

Supporting Information

Conditional cooperation with longer memory

Nikoleta E. Glynatsi, Martin A. Nowak, Christian Hilbe

Contents

1	Model and basic results	2
1.1	The repeated prisoner's dilemma	2
1.2	Finite-memory strategies	3
1.3	Computing the payoffs of finite-memory strategies	4
1.4	An Extension of Akin's Lemma	6
2	Characterizing the partner strategies among the reactive-n strategies	7
2.1	Partner strategies	7
2.2	Tit For Tat and Generous Tit For Tat with arbitrary memory lengths	8
2.3	An algorithm to check whether a reactive- n strategy is a Nash equilibrium	9
2.4	Reactive partner strategies in the donation game	12
2.5	Reactive partner strategies in the general prisoner's dilemma	14
2.6	Reactive defecting Nash strategies in the donation game	17
3	Evolutionary Simulations	18
4	Errors	19
5	Appendix: Proofs	20
5.1	Proof of Lemma 1: Akin's lemma	20
5.2	Proof of Lemma 2: Sufficiency of testing self-reactive strategies	21
5.3	Proof of Theorem 1: Sufficiency of pure self-reactive strategies	22
5.4	Proof of Theorem 2: Reactive-2 partner strategies in the donation game	24
5.5	Proof of Theorem 3: Reactive-3 partner strategies in the donation game	25
5.6	Proof of Theorem 4: Reactive- n counting strategies in the donation game	26
5.7	Proof of Theorem 5: Reactive-2 partner strategies in the prisoner's dilemma	28
5.8	Proof of Theorem 6: Reactive-3 partner strategies in the prisoner's dilemma	29
6	Supplementary References	30

This document provides further details on our methodology and our analytical results. Section 1 summarizes the model. In particular, we introduce all relevant strategy spaces, and we show how to compute long-term payoffs for strategies with more than one-round memory. Section 2 contains our key results. Here, we define partner strategies, we present an algorithm that allows us to verify whether a given reactive- n strategy is a partner, and we apply this algorithm to fully characterize the reactive- n partner strategies for $n=2$ and $n=3$. All proofs are presented in the Appendix in Section 5.

1 Model and basic results

1.1 The repeated prisoner's dilemma

We consider the infinitely repeated prisoner's dilemma between two players, player 1 and player 2. Each round, each player can either cooperate (C) or defect (D). The resulting payoffs are given by the matrix

$$\begin{array}{cc} & \begin{array}{cc} C & D \end{array} \\ \begin{array}{c} C \\ D \end{array} & \begin{pmatrix} R & S \\ T & P \end{pmatrix} \end{array}. \quad (1)$$

Here, R is the reward payoff of mutual cooperation, T is the temptation to defect, S is the sucker's payoff, and P is the punishment payoff for mutual defection. For the game to be a prisoner's dilemma, we require

$$T > R > P > S \quad \text{and} \quad 2R > T + S. \quad (2)$$

That is, mutual cooperation is the best outcome to maximize the players' total payoffs, but each player's dominant action is to defect. For some of our results, we focus on a special case of the prisoner's dilemma, the donation game. This game only depends on two free parameters, the benefit b and the cost c of cooperation. The payoff matrix of the donation game takes the form

$$\begin{array}{cc} & \begin{array}{cc} C & D \end{array} \\ \begin{array}{c} C \\ D \end{array} & \begin{pmatrix} b - c & -c \\ b & 0 \end{pmatrix} \end{array}. \quad (3)$$

For this game to satisfy the conditions (2) of a prisoner's dilemma, we assume $b > c > 0$ throughout.

