Uncertainty Quantification course

Lecture 8: Multidimensional SC and NISP

University of Amsterdam, fall 2023

1 November 2023

Outline

1 Homework assignment from last week

Stochastic collocation with dim>1

Oiscrete projection & NISP

Homework assignment from last week

Presentation by team 1

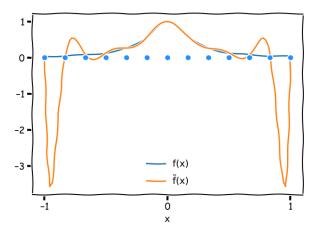
Stochastic Collocation: recap

PDE:
$$\mathcal{L}(u, Z) = \mathcal{F}(Z)$$
, $u = u(x, Z)$, dim $(Z)=1$

- (i) Select set of nodes in *Z*-space: $\theta_M := \{Z^{(j)}\}_{j=1}^M$
- (ii) Solve PDE on all nodes: $\mathcal{L}(u, Z^{(j)}) = \mathcal{F}(Z^{(j)})$ Solutions: $u^{(j)}(x)$ at node $Z^{(j)}$
- (iii) Interpolation: $\tilde{u}_M(x,Z) = \sum_{j=1}^M u^{(j)}(x) L_j(Z)$

with Lagrange basis functions $L_j(Z) := \prod_{i \neq j} \frac{Z - Z^{(i)}}{Z^{(j)} - Z^{(i)}}$

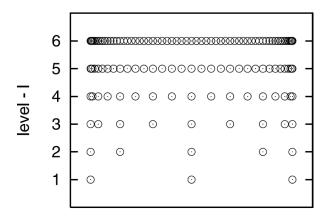
Uniform $\Theta_m := \{Z_i\}_{i=1}^m$: Runge phenomenon



standard SC in 1D

Better choice:

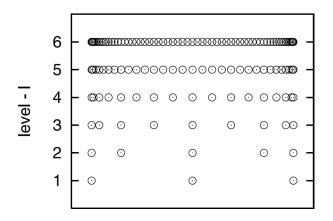
 non-uniform 1D quadrature points: building blocks SC method e.g. Clenshaw-Curtis (CC) nodes



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Can be a nested rule over 'levels'.

Example over [0, 1] (rescaled from standard CC domain of [-1, 1])

- Level I = 1: $x_i^{(1)} \in \{0.5\}$,
- Level I = 2: $x_i^{(2)} \in \{0.0, 0.5, 1.0\},\$
- Level I = 3: $x_i^{(3)} \in \{0.0, 0.146, 0.5, 0.854, 1.0\}$.

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Exponential increase in number of points per level;

$$m_{I} = \begin{cases} 2^{I-1} + 1 & I > 1 \\ 1 & I = 1 \end{cases}$$

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Weights are such that $\sum_i f(x_i)w_i \approx \int f(x)p(x)dx = \mathbb{E}[f]$.

Extension to higher dimensions

• Remember gPC with d = 2 uncertain inputs:

$$f(Z_1, Z_2) \approx \tilde{f}(Z_1, Z_2) = \sum_{i} \hat{f}_i \Phi_{i_1}(Z_1) \Phi_{i_2}(Z_2)$$

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- Here, **i** was a **multi index** (i_1, i_2) , determining the order of the basis functions.
- In SC, a multi index $I = (I_1, I_2)$ determines the order of 1D nodes used for each input.
- Example: level 1 nodes = $\{0\}$, level 2 nodes = $\{-1, 0, 1\}$

$$\textbf{I} = (1,2) \Rightarrow \{0\} \otimes \{-1,0,1\} = \{(0,-1),(0,0),(0,1)\}$$

Code is sampled at these 3 (Z_1, Z_2) via this multi index.

d=1: nodal set
$$\theta_M$$
, $\tilde{u}(Z)=\sum_{j=1}^M u(Z^{(j)})\,L_j(Z;\theta_M),\;\;Z^{(j)}\in\theta_M$

d>1: Tensor product, nodal set $\Theta_{\textit{M}} = \theta_{\textit{M}_1} \otimes \cdots \otimes \theta_{\textit{M}_d}$

