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#### 1 Common

#### Setup

- 1. Terminal: font Monospace 12
- 2. Gedit: Oblivion, font Monospace 12, auto indent, display line numbers, tab 4, highlight matching brackets, highlight current line, F9 (side panel)
- 3. /.bashrc: export CXXFLAGS='-Wall -Wshadow -Wextra -Wconversion -Wno-unused-result -Wno-deprecated-declarations -O2 -std=gnu++11 -g -DLOCAL'
- 4. for i in {A..K}; do mkdir \$i; cp main.cpp \$i/\$i.cpp; done

## **Template**

```
main.cpp:
#include <bits/stdc++.h>
using namespace std;
#define pb push_back
#define mp make_pair
#define fst first
#define snd second
#define sz(x) (int) ((x).size())
#define form(i, n) for (int i = 0; i < (n); ++i)
#define form (i, n) for (int i = (n) - 1; i \ge 0; --i)
#define forab(i, a, b) for (int i = (a); i < (b); ++i)
#define all(c) (c).begin(), (c).end()
using ll = long long;
using vi = vector<int>;
using pii = pair<int, int>;
#define FNAME ""
int main() {
#ifdef LOCAL
^^Ifreopen(FNAME".in", "r", stdin);
^^Ifreopen(FNAME".out", "w", stdout);
#endif
^^Icin.tie(0);
^^Iios_base::sync_with_stdio(0);
```

#### 3 Stress

^^Ireturn 0;

```
stress.sh:
#!/bin/bash
```

```
for ((i = 0;; i++)); do
^^I./gen $i >in || exit
^^I./main <in >out1 || exit
^^I./stupid <in >out2 || exit
^^Idiff out1 out2 || exit
^^Iecho $i OK
done
```

#### Java 4

```
Java template:
```

```
import java.io.BufferedReader;
import java.io.FileNotFoundException;
import java.io.FileReader;
```

```
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```

```
import java.io.IOException;
 import java.io.InputStreamReader;
 import java.io.PrintWriter;
 import java.util.*;
public class Main {
    ^IFastScanner in;
 ^^IPrintWriter out;
 ^^Ivoid solve() {
 ^^I^^Iint a = in.nextInt();
 ^^I^^Iint b = in.nextInt();
 ^^I^^Iout.print(a + b);
 ^^I}
 ^^Ivoid run() {
 ^^I^^Itry {
 ^^I^^I^^Iin = new FastScanner("input.txt");
 ^^I^^I^^Iout = new PrintWriter("output.txt");
 ^^I^^I^^Isolve();
^^I^^I^^Iout.flush();
^^I^^I^^Iout.close();
^^I^^I} catch (FileNotFoundException e) {
 ^^I^^I^^Ie.printStackTrace();
^^I^^I^^ISystem.exit(1);
^^I^^I}
^^I}
^^Iclass FastScanner {
^^I^^IBufferedReader br;
{\hat{\ }}{\hat{\  }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\  }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\  }}{\hat{\ }}{\hat{\ }}{\hat{\ }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}{\hat{\  }}
 ^^I^^Ipublic FastScanner() {
 ^^I^^I^^Ibr = new BufferedReader(new
 ^^I^^I}
^^I^^I^^Itry {
 ^^I^^I^^I^^Ibr = new BufferedReader(new FileReader(s));
 ^^I^^I^^I} catch (FileNotFoundException e) {
 ^^I^^I^^Ie.printStackTrace();
^^I^^I^^I}
^^I^^I}
^^I^^IString nextToken() {
^^I^^I^^Iwhile (st == null || !st.hasMoreElements()) {
 ^^I^^I^^I^^Itry {
 ^^I^^I^^I^^I^^Ist = new StringTokenizer(br.readLine());
 ^^I^^I^^I^^I} catch (IOException e) {
 ^^I^^I^^I^^I^^Ie.printStackTrace();
^^I^^I^^I^^I}
^^I^^I^^I}
^^I^^I^^Ireturn st.nextToken();
 ^^I^^I}
^^I^^Iint nextInt() {
 ^^I^^I^^Ireturn Integer.parseInt(nextToken());
 ^^I^^I}
^^I^^Ilong nextLong() {
 ^^I^^I^^Ireturn Long.parseLong(nextToken());
 ^^I^^I}
 ^^I^^Idouble nextDouble() {
^^I^^I^^Ireturn Double.parseDouble(nextToken());
^^I^^I}
 ^^I^^Ichar nextChar() {
 ^^I^^I^^Itry {
 ^^I^^I^^I^^Ireturn (char) (br.read());
```

```
^^I^^I^^I} catch (IOException e) {
^^I^^I^^I^^Ie.printStackTrace();
^^I^^I^^I}
^^I^^I^^Ireturn 0;
^^I^^I}
^^I^^IString nextLine() {
^^I^^I^^Itry {
^^I^^I^^I^^Ireturn br.readLine();
^^I^^I^^I} catch (IOException e) {
^^I^^I^^I^^Ie.printStackTrace();
^^I^^I^^I}
^^I^^I^^Ireturn "";
^^I^^I}
^^I}
^^Ipublic static void main(String[] args) {
^^I^^Inew Main().run();
^^I}
```

# 2 Big numbers

#### 5 Big Int

```
const int BASE_LEN = 9;
const int NUM_LEN = 50000 / BASE_LEN + 2; // LEN <= NUM_LEN *</pre>
→ BASE LEN
const int BASE = pow(10, BASE_LEN);
const 11 INF = 8e18, ADD = INF / BASE;
struct num {
 11 a[NUM_LEN];
  int len; // always > 0
  inline const ll& operator [] ( int i ) const { return a[i];
  → }
  inline ll& operator [] ( int i ) { return a[i]; }
  num\& operator = (const num \&x) { // copy}
   len = x.len;
   memcpy(a, x.a, sizeof(a[0]) * len);
   return *this;
  num( const num &x ) { *this = x; } // copy
  num() \{ len = 1, a[0] = 0; \} // 0
  num( 11 x ) { // x
   len = 0:
    while (!len | | x) {
     assert(len < NUM_LEN); // to catch overflow</pre>
      a[len++] = x \% BASE, x /= BASE;
   }
  num& cor() {
    while (a[len - 1] >= BASE) {
      assert(len < NUM_LEN); // to catch overflow</pre>
      if (a[len - 1] >= 2 * BASE)
       a[len] = a[len - 1] / BASE, a[len - 1] %= BASE;
       a[len] = 1, a[len - 1] -= BASE;
     len++;
   while (len > 1 && !a[len - 1])
     len--;
   return *this;
```

```
int length() {
 if (!len)
   return 0;
 int x = a[len - 1], res = 0;
 assert(x);
 while (x || !res)
   x /= 10, res++;
 return res + (len - 1) * BASE_LEN;
void out() const {
 int i = len - 1:
 printf("%d", (int)a[i--]);
 while (i >= 0)
   printf("%0*d", BASE_LEN, (int)a[i--]);
 puts("");
void init( const char *s ) {
 int sn = strlen(s);
 while (sn \&\& s[sn - 1] \le 32)
 len = (sn + BASE_LEN - 1) / BASE_LEN;
  memset(a, 0, sizeof(a[0]) * len);
 forn(i, sn) {
   11 &r = a[(sn - i - 1) / BASE_LEN];
   r = r * 10 + (s[i] - '0');
bool read() {
 static const int L = NUM_LEN * BASE_LEN + 1;
  static char s[L];
 if (!fgets(s, L, stdin))
   return 0;
 assert(!s[L - 2]);
 init(s);
 return 1;
void mul2() {
 forn(i, len)
   a[i] <<= 1;
  forn(i, len - 1)
   if (a[i] >= BASE)
     a[i + 1]++, a[i] -= BASE;
  cor();
void div2() {
 for (int i = len - 1; i >= 0; i--) {
   if (i && (a[i] & 1))
     a[i - 1] += BASE;
   a[i] >>= 1;
 }
  cor();
static ll cmp( ll *a, ll *b, int n ) {
 while (n--)
   if (a[n] != b[n])
     return a[n] - b[n];
 return 0;
int cmp( const num &b ) const {
 if (len != b.len)
   return len - b.len;
  for (int i = len - 1; i >= 0; i--)
   if (a[i] != b[i])
     return a[i] - b[i];
```

```
return 0:
                                                                      a.len = 0:
                                                                      return a;
                                                                    forn(i, a.len)
 bool zero() {
   return len == 1 && !a[0];
                                                                      a[i] *= k;
                                                                    forn(i, a.len - 1)
                                                                      if (a[i] >= BASE)
                                                                       a[i + 1] += a[i] / BASE, a[i] %= BASE;
  /** c = this/b, this %= b */
 num &div( num b, num &c ) {
                                                                    return a.cor();
   c.len = len - b.len;
    for (int i = c.len; i >= 0; i--) {
     int k = (1.0L * a[len - 1] * BASE + (len >= 2 ? a[len -
                                                                 num& operator /= ( num &a, int k ) {
      \rightarrow 2] : 0)) / (1.0L * b[b.len - 1] * BASE + (b.len >= 2
                                                                   if (k == 1)
      \rightarrow ? b[b.len - 2] + 1 : 0));
                                                                     return a;
      c[i] = 0;
                                                                    assert(k != 0);
     if (k > 0) {
                                                                    for (int i = a.len - 1; i > 0; i--)
       c[i] += k;
                                                                     a[i - 1] += (11)(a[i] \% k) * BASE, a[i] /= k;
        forn(j, b.len)
                                                                    a[0] /= k;
         a[i + j] = (11)b[j] * k;
                                                                   return a.cor():
       forn(j, b.len)
         if (a[i + j] < 0) {
           11 k = (-a[i + j] + BASE - 1) / BASE;
                                                                  num\&\ mul(\ const\ num\ \&a,\ const\ num\ \&b,\ num\ \&x\ ) {
           a[i + j] += k * BASE, a[i + j + 1] -= k;
                                                                   assert(a.len + b.len - 1 <= NUM_LEN);</pre>
                                                                    memset(x.a, 0, sizeof(x[0]) * (a.len + b.len - 1));
     }
                                                                    forn(i, a.len)
                                                                     forn(j, b.len)
        len--, a[len - 1] += a[len] * BASE, a[len] = 0;
                                                                        if ((x[i + j] += a[i] * b[j]) >= INF)
      else if (cmp(a, b.a, b.len) >= 0) {
                                                                         x[i + j + 1] += ADD, x[i + j] -= INF;
                                                                    x.len = a.len + b.len - 1;
       c[0]++;
       forn(j, b.len)
                                                                    forn(i, x.len - 1)
          if ((a[j] -= b[j]) < 0)
                                                                      if (x[i] >= BASE)
           a[j] += BASE, a[j + 1]--;
                                                                       x[i + 1] += x[i] / BASE, x[i] %= BASE;
                                                                    return x.cor():
   7
   if (c.len < 0)
     c[c.len = 0] = 0;
                                                                  bool operator == ( const num &a, const num &b ) { return
    forn(i, c.len)
                                                                  \hookrightarrow a.cmp(b) == 0; }
     if (c[i] >= BASE)
                                                                  bool operator != ( const num &a, const num &b ) { return
       c[i + 1] += c[i] / BASE, c[i] %= BASE;
                                                                  \rightarrow a.cmp(b) != 0; }
    c.len += (!c.len || c[c.len]);
                                                                  bool operator < ( const num &a, const num &b ) { return</pre>
                                                                  \rightarrow a.cmp(b) < 0; }
   return cor();
                                                                  bool operator > ( const num &a, const num &b ) { return
                                                                  \hookrightarrow a.cmp(b) > 0; }
                                                                  bool operator <= ( const num &a, const num &b ) { return</pre>
num& operator += ( num &a, const num &b ) {
                                                                  \hookrightarrow a.cmp(b) <= 0; }
 while (a.len < b.len)
                                                                  bool operator >= ( const num &a, const num &b ) { return
   a[a.len++] = 0;
                                                                  \hookrightarrow a.cmp(b) >= 0; }
 forn(i, b.len)
   a[i] += b[i];
                                                                  num& add( const num &a, const num &b, num &c ) { c = a; c +=
 forn(i, a.len - 1)
                                                                  if (a[i] >= BASE)
                                                                  num% sub( const num %a, const num %b, num %c ) { c = a; c -=
     a[i] -= BASE, a[i + 1]++;
                                                                  return a.cor();
                                                                  num& mul( const num &a, int k, num &c )
                                                                  num& div( const num &a, int k, num &c ) { c = a; c /=
num\& operator -= ( num \&a, const num \&b ) {
                                                                  while (a.len < b.len)
   a[a.len++] = 0;
                                                                  num\& operator *= ( num \&a, const num \&b ) {
 forn(i, b.len)
                                                                    static num tmp;
   a[i] -= b[i];
                                                                    mul(a, b, tmp);
 forn(i, a.len - 1)
                                                                   return a = tmp;
   if (a[i] < 0)
     a[i] += BASE, a[i + 1]--;
 assert(a[a.len - 1] >= 0); // a >= b
                                                                  num operator ^ ( const num &a, int k ) {
 return a.cor();
                                                                   num res(1);
                                                                   forn(i, k)
                                                                     res *= a;
num& operator *= ( num &a, int k ) {
                                                                   return res;
 if (k == 1)
   return a:
 if (k == 0) {
                                                                  num& gcd_binary( num &a, num &b ) {
```

