Squared distance between point p and segment [a..b]
dbl dist2(Pnt p, Pnt a, Pnt b){
        if ((p - a) * (b - a) < 0) return (p - a).len2();
if ((p - b) * (a - b) < 0) return (p - b).len2();
        dbl d = fabs((p - a) \% (b - a));
        return d * d / (b - a).len2();
31 GeometrySimple
int sign(dbl a) { return (a > EPS) - (a < -EPS); }</pre>
// Checks, if point is inside the segment
inline bool inSeg(const Pnt &p, const Pnt &a, const Pnt &b) {
        return eq((p - a) \% (p - b), 0) && leq((p - a) * (p - b), 0);
// Checks, if two intervals (segments without ends) intersect AND
\hookrightarrow do not lie on the same line
inline bool subIntr(const Pnt &a, const Pnt &b, const Pnt &c,
\hookrightarrow const Pnt &d){
        return
                         sign((b - a) \% (c - a)) * sign((b - a) \% (d - a)) ==
                           sign((d - c) \% (a - c)) * sign((d - c) \% (b - c)) ==
                          }
// Checks, if two seaments (ends are included) has an intersection
inline bool checkSegInter(const Pnt &a, const Pnt &b, const Pnt
\hookrightarrow &c, const Pnt &d){
         return inSeg(c, a, b) || inSeg(d, a, b) || inSeg(a, c, d) ||
         \hookrightarrow inSeg(b, c, d) || subIntr(a, b, c, d);
inline dbl area(vector<Pnt> p){
        dbls = 0;
        int n = sz(p);
        p.pb(p[0]);
        forn(i, n)
                s += p[i + 1] \% p[i];
        p.pop_back();
        return abs(s) / 2;
}
// Check if point p is inside polygon <n, q[]>
int containsSlow(Pnt p, Pnt *z, int n){
        int cnt = 0;
        forn(j, n){
                Pnt a = z[j], b = z[(j + 1) \% n];
                 if (inSeg(p, a, b))
                         return -1; // border
                 if (min(a.y, b.y) - EPS \le p.y \&\& p.y \le max(a.y, b.y) -
                         cnt += (p.x < a.x + (p.y - a.y) * (b.x - a.x) / (b.y
                          \rightarrow -a.y));
        }
        return cnt & 1; // O = outside, 1 = inside
}
//for convex polygon
//assume polygon is counterclockwise-ordered
bool containsFast(Pnt p, Pnt *z, int n) {
        Pnt o = z[0];
        if(gr((p - o) \% (z[1] - o), 0) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 1] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o)) || ls((p - o) \% (z[n - 0] - o
         \rightarrow o), 0))
                return 0;
        int 1 = 0, r = n - 1;
        while(r - 1 > 1){
                 int m = (1 + r) / 2;
                 if(gr((p - o) \% (z[m] - o), 0))
                         r = m;
                 else
```

1 = m:

```
SPb HSE (Labutin, Podguzov, Bogomolov)
                                                                   struct Edge {
    return leq((p - z[1]) % (z[r] - z[1]), 0);
// Checks, if point "p" is in the triangle "abc" IFF triangle in
inline int isInTr(const Pnt &p, const Pnt &a, const Pnt &b, const
→ Pnt &c){
    return
            gr((b - a) % (p - a), 0) &&
            gr((c - b) % (p - b), 0) &&
            gr((a - c) \% (p - c), 0);
}
                                                                     }
                                                                   }
    Graphs
    2-SAT
                                                                       return;
32
// MAXVAR - 2 * vars
                                                                     st.pb(v);
int cntVar = 0, val[MAXVAR], usedSat[MAXVAR], comp[MAXVAR];
vi topsortSat;
vi g[MAXVAR], rg[MAXVAR];
inline int newVar() {
 cntVar++;
 return (cntVar - 1) * 2;
inline int Not(int v) { return v ^ 1; }
inline void Implies(int v1, int v2) { g[v1].pb(v2),
                                                                   }
\rightarrow rg[v2].pb(v1); }
inline void Or(int v1, int v2) { Implies(Not(v1), v2),
inline void Nand(int v1, int v2) { Or(Not(v1), Not(v2)); }
inline void setTrue(int v) { Implies(Not(v), v); }
                                                                   }
void dfs1(int v) {
  usedSat[v] = 1;
  for (int to : g[v])
    if (!usedSat[to]) dfs1(to);
  topsortSat.pb(v);
                                                                   vi st:
void dfs2(int v, int c) {
  comp[v] = c;
 for (int to : rg[v])
    if (!comp[to]) dfs2(to, c);
                                                                   };
int getVal(int v) { return val[v]; }
// cntVar
bool solveSat() {
  forn(i, 2 * cntVar) usedSat[i] = 0;
  forn(i, 2 * cntVar)
   if (!usedSat[i]) dfs1(i);
  reverse(all(topsortSat));
  int c = 0;
  for (int v : topsortSat)
    if (!comp[v]) dfs2(v, ++c);
  forn(i, cntVar) {
   if (comp[2 * i] == comp[2 * i + 1]) return false;
    if (comp[2 * i] < comp[2 * i + 1]) val[2 * i + 1] = 1;
   else val[2 * i] = 1;
 return true;
                                                                             }
                                                                           }
33 Bridges
int up[MAX_N], tIn[MAX_N], timer;
```

```
vector<vi>vectors;
```

```
int to, id;
  Edge(int _to, int _id) : to(_to), id(_id) {}
vector<Edge> g[MAX_N];
void newComp(int size = 0) {
  comps.emplace_back(); //
  while (sz(st) > size) {
    comps.back().pb(st.back());
    st.pop_back();
void findBridges(int v, int parentEdge = -1) {
 if (up[v]) //
 up[v] = tIn[v] = ++timer;
 for (Edge e : g[v]) {
   if (e.id == parentEdge)
     continue:
    int u = e.to;
   if (!tIn[u]) {
      int size = sz(st);
      findBridges(u, e.id);
     if (up[u] > tIn[v])
       newComp(size);
    up[v] = min(up[v], up[u]);
          find_bridges newComp()
void run(int n) {
 forn (i, n)
   if (!up[i]) {
      findBridges(i);
      newComp();
34 Cut Points
bool used[MAX_M];
int tIn[MAX_N], timer, isCut[MAX_N], color[MAX_M], compCnt;
struct Edge {
 int to, id;
 Edge(int _to, int _id) : to(_to), id(_id) {}
vector<Edge> g[MAX_N];
int dfs(int v, int parent = -1) {
 tIn[v] = ++timer;
  int up = tIn[v], x = 0, y = (parent != -1);
  for (Edge p : g[v]) {
   int u = p.to, id = p.id;
   if (id != parent) {
     int t, size = sz(st);
      if (!used[id])
       used[id] = 1, st.push_back(id);
      if (!tIn[u]) { // not visited yet
       t = dfs(u, id);
        if (t >= tIn[v]) {
         ++x, ++compCnt;
          while (sz(st) != size) {
            color[st.back()] = compCnt;
            st.pop_back();
      } else
       t = tIn[u]:
      up = min(up, t);
   }
  if (x + y >= 2)
```

if (root->r) root->r->p = root;