Players interact in the repeated prisoner's dilemma for infinitely many rounds, and future payoffs are not discounted. A strategy σ^i for player i is a rule that tells the player what to do in any given round, depending on the outcome of all previous rounds. Given the player's strategies σ^1 and σ^2 , one can compute each player i 's expected payoff $\pi_{\sigma^1, \sigma^2}^i(t)$ in round t . For the entire repeated game, we define the players' payoffs as the expected payoff per round,

$$\pi^i(\sigma^1, \sigma^2) = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=1}^{\tau} \pi_{\sigma^1, \sigma^2}^i(t). \quad (4)$$

For general strategies σ^1 and σ^2 , the above limit may not always exist. Problems may arise, for example, if one of the players cooperates in the first round, defects in the two subsequent rounds, cooperates in the four rounds thereafter, etc., which prevents the time averages from converging. However, in the following, we focus on strategies with finite memory. When both players adopt such a strategy, the existence of the limit (4) is guaranteed, as we discuss further below.

1.2 Finite-memory strategies

In this study, we focus on strategies that ignore all events that happened more than n rounds ago. To define these strategies, we need some notation. An n -history for player i is a string $\mathbf{h}^i = (a_{-n}^i, \dots, a_{-1}^i) \in \{C, D\}^n$. We interpret the string's entry a_{-k}^i as player i 's action k rounds ago. We denote the space of all n -histories for player i as H^i . This space contains $|H^i| = 2^n$ elements. A pair $\mathbf{h} = (\mathbf{h}^1, \mathbf{h}^2)$ is called an n -history of the game. We use $H = H^1 \times H^2$ to denote the space of all such histories, which contains $|H| = 2^{2n}$ elements.

Memory- n strategies. Based on this notation, a *memory- n strategy* for player i as a tuple $\mathbf{m} = (m_{\mathbf{h}})_{\mathbf{h} \in H}$. Each input $\mathbf{h} = (h^i, h^{-i})$ refers to a possible n -history, where now \mathbf{h}^i and \mathbf{h}^{-i} refer to the n -histories of the focal player and the co-player, respectively. The corresponding output $m_{\mathbf{h}} \in [0, 1]$ is the focal player's cooperation probability in the next round, contingent on the outcome of the previous n rounds. We refer to the set of all memory- n strategies as

$$\mathcal{M}_n := \left\{ \mathbf{m} = (m_{\mathbf{h}})_{\mathbf{h} \in H} \mid 0 \leq m_{\mathbf{h}} \leq 1 \text{ for all } \mathbf{h} \in H \right\} = [0, 1]^{2^{2n}}. \quad (5)$$

This definition leaves the strategy's actions during the first n rounds unspecified, for which no complete n -history is yet available. However, because we consider infinitely repeated games without discounting, these first n rounds are usually irrelevant for the long-run dynamics, as we show further below. In the following, we therefore only specify a strategy's move during the first n rounds when necessary.

Among all memory- n spaces \mathcal{M}_n , the one with $n = 1$ is the most frequently studied. Memory-1 strategies take the form $\mathbf{m} = (m_{CC}, m_{CD}, m_{DC}, m_{DD})$. The first index refers to the focal player's last action (1-history) and the second index refers to the co-player's last action. As an example of a well-known memory-1 strategy, we mention Win-Stay Lose-Shift [1], $\mathbf{m} = (1, 0, 0, 1)$. However, there are many others [2].

Reactive- n strategies. For our following analysis, two particular subsets of memory- n strategies will play an important role. The first subset is the set of *reactive- n strategies*,

$$\mathcal{R}_n := \left\{ \mathbf{m} \in \mathcal{M}_n \mid m_{(\mathbf{h}^i, \mathbf{h}^{-i})} = m_{(\tilde{\mathbf{h}}^i, \mathbf{h}^{-i})} \text{ for all } \mathbf{h}^i, \tilde{\mathbf{h}}^i \in H^i \text{ and } \mathbf{h}^{-i} \in H^{-i} \right\}. \quad (6)$$

That is, reactive- n strategies are independent of the focal player's own n -history. The space of reactive- n strategies can be naturally identified with the space of all 2^n -dimensional vectors

$$\mathbf{p} = (p_{\mathbf{h}^{-i}})_{\mathbf{h}^{-i} \in H^{-i}} \text{ with } 0 \leq p_{\mathbf{h}^{-i}} \leq 1 \text{ for all } \mathbf{h}^{-i} \in H^{-i}. \quad (7)$$

