

**Exercise 1: Big-O Proofs**

For the following statements, either prove them or give a counterexample.

- (a)  $\max[|f(n)|, |g(n)|] \in \Theta(|f(n)| + |g(n)|)$ .
- (b) For eventually nonzero  $f(n)$  and  $g(n)$ ,  $f(n) \in O(g(n))$  implies  $g(n) \in O(f(n))$ .
- (c) For eventually nonzero  $f(n)$  and  $g(n)$ ,  $f(n) \in O(g(n))$  if and only if  $g(n) \in \Omega(f(n))$ .
- (d)  $f(n) \in O(f(2n))$ .

$f(n)$  being “eventually nonzero” here means that there is some  $n_0$  so that for all  $n > n_0$ ,  $f(n) \neq 0$ .

**Exercise 2: Big-O Order**

- (a) Partition the following functions into  $\Theta$ -equivalence classes, where  $f \sim g$  means  $f \in \Theta(g)$ . List the classes in increasing order of asymptotic growth rate, denoted by  $\prec$ , where  $C_i \prec C_j$  means  $f(n) \in o(g(n))$  for all  $f \in C_i$  and  $g \in C_j$ .

In all of these,  $\log$  is the natural logarithm.

- $f(n) = n^2$
- $g(n) = n^2 + \log n$
- $h(n) = n^n$
- $i(n) = \log n$
- $j(n) = (\log n)^2$
- $k(n) = \log(n^2)$
- $l(n) = 2^{2^n}$
- $m(n) = 2^n$
- $n(n) = n!$
- $p(n) = 2^{\log n}$

You may need to use the following approximation of  $n!$ , called Stirling’s approximation:

$$n! = \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \left(1 + O\left(\frac{1}{n}\right)\right)$$

- (b) In which of the cases in (a) does the base of the logarithm matter? I.e., in which case is a given function not in the  $\Theta$ -equivalence class of the same function using a logarithm of a different base?

**Exercise 3: Empirical Epsilon**

- (a) Write a function in R that computes, for a given positive floating point number  $x$ , the smallest number  $u(x)$  such that  $x + u(x) \neq x$  in machine arithmetic. Use a bisection method: start with candidates  $u_{\text{lower}}$ , for which you know that  $x + u_{\text{lower}} = x$ , and  $u_{\text{upper}}$ , for which you know that  $x + u_{\text{upper}} \neq x$ . Then, repeatedly check the midpoint of the interval  $[u_{\text{lower}}, u_{\text{upper}}]$  on whether adding it to  $x$  yields  $x$  or not, and replace either the lower or upper bound by the midpoint. Make sure that your function works with very small numbers: The midpoint between  $u_{\text{lower}}$  and  $u_{\text{upper}}$  might not be representable as a value distinct from either  $u_{\text{lower}}$  or  $u_{\text{upper}}$ , your function should not get stuck in an infinite loop in this case.

- (b) Use your function to calculate the value of  $u(x)$  for a reasonably dense grid of values of  $x$  between  $2^{-1024}$  and  $2^{-1010}$  on a logarithmic scale. (“reasonably dense” here means you should not only consider integer powers of 2. Notice that these are all very small numbers.) Plot the result on a log-log plot. What do you observe?

**Bonus:** (think about this, but don’t worry if you can’t solve it) Why is  $u(x)$  not monotonically non-decreasing? I.e., why are there apparently random fluctuations, with values of  $x_1 < x_2$  for which  $u(x_1) > u(x_2)$ ?