```
isCut[v] = 1; // v is cut vertex
                                                                          a = root, a \rightarrow p = a;
  return up;
                                                                          a->p = a;
                                                                        recalc(root);
     Eulerian Cycle
struct Edge {
                                                                      inline Node* rotate(Node* root, int k) {
                                                                       if (k == 0) return root;
  int to, used;
                                                                        Node *1, *r;
  Edge(): to(-1), used(0) {}
                                                                        split(root, 1, r, k);
  Edge(int v): to(v), used(0) {}
                                                                        return merge(r, 1);
vector<Edge> edges;
                                                                      inline pair<Node*, int> goUp(Node* root) {
vi g[MAX_N], res, ptr;
                                                                        int pos = getSize(root->1);
              ptr
                                                                        while (root->p != root)
                                                                          pos += (root->p->r == root ? getSize(root->p->l) + 1 : 0),
void dfs(int v) {
  for(; ptr[v] < sz(g[v]);) {</pre>
                                                                          \hookrightarrow root = root->p;
                                                                        return mp(root, pos);
    int id = g[v][ptr[v]++];
    if (!edges[id].used) {
      edges[id].used = edges[id ^ 1].used = 1;
                                                                      inline Node* deleteFirst(Node* root) {
      dfs(edges[id].to);
                                                                        Node* a:
      res.pb(id); //
    }
                                                                        split(root, a, root, 1);
                                                                        edges[a->e.v].erase(mp(a->e.u, a));
  }
 res.pb(v); // res
                                                                        return root;
                                                                      inline Node* getNode(int v, int u) {
    Euler Tour Tree
                                                                       return edges[v].lower_bound(mp(u, nullptr))->snd;
mt19937 rng(239);
                                                                      inline void cut(int v, int u) {
struct Edge {
                                                                        auto pV = goUp(getNode(v, u));
   int v. u:
                                                                        auto pU = goUp(getNode(u, v));
   Edge(int _v, int _u): v(_v), u(_u) {}
                                                                        int 1 = min(pV.snd, pU.snd), r = max(pV.snd, pU.snd);
                                                                        Node *a, *b, *c;
                                                                        split(pV.fst, a, b, 1);
struct Node {
                                                                        split(b, b, c, r - 1);
  Node *1, *r, *p;
                                                                        deleteFirst(b);
  Edge e;
                                                                        merge(a, deleteFirst(c));
  int y, size;
  Node(Edge _e): l(nullptr), r(nullptr), p(this), e(_e), y(rng()),
  \rightarrow size(1) {}
                                                                      inline pair<Node*, int> getRoot(int v) {
                                                                        return !sz(edges[v]) ? mp(nullptr, 0) :

    goUp(edges[v].begin()->snd);

inline int getSize(Node* root) { return root ? root->size : 0; }
inline void recalc(Node* root) { root->size = getSize(root->1) +
                                                                      inline Node* makeRoot(int v) {

    getSize(root->r) + 1; }

                                                                        auto root = getRoot(v);
                                                                        return rotate(root.fst, root.snd);
set<pair<int, Node*>> edges[MAX_N];
Node* merge(Node *a, Node *b) {
                                                                      inline Node* makeEdge(int v, int u) {
  if (!a) return b;
                                                                        Node* e = new Node(Edge(v, u));
  if (!b) return a;
                                                                        edges[v].insert(mp(u, e));
  if (a->y < b->y) {
                                                                       return e;
    a->r = merge(a->r, b);
    if (a->r) a->r->p = a;
    recalc(a):
                                                                      inline void link(int v, int u) {
    return a;
                                                                        Node *vN = makeRoot(v), *uN = makeRoot(u);
                                                                        merge(merge(vN, makeEdge(v, u)), uN), makeEdge(u, v));
  b->1 = merge(a, b->1);
  if (b->1) b->1->p = b;
  recalc(b);
  return b;
                                                                      37
                                                                          Hamilton Cycle
                                                                                             n*2^n
void split(Node *root, Node *&a, Node *&b, int size) {
                                                                      vi g[MAX_MASK];
                                                                      int adj[MAX_MASK], dp[1 << MAX_MASK];</pre>
  if (!root) {
    a = b = nullptr;
    return;
                                                                      vi hamiltonCycle(int n) {
                                                                        fill(dp, dp + (1 << n), 0);
  int lSize = getSize(root->1);
                                                                        forn (v, n) {
  if (lSize >= size) {
                                                                          adj[v] = 0;
                                                                          for (int to : g[v])
    split(root->1, a, root->1, size);
    if (root->1) root->1->p = root;
                                                                            adj[v] |= (1 << to);
    b = root, b->p = b;
  } else {
                                                                        dp[1] = 1;
                                                                        forn (mask, (1 \ll n))
    split(root->r, root->r, b, size - lSize - 1);
```

forn(v, n)

if (used[v]) return false;

```
if (mask & (1 << v) && dp[mask ^ (1 << v)] & adj[v])
                                                                      used[v] = 1;
        dp[mask] = (1 << v);
                                                                      for (int to : g[v])
                                                                        if (paired[to] == -1 || dfs(paired[to])) {
  vi ans;
  int mask = (1 << n) - 1, v;</pre>
                                                                          paired[to] = v, paired[v] = to;
  if (dp[mask] & adj[0]) {
                                                                          return true:
                                                                        }
    forab (i, 1, n)
      if ((1 << i) & (mask & adj[0]))
                                                                      return false;
       v = i;
    ans.pb(v);
   mask ^= (1 << v);
                                                                    int kuhn() {
    while(v) {
                                                                      int ans = 0;
                                                                      forn (i, n + m) paired[i] = -1;
     forn(i, n)
        if ((dp[mask] & (1 << i)) && (adj[i] & (1 << v))) {
                                                                      for (int run = 1; run;) {
                                                                       run = 0:
         v = i:
                                                                        fill(used, used + n + m, 0);
         break;
       }
                                                                        forn(i, n)
     mask ^= (1 << v);
                                                                          if (!used[i] && paired[i] == -1 && dfs(i))
      ans.pb(v);
                                                                            ans++, run = 1;
                                                                      }
 }
                                                                      return ans;
 return ans:
                                                                    // Max
                                                                                  -- A+, B-
     Karp with cycle
                                                                    // Min
                                                                                    -- A-, B+
int d[MAX_N][MAX_N], p[MAX_N][MAX_N];
vi g[MAX_N], ans;
                                                                    vi minCover, maxIndependent;
                                                                    void dfsCoverIndependent(int v) {
struct Edge {
                                                                      if (used[v]) return;
 int a, b, w;
                                                                      used[v] = 1;
 Edge(int _a, int _b, int _w): a(_a), b(_b), w(_w) {}
                                                                      for (int to : g[v])
                                                                        if (!used[to])
                                                                          used[to] = 1, dfsCoverIndependent(paired[to]);
vector<Edge> edges;
void fordBellman(int s, int n) {
 forn (i, n + 1)
   forn (j, n + 1)
                                                                    void findCoverIndependent() {
                                                                      fill(used, used + n + m, 0);
     d[i][j] = INF;
                                                                      forn (i, n)
 d[0][s] = 0;
                                                                        if (paired[i] == -1)
 forab (i, 1, n + 1)
    for (auto &e : edges)
                                                                          dfsCoverIndependent(i);
      forn (i, n)
        d[i][e.b] = d[i - 1][e.a] + e.w, p[i][e.b] = e.a;
                                                                        if (used[i]) maxIndependent.pb(i);
                                                                        else minCover.pb(i);
                                                                      forab (i, n, n + m)
ld karp(int n) {
                                                                        if (used[i]) minCover.pb(i);
                                                                        else maxIndependent.pb(i);
 int s = n++;
                                                                    }
  forn (i, n - 1)
    g[s].pb(sz(edges)), edges.pb(Edge(s, i, 0));
  fordBellman(s, n);
                                                                    40 LCA
 ld ansValue = INF;
  int curV = -1, dist = -1;
                                                                    int tin[MAX_N], tout[MAX_N], up[MAX_N][MAX_LOG];
  forn (v, n - 1)
                                                                    vi g[MAX N];
    if (d[n][v] != INF) {
                                                                    int curTime = 0;
     ld curAns = -INF;
     int curPos = -1:
                                                                    void dfs(int v, int p) {
      forn(k, n)
                                                                      up[v][0] = p;
       if (curAns \leftarrow (d[n][v] - d[k][v]) * (ld) (1) / (n - k))
                                                                      forn (i, MAX_LOG - 1)
         curAns = (d[n][v] - d[k][v]) * (ld) (1) / (n - k),
                                                                        up[v][i + 1] = up[up[v][i]][i];
          \hookrightarrow curPos = k;
                                                                      tin[v] = curTime++;
      if (ansValue > curAns)
                                                                      for (int u : g[v])
        ansValue = curAns, dist = curPos, curV = v;
                                                                        if (u != p)
                                                                          dfs(u, v);
  if (curV == -1) return ansValue;
                                                                      tout[v] = curTime++;
  for (int iter = n; iter != dist; iter--)
    ans.pb(curV), curV = p[iter][curV];
 reverse(all(ans));
                                                                    int isUpper(int v, int u) {
 return ansValue;
                                                                      return tin[v] <= tin[u] && tout[v] >= tout[u];
                                                                    int lca(int v, int u) {
39
     Kuhn's algorithm
                                                                      if (isUpper(u, v)) return u;
//
                                                                      fornr (i, MAX_LOG)
                                                                        if (!isUpper(up[u][i], v))
int n, m, paired[2 * MAX_N], used[2 * MAX_N];
                                                                          u = up[u][i];
vi g[MAX_N];
                                                                      return up[u][0];
bool dfs(int v) {
```

void init() {