In this reduced representation, each entry $p_{\mathbf{h}^{-i}}$ corresponds to the player's cooperation probability in the next round based on the co-player's actions in the previous n rounds. Again, the most studied case of reactive- n strategies is when $n=1$. Here, the reduced representation according to Eq. (7) takes the form $\mathbf{p} = (p_C, p_D)$. Probably the best-known example of a reactive-1 strategy is Tit-for-Tat, T_{FT} [3]. T_{FT} cooperates if and only if the co-player cooperated in the previous round. Hence, its memory-1 representation is $\mathbf{m} = (1, 0, 1, 0)$, whereas its reduced representation is $\mathbf{p} = (1, 0)$. Another example is the strategy Generous Tit-for-Tat, GT_{FT} [4, 5]. GT_{FT} occasionally cooperates even if the co-player defected. For that strategy, the memory-1 representation is $\mathbf{m} = (1, p_D^*, 1, p_D^*)$, and the reduced representation is $\mathbf{p} = (1, p_D^*)$, where

$$p_D^* := \min \left\{ 1 - (T - R)/(R - S), (R - P)/(T - P) \right\}. \quad (8)$$

In the special case that payoffs are given by the donation game, this condition simplifies to $p_D^* = 1 - c/b$.

Self-reactive- n strategies. The other important subspace of memory- n strategies is the set of self-reactive- n strategies,

$$\mathcal{S}_n := \left\{ \mathbf{m} \in \mathcal{M}_n \mid m_{(\mathbf{h}^i, \mathbf{h}^{-i})} = m_{(\mathbf{h}^i, \tilde{\mathbf{h}}^{-i})} \text{ for all } \mathbf{h}^i \in H^i \text{ and } \mathbf{h}^{-i}, \tilde{\mathbf{h}}^{-i} \in H^{-i} \right\}. \quad (9)$$

These strategies only depend on the focal player's own decisions during the last n rounds, independent of the co-player's decisions. Again, we can identify any self-reactive- n strategies with a 2^n -dimensional vector,

$$\tilde{\mathbf{p}} = (\tilde{p}_{\mathbf{h}^i})_{\mathbf{h}^i \in H^i} \text{ with } 0 \leq \tilde{p}_{\mathbf{h}^i} \leq 1 \text{ for all } \mathbf{h}^i \in H^i. \quad (10)$$

Each entry $\tilde{p}_{\mathbf{h}^i}$ corresponds to the player's cooperation probability in the next round, contingent on the player's own actions in the previous n rounds. A special subset of self-reactive strategies is given by the round- k -repeat strategies, for some $1 \leq k \leq n$. In any given round, players with a *round- k -repeat strategy* $\tilde{\mathbf{p}}^{k\text{-Rep}}$ choose the same action as they did k rounds ago. Formally, the entries of $\tilde{\mathbf{p}}^{k\text{-Rep}}$ are defined by

$$p_{\mathbf{h}^i}^{k\text{-Rep}} = \begin{cases} 1 & \text{if } a_{-k}^i = C \\ 0 & \text{if } a_{-k}^i = D. \end{cases} \quad (11)$$

From this point forward, we will use the notations \mathbf{m} , \mathbf{p} , and $\tilde{\mathbf{p}}$ to denote memory- n , reactive- n , and self-reactive- n strategies, respectively. We say these strategies are *pure* or *deterministic* if all conditional cooperation probabilities are either zero or one. If all cooperation probabilities are strictly between zero and one, we say the strategy is *strictly stochastic*. When it is convenient to represent the self-reactive repeat strategies as elements of the memory- n strategy space, we write $\mathbf{m}^{k\text{-Rep}} \in [0, 1]^{2^{2n}}$ instead of $\tilde{\mathbf{p}}^{k\text{-Rep}} \in [0, 1]^{2^n}$.