```
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```

```
int cnt = 0:
                                                                             const Num z = a[i + j + k] * rt[j + k];
                                                                             a[i + j + k] = a[i + j] - z;
  while (!a.zero() && !b.zero()) {
    while (!(b[0] & 1) && !(a[0] & 1))
                                                                             a[i + j] += z;
      cnt++, a.div2(), b.div2();
    while (!(b[0] & 1))
                                                                    }
     b.div2();
    while (!(a[0] & 1))
                                                                     void fftInv(Num *a, int n) {
     a.div2();
                                                                      fft(a, n):
    if (a.cmp(b) < 0)
                                                                       reverse(a + 1, a + n);
                                                                      forn (i, n)
     b -= a;
    else
                                                                         a[i] /= n;
      a -= b:
  if (a.zero())
                                                                    void doubleFft(Num *a, Num *fa, Num *fb, int n) { // only if
   std::swap(a, b);
                                                                     \hookrightarrow you need it
  while (cnt)
                                                                      fft(a, n);
   a.mul2(), cnt--;
                                                                      const int n1 = n - 1;
  return a;
                                                                       forn (i, n) {
                                                                        const Num &z0 = a[i], &z1 = a[(n - i) & n1];
                                                                         fa[i] = Num(z0.real() + z1.real(), z0.imag() - z1.imag())
num& gcd( num &a, num &b ) {
                                                                         → * 0.5:
                                                                        fb[i] = Num(z0.imag() + z1.imag(), z1.real() - z0.real())
  static num tmp:
                                                                        → * 0.5;
 return b.zero() ? a : gcd(b, a.div(b, tmp));
                                                                      }
                                                                    }
  FFT
                                                                    Num tmp[MAX_N];
                                                                    template<class T>
int rev[MAX_N];
                                                                     void mult(T *a, T *b, T *r, int n) { // n = 2 \hat{k}
                                                                      forn (i, n)
//typedef complex<dbl> Num;
                                                                         tmp[i] = Num((dbl) a[i], (dbl) b[i]);
struct Num {
                                                                       fft(tmp, n);
  dbl x, y;
                                                                       const int n1 = n - 1;
  Num() {}
                                                                       const Num c = Num(0, -0.25 / n);
  \label{eq:num} \mbox{Num(dbl $\underline{\ }$x, dbl $\underline{\ }$y): $x(\underline{\ }$x), $y(\underline{\ }$y) $\{$}
                                                                      fornr (i, n / 2 + 1) {
  inline dbl real() const { return x; }
                                                                        const int j = (n - i) & n1;
 inline dbl imag() const { return y; }
                                                                         const Num z0 = sqr(tmp[i]), z1 = sqr(tmp[j]);
  inline Num operator+(const Num &B) const { return Num(x +
                                                                         tmp[i] = (z1 - conj(z0)) * c;
  \hookrightarrow B.x, y + B.y); }
                                                                         tmp[j] = (z0 - conj(z1)) * c;
  inline Num operator-(const Num &B) const { return Num(x -
  \hookrightarrow B.x, y - B.y); }
                                                                       fft(tmp, n);
  inline Num operator*(dbl k) const { return Num(x * k, y *
                                                                       forn (i, n)
  \hookrightarrow k): }
                                                                        r[i] = (T) round(tmp[i].real());
  inline Num operator*(const Num &B) const { return Num(x *
  \hookrightarrow B.x - y * B.y, x * B.y + y * B.x); }
 inline void operator+=(const Num &B) { x += B.x, y += B.y; }
                                                                     void init() { // don't forget to init
 inline void operator/=(dbl k) { x \neq k, y \neq k; }
                                                                       forn(i, MAX_N)
  inline void operator*=(const Num &B) { *this = *this * B; }
                                                                         rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (LOG - 1));
};
                                                                       rt[1] = Num(1, 0);
Num rt[MAX_N];
                                                                       for (int k = 1, p = 2; k < LOG; k++, p *= 2) {
                                                                        const Num x(cos(PI / p), sin(PI / p));
inline Num sqr(const Num &x) { return x * x; }
                                                                         forab (i, p / 2, p)
inline Num conj(const Num &x) { return Num(x.real(),
                                                                           rt[2 * i] = rt[i], rt[2 * i + 1] = rt[i] * x;
\rightarrow -x.imag()); }
                                                                    7
inline int getN(int n) {
  int k = 1;
  while(k < n)
                                                                         FFT by mod and FFT with digits up to 10^6
   k <<= 1;
                                                                     Num ta[MAX_N], tb[MAX_N], tf[MAX_N], tg[MAX_N];
 return k:
                                                                     const int HALF = 15:
void fft(Num *a, int n) {
  assert(rev[1]); // don't forget to init
                                                                    void mult(int *a, int *b, int *r, int n, int mod) {
  int q = MAX_N / n;
                                                                      int tw = (1 << HALF) - 1;</pre>
  forn (i, n)
                                                                      forn (i, n) {
   if(i < rev[i] / q)
                                                                        int x = int(a[i] % mod);
      swap(a[i], a[rev[i] / q]);
                                                                         ta[i] = Num(x \& tw, x >> HALF);
  for (int k = 1; k < n; k <<= 1)
    for (int i = 0; i < n; i += 2 * k)
                                                                       forn (i, n) {
      forn (j, k) {
                                                                         int x = int(b[i] % mod);
```

```
tb[i] = Num(x \& tw, x >> HALF);
  fft(ta, n), fft(tb, n);
  forn (i, n) {
    int j = (n - i) & (n - 1);
    Num a1 = (ta[i] + conj(ta[j])) * Num(0.5, 0);
    Num a2 = (ta[i] - conj(ta[j])) * Num(0, -0.5);
    Num b1 = (tb[i] + conj(tb[j])) * Num(0.5 / n, 0);
    Num b2 = (tb[i] - conj(tb[j])) * Num(0, -0.5 / n);
    tf[j] = a1 * b1 + a2 * b2 * Num(0, 1);
    tg[j] = a1 * b2 + a2 * b1;
  fft(tf, n), fft(tg, n);
  forn (i, n) {
    11 aa = 11(tf[i].x + 0.5);
    11 bb = 11(tg[i].x + 0.5);
    11 cc = 11(tf[i].y + 0.5);
    r[i] = int((aa + ((bb \% mod) << HALF) + ((cc \% mod) << (2))

    * HALF))) % mod);
 }
}
int tc[MAX_N], td[MAX_N];
const int MOD1 = 1.5e9, MOD2 = MOD1 + 1;
void multLL(int *a, int *b, ll *r, int n){
  mult(a, b, tc, n, MOD1), mult(a, b, td, n, MOD2);
  forn(i, n)
    r[i] = tc[i] + (td[i] - tc[i] + (11)MOD2) * MOD1 % MOD2 *
    \hookrightarrow MOD1:
```

## **Data Structures**

#### **Centroid Decomposition**

```
vi g[MAX_N];
int d[MAX_N], par[MAX_N], centroid;
//d par -
int find(int v, int p, int total) {
 int size = 1, ok = 1;
  for (int to : g[v])
   if (d[to] == -1 \&\& to != p) {
      int s = find(to, v, total);
      if (s > total / 2) ok = 0;
      size += s;
   }
  if (ok && size > total / 2) centroid = v;
 return size:
void calcInComponent(int v, int p, int level) {
  // do something
  for (int to : g[v])
    if (d[to] == -1 \&\& to != p)
      calcInComponent(to, v, level);
//fill(d, d + n, -1)
//decompose(0, -1, 0)
void decompose(int root, int parent, int level) {
  find(root, -1, find(root, -1, INF));
  int c = centroid;
  par[c] = parent, d[c] = level;
  calcInComponent(centroid, -1, level);
  for (int to : g[c])
   if (d[to] == -1)
```

```
decompose(to, c, level + 1);
```

#### 9 Convex Hull Trick

}

```
struct Line {
  int k, b;
  Line() {}
  Line(int _k, int _b): k(_k), b(_b) {}
  ll get(int x) { return b + k * 111 * x; }
  bool operator<(const Line &1) const { return k < 1.k; } //</pre>
};
                    (a,b)
                             (a,c)
inline bool check(Line a, Line b, Line c) {
  return (a.b - b.b) * 111 * (c.k - a.k) < (a.b - c.b) * 111 *
  \hookrightarrow (b.k - a.k);
struct Convex {
  vector<Line> st;
  inline void add(Line 1) {
    while (sz(st) \ge 2 \&\& ! check(st[sz(st) - 2], st[sz(st) -
    \hookrightarrow 1], 1))
     st.pop_back();
    st.pb(1);
  int get(int x) {
    int l = 0, r = sz(st);
    while (r - 1 > 1) {
      int m = (1 + r) / 2; // >
      if (st[m - 1].get(x) < st[m].get(x))
        1 = m;
      else
    }
    return 1;
  }
  Convex() {}
  Convex(vector<Line> &lines) {
    st.clear();
    for(Line &1 : lines)
      add(1);
  Convex(Line line) { st.pb(line); }
  Convex(const Convex &a, const Convex &b) {
    vector<Line> lines:
    lines.resize(sz(a.st) + sz(b.st));
    merge(all(a.st), all(b.st), lines.begin());
    st.clear();
    for(Line &1 : lines)
      add(1);
  }
}:
10 DSU
```

```
int pr[MAX_N];
int get(int v) {
 return v == pr[v] ? v : pr[v] = get(pr[v]);
bool unite(int v, int u) {
 v = get(v), u = get(u);
 if (v == u) return 0;
 pr[u] = v;
 return 1;
```

int getST(int path, int v, int vl, int vr, int ind) {

pushST(path, v, vl, vr);

if (vl == vr - 1)
 return t[path][v];

```
void init(int n) {
                                                                   int vm = (vl + vr) / 2;
 forn (i, n) pr[i] = i;
                                                                   if (ind >= vm)
                                                                    return getST(path, 2 * v + 1, vm, vr, ind);
                                                                   return getST(path, 2 * v, v1, vm, ind);
     Fenwick Tree
11
                                                                 void setST(int path, int v, int vl, int vr, int l, int r, int
int t[MAX_N];

    val) {

                                                                   if (vl >= l && vr <= r) {
int get(int ind) {
                                                                     toPush[path][v] = val;
 int res = 0;
                                                                     pushST(path, v, vl, vr);
  for (; ind >= 0; ind \&= (ind + 1), ind--)
                                                                     return:
   res += t[ind];
  return res;
                                                                   pushST(path, v, v1, vr);
                                                                   if (vl >= r || l >= vr)
                                                                    return;
void add(int ind, int n, int val) {
                                                                   int vm = (v1 + vr) / 2;
  for (; ind < n; ind |= (ind + 1))
   t[ind] += val;
                                                                   setST(path, 2 * v, vl, vm, l, r, val);
                                                                   setST(path, 2 * v + 1, vm, vr, 1, r, val);
                                                                   t[path][v] = min(t[path][2 * v], t[path][2 * v + 1]);
int sum(int 1, int r) { // [l, r)
  return get(r - 1) - get(l - 1);
                                                                 bool isUpper(int v, int u) {
                                                                   return tin[v] <= tin[u] && tout[v] >= tout[u];
12 Hash Table
                                                                  int getHLD(int v) {
using H = 11;
                                                                  return getST(comp[v], 1, 0, sz(t[comp[v]]) / 2, num[v]);
const int HT_SIZE = 1<<20, HT_AND = HT_SIZE - 1, HT_SIZE_ADD =</pre>
→ HT SIZE / 100:
H ht[HT_SIZE + HT_SIZE_ADD];
                                                                 int setHLD(int v, int u, int val) {
int data[HT_SIZE + HT_SIZE_ADD];
                                                                   int ans = 0, w = 0;
                                                                   forn (i, 2) {
int get(const H &hash){
                                                                     while (!isUpper(w = top[comp[v]], u))
 int k = ((11) hash) & HT_AND;
                                                                       setST(comp[v], 1, 0, sz(t[comp[v]]) / 2, 0, num[v] + 1,
  while (ht[k] && ht[k] != hash) ++k;
                                                                        \hookrightarrow val), v = pr[w];
 return k:
                                                                     swap(v, u);
                                                                   setST(comp[v], 1, 0, sz(t[comp[v]]) / 2, min(num[v],
void insert(const H &hash, int x){
                                                                   \rightarrow num[u]), max(num[v], num[u]) + 1, val);
 int k = get(hash);
                                                                   return ans;
  if (!ht[k]) ht[k] = hash, data[k] = x;
                                                                 void dfs(int v, int p) {
bool count(const H &hash, int x){
                                                                   tin[v] = curTime++;
 int k = get(hash);
                                                                   size[v] = 1;
  return ht[k] != 0;
                                                                   pr[v] = p;
                                                                   for (int u : g[v])
                                                                     if (u != p) {
13 Heavy Light Decomposition
                                                                       dfs(u, v);
                                                                       size[v] += size[u];
vi g[MAX_N];
int size[MAX_N], comp[MAX_N], num[MAX_N], top[MAX_N],
                                                                   tout[v] = curTime++;

    pr[MAX_N], tin[MAX_N], tout[MAX_N];

vi t[MAX_N], toPush[MAX_N], lst[MAX_N];
int curPath = 0, curTime = 0;
                                                                 void build(int v) {
                                                                   if (v == 0 \mid \mid size[v] * 2 < size[pr[v]])
void pushST(int path, int v, int vl, int vr) {
                                                                     top[curPath] = v, comp[v] = curPath, num[v] = 0,
  if (toPush[path][v] != -1) {
                                                                     if (vl != vr - 1)
      forn (j, 2)
                                                                     comp[v] = comp[pr[v]], num[v] = num[pr[v]] + 1;
        toPush[path][2 * v + j] = toPush[path][v];
                                                                   lst[comp[v]].pb(v);
                                                                   for (int u : g[v])
      t[path][v] = toPush[path][v];
                                                                     if (u != pr[v])
    toPush[path][v] = -1;
                                                                       build(u);
```