```
dfs(0, 0);
    LCA offline (Tarjan)
vi g[MAX_N], q[MAX_N];
int pr[MAX_N], ancestor[MAX_N], used[MAX_N];
int get(int v) {
 return v == pr[v] ? v : pr[v] = get(pr[v]);
void unite(int v, int u, int anc) {
 v = get(v), u = get(u);
 pr[u] = v, ancestor[v] = anc;
void dfs(int v) {
 used[v] = 1;
 for (int u : g[v])
   if (!used[u])
     dfs(u), unite(v, u, v);
  for (int u : q[v])
   if (used[u])
      ancestor[get(u)]; // handle answer somehow
void init(int n) {
 forn (i, n) pr[i] = i, ancestor[i] = i;
 dfs(0);
42 2 Chinese
struct Edge {
   int fr, to, w, id;
    bool operator < (const Edge& o) const { return w < o.w; }
// find oriented mst (tree)
// there are no edge --> root (root is 0)
// 0 .. n - 1, weights and vertices will be changed, but ids are
\hookrightarrow ok
vector<Edge> work(const vector<vector<Edge>>& graph) {
    int n = (int) graph.size();
    vector<int> color(n), used(n, -1);
    for (int i = 0; i < n; i++)
       color[i] = i;
   vector<Edge> e(n);
    for (int i = 0; i < n; i++) {
        if (graph[i].empty()) {
           e[i] = \{-1, -1, -1, -1\};
       } else {
            e[i] = *min_element(graph[i].begin(),

    graph[i].end());

    }
    vector<vector<int>> cycles;
    used[0] = -2;
    for (int s = 0; s < n; s++) {
       if (used[s] != -1)
           continue;
        int x = s;
        while (used[x] == -1) {
           used[x] = s;
            x = e[x].fr;
       }
        if (used[x] != s)
            continue:
        vector<int> cycle = {x};
        for (int y = e[x].fr; y != x; y = e[y].fr)
            cycle.push_back(y), color[y] = x;
        cycles.push_back(cycle);
    }
    if (cycles.empty())
       return e:
    vector<vector<Edge>> next_graph(n);
    for (int s = 0; s < n; s++) {
        for (const Edge& edge : graph[s]) {
           if (color[edge.fr] != color[s])
```

```
next_graph[color[s]].push_back({
                color[edge.fr], color[s], edge.w - e[s].w,

→ edge.id

            });
    }
}
vector<Edge> tree = work(next_graph);
for (const auto& cycle : cycles) {
   int cl = color[cycle[0]];
   Edge next_out = tree[cl], out{};
    int from = -1;
    for (int v : cycle) {
        tree[v] = e[v];
        for (const Edge& edge : graph[v])
            if (edge.id == next_out.id)
                from = v, out = edge;
    tree[from] = out;
}
return tree;
 Math
```

}

# 43 CRT (KTO)

```
vi crt(vi a, vi mod) {
 int n = sz(a):
  vi x(n);
  forn (i, n) {
    x[i] = a[i];
    forn (j, i) {
      x[i] = inverse(mod[j], mod[i]) * (x[i] - x[j]) % mod[i];
      if (x[i] < 0) x[i] += mod[i];</pre>
   }
 }
 return x;
```

# 44 Discrete Logarithm

```
// Returns x: a^x = b \pmod{mod} or -1, if no such x exists
int discreteLogarithm(int a, int b, int mod) {
  int sq = sqrt(mod);
  int sq2 = mod / sq + (mod % sq ? 1 : 0);
  vector<pii> powers(sq2);
  forn (i, sq2)
   powers[i] = mp(power(a, (i + 1) * sq, mod), i + 1);
  sort(all(powers));
  forn (i, sq + 1) {
   int cur = power(a, i, mod);
    cur = (cur * 111 * b) % mod;
   auto it = lower_bound(all(powers), mp(cur, 0));
    if (it != powers.end() && it->fst == cur)
      return it->snd * sq - i;
  }
  return -1;
}
```

# 45 Discrete Root

```
// Returns x: x^k = a \mod mod, mod is prime
int discreteRoot(int a, int k, int mod) {
 if (a == 0)
   return 0;
 int g = primitiveRoot(mod);
 int y = discreteLogarithm(power(g, k, mod), a, mod);
 return power(g, y, mod);
```

# 46 Eratosthenes

```
vi eratosthenes(int n) {
 vi minDiv(n + 1, 0);
 minDiv[1] = 1;
 forab (i, 2, n + 1)
   if (minDiv[i] == 0)
     for (int j = i; j \le n; j += i)
       if (minDiv[j] == 0) minDiv[j] = i;
```

vi par(m, -1);