1.3 Computing the payoffs of finite-memory strategies

A Markov chain representation. The interaction between two players with memory- n strategies \mathbf{m}^1 and \mathbf{m}^2 can be represented as a Markov chain. The states of the Markov chain are the possible n -histories $\mathbf{h} \in H$. To compute the transition probabilities from one state to another within a single round, suppose players

currently have the n -history $\mathbf{h} = (\mathbf{h}^1, \mathbf{h}^2)$ in memory. Then the transition probability that the state after one round is $\tilde{\mathbf{h}} = (\tilde{\mathbf{h}}^1, \tilde{\mathbf{h}}^2)$ is a product of two factors,

$$M_{\mathbf{h}, \tilde{\mathbf{h}}} = x^1 \cdot x^2, \quad (12)$$

The two factors represent the (independent) decisions of the two players,

$$x^i = \begin{cases} m_{(\mathbf{h}^i, \mathbf{h}^{-i})}^i & \text{if } \tilde{a}_{-1}^i = C, \text{ and } \tilde{a}_{-t}^i = a_{-t+1}^i \text{ for } t \in \{2, \dots, n\} \\ 1 - m_{(\mathbf{h}^i, \mathbf{h}^{-i})}^i & \text{if } \tilde{a}_{-1}^i = D, \text{ and } \tilde{a}_{-t}^i = a_{-t+1}^i \text{ for } t \in \{2, \dots, n\} \\ 0 & \text{if } \tilde{a}_{-t}^i \neq a_{-t+1}^i \text{ for some } t \in \{2, \dots, n\}. \end{cases} \quad (13)$$

The resulting $2^{2n} \times 2^{2n}$ transition matrix $M = (M_{\mathbf{h}, \tilde{\mathbf{h}}})$ fully describes the dynamics among the two players after the first n rounds. More specifically, suppose $\mathbf{v}(t) = (v_{\mathbf{h}}(t))_{\mathbf{h} \in H}$ is the probability distribution of observing state \mathbf{h} after players made their decisions for round $t \geq n$. Then the respective probability distribution after round $t+1$ is given by $\mathbf{v}(t+1) = \mathbf{v}(t) \cdot M$. The long-run dynamics is particularly simple to describe when the matrix M is primitive (which happens, for example, when the two strategies $m_{\mathbf{h}}^i$ are strictly stochastic). In that case, it follows by the theorem of Perron and Frobenius that $\mathbf{v}(t)$ converges to some \mathbf{v} as $t \rightarrow \infty$. As a result, also the respective time average exists and converges to \mathbf{v} ,

$$\mathbf{v} = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} \mathbf{v}(t). \quad (14)$$

This limiting distribution \mathbf{v} can be computed as the unique solution of the system $\mathbf{v} = \mathbf{v}M$, with the additional constraint that the entries of \mathbf{v} need to sum up to one.

But even when M is not ergodic, $\mathbf{v}(t)$ still converges to an invariant distribution \mathbf{v} that satisfies $\mathbf{v} = \mathbf{v}M$. However, in that case, the system $\mathbf{v} = \mathbf{v}M$ no longer has a unique solution. Instead, the limiting distribution \mathbf{v} depends on the very first n -history after the first n rounds, $\mathbf{v}(n)$, which in turn depends on the players' moves during the first n rounds.

A formula for the payoffs among memory- n players. Based on the above considerations, we can derive an explicit formula for the payoffs according to Eq. (4) when players use memory- n strategies \mathbf{m}^1 and \mathbf{m}^2 . To this end, we introduce a 2^{2n} -dimensional vector $\mathbf{g}^i(k) = (g_{\mathbf{h}}^i(k))_{\mathbf{h} \in H}$, that takes an n -history \mathbf{h} as an input and returns player i 's payoff k rounds ago, for $k \leq n$. That is,

$$g_{\mathbf{h}}^i(k) = \begin{cases} R & \text{if } a_{-k}^i = C \text{ and } a_{-k}^{-i} = C \\ S & \text{if } a_{-k}^i = C \text{ and } a_{-k}^{-i} = D \\ T & \text{if } a_{-k}^i = D \text{ and } a_{-k}^{-i} = C \\ P & \text{if } a_{-k}^i = D \text{ and } a_{-k}^{-i} = D. \end{cases} \quad (15)$$