void initHLD() {

dfs(0, 0);

build(0):

```
forn (i, curPath) {
  int curSize = 1;
  while (curSize < sz(lst[i]))
    curSize *= 2;
  t[i].resize(curSize * 2);
  toPush[i] = vi(curSize * 2, -1);
  //initialize t[i]
}</pre>
```

#### 14 Next Greater in Segment Tree

#### 15 Sparse Table

```
int st[MAX_N][MAX_LOG];
int lg[MAX_N];

int get(int 1, int r) { // [l, r)
    int curLog = lg[r - 1];
    return min(st[l][curLog], st[r - (1 << curLog)][curLog]);
}

void initSparseTable(int *a, int n) {
    lg[1] = 0;
    forab (i, 2, n + 1) lg[i] = lg[i / 2] + 1;
    forn (i, n) st[i][0] = a[i];
    forn (j, lg[n])
        forn (i, n - (1 << (j + 1)) + 1)
            st[i][j + 1] = min(st[i][j], st[i + (1 << j)][j]);
}</pre>
```

#### 16 Fenwick Tree 2D

```
11 a[4][MAX_N][MAX_N];
int n, m;
inline int f(int x) { return x & ~(x - 1); }
inline void add(int k, int x, int y, ll val) {
  for (; x \le n; x += f(x))
   for (int j = y; j \le m; j += f(j))
      a[k][x][j] += val;
inline ll get(int k, int x, int y) {
 11 s = 0;
  for (; x > 0; x -= f(x))
    for (int j = y; j > 0; j -= f(j))
      s += a[k][x][j];
  return s;
inline ll get(int x, int y) {
 return 11(x + 1) * (y + 1) * get(0, x, y) - (y + 1) * get(1, y)
  \hookrightarrow x, y)
      -(x + 1) * get(2, x, y) + get(3, x, y);
inline void add(int x, int y, ll val) {
```

## 4 Dynamic Programming

#### **17** LIS

```
int longestIncreasingSubsequence(vi a) {
  int n = sz(a);
  vi d(n + 1, INF);
  d[0] = -INF;
  forn (i, n)
    *upper_bound(all(d), a[i]) = a[i];
  fornr (i, n + 1) if (d[i] != INF) return i;
  return 0;
}
```

#### 18 DP tree

```
int dp[MAX_N][MAX_N], a[MAX_N];
vi g[MAX_N];
int dfs(int v, int n) {
  forn (i, n + 1)
   dp[v][i] = -INF;
  dp[v][1] = a[v];
 int curSz = 1:
  for (int to : g[v]) {
   int toSz = dfs(to, n);
   for (int i = curSz; i >= 1; i--)
     fornr (j, toSz + 1)
       dp[v][i + j] = max(dp[v][i + j], dp[v][i] +
        \hookrightarrow dp[to][j]);
   curSz += toSz;
 7
 return curSz;
}
```

#### 19 Masks tricks

#### 5 Flows

#### 20 Utilities

#### 21 Ford-Fulkerson

```
int used[MAX_N], pr[MAX_N];
int curTime = 1;
int dfs(int v, int can, int toPush, int t) {
 if (v == t) return can;
  used[v] = curTime;
  for (int edge : g[v]) {
    auto &e = edges[edge]:
    if (used[e.u] != curTime && e.c - e.f >= toPush) {
      int flow = dfs(e.u, min(can, e.c - e.f), toPush, t);
      if (flow > 0) {
        addFlow(edge, flow), pr[e.u] = edge;
        return flow;
   }
 }
 return 0;
int fordFulkerson(int s, int t) {
  int ansFlow = 0, flow = 0;
  // Without scaling
 while ((flow = dfs(s, INF, 1, t)) > 0)
   ansFlow += flow, curTime++;
  // With scaling
  fornr (i, INF_LOG)
   for (curTime++; (flow = dfs(s, INF, (1 << i), t)) > 0;
    \hookrightarrow curTime++)
     ansFlow += flow;
 return ansFlow;
```

#### 22 Dinic

```
int pr[MAX_N], d[MAX_N], q[MAX_N], first[MAX_N];
int dfs(int v, int can, int toPush, int t) {
   if (v == t) return can;
   int sum = 0;
   for (; first[v] < (int) g[v].size(); first[v]++) {
      auto &e = edges[g[v][first[v]]];
      if (d[e.u] != d[v] + 1 || e.c - e.f < toPush) continue;
   int flow = dfs(e.u, min(can, e.c - e.f), toPush, t);
      addFlow(g[v][first[v]], flow);</pre>
```

```
can -= flow, sum += flow;
   if (!can)
   return sum;
 }
  return sum;
bool bfs(int n, int s, int t, int curPush) {
  forn (i, n) d[i] = INF, first[i] = 0;
  int head = 0, tail = 0;
  q[tail++] = s;
  d[s] = 0;
  while (tail - head > 0) {
   int v = q[head++];
   for (int edge : g[v]) {
     auto &e = edges[edge];
     if (d[e.u] > d[v] + 1 \&\& e.c - e.f >= curPush)
        d[e.u] = d[v] + 1, q[tail++] = e.u;
   }
  }
  return d[t] != INF;
int dinic(int n, int s, int t) {
 int ansFlow = 0;
  // Without scaling
  while (bfs(n, s, t, 1))
   ansFlow += dfs(s, INF, 1, t);
  // With scaling
 fornr (j, INF_LOG)
    while (bfs(n, s, t, 1 \ll j))
     ansFlow += dfs(s, INF, 1 << j, t);</pre>
  return ansFlow:
```

## 23 Hungarian

```
const int INF = 1e9;
int a[MAX_N][MAX_N];
// min = sum of a[pa[i],i]
// you may optimize speed by about 15%, just change all vectors
vi Hungarian(int n) {
  vi pa(n + 1, -1), row(n + 1, 0), col(n + 1, 0), la(n + 1);
 forn (k, n) {
   vi u(n + 1, 0), d(n + 1, INF);
   pa[n] = k;
   int 1 = n, x;
    while ((x = pa[1]) != -1) {
      u[1] = 1:
      int minn = INF, tmp, 10 = 1;
      forn (j, n)
       if (!u[j]) {
          if ((tmp = a[x][j] + row[x] + col[j]) < d[j])
           d[j] = tmp, la[j] = 10;
          if (d[j] < minn)</pre>
           minn = d[j], l = j;
       }
      forn (j, n + 1)
       if (u[i])
          col[j] += minn, row[pa[j]] -= minn;
          d[j] -= minn;
    while (1 != n)
      pa[1] = pa[la[1]], 1 = la[1];
  }
  return pa;
}
```

#### 24 Min Cost Max Flow

```
const int MAX_M = 1e4;
int pr[MAX_N], in[MAX_N], q[MAX_N * MAX_M], used[MAX_N],
\hookrightarrow d[MAX_N], pot[MAX_N];
vi g[MAX_N];
struct Edge {
  int v, u, c, f, w;
 Edge() {}
 Edge(int _v, int _u, int _c, int _w): v(_v), u(_u), c(_c),
 \hookrightarrow f(0), w(_w) {}
vector<Edge> edges;
inline void addFlow(int e, int flow) {
  edges[e].f += flow, edges[e ^ 1].f -= flow;
inline void addEdge(int v, int u, int c, int w) {
  g[v].pb(sz(edges)), edges.pb(Edge(v, u, c, w));
  g[u].pb(sz(edges)), edges.pb(Edge(u, v, 0, -w));
int dijkstra(int n, int s, int t) {
  forn (i, n) used[i] = 0, d[i] = INF;
  d[s] = 0;
  while (1) {
   int v = -1;
    forn (i, n)
     if (!used[i] && (v == -1 \mid \mid d[v] > d[i]))
        v = i;
    if (v == -1 \mid \mid d[v] == INF) break;
    used[v] = 1;
    for (int edge : g[v]) {
      auto &e = edges[edge];
      int w = e.w + pot[v] - pot[e.u];
     if (e.c > e.f \&\& d[e.u] > d[v] + w)
        d[e.u] = d[v] + w, pr[e.u] = edge;
   }
 }
  if (d[t] == INF) return d[t];
  forn (i, n) pot[i] += d[i];
  return pot[t];
int fordBellman(int n, int s, int t) {
  forn (i, n) d[i] = INF;
  int head = 0, tail = 0;
  d[s] = 0, q[tail++] = s, in[s] = 1;
  while (tail - head > 0) {
    int v = q[head++];
   in[v] = 0;
   for (int edge : g[v]) {
      auto &e = edges[edge];
      if (e.c > e.f && d[e.u] > d[v] + e.w) {
       d[e.u] = d[v] + e.w;
       pr[e.u] = edge;
       if (!in[e.u])
          in[e.u] = 1, q[tail++] = e.u;
   }
  }
  return d[t];
int minCostMaxFlow(int n, int s, int t) {
  int ansFlow = 0, ansCost = 0, dist;
  while ((dist = dijkstra(n, s, t)) != INF) {
    int curFlow = INF;
    for (int cur = t; cur != s; cur = edges[pr[cur]].v)
```

#### 6 Games

#### 25 Retrograde Analysis

```
int win[MAX_N], lose[MAX_N], outDeg[MAX_N];
vi rg[MAX_N];
void retro(int n) {
  queue<int> q;
 forn (i, n)
   if (!outDeg[i])
     lose[i] = 1, q.push(i);
  while (!q.empty()) {
   int v = q.front();
   q.pop();
   for (int to : rg[v])
     if (lose[v]) {
        if (!win[to])
          win[to] = 1, q.push(to);
      } else {
        outDeg[to]--;
        if (!outDeg[to])
          lose[to] = 1, q.push(to);
 }
```

# 7 Geometry

## 26 ClosestPoints (SweepLine)

```
#include "header.h"
const int N = 2e5:
struct Pnt {
 int x, y, i;
  bool operator <(const Pnt &p) const { return mp(y, i) <</pre>
  \hookrightarrow mp(p.y, p.i); }
11 d2 = 8e18, d = (11) sqrt(d2) + 1;
Pnt p[N];
inline ll sqr(int x){
 return (LL)x * x;
inline void relax(const Pnt &a, const Pnt &b){
 11 tmp = sqr(a.x - b.x) + sqr(a.y - b.y);
  if (tmp < d2)
    d2 = tmp, d = (LL)(sqrt(d2) + 1 - 1e-9); // round up
inline bool xless(const Pnt &a, const Pnt &b){
  return a.x < b.x;
int main() {
```

```
int n:
scanf("%d", &n);
forn(i, n)
 scanf("%d%d", &p[i].x, &p[i].y), p[i].i = i;
sort(p, p + n, xless);
set <Pnt> s;
int 1 = 0;
forn(r, n){
  set<Pnt>::iterator it_r = s.lower_bound(p[r]), it_l =
 for (; it_r != s.end() && it_r->y - p[r].y < d; ++it_r)
   relax(*it_r, p[r]);
  while (it_1 != s.begin() && p[r].y - (--it_1)->y < d)
   relax(*it_l, p[r]);
  s.insert(p[r]);
 while (1 <= r \&\& p[r].x - p[1].x >= d)
    s.erase(p[l++]);
printf("%.9f\n", sqrt(d2));
return 0;
```