```
return minDiv;
                                                                       for (col = 0; col < m && row < n; col++) {
                                                                         int best = row;
                                                                         for (int i = row; i < n; i++)</pre>
vi eratosthenesLinear(int n) {
                                                                           if (a[i][col] > a[best][col])
 vi minDiv(n + 1, 0), primes;
                                                                             best = i:
                                                                         if (a[best][col] == 0)
 minDiv[1] = 1;
 forab (i, 2, n + 1) {
                                                                           continue;
   if (minDiv[i] == 0)
                                                                         par[col] = row;
     minDiv[i] = i, primes.pb(i);
                                                                         swap(a[row], a[best]);
   for (int j = 0; j < sz(primes) && primes[j] <= minDiv[i] && i
                                                                         forn (i, n)
    \hookrightarrow * primes[j] <= n; j++)
                                                                           if (i != row && a[i][col])
     minDiv[i * primes[j]] = primes[j];
                                                                              a[i] ^= a[row];
                                                                       }
 return minDiv:
                                                                       vi ans(m, 0);
                                                                       forn (i, m)
                                                                         if (par[i] != -1)
    Factorial
                                                                           ans[i] = a[par[i]][n] / a[par[i]][i];
// Returns pair (rem, deg), where rem = n! % mod,
                                                                       bool ok = 1;
// deg = k: mod^k / n!, mod is prime, O(mod log mod)
                                                                       forn (i, n) {
                                                                         int cur = 0:
pii fact(int n, int mod) {
                                                                         forn (j, m) cur ^= (ans[j] & a[i][j]);
  int rem = 1, deg = 0, nCopy = n;
                                                                         if (cur != a[i][n]) ok = 0;
 while (nCopy) nCopy /= mod, deg += nCopy;
                                                                       }
 while (n > 1) {
   rem = (rem * ((n / mod) % 2 ? -1 : 1) + mod) % mod;
                                                                       return ok:
   for (int i = 2; i <= n % mod; i++)
     rem = (rem * 111 * i) % mod;
   n /= mod:
 }
                                                                     50 Gcd
 return mp(rem % mod, deg);
                                                                     int gcd(int a, int b) {
                                                                       return b ? gcd(b, a % b) : a;
48
     Gauss
const double EPS = 1e-9;
                                                                     int gcd(int a, int b, int &x, int &y) {
                                                                       if (b == 0) {
int gauss(double **a, int n, int m) { // n is number of equations,
                                                                        x = 1, y = 0;
\rightarrow m is number of variables
                                                                         return a;
 int row = 0, col = 0;
 vi par(m, −1);
                                                                       int g = gcd(b, a % b, x, y), newX = y;
 vector<double> ans(m, 0);
                                                                       y = x - a / b * y;
 for (col = 0; col < m && row < n; col++) {
                                                                       x = newX;
    int best = row;
                                                                       return g;
    for (int i = row; i < n; i++)
     if (abs(a[i][col]) > abs(a[best][col]))
       best = i:
                                                                     void diophant(int a, int b, int c, int &x, int &y) {
    if (abs(a[best][col]) < EPS) continue;</pre>
                                                                       int g = gcd(a, b, x, y);
    par[col] = row;
                                                                       if (c % g != 0) return;
    forn (i, m + 1) swap(a[row][i], a[best][i]);
                                                                       x *= c / g, y *= c / g;
    forn (i, n)
                                                                       // next solutions: x += b / g, y -= a / g
     if (i != row) {
       double k = a[i][col] / a[row][col];
       for (int j = col; j <= m; j++)</pre>
                                                                     int inverse(int a, int mod) { // Returns -1, if a and mod are not
          a[i][j] -= k * a[row][j];
                                                                     }
                                                                      int x, y;
   row++;
                                                                       int g = gcd(a, mod, x, y);
 }
                                                                       return g == 1 ? (x \% mod + mod) \% mod : -1;
  int single = 1;
  forn (i, m)
    if (par[i] != -1) ans[i] = a[par[i]][m] / a[par[i]][i];
                                                                     vi inverseForAll(int mod) {
    else single = 0;
                                                                       vi r(mod, 0);
 forn (i, n) {
                                                                       r[1] = 1:
    double cur = 0;
                                                                       for (int i = 2; i < mod; i++)
   for (int j = 0; j < m; j++)
                                                                         r[i] = (mod - r[mod \% i]) * (mod / i) \% mod;
     cur += ans[j] * a[i][j];
                                                                       return r:
    if (abs(cur - a[i][m]) > EPS)
     return 0;
 if (!single)
                                                                     51 Gray
   return 2:
 return 1:
                                                                     int gray(int n) {
                                                                       return n ^ (n >> 1);
49
     Gauss binary
const int MAX = 1024;
                                                                     int revGray(int n) {
                                                                       int k = 0;
                                                                      for (; n; n >>= 1) k ^= n;
int gaussBinary(vector<bitset<MAX>> a, int n, int m) {
 int row = 0, col = 0;
                                                                       return k;
```

}

```
SPb HSE (Labutin, Podguzov, Bogomolov)
     Miller-Rabin Test
52
vector <int> primes = {2, 3, 5, 7, 11, 13, 17, 19, 23};
bool isPrimeMillerRabin(ll n) {
 int k = 0:
  11 t = n - 1;
  while (t \% 2 == 0) k++, t /= 2;
  for (auto p : primes) {
    11 g = \_gcd(n, (11) p);
    if (g > 1 && g < n) return 0;
    if (g == n) return 1;
    ll b = powerLL(p, t, n), last = n - 1;
    bool was = 0;
    forn (i, k + 1) {
     if (b == 1 && last != n - 1)
       return 0;
      if (b == 1) {
       was = 1;
        break:
     }
     last = b, b = mul(b, b, n);
   if (!was) return 0;
  }
 return 1;
53 Phi
int phi(int n) {
  int result = n;
  for (int i = 2; i * i <= n; i++)
   if (n % i == 0) {
      while (n \% i == 0) n /= i;
      result -= result / i;
   }
  if (n > 1) result -= result / n;
  return result;
int inversePhi(int a, int mod) {
 return power(a, phi(mod) - 1, mod);
54 Pollard
inline void pollardFoo(ll& x, ll mod) {
 x = (mul(x, x, mod) + 1) \% mod;
vector<pair<11, int>> factorize(11 n) {
  if (n == 1) return {};
  if (isPrimeMillerRabin(n)) return {mp(n, 1)};
  if (n <= 100) {
   vector<pair<ll, int>> ans;
    for (int i = 2; i * i <= n; i++)
      if (n % i == 0) {
        int cnt = 0;
        while (n \% i == 0) n /= i, cnt++;
        ans.pb(mp(i, cnt));
      }
    if (n != 1) ans.pb(mp(n, 1));
    sort(all(ans));
   return ans;
  }
  while (1) {
    11 a = rand() \% n, b = a;
    while (1) {
```

pollardFoo(a, n), pollardFoo(b, n), pollardFoo(b, n);

 $11 g = \_gcd(abs(a-b), n);$ 

auto ans1 = factorize(g);

vector<pair<11, int>> ans;

auto ans2 = factorize(n / g);

ans1.insert(ans1.end(), all(ans2));

if (g != 1) {

break;

if (g == n)

sort(all(ans1));

for (auto np : ans1)

```
if (sz(ans) == 0 || np.fst != ans.back().fst)
           ans.pb(np);
         else
           ans.back().snd += np.snd;
       return ans:
     }
   }
 }
 assert(0);
55 Power And Mul
```

