# TECHNICAL UNIVERSITY OF MOLDOVA



# FACULTY OF COMPUTERS, INFORMATICS AND MICROELECTRONICS SOFTWARE ENGINEERING AND AUTOMATION DEPARTMENT

# Homework 1

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# $\mathbf{E}\mathbf{x}\mathbf{1}$

Present the language L(G) generated by the given grammar  $G = (V_n, V_n, S, P)$ 

• 
$$V_n = \{S\}; V_t = \{a\}; P = \{S \to aS, S \to \varepsilon \}$$

$$L = \{a^n | n \ge 0\}$$

$$(1)$$

• 
$$V_n = \{S\}; V_t = \{a\}; P = \{S \to aS, S \to a \}$$

$$L = \{a^n | n \ge 1\}$$
(2)

• 
$$V_n = \{A\}; V_t = \{0,1\}; P = \{A \to 0A1, A \to 01\}$$

$$L = \{0^n, 1^n | n \ge 1\}$$
(3)

## $\mathbf{Ex2}$

Identify the grammar for the following languages:

 $\bullet \ L = \{a^n, b^n | n \ge 0\}$ 

$$V_n = \{S\}; V_t = \{a, b\}; P = \{S \to \varepsilon, S \to aSb\}$$

$$\tag{4}$$

•  $L = \{a^n, b^n | n > 0\}$ 

$$V_n = \{S\}; V_t = \{a, b\}; P = \{S \to aSb, S \to ab\}$$
 (5)

•  $L = \{a^n, b^{n+1} | n \ge 1\}$ 

$$V_n = \{S\}; V_t = \{a, b\}; P = \{S \to aSb, S \to abb\}$$
 (6)

 $\bullet \ L = \{a^n, b^n, c^m, d^m | n > 1, m > 1\}$ 

$$V_n = \{S, A\}; V_t = \{a, b, c, d\}; P = \{S \to aSb|aAb; A \to dAc|dc\}$$
 (7)

•  $L = \{a^n, b^n, c^m, d^m | n \ge 0, m \ge 0\}$ 

$$V_n = \{S,A\}; V_t = \{a,b,c,d\}; P = \{S \rightarrow aSb|aAb|\varepsilon; A \rightarrow dAc|dc|\varepsilon\}$$
 (8)

 $\bullet \ L = \{a^n, b^m, c^m, d^n | n > 1, m > 1\}$ 

$$V_n = \{S, A\}; V_t = \{a, b, c, d\}; P = \{S \to aSd | aAd; A \to bAc | bc\}$$
 (9)

•  $L = \{a^m, b^n, c^{m+n} | n \ge 1, m \ge 1\}$ 

$$V_n = \{S, A, B\}; V_t = \{a, b, c\}; P = \{S \to AB; A \to aAc | ac; B \to bBc | bc\}$$
 (10)

•  $L = \{a^m, b^{m+n}, c^n | n \ge 1, m \ge 1\}$ 

$$V_n = \{S, A, B\}; V_t = \{a, b, c\}; P = \{S \to AB; A \to aAb|ab; B \to bBc|bc\}$$
 (11)

### Ex3

For the given grammar identify the generated word:

```
G=(V_n, V_n, S, P) V_n = {programm>,<set of afirmation>, <afirmation>,
                     <assigment>,<test>,<variable>,<number>, jalpha>};
V_t = \{ \text{begin, end, succ, pred, while, do, } :=, \neq, ;, (, ), 0,1....9, A, B,....Z \}
                           P = \{ < program > \rightarrow begin end \}
                  \langle program \rangle \rightarrow begin \langle set of a firmation \rangle end
                        \langleset of afirmation\rangle \rightarrow \langleafirmation\rangle
          \langleset of afirmation\rangle \rightarrow \langleset of afirmation\rangle \langleafirmation\rangle
                            \langle afirmation \rangle \rightarrow \langle assignment \rangle
        <afirmation> → while<test>do<afirmation>—                                                                                                                                                                                                                                                                                                                                              
                             \langle afirmation \rangle \rightarrow \langle program \rangle
                        \langle \text{test} \rangle \rightarrow \langle \text{variable} \rangle \neq \langle \text{variable} \rangle
                           \langle assigment \rangle \rightarrow \langle variable \rangle := 0
                 \langle assigment \rangle \rightarrow \langle variable \rangle := succ(\langle variable \rangle)
                 \langle assigment \rangle \rightarrow \langle variable \rangle := pred(\langle variable \rangle)
                                \langle \text{variable} \rangle \rightarrow \langle \text{alpha} \rangle
                          \langle variable \rangle \rightarrow \langle variable \rangle \langle alpha \rangle
                        \langle variable \rangle \rightarrow \langle variable \rangle \langle number \rangle
                  <number> \rightarrow 0-1-2-3-4-5-6-7-8-9
```

#### Solution

```
<program> \rightarrow begin <set of afirmation > end \rightarrow begin <afirmation> end \rightarrow begin <assignment> end \rightarrow begin <variable> := 0 end \rightarrow begin <alpha> := 0 end \rightarrow begin A := 0 end \rightarrow end
```

# $\mathbf{Ex4}$

Define grammar tha generates the variable identifiers from Java.

$$V_n = \{<\text{start}>, <\text{next}>\}$$

$$V_t = \{0,1,2,...,9,A,B,C,...,Z,a,b,c,...,z, \$, \_\}$$

$$P = \{<\text{start}> \to A|B|...|Z|a|b|...|z|o|1|...|g|\$|...<\text{next}>,$$

$$<\text{next}> \to E\}$$

## Ex5

Define the grammar that generates all real literals in Java.

$$V_{n} = \{<\text{float}>, <\text{decimal}>\}$$

$$V_{t} = \{0,1,2,...,9,A,B,C,...,Z,a,b,c,...,z,..\}$$

$$P = \{<\text{float}> \rightarrow A|B|...|Z|a|b|...|z<\text{decimal}>,$$

$$<\text{decimal}> \rightarrow .|A|B|...|Z|a|b|...|z|0|1|...|9|$|...<\text{decimal}>,$$

$$<\text{decimal}> \rightarrow \varepsilon \}$$

## Ex6

Define the grammar that generates the strings that correspond to valid currency amounts. A valid string is either a dollar sign followed by a number which has no leading 0's, and may have a decimal point in which case it must be followed by exactly two decimal digits, OR a one or two-digit amount followed by the cent sign c. The single exception to this rule is that strings which begin with "\$0." and are followed by exactly two digits are also acceptable. Thus, \$432.63, 1,0.29, 47c, 2c are all accepted, but \$021, \$4.3, \$8.63c, \$0.0 are not accepted.

#### Solution

$$V_n = \{ S,A \}$$

$$V_t = \{ 0,1,2,...,9,...,\$,c \}$$

$$P = \{ S \rightarrow \$1|2|...|9A, S \rightarrow Dc, S \rightarrow DDc, S \rightarrow DDc, S \rightarrow \$0.DD, A \rightarrow 0|1|...|9A, A \rightarrow .DD, D \rightarrow 0|1|...|9 \}$$