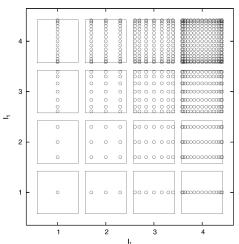
$$\tilde{u}(Z) = \sum_{j_1=1}^{M_1} \cdots \sum_{j_d=1}^{M_d} u(Z^{(j_1,\dots,j_d)}) L_{j_1,\dots,j_d}(Z;\Theta_M)$$

nodes
$$Z^{(j_1,...,j_d)} = (Z_1^{(j_1)},...,Z_d^{(j_d)}) \in \Theta_M$$

and basis functions
$$L_{j_1,...,j_d}(Z;\Theta_M) = \prod_{n=1}^d L_{j_n}(Z_n;\theta_{M_n})$$

standard SC in multiple dimension

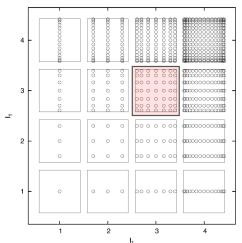
- More than 1 input: tensor products of 1D quad rules
- In 2D:



• Standard SC: user picks multi index $I = (l_1, l_2)$

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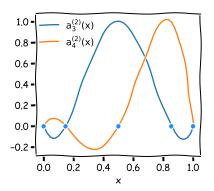
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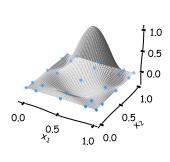
standard SC in multiple dimension

Surrogate in 2D:

$$\tilde{f}(\mathbf{Z}) = \sum_{j_1=1}^{m_{l_1}} \sum_{j_2=1}^{m_{l_2}} f\left(Z_1^{(j_1)}, Z_2^{(j_2)}\right) L_{j_1}(Z_1) \otimes L_{j_2}(Z_2)$$

2D basis functions:





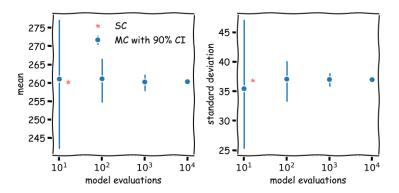
Example: wing-weight function

$$f(\mathbf{Z}) = 0.036 S_w^{0.758} W_{fw}^{0.0035} \left(\frac{A}{\cos^2(\Lambda)}\right)^{0.6} q^{0.006} \lambda^{0.04} \left(\frac{100 t_c}{\cos(\Lambda)}\right)^{-0.3} \left(\frac{N_z}{W_{dg}}\right)^{0.49} + S_w W_D$$

| S _w ∈ [150, 200] | wing area (ft ²) |
|--------------------------------|--|
| W _{fw} ∈ [220, 300] | weight of fuel in the wing (lb) |
| A ∈ [6, 10] | aspect ratio |
| Λ ∈ [-10, 10] | quarter-chord sweep (degrees) |
| q ∈ [16, 45] | dynamic pressure at cruise (lb/ft ²) |
| $\lambda \in [0.5, 1]$ | taper ratio |
| t _c ∈ [0.08, 0.18] | aerofoil thickness to chord ratio |
| N _z ∈ [2.5, 6] | ultimate load factor |
| W _{dg} ∈ [1700, 2500] | flight design gross weight (lb) |
| $W_p \in [0.025, 0.08]$ | paint weight (lb/ft ²) |

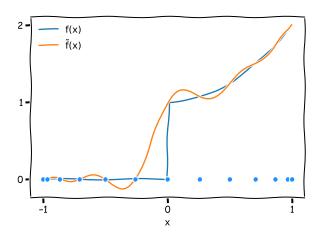
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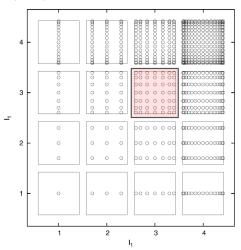
SC can outperform MC

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- Caveat 1: Accuracy poor if $f(\mathbf{x})$ is not smooth



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- Caveat 2: $\mathbf{Z} = \{Z_1, \dots, Z_d\}$, d must be small (e.g. \leq 5) 2D: given $\mathbf{I} = (I_1, I_2)$, there are $M = M_1 \times M_2$ code evaluations.



With tensor product $\Theta_M = \theta_{M_1} \otimes \cdots \otimes \theta_{M_d}$, total number of nodes is $M_1 \times M_2 \times \cdots \times M_d$.

Grows very rapidly with d: curse of dimension

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- 6 parameters, simulation time \approx 1.5 weeks.

Curse of dimension example:

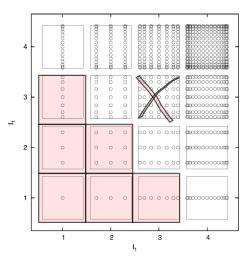
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- 6 parameters, simulation time \approx 1.5 weeks.