#### 27 ConvexHull

```
typedef vector<Pnt> vpnt;
inline bool by Angle (const Pnt &a, const Pnt &b) {
 dbl x = a \% b;
  return eq(x, 0) ? a.len2() < b.len2() : x < 0;
vpnt convexHull(vpnt p) {
  int n = sz(p);
  assert(n > 0);
  swap(p[0], *min_element(all(p)));
 forab(i, 1, n)
 p[i] = p[i] - p[0];
  sort(p.begin() + 1, p.end(), byAngle);
                            (1) (2)
  (1):
  int k = p.size() - 1;
  while(k > 0 \& eq((p[k-1]-p.back()) \% p.back(), 0))
   --k;
  reverse(pi.begin() + k, pi.end());*/
  int rn = 0:
  vpnt r(n);
  r[rn++] = p[0];
  forab(i, 1, n){
   Pnt q = p[i] + p[0];
   while(rn >= 2 && geq((r[rn - 1] - r[rn - 2]) % (q - r[rn -
    \hookrightarrow 2]), 0)) // (2) ge
     --rn;
   r[rn++] = q;
 }
 r.resize(rn);
 return r:
}
```

## 28 GeometryBase

```
#include<bits/stdc++.h>
using namespace std;

typedef long long ll;
typedef long double ld;
typedef double dbl;

const dbl EPS = 1e-9;
```

```
const int PREC = 20:
inline bool eq(dbl a, dbl b) { return abs(a-b)<=EPS; }</pre>
inline bool gr(dbl a, dbl b) { return a>b+EPS; }
inline bool geq(dbl a, dbl b) { return a>=b-EPS; }
inline bool ls(dbl a, dbl b) { return a < b - EPS; }</pre>
inline bool leq(dbl a, dbl b) { return a <= b + EPS; }</pre>
struct Pnt {
    dbl x,y;
    Pnt(): x(0), y(0) {}
    Pnt(dbl xx, dbl yy): x(xx), y(yy) {}
    inline Pnt operator +(const Pnt &p) const { return Pnt(x +
    \rightarrow p.x, y + p.y); }
    inline Pnt operator -(const Pnt &p) const { return Pnt(x -
    \rightarrow p.x, y - p.y); }
    inline dbl operator *(const Pnt &p) const { return x * p.x
    \rightarrow + y * p.y; } // ll
    inline dbl operator %(const Pnt &p) const { return x * p.y
    \hookrightarrow - y * p.x; } // ll
    inline Pnt operator *(dbl k) const { return Pnt(x * k, y *
    \rightarrow k): }
    inline Pnt operator /(dbl k) const { return Pnt(x / k, y /
    \hookrightarrow k); }
    inline Pnt operator -() const { return Pnt(-x, -y); }
    inline void operator +=(const Pnt &p) { x += p.x, y +=
    \hookrightarrow p.y; }
    inline void operator -= (const Pnt &p) { x -= p.x, y -=
    \rightarrow p.y; }
    inline void operator *=(dbl k) { x*=k, y*=k; }
    inline bool operator ==(const Pnt &p) { return
     \rightarrow abs(x-p.x)<=EPS && abs(y-p.y)<=EPS; }
    inline bool operator !=(const Pnt &p) { return
    \rightarrow abs(x-p.x)>EPS || abs(y-p.y)>EPS; }
    inline bool operator <(const Pnt &p) { return
    \rightarrow abs(x-p.x)<=EPS ? y<p.y-EPS : x<p.x; }
    inline dbl angle() const { return atan2(y, x); } // ld
    inline dbl len2() const { return x*x+y*y; } // ll
    inline dbl len() const { return sqrt(x*x+y*y); } // ll, ld
    inline Pnt getNorm() const {
        auto 1 = len();
        return Pnt(x/1, y/1);
    }
    inline void normalize() {
        auto 1 = len();
        x/=1, y/=1;
    inline Pnt getRot90() const { //counter-clockwise
        return Pnt(-y, x);
    inline Pnt getRot(dbl a) const { // ld
        dbl si = sin(a), co = cos(a);
        return Pnt(x*co - y*si, x*si + y*co);
    inline void read() {
        int xx, yy;
    cin >> xx >> yy;
        x = xx, y = yy;
    inline void write() const{
        cout << fixed << (double)x << (double)y;</pre>
};
struct Line{
```

```
dbl a, b, c;
    Line(): a(0), b(0), c(0) {}
    Line(dbl aa, dbl bb, dbl cc): a(aa), b(bb), c(cc) {}
    Line(const Pnt &A, const Pnt &p){ // it normalizes (a,b),
    \rightarrow important in d(), normalToP()
       Pnt n = (p-A).getRot90().getNorm();
        a = n.x, b = n.y, c = -(a * A.x + b * A.y);
    inline dbl d(const Pnt &p) const { return a*p.x + b*p.y +
    inline Pnt no() const {return Pnt(a, b);}
    inline Pnt normalToP(const Pnt &p) const { return Pnt(a,b)
    \rightarrow * (a*p.x + b*p.y + c); }
    inline void write() const{
      cout << fixed << (double)a << " " << (double)b << " " <<
      }:
```

## 29 GeometryInterTangent

```
void buildTangent(Pnt p1, dbl r1, Pnt p2, dbl r2, Line &1) {
\leftrightarrow // r1, r2 = radius with sign
    Pnt p = p2 - p1;
    1.c = r1;
    dbl c2 = p.len2(), c1 = sqrt(c2 - sqr(r2));
    1.a = (-p.x * (r1 - r2) + p.y * c1) / c2;
    1.b = (-p.y * (r1 - r2) - p.x * c1) / c2;
    1.c -= 1.no() * p1;
    assert(eq(l.d(p1), r1));
    assert(eq(1.d(p2), r2));
}
struct Circle {
    Pnt p;
    dbl r;
vector<Pnt> v; // to store intersection
// Intersection of two lines
int line_line(const Line &1, const Line &m){
    dbl z = m.a * 1.b - 1.a * m.b;
  dbl x = m.c * 1.b - 1.c * m.b;
  dbl y = m.c * l.a - l.c * m.a;
    if(fabs(z) > EPS){
        v.pb(Pnt(-x/z, y/z));
        return 1:
    }else if(fabs(x) > EPS || fabs(y) > EPS)
       return 0; // parallel lines
    else
        return 2; // same lines
}
// Intersection of Circle and line
void circle_line(const Circle &c, const Line &l){
    dbl d = 1.d(c.p);
    if(fabs(d) > c.r + EPS)
    if(fabs(fabs(d) / c.r - 1) < EPS)
       v.pb(c.p - 1.no() * d);
    else{
        dbl s = sqrt(fabs(sqr(c.r) - sqr(d)));
        v.pb(c.p - 1.no() * d + 1.no().getRot90() * s);
        v.pb(c.p - 1.no() * d - 1.no().getRot90() * s);
    }
}
```

#### 30 GeometrySimple

//for convex polygon

```
int sign(dbl a) { return (a > EPS) - (a < -EPS); }</pre>
// Checks, if point is inside the segment
inline bool inSeg(const Pnt &p, const Pnt &a, const Pnt &b) {
   return eq((p - a) % (p - b), 0) && leq((p - a) * (p - b),
}
// Checks, if two intervals (segments without ends) intersect
\hookrightarrow AND do not lie on the same line
inline bool subIntr(const Pnt &a, const Pnt &b, const Pnt &c,
return
            sign((b - a) \% (c - a)) * sign((b - a) \% (d - a))
            sign((d - c) \% (a - c)) * sign((d - c) \% (b - c))
}
// Checks, if two segments (ends are included) has an
\hookrightarrow intersection
inline bool checkSegInter(const Pnt &a, const Pnt &b, const
→ Pnt &c, const Pnt &d){
    return inSeg(c, a, b) || inSeg(d, a, b) || inSeg(a, c, d)
    \rightarrow || inSeg(b, c, d) || subIntr(a, b, c, d);
inline dbl area(vector<Pnt> p){
    dbl s = 0;
    int n = sz(p);
    p.pb(p[0]);
    forn(i, n)
        s += p[i + 1] \% p[i];
    p.pop_back();
    return abs(s) / 2;
}
// Check if point p is inside polygon <n, q[]>
int containsSlow(Pnt p, Pnt *z, int n){
    int cnt = 0;
    forn(j, n){
        Pnt a = z[j], b = z[(j + 1) \% n];
        if (inSeg(p, a, b))
            return -1; // border
        if (min(a.y, b.y) - EPS \le p.y \&\& p.y \le max(a.y, b.y)

→ EPS)

           cnt += (p.x < a.x + (p.y - a.y) * (b.x - a.x) /
            \hookrightarrow (b.y - a.y));
    }
    return cnt & 1; // 0 = outside, 1 = inside
}
```

```
//assume polygon is counterclockwise-ordered
bool containsFast(Pnt p, Pnt *z, int n) {
    Pnt o = z[0];
    if(gr((p - o) \% (z[1] - o), 0) || ls((p - o) \% (z[n - 1] - o))
    \rightarrow o), 0))
        return 0;
    int 1 = 0, r = n - 1;
    while(r - 1 > 1){
        int m = (1 + r) / 2;
        if(gr((p - o) % (z[m] - o), 0))
            r = m;
        else
            1 = m;
    }
    return leq((p - z[1]) % (z[r] - z[1]), 0);
}
// Checks, if point "i" is in the triangle "abc" IFF triangle
\hookrightarrow in CCW order
inline int isInTr(int i, int a, int b, int c){
    return
            gr((p[b] - p[a]) % (p[i] - p[a]), 0) &&
            gr((p[c] - p[b]) \% (p[i] - p[b]), 0) &&
            gr((p[a] - p[c]) % (p[i] - p[c]), 0);
}
```

#### 31 Halfplanes Intersection

```
const int maxn = (int)4e5 + 9;
const dbl eps = 1e-12;
dbl sqr( dbl x ) { return x * x; }
struct pnt{
 LL operator * ( pnt p ) { return (LL)x * p.y - (LL)y * p.x;
 LL operator ^ ( pnt p ) { return (LL)x * p.x + (LL)y * p.y;
  → }
  pnt ort() { return pnt(-y, x); }
  dbl ang() { return atan2(y, x); }
  LL d2() { return x * x + y * y; }
pnt st, v, p[maxn];
int n, sp, ss[maxn], ind[maxn], no[maxn], cnt[maxn], k = 0,
\hookrightarrow a[maxn], b[maxn];
dbl ang[maxn];
pnt Norm( int k ){ return (p[a[k]] - p[b[k]]).ort();}
void AddPlane( int i, int j ){
  a[k] = i, b[k] = j, ind[k] = k;
  ang[k] = Norm(k).ang();
}
bool angLess( int i, int j ){ return ang[i] < ang[j];}</pre>
void Unique()
{
  int i = 0, k2 = 0;
  while (i < k)
   int ma = ind[i], st = i;
    pnt no = Norm(ma);
    for (i++; i < k && fabs(ang[ind[st]] - ang[ind[i]]) < eps;</pre>
     if ((no ^ p[a[ma]]) < (no ^ p[a[ind[i]]]))</pre>
        ma = ind[i];
    ind[k2++] = ma;
```

```
k = k2;
dbl xx, yy, tmp;
#define BUILD(a1, b1, c1, i) \
 tmp = sqrt(a1 * a1 + b1 * b1); \
  a1 /= tmp, b1 /= tmp; \
 dbl\ c1 = -(a1 * p[a[i]].x + b1 * p[a[i]].y);
void FindPoint( int i, int j, dbl step = 0.0 ){
  BUILD(a1, b1, c1, i);
  BUILD(a2, b2, c2, j);
 xx = -(c1 * b2 - c2 * b1) / (a1 * b2 - a2 * b1);
 yy = (c1 * a2 - c2 * a1) / (a1 * b2 - a2 * b1);
 dbl no = sqrt(sqr(a1 + a2) + sqr(b1 + b2));
 xx += (a1 + a2) * step / no;
  yy += (b1 + b2) * step / no;
void TryShiftPoint( int i, int j, dbl step )
 FindPoint(i, j, step);
 forn(i, k){
    BUILD(a1, b1, c1, ind[i]);
   if (a1 * xx + b1 * yy + c1 < eps)
     return:
  puts("Possible");
  printf("%.201f %.201f\n", (double)xx, (double)yy);
  exit(0);
void PushPlaneIntoStack( int i )
  while (sp \ge 2 \&\& ang[i] - ang[ss[sp - 2]] + eps < M_PI){
   FindPoint(i, ss[sp - 2]);
    BUILD(a1, b1, c1, ss[sp - 1]);
   if ((a1 * xx + b1 * yy + c1) < -eps)
     break;
   sp--;
  ss[sp++] = i;
int main()
 scanf("%d", &n);
 forn(i, n)
   scanf("%d%d", &p[i].x, &p[i].y);
 p[n] = p[0];
  // Find set of planes
  forn(i, sp)
   AddPlane(max(ss[i], ss[i + 1]), min(ss[i], ss[i + 1]));
  forn(i, n - 1)
   AddPlane(i + 1, i);
  sort(ind, ind + k, angLess);
  int oldK = k;
  Unique();
```

```
forn(i, oldK)
  no[i] = i;
forn(i, k){
  int j = oldK + i, x = ind[i];
  ang[j] = ang[x] + 2 * M_PI;
  a[j] = a[x];
  b[j] = b[x];
  ind[i + k] = j, no[j] = x;
sp = 0;
forn(i, 2 * k)
 PushPlaneIntoStack(ind[i]);
forn(t, sp)
 if (++cnt[no[ss[t]]] > 1){
   TryShiftPoint(ss[t], ss[t - 1], 1e-5);
    break;
  }
return 0:
```