Now for a given $t \geq n$, given that $\mathbf{v}(t)$ captures the state of the system after round t , we can write player i 's

expected payoff in that round as

$$\pi_{\mathbf{m}^1, \mathbf{m}^2}^i(t) = \langle \mathbf{v}(t), \mathbf{g}^i(1) \rangle = \sum_{\mathbf{h} \in H} v_{\mathbf{h}}(t) \cdot g_{\mathbf{h}}^i(1). \quad (16)$$

As a result, we obtain for the player's average payoff across all rounds

$$\begin{aligned} \pi^i(\mathbf{m}^1, \mathbf{m}^2) &\stackrel{(4)}{=} \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=1}^{\tau} \pi_{\mathbf{m}^1, \mathbf{m}^2}^i(t) = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} \pi_{\mathbf{m}^1, \mathbf{m}^2}^i(t) \\ &\stackrel{(16)}{=} \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} \langle \mathbf{v}(t), \mathbf{g}^i(1) \rangle = \left\langle \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} \mathbf{v}(t), \mathbf{g}^i(1) \right\rangle \\ &\stackrel{(14)}{=} \langle \mathbf{v}, \mathbf{g}^i(1) \rangle. \end{aligned} \quad (17)$$

That is, given we know the invariant distribution \mathbf{v} that captures the game's long-run dynamics, it is straightforward to compute payoffs by taking the scalar product with the vector $\mathbf{g}^i(1)$. With a similar approach as in Eq. (17), one can also show

$$\langle \mathbf{v}, \mathbf{g}^i(1) \rangle = \langle \mathbf{v}, \mathbf{g}^i(2) \rangle = \dots = \langle \mathbf{v}, \mathbf{g}^i(n) \rangle. \quad (18)$$

That is, to compute player i 's expected payoff, it does not matter whether one refers to the last round of an n -history or to an earlier round of an n -history. All rounds k with $1 \leq k \leq n$ are equivalent.

1.4 An Extension of Akin's Lemma

The above Markov chain approach allows us to analyze games when both players adopt memory- n strategies. But even if only one player adopts a memory- n strategy (and the other player's strategy is arbitrary), one can still derive certain constraints on the game's long-run dynamics. One such constraint was first described by Akin [6]: if player 1 adopts a memory-1 strategy \mathbf{m} against an arbitrary opponent, and if the time average \mathbf{v} defined by the right hand side of Eq. (14) exists, then

$$\langle \mathbf{v}, \mathbf{m} - \mathbf{m}^{1-\text{Rep}} \rangle = 0. \quad (19)$$

That is, the limiting distribution \mathbf{v} needs to be orthogonal to the vector $\mathbf{m} - \mathbf{m}^{1-\text{Rep}}$. This result has been termed *Akin's Lemma* [7]. With similar methods as in Ref. [6], one can generalize this result to the context of memory- n strategies.

Lemma 1 (A generalized version of Akin's Lemma)

Let player 1 use a memory- n strategy, and let player 2 use an arbitrary strategy. For the resulting game and all $t \geq n$, let $\mathbf{v}(t) = (v_{\mathbf{h}}(t))_{\mathbf{h} \in H}$ denote the probability distribution of observing each possible n -history $\mathbf{h} \in H$ after players made their decisions for round t . Moreover, suppose the respective time average \mathbf{v}

according to Eq. (14) exists. Then for each k with $1 \leq k \leq n$, we obtain

$$\langle \mathbf{v}, \mathbf{m} - \mathbf{m}^{k-Rep} \rangle = 0. \quad (20)$$

All proofs are presented in the Appendix. Here we provide an intuition. The expression $\langle \mathbf{v}, \mathbf{m} \rangle = \sum_{\mathbf{h}} v_{\mathbf{h}} m_{\mathbf{h}}$ can be interpreted as player 1's average cooperation rate across all rounds of the repeated game. To compute that average cooperation rate, one first draws an n -history \mathbf{h} (with probability $v_{\mathbf{h}}$), and then one computes how likely player 1 would cooperate in the subsequent round (with probability $m_{\mathbf{h}}$). Alternatively, one could compute the average cooperation rate by drawing an n -history \mathbf{h} and then checking how likely player 1 was to cooperate k rounds ago, according to that n -history. That second interpretation leads to the expression $\langle \mathbf{v}, \mathbf{m}^{k-Rep} \rangle$. According to Eq. (20), both interpretations are equivalent.