...

20 parameters, simulation time $\approx 3.17 \cdot 10^{12}$ years, 226 times the age of the universe.

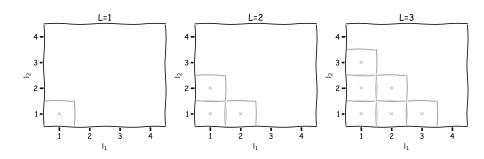
No supercomputer can beat this, need smarter algorithms.

• Postpone curse: sparse grids $\text{Select index set } \Lambda = \{\textbf{I}_1, \textbf{I}_2, \cdots \} \text{ instead of a single } \textbf{I}$



Smolyak sparse grids

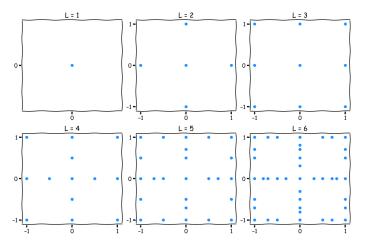
Visually, a standard sparse grid with $\{I \mid |I| \le L + d - 1\}$ selects a 'triangle' of multi indices when d = 2 (2 uncertain inputs):



Remember: $|\mathbf{I}| = I_1 + \cdots + I_d$.

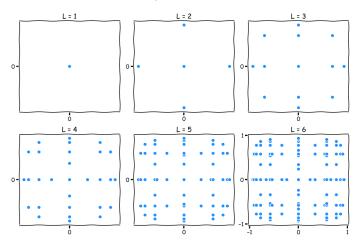
Isotropic sparse grids: sampling plan construction

- Given $\Lambda = \{(1,1),(1,2),\cdots\}$, create sampling plan by taking combination of tensor products for each $I \in \Lambda$
- For nested Clenshaw Curtis quad rule:



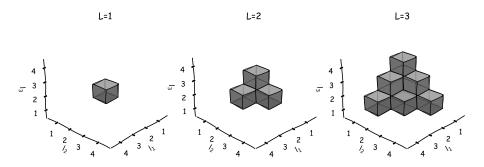
Isotropic sparse grids: sampling plan construction

- Given $\Lambda = \{(1,1),(1,2),\cdots\}$, create sampling plan by taking combination of tensor products for each $I \in \Lambda$
- For non-nested Gaussian quad rule:



Smolyak sparse grids

For d = 3, $\{I \mid |I| \le L + d - 1\}$ selects a 'simplex' of multi indices:



Smolyak sparse grids

• No explicit formula available to calculate M, total #nodes in sparse grid. Can be calculated via algorithm¹. For Clenshaw-Curtis:

| DIM: | 1 | 2 | 3 | 4 | 5 | |
|-------------|---------|-----------|-----------|------------|------------|--------|
| CFN_E LEVEL | | | | | | |
| 0 | 1 | 1 | 1 | 1 | 1 | |
| 1 | 3 | 5 | 7 | 9 | 11 | |
| 2 | 5 | 13 | 25 | 41 | 61 | |
| 3 | 9 | 29 | 69 | 137 | 241 | |
| 4 | 17 | 65 | 177 | 401 | 801 | |
| 5 | 33 | 145 | 441 | 1,105 | 2,433 | |
| 6 | 65 | 321 | 1,073 | 2,929 | 6,993 | |
| 7 | 129 | 705 | 2,561 | 7,537 | 19,313 | |
| 8 | 257 | 1,537 | 6,017 | 18,945 | 51,713 | |
| 9 | 513 | 3,329 | 13,953 | 46,721 | 135,073 | |
| 10 | 1,025 | 7,169 | 32,001 | 113,409 | 345,665 | |
| DIM: | 6 | 7 | 8 | 9 | 10 | ĺ |
| CFN_E LEVEL | | | | | | 1 |
| 0 | 1 | 1 | 1 | 1 | 1 | |
| 1 | 13 | 15 | 17 | 19 | 21 | |
| 2 | 85 | 113 | 145 | 181 | 221 | |
| 3 | 389 | 589 | 849 | 1,177 | 1,581 | |
| 4 | 1,457 | 2,465 | 3,937 | 6,001 | 8,801 | Full t |
| 5 | 4,865 | 9,017 | 15,713 | 26,017 | 41,265 | equiva |
| 6 | 15,121 | 30,241 | 56,737 | 100,897 | 171,425 | |
| 7 | 44,689 | 95,441 | 190,881 | 361,249 | 652,065 | be |
| 8 | 127,105 | 287,745 | 609,025 | 1,218,049 | 2,320,385 | 4 |
| 9 | 350,657 | 836,769 | 1,863,937 | 3,918,273 | 7,836,545 | |
| 10 | 943,553 | 2,362,881 | 5,515,265 | 12,133,761 | 25,370,753 | |
| | | | | | | |