## 8 Graphs

#### 32 2-SAT

```
// MAXVAR - 2 * vars
int cntVar = 0, val[MAXVAR], usedSat[MAXVAR], comp[MAXVAR];
vi topsortSat:
vi g[MAXVAR], rg[MAXVAR];
inline int newVar() {
  cntVar++;
  return (cntVar - 1) * 2;
inline int Not(int v) { return v ^ 1; }
inline void Implies(int v1, int v2) { g[v1].pb(v2),
\rightarrow rg[v2].pb(v1); }
inline void Or(int v1, int v2) { Implies(Not(v1), v2),

    Implies(Not(v2), v1); }

inline void Nand(int v1, int v2) { Or(Not(v1), Not(v2)); }
inline void setTrue(int v) { Implies(Not(v), v); }
void dfs1(int v) {
 usedSat[v] = 1;
  for (int to : g[v])
   if (!usedSat[to]) dfs1(to);
  topsortSat.pb(v);
void dfs2(int v, int c) {
 comp[v] = c;
  for (int to : rg[v])
   if (!comp[to]) dfs2(to, c);
int getVal(int v) { return val[v]; }
// cntVar
bool solveSat() {
 forn(i, 2 * cntVar) usedSat[i] = 0;
  forn(i, 2 * cntVar)
   if (!usedSat[i]) dfs1(i);
  reverse(all(topsortSat));
 int c = 0;
```

```
for (int v : topsortSat)
  if (!comp[v]) dfs2(v, ++c);
forn(i, cntVar) {
  if (comp[2 * i] == comp[2 * i + 1]) return false;
  if (comp[2 * i] < comp[2 * i + 1]) val[2 * i + 1] = 1;
  else val[2 * i] = 1;
}
return true;
}</pre>
```

## 33 Bridges

```
int up[MAX_N], tIn[MAX_N], timer;
vector<vi> comps;
vi st:
struct Edge {
 int to, id;
  Edge(int _to, int _id) : to(_to), id(_id) {}
vector<Edge> g[MAX_N];
void newComp(int size = 0) {
  comps.emplace_back(); //
  while (sz(st) > size) {
   comps.back().pb(st.back());
    st.pop_back();
 }
}
void findBridges(int v, int parentEdge = -1) {
 if (up[v]) //
   return;
  up[v] = tIn[v] = ++timer;
  st.pb(v);
 for (Edge e : g[v]) {
   if (e.id == parentEdge)
     continue;
    int u = e.to;
   if (!tIn[u]) {
     int size = sz(st);
      findBridges(u, e.id);
     if (up[u] > tIn[v])
       newComp(size);
   }
   up[v] = min(up[v], up[u]);
 }
          find bridges newComp()
void run(int n) {
 forn (i, n)
   if (!up[i]) {
     findBridges(i):
      newComp();
}
```

#### 34 Cut Points

```
bool used[MAX_M];
int tIn[MAX_N], timer, isCut[MAX_N], color[MAX_M], compCnt;
vi st;

struct Edge {
  int to, id;
  Edge(int _to, int _id) : to(_to), id(_id) {}
};

vector<Edge> g[MAX_N];
```

```
int dfs(int v, int parent = -1) {
 tIn[v] = ++timer;
 int up = tIn[v], x = 0, y = (parent != -1);
 for (Edge p : g[v]) {
   int u = p.to, id = p.id;
   if (id != parent) {
     int t, size = sz(st);
     if (!used[id])
       used[id] = 1, st.push_back(id);
     if (!tIn[u]) { // not visited yet
       t = dfs(u, id);
       if (t >= tIn[v]) {
         ++x, ++compCnt;
         while (sz(st) != size) {
           color[st.back()] = compCnt;
           st.pop_back();
         }
       }
     } else
       t = tIn[u]:
     up = min(up, t);
   }
 }
 if (x + y >= 2)
   isCut[v] = 1; // v is cut vertex
 return up;
```

#### 35 Eulerian Cycle

```
struct Edge {
   int to, used;
   Edge(): to(-1), used(0) {}
   Edge(int v): to(v), used(0) {}
};

vector<Edge> edges;
vi g[MAX_N], res, ptr;
// ptr

void dfs(int v) {
   for(; ptr[v] < sz(g[v]);) {
     int id = g[v][ptr[v]++];
     if (!edges[id].used) {
       edges[id].used = edges[id ^ 1].used = 1;
       dfs(edges[id].to);
       res.pb(id); //
   }
}
   res.pb(v); // res
}</pre>
```

#### 36 Euler Tour Tree

```
inline int getSize(Node* root) { return root ? root->size : 0;
inline void recalc(Node* root) { root->size = getSize(root->1)
\hookrightarrow + getSize(root->r) + 1; }
set<pair<int, Node*>> edges[MAX_N];
Node* merge(Node *a, Node *b) {
  if (!a) return b;
  if (!b) return a;
  if (a->y < b->y) {
    a->r = merge(a->r, b);
    if (a->r) a->r->p = a;
    recalc(a);
    return a;
  b->1 = merge(a, b->1);
  if (b->1) b->1->p = b;
  recalc(b);
  return b:
void split(Node *root, Node *&a, Node *&b, int size) {
  if (!root) {
    a = b = nullptr;
    return;
  int lSize = getSize(root->1);
  if (lSize >= size) {
    split(root->1, a, root->1, size);
    if (root->l) root->l->p = root;
    b = root, b \rightarrow p = b;
  } else {
   split(root->r, root->r, b, size - lSize - 1);
    if (root->r) root->r->p = root;
    a = root, a \rightarrow p = a;
    a->p = a;
  }
  recalc(root);
}
inline Node* rotate(Node* root, int k) {
  if (k == 0) return root;
  Node *1, *r;
  split(root, 1, r, k);
  return merge(r, 1);
inline pair<Node*, int> goUp(Node* root) {
  int pos = getSize(root->1);
  while (root->p != root)
   pos += (root->p->r == root ? getSize(root->p->1) + 1 : 0),

→ root = root->p;

  return mp(root, pos);
inline Node* deleteFirst(Node* root) {
  Node* a;
  split(root, a, root, 1);
  edges[a->e.v].erase(mp(a->e.u, a));
  return root;
inline Node* getNode(int v, int u) {
 return edges[v].lower_bound(mp(u, nullptr))->snd;
inline void cut(int v, int u) {
  auto pV = goUp(getNode(v, u));
  auto pU = goUp(getNode(u, v));
```

```
int 1 = min(pV.snd, pU.snd), r = max(pV.snd, pU.snd);
 Node *a, *b, *c;
 split(pV.fst, a, b, 1);
 split(b, b, c, r - 1);
 deleteFirst(b);
 merge(a, deleteFirst(c));
inline pair<Node*, int> getRoot(int v) {
 return !sz(edges[v]) ? mp(nullptr, 0) :

    goUp(edges[v].begin()->snd);

inline Node* makeRoot(int v) {
 auto root = getRoot(v);
 return rotate(root.fst, root.snd);
inline Node* makeEdge(int v, int u) {
 Node* e = new Node(Edge(v, u));
 edges[v].insert(mp(u, e));
inline void link(int v, int u) {
 Node *vN = makeRoot(v), *uN = makeRoot(u);
 merge(merge(vN, makeEdge(v, u)), uN), makeEdge(u, v));
```

## 37 Hamilton Cycle

```
n*2^n
vi g[MAX_MASK];
int adj[MAX_MASK], dp[1 << MAX_MASK];</pre>
vi hamiltonCycle(int n) {
  fill(dp, dp + (1 << n), 0);
  forn (v, n) {
    adj[v] = 0;
    for (int to : g[v])
      adj[v] |= (1 << to);
  }
  dp[1] = 1;
  forn (mask, (1 << n))
   forn(v. n)
      if (mask & (1 << v) && dp[mask \hat{} (1 << v)] & adj[v])
        dp[mask] = (1 << v);
  vi ans:
  int mask = (1 << n) - 1, v;</pre>
  if (dp[mask] & adj[0]) {
    forab (i, 1, n)
      if ((1 << i) & (mask & adj[0]))</pre>
       v = i;
    ans.pb(v);
    mask ^= (1 << v);
    while(v) {
        if ((dp[mask] & (1 << i)) && (adj[i] & (1 << v))) {
          v = i:
          break;
      mask \hat{} = (1 << v);
      ans.pb(v);
   }
 }
  return ans;
```

#### 38 Karp with cycle

```
int d[MAX_N][MAX_N], p[MAX_N][MAX_N];
vi g[MAX_N], ans;
struct Edge {
  int a, b, w;
  Edge(int _a, int _b, int _w): a(_a), b(_b), w(_w) {}
vector<Edge> edges;
void fordBellman(int s, int n) {
  form (i, n + 1)
    forn (j, n + 1)
      d[i][j] = INF;
  d[0][s] = 0;
  forab (i, 1, n + 1)
    for (auto &e : edges)
     if (d[i-1][e.a] < INF && d[i][e.b] > d[i-1][e.a] +
      \hookrightarrow e.w)
       d[i][e.b] = d[i - 1][e.a] + e.w, p[i][e.b] = e.a;
ld karp(int n) {
  int s = n++:
  forn (i, n - 1)
    g[s].pb(sz(edges)), edges.pb(Edge(s, i, 0));
  fordBellman(s, n);
  ld ansValue = INF:
  int curV = -1, dist = -1;
  forn (v, n - 1)
    if (d[n][v] != INF) {
      ld curAns = -INF;
      int curPos = -1;
      forn(k, n)
       if (curAns <= (d[n][v] - d[k][v]) * (ld) (1) / (n -
        \hookrightarrow k))
          curAns = (d[n][v] - d[k][v]) * (ld) (1) / (n - k),
          \hookrightarrow curPos = k;
      if (ansValue > curAns)
        ansValue = curAns, dist = curPos, curV = v;
  if (curV == -1) return ansValue;
  for (int iter = n; iter != dist; iter--)
    ans.pb(curV), curV = p[iter][curV];
  reverse(all(ans));
  return ansValue;
39 Kuhn's algorithm
//
int n, m, paired[2 * MAX_N], used[2 * MAX_N];
vi g[MAX_N];
bool dfs(int v) {
 if (used[v]) return false;
  used[v] = 1;
  for (int to : g[v])
```

if (paired[to] == -1 || dfs(paired[to])) {
 paired[to] = v, paired[v] = to;

return true:

forn (i, n + m) paired[i] = -1;
for (int run = 1; run;) {

7

return false;

int ans = 0;

int kuhn() {

```
run = 0:
   fill(used, used + n + m, 0);
   forn(i, n)
     if (!used[i] && paired[i] == -1 && dfs(i))
       ans++, run = 1;
 return ans;
//
// Max
              -- A+. B-
// Min
              -- A-, B+
vi minCover, maxIndependent;
void dfsCoverIndependent(int v) {
 if (used[v]) return;
 used[v] = 1;
 for (int to : g[v])
   if (!used[to])
     used[to] = 1, dfsCoverIndependent(paired[to]);
void findCoverIndependent() {
 fill(used, used + n + m, 0);
 forn (i, n)
   if (paired[i] == -1)
     dfsCoverIndependent(i);
   if (used[i]) maxIndependent.pb(i);
   else minCover.pb(i);
 forab (i, n, n + m)
   if (used[i]) minCover.pb(i);
    else maxIndependent.pb(i);
```