```
inline ll fix(ll a, ll mod) { // a in [0, 2 * mod)
  if (a \ge mod) a -= mod;
  return a;
// Returns (a * b) % mod, 0 <= a < mod, 0 <= b < mod
11 mulSlow(11 a, 11 b, 11 mod) {
 if (!b) return 0;
  ll c = fix(mulSlow(a, b / 2, mod) * 2, mod);
  return b & 1 ? fix(c + a, mod) : c;
11 mul(11 a, 11 b, 11 mod) {
  11 q = (1d) a * b / mod;
  11 r = a * b - mod * q;
  while (r < 0) r += mod;
  while (r >= mod) r -= mod;
  return r:
int power(int a, int n, int mod) {
 if (!n) return 1;
  int b = power(a, n / 2, mod);
 b = (b * 111 * b) \% mod;
  return n & 1 ? (a * 111 * b) % mod : b;
11 powerLL(11 a, 11 n, 11 mod) {
 if (!n) return 1;
  11 b = powerLL(a, n / 2, mod);
  b = mul(b, b, mod);
  return n & 1 ? mul(a, b, mod) : b;
int powerFast(int a, int n, int mod) {
  int res = 1;
  while (n) {
   if (n & 1)
     res = (res * 111 * a) % mod;
    a = (a * 111 * a) % mod;
    n /= 2;
  }
  return res;
```

# **56** Primitive Root

```
int primitiveRoot(int mod) { // Returns -1 if no primitive root
\hookrightarrow exists
 vi fact;
 int ph = phi(mod);
  int n = mod;
  for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) {
      fact.pb(i);
      while (n \% i == 0) n /= i;
   }
  }
  if (n > 1) fact.pb(n);
  forab (i, 2, mod + 1) {
   bool ok = 1;
   for (int j = 0; j < sz(fact) && ok; <math>j++)
     ok &= power(i, ph / fact[j], mod) != 1;
    if (ok) return i;
```

```
57 Simpson
```

return -1:

# 10 Strings

# 58 Aho-Corasick

```
const int ALPHA = 26;
const int MAX_N = 1e5;
struct Node {
  int next[ALPHA], term; //
  int go[ALPHA], suf, p, pCh; //
  Node(): term(0), suf(-1), p(-1) {
   fill(next, next + ALPHA, -1);
    fill(go, go + ALPHA, -1);
 }
};
Node g[MAX_N];
int last;
void add(const string &s) {
 int now = 0:
  for(char x : s) {
    if (g[now].next[x - 'a'] == -1) {
      g[now].next[x - 'a'] = ++last;
      g[last].p = now, g[last].pCh = x;
   now = g[now].next[x - 'a'];
  }
  g[now].term = 1;
int go(int v, int c);
int getLink(int v) {
  if (g[v].suf == -1) {
   if (!v || !g[v].p) g[v].suf = 0;
    else g[v].suf = go(getLink(g[v].p), g[v].pCh);
  }
 return g[v].suf;
int go(int v, int c) {
  if (g[v].go[c] == -1) {
    if (g[v].next[c] != -1) g[v].go[c] = g[v].next[c];
    else g[v].go[c] = !v ? 0 : go(getLink(v), c);
  }
 return g[v].go[c];
```

# 59 Prefix-function

```
vi prefix(const string &s) {
  int n = sz(s);
  vi pr(n);
  forab (i, 1, n + 1) {
    int j = pr[i - 1];
    while (j > 0 && s[i] != s[j]) j = pr[j - 1];
    if (s[i] == s[j]) j++;
    pr[i] = j;
  }
  return pr;
}
```

### 60 Z-function

```
vi z(const string& s) {
   int n = sz(s);
   vi z(n);
   for (int i = 1, l = 0, r = 0; i < n; i++) {
      if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
      while (i + z[i] < n && s[z[i]] == s[i + z[i]]) z[i]++;
      if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
   }
   return z;
}
```

### 61 Hashes

```
const int P = 239017;
inline int add(int a, int b, int m) {
  a += b;
 return a >= m ? a - m : a;
inline int sub(int a, int b, int m) {
 a -= b:
 return a < 0? a + m : a;
const int MOD_X = 1e9 + 9, MOD_Y = 1e9 + 7;
// using H = unsigned long long;
struct H {
  int x, y;
 H(): x(0), y(0) {}
 H(int _x): x(_x), y(_x) {}
 H(int _x, int _y): x(_x), y(_y) {}
 inline H operator+(const H& h) const { return H(add(x, h.x,

    MOD_X), add(y, h.y, MOD_Y)); }

  inline H operator-(const H& h) const { return H(sub(x, h.x,
  \hookrightarrow MOD_X), sub(y, h.y, MOD_Y)); }
 inline H operator*(11 k) const { return H(int((x * k) % MOD_X),

    int((y * k) % MOD_Y)); }

 inline H operator*(const H& h) const{ return H(int((ll(x) * h.x)
  inline bool operator==(const H\& h) const { return x == h.x \&\& y
  \hookrightarrow == h.y; }
 inline bool operator!=(const H& h) const { return x != h.x || y
 inline bool operator<(const H& h) const { return x < h.x || (x
  \hookrightarrow == h.x && y < h.y); }
  explicit inline operator ll() const { return ll(x) * MOD_Y + y + y

→ 1; } // > 0

H deg[MAX_N], h[MAX_N];
inline H get(int 1, int r) { return h[r] - h[l] * deg[r - l]; }
void init(const string& s) {
 int n = sz(s);
  deg[0] = 1;
 forn (i, n)
   h[i + 1] = h[i] * P + s[i], deg[i + 1] = deg[i] * P;
```

# 62 Manaker

```
void manaker(const string& s, int *z0, int *z1) {
  int n = sz(s);
  forn (t, 2) {
    int *z = t ? z1 : z0, l = -1, r = -1; // [l..r]
    forn (i, n - t) {
      int k = 0;
      if (r > i + t) {
        int j = l + (r - i - t);
        k = min(z[j], j - l);
      }
    while (i - k >= 0 && i + k + t < n && s[i - k] == s[i + k + condition of the condition
```

```
SPb HSE (Labutin, Podguzov, Bogomolov)
      if (k \&\& i + k + t > r)
        l = i - k + 1, r = i + k + t - 1;
 }
    Palindromic Tree
const int ALPHA = 26;
struct Vertex {
  int suf, len, next[ALPHA];
  Vertex() { fill(next, next + ALPHA, 0); }
int vn, v;
Vertex t[MAX_N + 2];
int n, s[MAX_N];
int get(int i) { return i < 0 ? -1 : s[i]; }</pre>
void init() {
 t[0].len = -1, vn = 2, v = 0, n = 0;
void add(int ch) {
  s[n++] = ch;
  while (v != 0 && ch != get(n - t[v].len - 2))
   v = t[v].suf;
  int& r = t[v].next[ch];
  if (!r) {
    t[vn].len = t[v].len + 2;
    if (!v) t[vn].suf = 1;
     v = t[v].suf;
      while (v != 0 \&\& ch != get(n - t[v].len - 2))
       v = t[v].suf;
      t[vn].suf = t[v].next[ch];
    }
   r = vn++;
 }
 v = r;
    Suffix Array (+stable)
int sLen, num[MAX_N + 1];
char s[MAX_N + 1];
int p[MAX_N], col[MAX_N], inv[MAX_N], lcp[MAX_N];
inline int add(int a, int b) {
  return a >= sLen ? a - sLen : a;
inline int sub(int a, int b) {
 a -= b;
 return a < 0 ? a + sLen : a;
void buildArray(int n) {
  sLen = n;
  int ma = max(n, 256);
  forn (i, n)
   col[i] = s[i], p[i] = i;
  for (int k2 = 1; k2 / 2 < n; k2 *= 2) {
    int k = k2 / 2;
    memset(num, 0, sizeof(num));
    forn (i, n) num[col[i] + 1]++;
    forn (i, ma) num[i + 1] += num[i];
    forn (i, n)
     inv[num[col[sub(p[i], k)]]++] = sub(p[i], k);
    int cc = 0;
      bool flag = col[inv[i]] != col[inv[i - 1]];
      flag |= col[add(inv[i], k)] != col[add(inv[i - 1], k)];
     if (i && flag) cc++;
     num[inv[i]] = cc;
    7
```