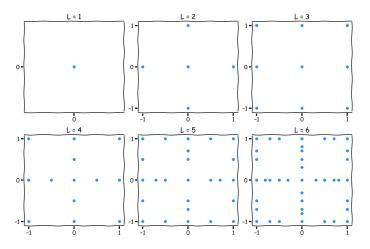
UQ course

tensor grid alent would $e 1025^{10}$

Burkardt, J. (2014). Counting Abscissas in Sparse Grids.

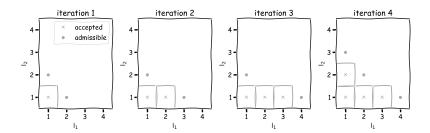
• Isotopic sparse grids: less points but all inputs are equal.

$$\Lambda = \{ \mathbf{I} \mid |\mathbf{I}|_1 - d + 1 \le L \}$$

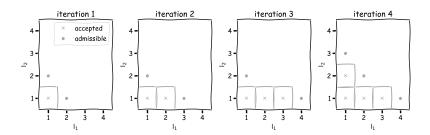


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- Dimension-adaptive sampling: find out 'on-the-fly' which ones
 - \rightarrow start with $\Lambda = \{1, 1, \cdots, 1\}$
 - \rightarrow adaptively refine Λ , add only 'important' $\mathbf{I} = (I_1, \cdots, I_d)$

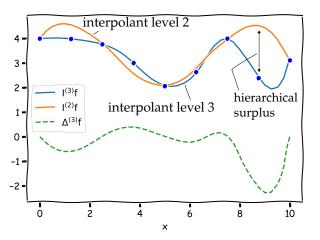


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Which I are important?

One possibility: **hierarchical surplus**, error between **new** code value and **old** interpolant prediction.



From all candidate I, select the one with the highest surplus.

Example:

$$f(\mathbf{x}) = 6x_1 + 4x_2 + 5.5x_3 + 3x_1x_2 + 2.2x_1x_3 + 1.4x_2x_3 + x_4 + 0.5x_5 + 0.2x_6 + 0.1x_7$$

Polynomial function, adaptive algorithm should be able to **exactly** find it in 10 iterations

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Polynomial function, adaptive algorithm should be able to **exactly** find it in 10 iterations

I = (1, 0, 0, 0, 0, 0, 0) refines x_1 line in 7-dimensional hypercube, enough to exactly describe $6x_1$ term.

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I = (1, 0, 0, 0, 0, 0, 0) refines x_1 line in 7-dimensional hypercube, enough to exactly describe $6x_1$ term.

I = (1, 1, 0, 0, 0, 0, 0) refines (x_1, x_2) plane in 7-dimensional hypercube, enough to exactly describe $3x_1x_2$ term.

aka Non-Intrusive Spectral Projection (NISP) or pseudospectral method

Recall (lecture 2) approximation of function f(x) by projection:

$$f(x) pprox f_N(x) = \sum_{n=0}^N \hat{f}_n \, \phi_n(x)$$
 with orthogonal basis functions $\{\phi_n(x)\}$, and $\hat{f}_n = \frac{\langle f, \phi_n \rangle}{\langle \phi_n, \phi_n \rangle}$ with appropriate (weighted) inner product $\langle ., . \rangle$

Discrete projection: approximate inner products with quadrature

Quadrature: numerical integration, $\int f(x)w(x) dx \approx \sum_{j=1}^{q} f(x^{(j)})\alpha^{(j)}$ with nodes $\{x^{(j)}\}$ and weights $\{\alpha^{(j)}\}$ (density w(x) absorbed into weights)

Quadrature in dim>1 is called cubature

Inner products for projection with *q*-point quadrature/cubature rule:

$$\langle f, \phi_n \rangle = \int f(x) \, \phi_n(x) \, w(x) \, dx \approx \sum_{j=1}^q f(x^{(j)}) \, \phi_n(x^{(j)}) \, \alpha^{(j)}$$

Stochastic system (e.g. PDE) with exact solution u(t, x, Z), $Z \in \mathbb{R}^d$