#### 40 LCA

```
int tin[MAX_N], tout[MAX_N], up[MAX_N][MAX_LOG];
vi g[MAX_N];
int curTime = 0;
void dfs(int v, int p) {
 up[v][0] = p;
 forn (i, MAX_LOG - 1)
   up[v][i + 1] = up[up[v][i]][i];
 tin[v] = curTime++;
 for (int u : g[v])
   if (u != p)
     dfs(u, v);
  tout[v] = curTime++;
int isUpper(int v, int u) {
 return tin[v] <= tin[u] && tout[v] >= tout[u];
int lca(int v, int u) {
 if (isUpper(u, v)) return u;
 fornr (i, MAX_LOG)
   if (!isUpper(up[u][i], v))
     u = up[u][i];
 return up[u][0];
void init() {
 dfs(0, 0);
```

## 41 LCA offline (Tarjan)

```
vi g[MAX_N], q[MAX_N];
int pr[MAX_N], ancestor[MAX_N], used[MAX_N];
int get(int v) {
 return v == pr[v] ? v : pr[v] = get(pr[v]);
void unite(int v, int u, int anc) {
 v = get(v), u = get(u);
 pr[u] = v, ancestor[v] = anc;
void dfs(int v) {
 used[v] = 1;
 for (int u : g[v])
   if (!used[u])
     dfs(u), unite(v, u, v);
 for (int u : q[v])
   if (used[u])
     ancestor[get(u)]; // handle answer somehow
}
void init(int n) {
 forn (i, n) pr[i] = i, ancestor[i] = i;
 dfs(0);
```

#### 9 Math

#### **42 CRT** (**KTO**)

```
vi crt(vi a, vi mod) {
   int n = sz(a);
   vi x(n);
   forn (i, n) {
      x[i] = a[i];
      forn (j, i) {
       x[i] = inverse(mod[j], mod[i]) * (x[i] - x[j]) % mod[i];
      if (x[i] < 0) x[i] += mod[i];
    }
}
return x;
}</pre>
```

#### 43 Discrete Logarithm

```
// Returns x: a^x = b \pmod{mod} or -1, if no such x exists
int discreteLogarithm(int a, int b, int mod) {
 int sq = sqrt(mod);
 int sq2 = mod / sq + (mod % sq ? 1 : 0);
 vector<pii> powers(sq2);
 forn (i, sq2)
   powers[i] = mp(power(a, (i + 1) * sq, mod), i + 1);
 sort(all(powers));
 forn (i, sq + 1) {
   int cur = power(a, i, mod);
   cur = (cur * 111 * b) % mod;
   auto it = lower_bound(all(powers), mp(cur, 0));
   if (it != powers.end() && it->fst == cur)
     return it->snd * sq - i;
 }
 return -1;
```

#### 44 Discrete Root

```
// Returns x: x~k = a mod mod, mod is prime
int discreteRoot(int a, int k, int mod) {
  if (a == 0)
```

}

row++;

```
return 0:
  int g = primitiveRoot(mod);
                                                                  int single = 1;
 int y = discreteLogarithm(power(g, k, mod), a, mod);
                                                                  forn (i, m)
                                                                    if (par[i] != -1) ans[i] = a[par[i]][m] / a[par[i]][i];
 return power(g, y, mod);
                                                                    else single = 0;
                                                                  forn (i, n) {
                                                                    double cur = 0;
45
    Eratosthenes
                                                                    for (int j = 0; j < m; j++)
                                                                       cur += ans[j] * a[i][j];
vi eratosthenes(int n) {
                                                                    if (abs(cur - a[i][m]) > EPS)
 vi minDiv(n + 1, 0);
                                                                       return 0;
 minDiv[1] = 1;
 forab (i, 2, n + 1)
                                                                  if (!single)
    if (minDiv[i] == 0)
                                                                    return 2;
     for (int j = i; j \le n; j += i)
                                                                  return 1;
       if (minDiv[j] == 0) minDiv[j] = i;
 return minDiv:
                                                                    Gauss binary
vi eratosthenesLinear(int n) {
 vi minDiv(n + 1, 0), primes;
                                                                const int MAX = 1024;
 minDiv[1] = 1;
 forab (i, 2, n + 1) {
                                                                int gaussBinary(vector<bitset<MAX>> a, int n, int m) {
   if (minDiv[i] == 0)
                                                                  int row = 0, col = 0;
     minDiv[i] = i, primes.pb(i);
                                                                  vi par(m, -1);
   for (int j = 0; j < sz(primes) && primes[j] <= minDiv[i]
                                                                  for (col = 0; col < m && row < n; col++) {
    int best = row;
     minDiv[i * primes[j]] = primes[j];
                                                                     for (int i = row; i < n; i++)
 }
                                                                      if (a[i][col] > a[best][col])
 return minDiv;
                                                                        best = i;
                                                                    if (a[best][col] == 0)
                                                                      continue;
                                                                    par[col] = row;
46 Factorial
                                                                    swap(a[row], a[best]);
                                                                    forn (i, n)
// Returns pair (rem, deg), where rem = n! % mod,
                                                                     if (i != row && a[i][col])
// deg = k: mod % / n!, mod is prime, O(mod log mod)
pii fact(int n, int mod) {
                                                                          a[i] ^= a[row];
                                                                    row++:
 int rem = 1, deg = 0, nCopy = n;
 while (nCopy) nCopy /= mod, deg += nCopy;
                                                                  }
                                                                  vi ans(m, 0);
 while (n > 1) {
                                                                  forn (i, m)
   rem = (rem * ((n / mod) % 2 ? -1 : 1) + mod) % mod;
                                                                    if (par[i] != -1)
   for (int i = 2; i <= n % mod; i++)
                                                                      ans[i] = a[par[i]][n] / a[par[i]][i];
    rem = (rem * 111 * i) % mod;
                                                                  bool ok = 1;
   n /= mod;
                                                                  forn (i, n) {
 7
                                                                    int cur = 0;
 return mp(rem % mod, deg);
                                                                    forn (j, m) cur ^= (ans[j] & a[i][j]);
                                                                    if (cur != a[i][n]) ok = 0;
                                                                  }
47
     Gauss
                                                                  return ok;
const double EPS = 1e-9;
int gauss(double **a, int n, int m) { // n is number of
                                                                49 Gcd
\hookrightarrow equations, m is number of variables
 int row = 0, col = 0;
                                                                int gcd(int a, int b) {
 vi par(m, -1);
                                                                  return b ? gcd(b, a % b) : a;
 vector<double> ans(m, 0);
 for (col = 0; col < m && row < n; col++) {
   int best = row;
                                                                int gcd(int a, int b, int &x, int &y) {
    for (int i = row; i < n; i++)</pre>
                                                                  if (b == 0) {
     if (abs(a[i][col]) > abs(a[best][col]))
                                                                   x = 1, y = 0;
       best = i;
                                                                    return a;
    if (abs(a[best][col]) < EPS) continue;</pre>
    par[col] = row:
                                                                  int g = gcd(b, a \% b, x, y), newX = y;
    forn (i, m + 1) swap(a[row][i], a[best][i]);
                                                                  y = x - a / b * y;
    forn (i, n)
                                                                  x = newX;
     if (i != row) {
                                                                  return g;
       double k = a[i][col] / a[row][col];
       for (int j = col; j \le m; j++)
         a[i][j] -= k * a[row][j];
                                                                void diophant(int a, int b, int c, int &x, int &y) {
```

int g = gcd(a, b, x, y);
if (c % g != 0) return;

```
x *= c / g, y *= c / g;
 // next solutions: x += b / g, y -= a / g
                                                                 int inversePhi(int a, int mod) {
                                                                  return power(a, phi(mod) - 1, mod);
int inverse(int a, int mod) { // Returns -1, if a and mod are
\hookrightarrow not coprime
                                                                 53 Pollard
 int x, y;
 int g = gcd(a, mod, x, y);
                                                                 inline void pollardFoo(ll& x, ll mod) {
 return g == 1 ? (x \% mod + mod) \% mod : -1;
                                                                  x = (mul(x, x, mod) + 1) \% mod;
vi inverseForAll(int mod) {
                                                                 vector<pair<11, int>> factorize(11 n) {
 vi r(mod, 0);
                                                                   if (n == 1) return {};
 r[1] = 1;
                                                                   if (isPrimeMillerRabin(n)) return {mp(n, 1)};
 for (int i = 2; i < mod; i++)
                                                                   if (n <= 100) {
   r[i] = (mod - r[mod % i]) * (mod / i) % mod;
                                                                     vector<pair<11, int>> ans;
 return r;
                                                                     for (int i = 2; i * i <= n; i++)
                                                                      if (n % i == 0) {
                                                                        int cnt = 0;
                                                                        while (n \% i == 0) n /= i, cnt++;
50 Gray
                                                                        ans.pb(mp(i, cnt));
int gray(int n) {
 return n ^ (n >> 1);
                                                                     if (n != 1) ans.pb(mp(n, 1));
                                                                     sort(all(ans));
                                                                     return ans;
int revGray(int n) {
 int k = 0;
                                                                   while (1) {
 for (; n; n >>= 1) k ^= n;
                                                                    ll a = rand() % n, b = a;
 return k;
                                                                     while (1) {
                                                                       pollardFoo(a, n), pollardFoo(b, n), pollardFoo(b, n);
                                                                       11 g = \_gcd(abs(a-b), n);
                                                                       if (g != 1) {
    Miller-Rabin Test
                                                                        if (g == n)
vector <int> primes = {2, 3, 5, 7, 11, 13, 17, 19, 23};
                                                                           break;
                                                                         auto ans1 = factorize(g);
bool isPrimeMillerRabin(ll n) {
                                                                         auto ans2 = factorize(n / g);
                                                                         vector<pair<11, int>> ans;
 int k = 0;
                                                                         ans1.insert(ans1.end(), all(ans2));
 11 t = n - 1;
 while (t \% 2 == 0) k++, t /= 2;
                                                                         sort(all(ans1));
                                                                         for (auto np : ans1)
 for (auto p : primes) {
                                                                          if (sz(ans) == 0 || np.fst != ans.back().fst)
   11 g = _{gcd(n, (11) p)};
   if (g > 1 && g < n) return 0;
                                                                            ans.pb(np);
   if (g == n) return 1;
                                                                           else
                                                                            ans.back().snd += np.snd:
   ll b = power(p, t, n), last = n - 1;
   bool was = 0;
                                                                        return ans;
                                                                       }
   forn (i, k + 1) {
                                                                     }
     if (b == 1 && last != n - 1)
                                                                   }
       return 0;
                                                                   assert(0);
     if (b == 1) {
       was = 1;
       break;
                                                                 54 Power And Mul
     last = b, b = mul(b, b, n);
                                                                 inline 11 fix(11 a, 11 mod) { // a in [0, 2 * mod)
   if (!was) return 0;
                                                                  if (a >= mod) a -= mod;
                                                                   return a;
 return 1;
                                                                 }
                                                                 // Returns (a * b) \% mod, 0 <= a < mod, 0 <= b < mod
                                                                 11 mulSlow(11 a, 11 b, 11 mod) {
52 Phi
                                                                   if (!b) return 0;
                                                                  ll c = fix(mulSlow(a, b / 2, mod) * 2, mod);
int phi(int n) {
                                                                   return b & 1 ? fix(c + a, mod) : c;
 int result = n:
 for (int i = 2; i * i <= n; i++)
   if (n % i == 0) {
     while (n \% i == 0) n /= i;
                                                                 11 mul(11 a, 11 b, 11 mod) {
     result -= result / i;
                                                                  11 q = (1d) a * b / mod;
   }
                                                                   11 r = a * b - mod * q;
 if (n > 1) result -= result / n;
                                                                   while (r < 0) r += mod;
 return result;
                                                                   while (r \ge mod) r -= mod;
                                                                   return r;
```

```
int power(int a, int n, int mod) {
 if (!n) return 1;
 int b = power(a, n / 2, mod);
 b = (b * 111 * b) \% mod;
 return n & 1 ? (a * 111 * b) % mod : b;
11 powerLL(11 a, 11 n, 11 mod) {
 if (!n) return 1;
 11 b = power(a, n / 2, mod);
 b = mul(b, b, mod);
 return n & 1 ? mul(a, b, mod) : b;
int powerFast(int a, int n, int mod) {
 int res = 1;
 while (n) {
   if (n & 1)
    res = (res * 111 * a) % mod;
   a = (a * 111 * a) \% mod;
   n /= 2;
 return res;
```

#### 55 Primitive Root

```
int primitiveRoot(int mod) { // Returns -1 if no primitive

→ root exists

 vi fact;
 int ph = phi(mod);
 int n = mod;
 for (int i = 2; i * i <= n; i++) {
   if (n % i == 0) {
     fact.pb(i);
     while (n \% i == 0) n /= i;
 }
 if (n > 1) fact.pb(n);
 forab (i, 2, mod + 1) {
   bool ok = 1;
   for (int j = 0; j < sz(fact) && ok; j++)
     ok &= power(i, ph / fact[j], mod) != 1;
   if (ok) return i;
 }
 return -1;
```