```
forn (i, n) p[i] = inv[i], col[i] = num[i];
  memset(num, 0, sizeof(num));
 forn (i, n) num[col[i] + 1]++;
  forn (i, ma) num[i + 1] += num[i];
 forn (i, n) inv[num[col[i]]++] = i;
 forn (i, n) p[i] = inv[i];
  forn (i, n) inv[p[i]] = i;
void buildLCP(int n) {
 int len = 0;
 forn (ind, n){
   int i = inv[ind];
   len = max(0, len - 1);
   if (i != n - 1)
     while (len < n && s[add(p[i], len)] == s[add(p[i + 1],
      \rightarrow len)])
       len++;
   lcp[i] = len;
   if (i != n - 1 & p[i + 1] == n - 1) len = 0;
}
65 Suffix Automaton
struct Vx {
   static const int AL = 26;
   int len, suf;
   int next[AL];
    Vx() {}
    Vx(int 1, int s): len(1), suf(s) {}
}:
struct SA {
   static const int MAX_LEN = 1e5 + 100, MAX_V = 2 * MAX_LEN;
    int last, vcnt;
   Vx v[MAX_V];
   SA() { vcnt = 1, last = newV(0, 0); } // root = vertex with
    \hookrightarrow number 1
    int newV(int len, int suf){
       v[vcnt] = Vx(len, suf);
        return vcnt++;
    int add(char ch) {
       int p = last, c = ch - 'a';
        last = newV(v[last].len + 1, 0);
        while (p && !v[p].next[c]) // added p &&
           v[p].next[c] = last, p = v[p].suf;
        if (!p)
           v[last].suf = 1;
        else {
           int q = v[p].next[c];
            if (v[q].len == v[p].len + 1) v[last].suf = q;
                int r = newV(v[p].len + 1, v[q].suf);
                v[last].suf = v[q].suf = r;
                memcpy(v[r].next, v[q].next, sizeof(v[r].next));
                while (p \&\& v[p].next[c] == q)
                    v[p].next[c] = r, p = v[p].suf;
            }
       }
        return last;
   }
};
66 Suffix Tree
const int MAX L=1e5+10;
char S[MAX_L];
int L;
struct Node;
struct Pos:
typedef Node *pNode;
typedef map<char,pNode> mapt;
```

struct Node{

```
pNode P,link;
  int L,R;
  mapt next;
  Node():P(NULL).link(this).L(0).R(0){}
  Node(pNode P,int L,int R):P(P),link(NULL),L(L),R(R){}
  inline int elen() const{return R-L;}
  inline pNode add_edge(int L,int R){return next[S[L]]=new
  struct Pos{
  pNode V;
  int up;
  Pos():V(NULL),up(0){}
  Pos(pNode V,int up):V(V),up(up){}
  pNode split_edge() const{
    if(!up)
     return V:
    int L=V->L, M=V->R-up;
    pNode P=V->P, n=new Node(P,L,M);
    P->next[S[L]]=n;
   n->next[S[M]]=V:
    V->P=n, V->L=M;
   return n;
 Pos next_char(char c) const{
    if(up)
      return S[V\rightarrow R-up]==c ? Pos(V,up-1) : Pos();
      mapt::iterator it=V->next.find(c);
      return it==V->next.end() ? Pos() :
      → Pos(it->snd,it->snd->elen()-1);
   }
 }
};
Pos go_down(pNode V,int L,int R){
  if(L==R)
    return Pos(V,0);
  while(1){
    V=V->next[S[L]];
   L+=V->elen();
    if(L>=R)
      return Pos(V,L-R);
}
inline pNode calc_link(pNode &V){
  if(!V->link)
    V->link=go\_down(V->P->link,V->L+!V->P->P,V->R).split\_edge();
  return V->link;
Pos add_char(Pos P,int k){
  while(1){
    Pos p=P.next_char(S[k]);
    if(p.V)
     return p;
    pNode n=P.split_edge();
    n->add_edge(k,MAX_L);
    if(!n->P)
      return Pos(n,0);
    P=Pos(calc_link(n),0);
 }
pNode Root;
void make_tree(){
 Root=new Node();
  Pos P(Root, 0);
 forn(i,L)
    P=add_char(P,i);
```

#### C++ Tricks 11

inline int readInt() {

if (c == '-')

int s = 0, c = readChar(), x = 0;

s = 1, c = readChar();

```
67
     Fast allocation
const int MAX_MEM = 1e8;
int mpos = 0;
char mem[MAX_MEM];
inline void* operator new(size_t n) {
 char *res = mem + mpos;
  mpos += n;
  assert(mpos <= MAX_MEM);</pre>
  return (void*) res;
}
inline void operator delete(void*) {}
inline void* operator new[](size_t) { assert(0); }
inline void operator delete[](void*) { assert(0); }
68 Hash of pair
struct PairHasher {
  size_t operator()(const pair<int, int>& p) const { return p.fst
  Ordered Set
69
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
typedef tree<int, null_type, less<int>, rb_tree_tag,

    tree_order_statistics_node_update> ordered_set;
void example() {
  ordered_set X;
  X.insert(1);
  cout << *X.find_by_order(1) << " " << X.order_of_key(1) <<</pre>
     "\n";
}
70 Hash Map
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
struct chash { // To use most bits rather than just the lowest

→ ones:

  const uint64_t C = 11(2e18 * PI) + 71; // large odd number
  const int RANDOM = 912387491;
  11 operator()(11 x) const { return __builtin_bswap64((x ^
  \hookrightarrow RANDOM) * C); }
};
template<class K, class V> using ht = gp_hash_table<K, V, chash>;
template<class K, class V> V get(ht<K, V>& u, K x) {
 return u.find(x) == end(u) ? 0 : u[x];
}
ht<11, int> h({}, {}, {}, {}, {1<<20});
71 Fast I/O (short)
inline int readChar();
inline int readInt();
template <class T> inline void writeInt(T x);
inline int readChar() {
 int c = getchar();
  while (c <= 32)
    c = getchar();
 return c;
```