2 Characterizing the partner strategies among the reactive- n strategies

2.1 Partner strategies

In this study, we are interested in identifying strategies that can sustain full cooperation in a Nash equilibrium. Strategies with these properties have been termed as being of *Nash type* by Akin [6], or as *partner strategies* by Hilbe *et al* [8]. In the following, we formally define them.

Definition (Partner strategies)

- (i) A strategy σ for the repeated prisoner's dilemma is a *Nash equilibrium* if it is a best response to itself. That is, we require $\pi^1(\sigma, \sigma)$ to exist and

$$\pi^1(\sigma, \sigma) \geq \pi^1(\sigma', \sigma) \text{ for all other strategies } \sigma' \text{ for which } \pi^1(\sigma', \sigma) \text{ exists.} \quad (21)$$

- (ii) A player's strategy is *nice*, if the player is never the first to defect.
- (iii) A *partner strategy* is a strategy that is both nice and a Nash equilibrium.

Several remarks are in order. First, we note that when two players with nice strategies interact, they both cooperate in every round. Partner strategies thus sustain mutual cooperation in a Nash equilibrium. Second, if a memory- n strategy $\mathbf{m} = (m_{\mathbf{h}})_{\mathbf{h} \in H}$ is to be nice, it needs to cooperate after n rounds of mutual cooperation. In other words, if $\mathbf{h}_C = (\mathbf{h}_C^i, \mathbf{h}_C^{-i})$ is the n -history that consists of mutual cooperation for the past n rounds, then the strategy needs to respond by cooperating with certainty, $m_{\mathbf{h}_C} = 1$. Similarly, a nice reactive- n strategy needs to satisfy $p_{\mathbf{h}_C^{-i}} = 1$. Third, we note that our definition of Nash equilibria only requires that players cannot profitably deviate towards strategies *for which a payoff can be defined*. If the strategy σ is a memory- n strategy, in the following we make the slightly looser requirement that the strategy is a best response among all σ' for which the limit (14) exists. Fourth, in general it is a difficult task to verify that any given strategy σ is a Nash equilibrium. After all, one needs to verify that it yields the highest payoff according to Eq. (21) among all (uncountably) many alternative strategies σ' . Fortunately, the situation is

somewhat simpler if the strategy under consideration is a memory- n strategy. In that case, it follows from an argument by Press and Dyson [9] that one only needs to compare the strategy to all other memory- n strategies. However, this still leaves us with uncountably many strategies to check. In fact, it is one aim of this paper to show that for reactive- n strategies, it suffices to check finitely many alternative strategies.

2.2 Tit For Tat and Generous Tit For Tat with arbitrary memory lengths

Zero-determinant strategies with n rounds memory. Before we provide a general algorithm to identify reactive- n partner strategies, we first generalize some of the well-known reactive-1 partner strategies, TFT and GTFT, to the case of memory- n . To this end, we use Lemma 1 to develop a theory of zero-determinant strategies within the class of memory- n strategies, see also Refs. [10, 11]. In the following, we say a memory- n strategy \mathbf{m} is a *zero-determinant strategy* if there are integers $k_1, k_2, k_3 \leq n$ and real numbers α, β, γ such that \mathbf{m}^i can be written as

$$\mathbf{m}^i = \alpha \mathbf{g}^i(k_1) + \beta \mathbf{g}^{-i}(k_2) + \gamma \mathbf{1} + \mathbf{m}^{k_3-\text{Rep}}. \quad (22)$$