## 56 Simpson

# 10 Strings

#### 57 Aho-Corasick

```
const int ALPHA = 26;
const int MAX_N = 1e5;
```

```
struct Node {
  int next[ALPHA], term; //
  int go[ALPHA], suf, p, pCh; //
  Node(): term(0), suf(-1), p(-1) {
   fill(next, next + ALPHA, -1);
    fill(go, go + ALPHA, -1);
 }
};
Node g[MAX_N];
int last;
void add(const string &s) {
 int now = 0;
  for(char x : s) {
   if (g[now].next[x - 'a'] == -1) {
     g[now].next[x - 'a'] = ++last;
      g[last].p = now, g[last].pCh = x;
   now = g[now].next[x - 'a'];
  g[now].term = 1;
int go(int v, int c);
int getLink(int v) {
 if (g[v].suf == -1) {
   if (!v || !g[v].p) g[v].suf = 0;
   else g[v].suf = go(getLink(g[v].p), g[v].pCh);
  return g[v].suf;
int go(int v, int c) {
  if (g[v].go[c] == -1) {
    if (g[v].next[c] != -1) g[v].go[c] = g[v].next[c];
    else g[v].go[c] = !v ? 0 : go(getLink(v), c);
  return g[v].go[c];
```

#### 58 Prefix-function

```
vi prefix(const string &s) {
   int n = sz(s);
   vi pr(n);
   forab (i, 1, n + 1) {
      int j = pr[i - 1];
      while (j > 0 && s[i] != s[j]) j = pr[j - 1];
      if (s[i] == s[j]) j++;
      pr[i] = j;
   }
   return pr;
}
```

#### 59 Z-function

```
vi z(const string% s) {
   int n = sz(s);
   vi z(n);
   for (int i = 1, 1 = 0, r = 0; i < n; i++) {
      if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
      while (i + z[i] < n && s[z[i]] == s[i + z[i]]) z[i]++;
      if (i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
   }
   return z;
}
```

#### 60 Hashes

```
const int P = 239017;
inline int add(int a, int b, int m) {
 a += b;
 return a >= m ? a - m : a;
inline int sub(int a, int b, int m) {
 a -= b;
 return a < 0 ? a + m : a;
const int MOD_X = 1e9 + 9, MOD_Y = 1e9 + 7;
// using H = unsigned long long;
struct H {
 int x, y;
 H(): x(0), y(0) {}
 H(int _x): x(_x), y(_x) {}
 H(int _x, int _y): x(_x), y(_y) {}
  inline H operator+(const H& h) const { return H(add(x, h.x,
  \hookrightarrow MOD_X), add(y, h.y, MOD_Y)); }
  inline H operator-(const H& h) const { return H(sub(x, h.x,

    MOD_X), sub(y, h.y, MOD_Y)); }

  inline H operator*(ll k) const { return H(int((x * k) %
  \hookrightarrow MOD_X), int((y * k) % MOD_Y)); }
  inline H operator*(const H& h) const{ return H(int((11(x) *

    h.x) % MOD_X), int((11(y) * h.y) % MOD_Y)); }

 inline bool operator==(const H& h) const { return x == h.x
  inline bool operator!=(const H& h) const { return x != h.x
  \rightarrow || y != h.y; }
  \hookrightarrow (x == h.x \&\& y < h.y); }
  explicit inline operator ll() const { return ll(x) * MOD_Y +

    y + 1; } // > 0

H deg[MAX_N], h[MAX_N];
inline H get(int 1, int r) { return h[r] - h[1] * deg[r - 1];
void init(const string& s) {
 int n = sz(s);
  deg[0] = 1;
 forn (i, n)
    h[i + 1] = h[i] * P + s[i], deg[i + 1] = deg[i] * P;
```

#### 61 Manaker

```
void manaker(const string& s, int *z0, int *z1) {
 int n = sz(s);
 forn (t, 2) {
   int *z = t ? z1 : z0, 1 = -1, r = -1; // [l..r]
   forn (i, n - t) {
     int k = 0;
     if (r > i + t) {
       int j = 1 + (r - i - t);
       k = min(z[j], j - 1);
     while (i - k >= 0 \&\& i + k + t < n \&\& s[i - k] == s[i +
      \hookrightarrow k + t])
      k++;
     z[i] = k;
     if (k \&\& i + k + t > r)
        1 = i - k + 1, r = i + k + t - 1;
 }
```

#### **62** Palindromic Tree

}

```
const int ALPHA = 26;
struct Vertex {
 int suf, len, next[ALPHA];
  Vertex() { fill(next, next + ALPHA, 0); }
int vn, v;
Vertex t[MAX_N + 2];
int n, s[MAX_N];
int get(int i) { return i < 0 ? -1 : s[i]; }</pre>
void init() {
 t[0].len = -1, vn = 2, v = 0, n = 0;
void add(int ch) {
  s[n++] = ch;
 while (v != 0 && ch != get(n - t[v].len - 2))
   v = t[v].suf;
  int& r = t[v].next[ch];
  if (!r) {
   t[vn].len = t[v].len + 2;
   if (!v) t[vn].suf = 1;
   else {
      v = t[v].suf;
     while (v != 0 && ch != get(n - t[v].len - 2))
       v = t[v].suf;
      t[vn].suf = t[v].next[ch];
    r = vn++;
  }
  v = r:
```

#### 63 Suffix Array (+stable)

```
int sLen, num[MAX_N + 1];
char s[MAX_N + 1];
int p[MAX_N], col[MAX_N], inv[MAX_N], lcp[MAX_N];
inline int mod(int x) {
 return (x + sLen) % sLen;
void buildArray(int n) {
  sLen = n:
 int ma = max(n, 256);
 forn (i, n)
   col[i] = s[i], p[i] = i;
 for (int k2 = 1; k2 / 2 < n; k2 *= 2) {
   int k = k2 / 2;
    memset(num, 0, sizeof(num));
   forn (i, n) num[col[i] + 1]++:
    forn (i, ma) num[i + 1] += num[i];
   forn (i, n)
     inv[num[col[mod(p[i] - k)]]++] = mod(p[i] - k);
   int cc = 0;
    forn (i, n) {
     bool add = col[inv[i]] != col[inv[i - 1]];
      add |= col[mod(inv[i] + k)] != col[mod(inv[i - 1] + k)];
     if (i && add) cc++;
     num[inv[i]] = cc;
    forn (i, n) p[i] = inv[i], col[i] = num[i];
```

```
}
 memset(num, 0, sizeof(num));
  forn (i, n) num[col[i] + 1]++;
  forn (i, ma) num[i + 1] += num[i];
  forn (i, n) inv[num[col[i]]++] = i;
  forn (i, n) p[i] = inv[i];
  forn (i, n) inv[p[i]] = i;
void buildLCP(int n) {
 int len = 0:
  forn (ind, n){
   int i = inv[ind];
    len = max(0, len - 1);
    if (i != n - 1)
      while (len < n && s[mod(p[i] + len)] == s[mod(p[i + 1] +
      \hookrightarrow len)])
        len++:
    lcp[i] = len;
    if (i != n - 1 && p[i + 1] == n - 1) len = 0;
```

#### 64 Suffix Automaton

```
struct Vx {
    static const int AL = 26;
    int len, suf;
    int next[AL];
    Vx() {}
    Vx(int 1, int s): len(1), suf(s) {}
};
    static const int MAX_LEN = 1e5 + 100, MAX_V = 2 * MAX_LEN;
    int last, vcnt;
    Vx v[MAX_V];
    SA() { vcnt = 1, last = newV(0, 0); } // root = vertex
    \hookrightarrow with number 1
    int newV(int len, int suf){
        v[vcnt] = Vx(len, suf);
        return vcnt++;
    }
    int add(char ch) {
        int p = last, c = ch - 'a';
        last = newV(v[last].len + 1, 0);
        while (p && !v[p].next[c]) // added p &&
            v[p].next[c] = last, p = v[p].suf;
        if (!p)
            v[last].suf = 1;
        else {
            int q = v[p].next[c];
            if (v[q].len == v[p].len + 1) v[last].suf = q;
                int r = newV(v[p].len + 1, v[q].suf);
                v[last].suf = v[q].suf = r;
                memcpy(v[r].next, v[q].next,

    sizeof(v[r].next));
                while (p \&\& v[p].next[c] == q)
                    v[p].next[c] = r, p = v[p].suf;
        }
        return last;
    }
};
```

#### 65 Suffix Tree

const int MAX L=1e5+10:

```
char S[MAX_L];
int L:
struct Node;
struct Pos;
typedef Node *pNode;
typedef map<char,pNode> mapt;
struct Node{
  pNode P,link;
  int L,R;
  mapt next;
  Node():P(NULL),link(this),L(0),R(0){}
  Node(pNode P,int L,int R):P(P),link(NULL),L(L),R(R){}
  inline int elen() const{return R-L;}
  inline pNode add_edge(int L,int R){return next[S[L]]=new
  → Node(this,L,R);}
}:
struct Pos{
  pNode V;
  int up;
  Pos():V(NULL),up(0){}
  Pos(pNode V, int up):V(V),up(up){}
  pNode split_edge() const{
   if(!up)
     return V;
    int L=V->L, M=V->R-up;
    pNode P=V->P, n=new Node(P,L,M);
    P->next[S[L]]=n;
    n->next[S[M]]=V;
    V->P=n, V->L=M;
    return n:
  Pos next_char(char c) const{
      return S[V->R-up]==c ? Pos(V,up-1) : Pos();
      mapt::iterator it=V->next.find(c);
      return it==V->next.end() ? Pos() :
      \hookrightarrow Pos(it->snd,it->snd->elen()-1);
    }
 }
};
Pos go_down(pNode V,int L,int R){
  if(L==R)
   return Pos(V,0);
  while(1){
    V=V->next[S[L]];
    L+=V->elen();
    if(L>=R)
      return Pos(V,L-R);
  }
}
inline pNode calc_link(pNode &V){
 if(!V->link)
    V->link=go\_down(V->P->link,V->L+!V->P->P,V->R).split\_edge();
  return V->link;
Pos add_char(Pos P,int k){
  while(1){
    Pos p=P.next_char(S[k]);
    if(p.V)
```

```
return p;
pNode n=P.split_edge();
n->add_edge(k,MAX_L);
if(!n->P)
    return Pos(n,0);
P=Pos(calc_link(n),0);
}

pNode Root;
void make_tree(){
    Root=new Node();
    Pos P(Root,0);
    forn(i,L)
        P=add_char(P,i);
}
```

#### 11 C++ Tricks

#### 66 Fast allocation

```
const int MAX_MEM = 1e8;
int mpos = 0;
char mem[MAX_MEM];
inline void* operator new(size_t n) {
   char *res = mem + mpos;
   mpos += n;
   assert(mpos <= MAX_MEM);
   return (void*) res;
}
inline void operator delete(void*) {}
inline void* operator new[](size_t) { assert(0); }
inline void operator delete[](void*) { assert(0); }</pre>
```

#### 67 Hash of pair

#### 68 Ordered Set

#### 69 Hash Map

#### 70 Fast I/O (short)

```
inline int readChar();
inline int readInt();
template <class T> inline void writeInt(T x);
inline int readChar() {
 int c = getchar();
  while (c <= 32)
   c = getchar();
 return c;
inline int readInt() {
 int s = 0, c = readChar(), x = 0;
  if (c == '-')
   s = 1, c = readChar();
  while ('0' <= c && c <= '9')
   x = x * 10 + c - '0', c = readChar();
 return s ? -x : x;
template <class T> inline void writeInt(T x) {
 if (x < 0)
   putchar('-'), x = -x;
  char s[24]:
  int n = 0:
  while (x \mid | \mid !n)
    s[n++] = '0' + x \% 10, x /= 10;
  while (n--)
    putchar(s[n]);
```