```
while ('0' \le c \&\& c \le '9')
   x = x * 10 + c - '0', c = readChar();
                                                                      template <class T>
 return s ? -x : x;
                                                                      inline T readInt() {
                                                                        int s = 1, c = readChar();
                                                                        T x = 0;
template <class T> inline void writeInt(T x) {
                                                                        if (c == '-')
                                                                          s = -1, c = getChar();
  if (x < 0)
   putchar('-'), x = -x;
                                                                        while ('^{0}' <= c && c <= '^{9}')
  char s[24];
                                                                          x = x * 10 + c - '0', c = getChar();
  int n = 0;
                                                                        return s == 1 ? x : -x;
  while (x \mid \mid !n)
   s[n++] = '0' + x \% 10, x /= 10;
  while (n--)
                                                                      inline double readDouble() {
                                                                        int s = 1. c = readChar():
    putchar(s[n]):
                                                                        double x = 0;
                                                                        if (c == '-')
                                                                         s = -1, c = getChar();
     Fast I/O (long)
                                                                        while ('0' <= c && c <= '9')
                                                                        x = x * 10 + c - '0', c = getChar();
template <class T = int> inline T readInt();
                                                                        if (c == '.') {
inline double readDouble();
                                                                         c = getChar();
inline int readUInt();
                                                                          double coef = 1;
inline int readChar();
                                                                          while ('0' <= c && c <= '9')
inline void readWord(char *s);
                                                                            x += (c - '0') * (coef *= 1e-1), c = getChar();
inline bool readLine(char *s); // do not save '\n'
inline bool isEof();
                                                                        return s == 1 ? x : -x;
inline int peekChar();
inline bool seekEof();
                                                                      inline void readWord(char *s) {
template <class T> inline void writeInt(T x, int len);
                                                                        int c = readChar();
template <class T> inline void writeUInt(T x, int len);
                                                                        while (c > 32)
template <class T> inline void writeInt(T x) { writeInt(x, -1); };
                                                                         *s++ = c, c = getChar();
template <class T> inline void writeUInt(T x) { writeUInt(x, -1);
→ };
inline void writeChar(int x);
inline void writeWord(const char *s);
                                                                      inline bool readLine(char *s) {
inline void writeDouble(double x, int len = 0);
                                                                        int c = getChar();
inline void flush();
                                                                        while (c != '\n' && c != -1)
                                                                         *s++ = c, c = getChar();
const int BUF_SIZE = 4096;
                                                                        *s = 0;
                                                                        return c != -1;
char buf[BUF_SIZE];
int bufLen = 0, pos = 0;
inline bool isEof() {
                                                                      int writePos = 0;
 if (pos == bufLen) {
                                                                      char writeBuf[BUF_SIZE];
   pos = 0, bufLen = fread(buf, 1, BUF_SIZE, stdin);
   if (pos == bufLen)
                                                                      inline void writeChar(int x) {
      return 1:
                                                                        if (writePos == BUF_SIZE)
 }
                                                                          fwrite(writeBuf, 1, BUF_SIZE, stdout), writePos = 0;
 return 0;
                                                                        writeBuf[writePos++] = x;
inline int getChar() {
                                                                      inline void flush() {
 return isEof() ? -1 : buf[pos++];
                                                                          fwrite(writeBuf, 1, writePos, stdout), writePos = 0;
inline int peekChar() {
 return isEof() ? -1 : buf[pos];
                                                                      template <class T>
                                                                      inline void writeInt(T x, int outputLen) {
                                                                       if (x < 0)
inline bool seekEof() {
                                                                          writeChar('-'), x = -x;
 int c;
 while ((c = peekChar()) != -1 && c <= 32)
                                                                        char s[24];
   pos++;
 return c == -1;
                                                                        int n = 0;
                                                                        while (x \mid | \mid n)
                                                                         s[n++] = '0' + x \% 10, x /= 10;
                                                                        while (n < outputLen)</pre>
inline int readChar() {
                                                                         s[n++] = '0';
 int c = getChar();
                                                                        while (n--)
  while (c !=-1 \&\& c <= 32)
                                                                          writeChar(s[n]);
   c = getChar();
 return c:
                                                                      template <class T>
                                                                      inline void writeUInt(T x, int outputLen) {
inline int readUInt() {
                                                                        char s[24];
 int c = readChar(), x = 0;
                                                                        int n = 0;
  while ('0' <= c && c <= '9')
   x = x * 10 + c - '0', c = getChar();
                                                                        while (x \mid | \mid n)
                                                                         s[n++] = '0' + char(x \% 10), x /= 10;
 return x;
                                                                        while (n < outputLen)</pre>
```

```
s[n++] = '0';
  while (n--)
    writeChar(s[n]);
inline void writeWord(const char *s) {
 while (*s)
    writeChar(*s++);
inline void writeDouble(double x, int outputLen) {
  if (x < 0)
    writeChar('-'), x = -x;
  int t = (int) x;
  writeUInt(t), x -= t;
  writeChar('.');
  for (int i = outputLen - 1; i > 0; i--) {
   x *= 10;
    t = std::min(9, (int) x);
    writeChar('0' + t), x = t;
 }
 x *= 10;
 t = std::min(9, (int)(x + 0.5));
  writeChar('0' + t);
```

#### **12 Notes**

#### **73** Работа с деревьями

Приемы для работы с деревьями:

- 1. Двоичные подъемы
- 2. Поддеревья как отрезки Эйлерова обхода
- 3. Вертикальные пути в Эйлеровом обходе (на ребрах вниз +k, на ребрах вверх
- 4. Храним в вершине значение функции на пути от корня до нее, дальше LCA.
- 5. Спуск с DFS, поддерживаем ДО на пути до текущей вершины.
- 6. Heavy-light decomposition
- 7. Centroid decomposition
- 8. Корневая по запросам
- 9. Тяжелые/легкие вершины
- 10. DFS  $\rightarrow$  дерево блоков, размеры  $\in [K..2K]$
- 11. У вершины не более  $O(\sqrt{N})$  разных поддеревьев
- 12. Сумма размеров поддеревьев без тяжелого ребенка  $O(n \log n)$
- 13. Сумма глубин поддеревьев без глубокого ребенка O(n)

#### Маски

Считаем динамику по маскам за  $O(2^n \cdot n)$  f[mask] = sum по submaskg[submask].

dp[mask][i]— значение динамики для маски mask, если младшие i бит в ней зафиксированы (то есть мы не можем удалять оттуда).

Ответ в dp[mask][0].

dp[mask][len] = g[mask]. Если i-ый бит 0, то dp[mask][i] = dp[mask][i+1], иначе dp[mask][i] = dp[mask][i+1] + dp[mask(1 << i)][i+1].

Старший бит: предподсчет.

Младший бит:  $x \& \sim (-x)$ 

Чтобы по степени двойки получить логарифм, можно воспользоваться тем, что все степени двойки имеют разный остаток по модулю 67.