In this expression, $\mathbf{g}^i(k)$ is the vector that returns player i 's payoff k rounds ago, as defined by Eq. (15), $\mathbf{m}^{k-\text{Rep}}$ is the memory- n strategy that repeats player i 's own move k rounds ago, and $\mathbf{1}$ is the 2^{2n} -dimensional vector for which every entry is one. Using the generalized version of Akin's Lemma, we obtain

$$\begin{aligned} 0 &\stackrel{(20)}{=} \langle \mathbf{v}, \mathbf{m} - \mathbf{m}^{k_3-\text{Rep}} \rangle \\ &\stackrel{(22)}{=} \langle \mathbf{v}, \alpha \mathbf{g}^i(k_1) + \beta \mathbf{g}^{-i}(k_2) + \gamma \mathbf{1} \rangle \\ &= \alpha \langle \mathbf{v}, \mathbf{g}^i(k_1) \rangle + \beta \langle \mathbf{v}, \mathbf{g}^{-i}(k_2) \rangle + \gamma \langle \mathbf{v}, \mathbf{1} \rangle \\ &\stackrel{(17),(18)}{=} \alpha \pi^i(\mathbf{m}^i, \sigma^{-i}) + \beta \pi^{-i}(\mathbf{m}^i, \sigma^{-i}) + \gamma. \end{aligned} \quad (23)$$

That is, a player with a zero-determinant strategy enforces a linear relationship between the players' payoffs, irrespective of the co-player's strategy. Remarkably, the parameters α, β , and γ of that linear relationship are entirely under player i 's control.

Generalized versions of Tit-for-tat. One interesting special case arises if $k_1 = k_2 = k_3 =: k$ and $\alpha = -\beta = 1/(T-S)$, $\gamma = 0$. In that case, formula (22) yields the strategy with entries

$$m_{\mathbf{h}} = \begin{cases} 1 & \text{if } a_{-k}^{-i} = C \\ 0 & \text{if } a_{-k}^{-i} = D \end{cases}$$

Therefore, a player with that strategy cooperates if and only if the co-player cooperated k rounds ago. Thus, the strategy implements TFT (for $k = 1$) or delayed versions thereof (for $k > 1$). By Eq. (23), the strategy enforces equal payoffs against any co-player,

$$\pi^i(\mathbf{m}^i, \sigma^{-i}) = \pi^{-i}(\mathbf{m}^i, \sigma^{-i}). \quad (24)$$

Moreover, this strategy is nice if we additionally require it to unconditionally cooperate during the first k rounds. Given this additional requirement, the payoff of \mathbf{m}^i against itself is R . Moreover, the strategy is a Nash equilibrium. To see why, suppose to the contrary that there is a strategy σ^{-i} with $\pi^{-i}(\mathbf{m}^i, \sigma^{-i}) > R$. Then it follows from (24) that $\pi^i(\mathbf{m}^i, \sigma^{-i}) + \pi^{-i}(\mathbf{m}^i, \sigma^{-i}) > 2R$. That is, the total payoff per round exceeds $2R$, which is incompatible with the basic assumptions on a prisoner's dilemma, Eq. (2). We conclude that all these versions of TFT are nice and they are Nash equilibria. Hence, they are partner strategies.

Generalized versions of Generous Tit-for-Tat. Another interesting special case arises in the donation game if $k_1 = k_2 = k_3 =: k$ and $\alpha = 0$, $\beta = -1/b$, $\gamma = 1 - c/b$. In that case Eq. (22) yields the strategy with entries

$$m_{\mathbf{h}} = \begin{cases} 1 & \text{if } a_{-k}^{-i} = C \\ 1 - c/b & \text{if } a_{-k}^{-i} = D \end{cases}$$

That is, the generated strategy is GTFT (if $k = 1$), or a delayed version thereof (for $k > 1$). By Eq. (23), the enforced payoff relationship is $\pi^{-i}(\mathbf{m}^i, \sigma^{-i}) = b - c$. That is, the co-player always obtains the mutual cooperation payoff, irrespective of the co-player's strategy. In particular, all these versions of GTFT are Nash equilibria (independent of how they act during the first n rounds). If we additionally require them to cooperate during the first n rounds, they are also nice. Hence, they are partner strategies.