Spectral (gPC) approximation:
$$u_N(t, x, Z) = \sum_{n=0}^{N} \hat{u}_n(t, x) \Phi_n(Z)$$

$$\hat{u}_n(t,x) = \gamma_n^{-1} \langle u, \Phi_n \rangle = \gamma_n^{-1} \mathbb{E}[u(t,x,Z), \Phi_n(Z)]$$

$$\approx \tilde{u}_n(t,x) = \gamma_n^{-1} \sum_{j=1}^q u(t,x,z^{(j)}) \, \Phi_n(z^{(j)}) \, \alpha^{(j)}$$

Needed: q solutions of original (non-stochastic) system, $u(t, x, z^{(j)})$ with j = 1, ..., q

→ NISP is a non-intrusive method

dim(Z)>1: Cubature

From quadrature if dim(Z)=1 to cubature if dim(Z)>1:

Tensor product construction, based on 1-d quadrature and nodal sets

Again, sparse grids to reduce no. of nodes

| 888 | 8 | 8 | 8 | 8 | 8 | 8 | 8 | 8 | 8 | 888 | | 0 | | | | 8 | | | | 0 | |
|-----|---|---|---|---|---|---|---|---|---|------|------|---|---|---|---|---|---|---|---|---|-----|
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| 000 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 00 | 00 o | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 000 |
| 000 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 000 | | | | | | 0 | | | | | |
| 000 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 000 | | 0 | | | | 0 | | | | 0 | |
| 000 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 000 | | | | | | 0 | | | | | |
| 000 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 000 | 0 | 0 | | 0 | | 0 | | 0 | | 0 | 0 |
| 888 | 8 | 8 | 8 | 8 | 8 | 0 | 8 | 8 | 0 | 888 | | 0 | | | | 8 | | | | 0 | |

dim(Z)>1: Cubature

Example: your homework assignment from lecture 6!

Consider the following function with two random inputs:

$$f(Z_1, Z_2) = -(Z_1 - 2)(Z_2 - 1)^3 + \exp\left[-\frac{1}{2}(Z_1 - 2)^2 - \frac{1}{10}(Z_2 - 1)^2\right]$$

One input has a uniform distribution on [1,3], the other has a normal distribution with mean 1 and variance 1. Thus, $Z_1 \sim \mathcal{U}[1,3]$ and $Z_2 \sim \mathcal{N}(1,1)$.

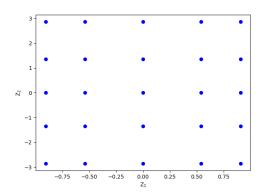
- If $f(Z_1, Z_2)$ is considered as your 'code'.
- This involved computing $\mathbb{E}\left[f(Z_1,Z_1)\Phi_{i_1}^{(1)}(Z_1)\Phi_{i_2}^{(2)}(Z_2)\right]/\gamma_i$.
- This can be done with in-built (SciPy) integration routines.

dim(Z)>1: Cubature

But also with cubature:

 $p=1 \ q=1$

$$\mathbb{E}\left[f(Z_1,Z_1)\Phi_{i_1}^{(1)}(Z_1)\Phi_{i_2}^{(2)}(Z_2)\right] \approx \sum_{p=1}^{m_p} \sum_{q=1}^{m_q} f(Z_{1,p}Z_{2,q})\Phi_{i_1}^{(1)}(Z_1)\Phi_{i_2}^{(2)}(Z_2)\alpha_{1,p}\alpha_{2,q}$$



Multidim SC & NISP **UQ** course 1 Nov 2023 31/33 Solution u(t, x, Z) of PDE system, approximated by $\tilde{u}(t, x, Z) = \sum_{n=0}^{N} \hat{u}_n(t, x) \Psi_n(Z)$

- Galerkin method: projection-based, leading to coupled equations for coefficients $\{\hat{u}_n(t,x)\}_{n=0}^N$. Intrusive.
- Collocation: interpolation in Z-space of deterministic solutions $\{u(t,x,z^{(j)})\}_{j=1}^q$. Non-intrusive.
- NISP: projection coefficients $\{\hat{u}_n(t,x)\}_{n=0}^N$ approximated with cubature using deterministic solutions $\{u(t,x,z^{(j)})\}_{j=1}^q$. Non-intrusive.

- Inverse UQ: Bayesian calibration and MCMC.
 Xiu section 8.2, Smith section 8.1 8.3 (on Canvas)
- Presentation about homework by TEAM 2?
- Homework: Multi-dimensional SC (more details on Canvas)