```
for (int mask = 0; mask < (1 << n); mask++)</pre>
^{\text{lsupmask}}: for (int s = mask; s < (1 << n); s = (s + 1) | mask)
```

#### Гранди 75

Теорема Шпрага-Гранди: берем mex всех значений функции Гранди по состояниям, в которые можем перейти из данного.

Если сумма независимых игр, то значение функции Гранди равно хог значений функций Гранди по всем играм.

Бывает полезно вывести первые п значений и поискать закономерность.

Часто сводится к xor по чему-нибудь.

# 76 Потоки

Потоки:

Name	Asympthotic
Ford-Fulkerson	$O( f  \cdot E)$
Ford-Fulkerson with scaling	$O(\log  f  \cdot E^2)$
Edmonds-Karp	$O(V \cdot E^2)$
Dinic	$O(V^2 \cdot E)$
Dinic with scaling	$O(V \cdot E \cdot \log C)$
Dinic on bipartite graph	$O(E\sqrt{V})$
Dinic on unit network	$O(E\sqrt{E})$

#### L-R потоки:

Есть граф с недостатками или избытками в каждой вершине. Создаем фиктивные исток и сток (из истока все ребра в избытки, из недостатков все ребра в сток).

Теперь пусть у нас есть L-R граф, для каждого ребра  $e\ (v \to u)$  известны  $L_e$ и  $R_e$ . Добавим в v избыток  $L_e$ , в u недостаток  $L_e$ , а пропускную способность сделаем  $R_e - L_e$ .

Получили решение задачи о LR-циркуляции.

Если у нас обычный граф с истоком и стоком, то добавляем бесконечное ребро из стока в сток и ищем циркуляцию.

Таким образом нашли удовлетворяющий условиям LR-поток. Если хотим максимальный поток, то на остаточной сети запускаем поиск максимального потока.

В новом графе в прямую сторону пропускная способность равна  $R_e - f_e$ , в обратную  $f_e - L_e$ .

MinCostCirculation:

Пока есть шикл отрицательного веса, запускаем алгоритм Карпа и пускаем максимальный поток по найденному циклу.

#### ДП 77

Табличка с оптимизациями для линамики:

Name	Original recurrence	From
	Sufficient Condition	То
CHT1	$dp[i] = \min_{j < i} dp[j] + b[j] \cdot a[i]$	$O(n^2)$
	$ b[j] \geqslant b[j+1] \mid  a[i] \leqslant a[i+1]$	O(n)
CHT2	$dp[i][j] = \min_{k < j} dp[i-1][k] + b[k] \cdot a[j]$	$O(kn^2)$
	$b[k] \geqslant b[k+1] \mid\mid a[j] \leqslant a[j+1]$	O(kn)
D&C	$dp[i][j] = \min_{k < j} dp[i-1][k] + c[k][j]$	$O(kn^2)$
	$p[i,j] \leqslant p[i,j+1]$	$O(kn\log n)$
Knuth	$dp[i][j] = \min_{i < k < j} dp[i][k] + dp[k][j] + c[i][j]$	$O(n^3)$
	$p[i, j-1] \leqslant p[i, j] \leqslant p[i+1, j]$	$O(n^2)$
IOI	$f_n(k)$ — best for fixed k	$O(k^{(2)}n)$
	$f_n$ — convex, add penalty $\lambda \cdot k$	$O(n \log C)$

#### **78** Комбинаторика

Биномиальные коэффициенты:

Теорема Люка для биномиальных коэффициентов: Хотим посчитать  $C_n^k$ , разложим в р-ичной системе счисления,  $n=(n_0,n_1,\dots), k=(k_0,k_1,\dots).$  ans= $C_{n_0}^{k_0} \cdot C_{n_1}^{k_1} \cdot \dots$  Способы вычисления  $C_n^k$ :

- 1.  $C_n^k = C_{n-1}^k + C_{n-1}^{k-1}$ precalc:  $O(n^2)$ , query: O(1).
- 2.  $C_n^k = \frac{n!}{k!(n-k)!}$ , предподсчитываем факториалы precalc:  $O(n \log n)$ , query:  $O(\log n)$
- 3. Теорема Люка precalc:  $O(p \log p)$ , query: O(log p).
- 4.  $C_n^k = C_n^{k-1} \cdot \frac{n-k+1}{k}$
- 5.  $C_n^k = \frac{n!}{k!(n-k)!}$ , для каждого факториала считаем степень вхождения и оста-

precalc:  $O(p \log p)$ , query: O(log p).

$$C_n^{\frac{n}{2}} = \frac{2^n}{\sqrt{\frac{\pi n}{2}}}$$

# 79 Делители

•  $\leq 20: d(12) = 6$ 

•  $\leq 50 : d(48) = 10$ 

•  $\leq 100 : d(60) = 12$ 

•  $\leq 1000 : d(840) = 32$ 

 $\bullet \le 10^4 : d(9\ 240) = 64$ 

 $\bullet \le 10^5 : d(83\ 160) = 128$ 

 $\bullet \ \le 10^6: d(720\ 720) = 240$ 

 $\bullet \le 10^7 : d(8\ 648\ 640) = 338$ 

 $\bullet \ \leq 10^8: d(91\,891\,800) = 768$ 

 $\bullet \le 10^9 : d(931\ 170\ 240) = 1344$ 

 $\bullet \le 10^{11} : d(97772875200) = 4032$ 

 $\bullet \ \le 10^{12} : d(963\ 761\ 198\ 400) = 6720$ 

 $\bullet \ \le 10^{15} : d(866\ 421\ 317\ 361\ 600) = 15360$ 

 $\bullet \le 10^{18} : d(897\ 612\ 484\ 786\ 617\ 600) = 103680$ 

# 80 Числа Белла

i	$B_i$	i	$B_i$
0	1	12	4,213,597
1	1	13	27,644,437
2	2	14	190,899,322
3	5	15	1,382,958,545
4	15	16	10,480,142,147
5	52	17	82,864,869,804
6	203	18	682,076,806,159
7	877	19	5,832,742,205,057
8	4,140	20	51,724,158,235,372
9	21,147	21	474,869,816,156,751
10	115,975	22	4,506,715,738,447,323
11	678,570	23	44,152,005,855,084,346

# 81 Разбиения

Число неупорядоченных разбиений n на положительные слагаемые.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$
$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$
$$\frac{n \quad | \ 0.1234567892050100}{p(n) \quad | \ 1.1235711152230627 \sim 2e5 \sim 2e8}$$

# 82 Матричные игры

Пишем матрицу стратегий  $A_{i,j}$  это выигрыш первого и проигрыш второго, i стратегия 1-го. Седловая точка есть для несмешанной стратегии если  $\max_i \min A_{i,*} = \min_j \max A_{*,j}$ . Иначе:

$$f(x) = sum(x_i) \to max, \ Ans = 1/f(x)$$
 
$$Ax \le 1_n, \ x_i \ge 0$$

Для  $2 \times 2$ , p первый игрок, q — второй:

$$p^* = \left(\frac{a_{22} - a_{21}}{a_{22} - a_{12} + a_{11} - a_{21}}; \frac{a_{11} - a_{12}}{a_{22} - a_{12} + a_{11} - a_{21}}\right)$$
$$q^* = \left(\frac{a_{22} - a_{12}}{a_{22} - a_{12} + a_{11} - a_{21}}; \frac{a_{11} - a_{21}}{a_{22} - a_{12} + a_{11} - a_{21}}\right)$$
$$Ans = \frac{a_{22}a_{11} - a_{12}a_{21}}{a_{11} + a_{22} - a_{12} - a_{21}}$$

# 83 Mixed

- ullet Формула Пика: S=Inside+Edge/2-1
- Теорема Люка:  $0 \le n, m \in \mathbb{Z}, p$  простое.  $n = n_k p^k + ... + n_1 p + n_0$  и  $m = m_k p^k + ... + m_1 p + m_0$ . Тогда  $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$ .
- Лемма Бернсайда: |X/G| число орбит G.  $X^g = \{x \in X | gx = x\}$

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$