#### **71 Fast I/O (long)**

```
template <class T = int> inline T readInt();
inline double readDouble();
inline int readUInt();
inline int readChar();
inline void readWord(char *s);
inline bool readLine(char *s); // do not save '\n'
inline bool isEof():
inline int peekChar();
inline bool seekEof();
template <class T> inline void writeInt(T x, int len);
template <class T> inline void writeUInt(T x, int len);
template <class T> inline void writeInt(T x) { writeInt(x,

→ -1): }:
template <class T> inline void writeUInt(T x) { writeUInt(x,
→ -1); };
inline void writeChar(int x);
inline void writeWord(const char *s);
inline void writeDouble(double x, int len = 0);
inline void flush();
const int BUF_SIZE = 4096;
char buf[BUF_SIZE];
int bufLen = 0, pos = 0;
```

```
inline bool isEof() {
 if (pos == bufLen) {
   pos = 0, bufLen = fread(buf, 1, BUF_SIZE, stdin);
   if (pos == bufLen)
     return 1;
 7
 return 0:
inline int getChar() {
 return isEof() ? -1 : buf[pos++];
}
inline int peekChar() {
 return isEof() ? -1 : buf[pos];
inline bool seekEof() {
  while ((c = peekChar()) != -1 \&\& c <= 32)
 return c == -1;
inline int readChar() {
 int c = getChar();
  while (c != -1 \&\& c <= 32)
   c = getChar();
 return c:
inline int readUInt() {
 int c = readChar(), x = 0;
  while ('0' <= c && c <= '9')
   x = x * 10 + c - '0', c = getChar();
 return x;
}
template <class T>
inline T readInt() {
  int s = 1, c = readChar();
 T x = 0;
 if (c == '-')
   s = -1, c = getChar();
  while ('0' <= c && c <= '9')
   x = x * 10 + c - '0', c = getChar();
 return s == 1 ? x : -x;
}
inline double readDouble() {
 int s = 1, c = readChar();
  double x = 0;
 if (c == '-')
   s = -1, c = getChar();
  while ('0' <= c && c <= '9')
 x = x * 10 + c - '0', c = getChar();
  if (c == '.') {
   c = getChar();
    double coef = 1;
   while ('0' <= c && c <= '9')
     x += (c - '0') * (coef *= 1e-1), c = getChar();
 }
  return s == 1 ? x : -x;
inline void readWord(char *s) {
 int c = readChar();
  while (c > 32)
   *s++ = c, c = getChar();
  *s = 0;
```

```
inline bool readLine(char *s) {
 int c = getChar();
  while (c != '\n' \&\& c != -1)
   *s++ = c, c = getChar();
  *s = 0;
 return c != -1:
int writePos = 0:
char writeBuf[BUF_SIZE];
inline void writeChar(int x) {
 if (writePos == BUF_SIZE)
   fwrite(writeBuf, 1, BUF_SIZE, stdout), writePos = 0;
  writeBuf[writePos++] = x;
inline void flush() {
 if (writePos)
    fwrite(writeBuf, 1, writePos, stdout), writePos = 0;
template <class T>
inline void writeInt(T x, int outputLen) {
 if (x < 0)
    writeChar('-'), x = -x;
 char s[24];
  int n = 0;
  while (x \mid | \mid !n)
   s[n++] = '0' + x \% 10, x /= 10;
  while (n < outputLen)</pre>
   s[n++] = '0';
  while (n--)
    writeChar(s[n]);
template <class T>
inline void writeUInt(T x, int outputLen) {
 char s[24];
  int n = 0;
 while (x \mid | !n)
   s[n++] = '0' + char(x \% 10), x /= 10;
  while (n < outputLen)</pre>
   s[n++] = '0';
  while (n--)
    writeChar(s[n]);
inline void writeWord(const char *s) {
  while (*s)
    writeChar(*s++);
inline void writeDouble(double x, int outputLen) {
 if (x < 0)
    writeChar('-'), x = -x;
  int t = (int) x;
  writeUInt(t), x -= t;
  writeChar('.');
  for (int i = outputLen - 1; i > 0; i--) {
   x *= 10;
    t = std::min(9, (int) x);
    writeChar('0' + t), x -= t;
  t = std::min(9, (int)(x + 0.5));
  writeChar('0' + t);
```

\_

## 12 Notes

#### 72 Работа с деревьями

Приемы для работы с деревьями:

- 1. Двоичные подъемы
- 2. Поддеревья как отрезки Эйлерова обхода
- 3. Вертикальные пути в Эйлеровом обходе (на ребрах вниз +k, на ребрах вверх -k).
- Храним в вершине значение функции на пути от корня до нее, дальше LCA.
- 5. Спуск с DFS, поддерживаем ДО на пути до текущей вершины.
- 6. Heavy-light decomposition
- 7. Centroid decomposition
- 8. Корневая по запросам
- 9. Тяжелые/легкие вершины
- 10. DFS  $\rightarrow$  дерево блоков, размеры  $\in [K..2K]$
- 11. У вершины не более  $O(\sqrt{N})$  разных поддеревьев
- 12. Сумма размеров поддеревьев без тяжелого ребенка  $O(n \log n)$
- 13. Сумма глубин поддеревьев без глубокого ребенка O(n)

#### 73 Маски

Считаем динамику по маскам за  $O(2^n \cdot n)$  f[mask] = sum по submask g[submask]. dp[mask][i] — значение динамики для маски mask, если младшие i бит в ней зафиксированы (то есть мы не можем удалять оттуда). Ответ в dp[mask][0]. dp[mask][len] = g[mask]. Если i-ый бит 0, то dp[mask][i] = dp[mask][i+1], иначе dp[mask][i] = dp[mask][i+1] + dp[mask](1 << i)[i+1].

Старший бит: предподсчет.

Младший бит:  $x \& \sim (-x)$ 

Чтобы по степени двойки получить логарифм, можно воспользоваться тем, что все степени двойки имеют разный остаток по модулю 67.

#### 74 Гранди

Теорема Шпрага-Гранди: берем mex всех значений функции Гранди по состояниям, в которые можем перейти из данного. Если сумма независимых игр, то значение функции Гранди равно хог значений функций Гранди по всем играм. Бывает полезно вывести первые п значений и поискать закономерность. Часто сводится к xor по чему-нибудь.

#### 75 Потоки

Потоки:

Name	Asympthotic
Ford-Fulkerson	$O( f  \cdot E)$
Ford-Fulkerson with scaling	$O(\log  f  \cdot E^2)$
Edmonds-Karp	$O(V \cdot E^2)$
Dinic	$O(V^2 \cdot E)$
Dinic with scaling	$O(V \cdot E \cdot \log C)$
Dinic on bipartite graph	$O(E\sqrt{V})$
Dinic on unit network	$O(E\sqrt{E})$

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L-R потоки: Есть граф с недостатками или избытками в каждой вершине. Создаем фиктивные исток и сток (из истока все ребра в избытки, из недостатков все ребра в сток). Теперь пусть у нас есть L-R граф, для каждого ребра  $e\ (v \to u)$  известны  $L_e$  и  $R_e$ . Добавим в v избыток  $L_e$ , в u недостаток  $L_e$ , а пропускную способность сделаем  $R_e-L_e$ . Получили решение задачи о LR-циркуляции. Если у нас обычный граф с истоком и стоком, то добавляем бесконечное ребро из стока в сток и ищем циркуляцию. Таким образом нашли удовлетворяющий условиям LR-поток. Если хотим максимальный поток, то на остаточной сети запускаем поиск максимального потока. В новом графе в прямую сторону пропускная способность равна  $R_e-f_e$ , в обратную  $f_e-L_e$ .

MinCostCirculation: Пока есть цикл отрицательного веса, запускаем алгоритм Карпа и пускаем максимальный поток по найденному циклу.

#### 76 ДП

Табличка с оптимизациями для динамики:

Name	Original recurrence	From
	Sufficient Condition	То
CHT1	$dp[i] = \min_{i < i} dp[j] + b[j] \cdot a[i]$	$O(n^2)$
	$b[j] \geqslant b[j+1] \mid\mid a[i] \leqslant a[i+1]$	O(n)
CHT2	$dp[i][j] = \min_{k < j} dp[i-1][k] + b[k] \cdot a[j]$	$O(kn^2)$
	$b[k] \geqslant b[k+1] \mid\mid a[j] \leqslant a[j+1]$	O(kn)
D&C	$dp[i][j] = \min_{k < j} dp[i-1][k] + c[k][j]$	$O(kn^2)$
	$p[i,j] \leqslant p[i,j+1]$	$O(kn\log n)$
Knuth	$dp[i][j] = \min_{i < k < j} dp[i][k] + dp[k][j] + c[i][j]$	$O(n^3)$
	$p[i, j-1] \leqslant p[i, j] \leqslant p[i+1, j]$	$O(n^2)$
IOI	$f_n(k)$ — best for fixed k	$O(k^{(2)}n)$
	$f_n$ — convex, add penalty $\lambda \cdot k$	$O(n \log C)$

#### 77 Комбинаторика

Биномиальные коэффициенты:

Теорема Люка для биномиальных коэффициентов: Хотим посчитать  $C_n^k$ , разложим в р-ичной системе счисления,  $n=(n_0,n_1,\dots), k=(k_0,k_1,\dots)$ .  $ans=C_{n_0}^{k_0}\cdot C_{n_1}^{k_1}\cdot\dots$ 

Способы вычисления  $C_n^k$ :

1.  $C_n^k = C_{n-1}^k + C_{n-1}^{k-1}$ precalc:  $O(n^2)$ , query: O(1).

2.  $C_n^k = \frac{n!}{k!(n-k)!}$ , предподсчитываем факториалы precalc:  $O(n \log n)$ , query:  $O(\log n)$ 

3. Теорема Люка

precalc:  $O(p \log p)$ , query: O(log p).

4.  $C_n^k = C_n^{k-1} \cdot \frac{n-k+1}{k}$ 

5.  $C_n^k = \frac{n!}{k!(n-k)!}$ , для каждого факториала считаем степень вхождения и остаток

precalc:  $O(p \log p)$ , query: O(log p).

$$C_n^{\frac{n}{2}} = \frac{2^n}{\sqrt{\frac{\pi n}{2}}}$$

#### 78 Делители

•  $\leq 20: d(12) = 6$ 

•  $\leq 50 : d(48) = 10$ 

•  $\leq 100 : d(60) = 12$ 

 $\bullet \le 1000 : d(840) = 32$ 

#### SPb HSE (Labutin, Podguzov, Bogomolov)

 $\bullet$  < 10<sup>4</sup> :  $d(9\ 240) = 64$ 

 $\bullet \le 10^5 : d(83\ 160) = 128$ 

 $\bullet \le 10^6 : d(720720) = 240$ 

 $\bullet \le 10^7 : d(8\ 648\ 640) = 338$ 

 $\bullet \le 10^8 : d(91\,891\,800) = 768$ 

 $\bullet \le 10^9 : d(931\ 170\ 240) = 1344$ 

 $\bullet \ \leq 10^{11}: d(97\ 772\ 875\ 200) = 4032$ 

 $\bullet \ \leq 10^{12} : d(963\ 761\ 198\ 400) = 6720$ 

•  $\leq 10^{15} : d(866\ 421\ 317\ 361\ 600) = 15360$ 

•  $\leq 10^{18} : d(897612484786617600) = 103680$ 

#### 79 Числа Белла

i	$B_i$	i	$B_i$
0	1	12	4,213,597
1	1	13	27,644,437
2	2	14	190,899,322
3	5	15	1,382,958,545
4	15	16	10,480,142,147
5	52	17	82,864,869,804
6	203	18	682,076,806,159
7	877	19	5,832,742,205,057
8	4,140	20	51,724,158,235,372
9	21,147	21	474,869,816,156,751
10	115,975	22	4,506,715,738,447,323
11	678,570	23	44,152,005,855,084,346

#### 80 Разбиения

Число неупорядоченных разбиений n на положительные слагаемые.

#### 81 Матричные игры

Пишем матрицу стратегий  $A_{i,j}$  это выигрыш первого и проигрыш второго, i стратегия 1-го. Седловая точка есть для несмешанной стратегии если  $\max_i \min A_{i,*} = \min_j \max A_{*,j}$ . Иначе:

$$f(x) = sum(x_i) \to max, \ Ans = 1/f(x)$$

$$Ax \le 1_n, \ x_i \ge 0$$

Для  $2 \times 2$ , p первый игрок, q — второй:

$$p^* = \left(\frac{a_{22} - a_{21}}{a_{22} - a_{12} + a_{11} - a_{21}}; \frac{a_{11} - a_{12}}{a_{22} - a_{12} + a_{11} - a_{21}}\right)$$
$$q^* = \left(\frac{a_{22} - a_{12}}{a_{22} - a_{12} + a_{11} - a_{21}}; \frac{a_{11} - a_{21}}{a_{22} - a_{12} + a_{11} - a_{21}}\right)$$
$$Ans = \frac{a_{22}a_{11} - a_{12}a_{21}}{a_{11} + a_{22} - a_{12} - a_{21}}$$

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#### 82 Mixed

- ullet Формула Пика: S = Inside + Edge/2 1
- Теорема Люка:  $0 \le n, m \in \mathbb{Z}, p$  простое.  $n = n_k p^k + \ldots + n_1 p + n_0$  и  $m = m_k p^k + \ldots + m_1 p + m_0$ . Тогда  $\binom{n}{m} \equiv \prod\limits_{i=0}^k \binom{n_i}{m_i} \pmod{p}$ .
- Лемма Бернсайда: |X/G| число орбит G.  $X^g=\{x\in X|gx=x\}$

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$