2.3 An algorithm to check whether a reactive- n strategy is a Nash equilibrium

Sufficiency of checking pure self-reactive strategies. After discussing these particular cases, we would like to derive a general algorithm that allows us to verify whether a given reactive- n strategy is a Nash equilibrium. In principle, this requires us to check the payoff of any other strategy (including strategies that have a much longer memory length than n). Fortunately, however, some simplifications are possible when we use an insight by Press and Dyson [9]. They discussed the case where one player uses a memory-1 strategy and the other player employs a longer memory strategy. They demonstrated that the payoff of the player with the longer memory is exactly the same as if the player had employed a specific shorter-memory strategy, disregarding any history beyond what is shared with the short-memory player. Here we show a result that follows a similar intuition. If there is a part of the game's history that one player does not take into account, then the co-player gains nothing by considering that part of the history.

Lemma 2 (Against reactive strategies, any feasible payoff can be generated with self-reactive strategies)

Let $\mathbf{p} \in \mathcal{R}_n$ be a reactive strategy for player 2. Moreover, suppose player 1 adopts some strategy σ such that for the resulting game, the time average \mathbf{v} according to Eq. (14) exists. Then there is a self-reactive- n strategy $\tilde{\mathbf{p}} \in \mathcal{S}_n$ such that $\pi^i(\sigma, \mathbf{p}) = \pi^i(\tilde{\mathbf{p}}, \mathbf{p})$ for $i \in \{1, 2\}$.

For an illustration of this result, see Figure S1. It shows that against a reactive-2 player, any payoff that can be achieved with a memory-2 strategy can already be achieved with a self-reactive-2 strategy.

If we are to verify that some given reactive- n strategy \mathbf{p} is a Nash equilibrium, Lemma 2 simplifies our task considerably. Instead of checking condition (21) for all possible strategies σ' , we only need to check it

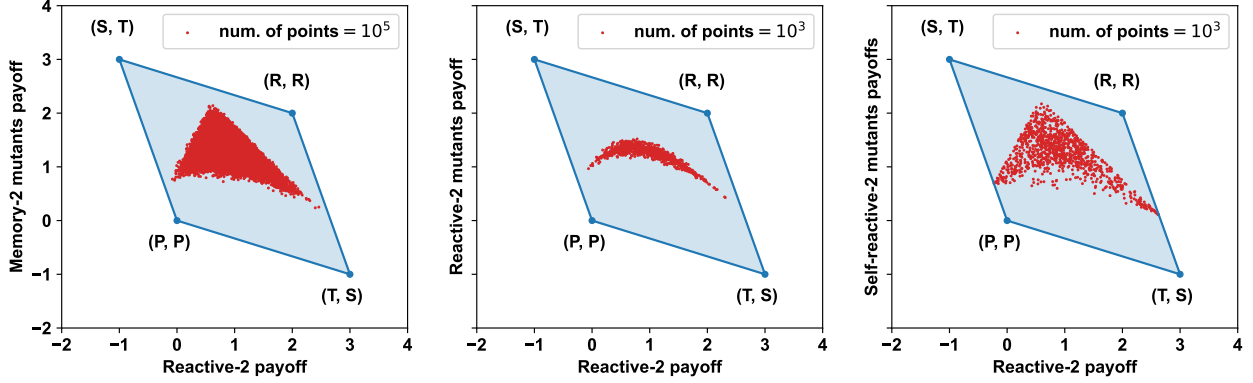


Figure S1: Feasible payoffs for a reactive-2 strategy. We consider a player with reactive-2 strategy $\mathbf{p} = (0.37, 0.89, 0.95, 0.23)$. The player interacts with many other players (referred to as ‘mutants’) who adopt either some random memory-2 strategy (left), a random reactive-2 strategy (middle), or a random self-reactive-2 strategy (right panel). The panels show the resulting payoffs to the two players as red dots, with the x -axis showing the payoff of the focal player, and the y -axis showing the payoff of the mutants. We observe that when mutants use memory-2 strategies and self-reactive-2 strategies, we obtain the same region of feasible payoffs, in line with Lemma 2. In contrast, if mutants are restricted to reactive-2 strategies, the set of feasible payoffs is strictly smaller. Here, we consider a donation game with $b=3$ and $c=1$.

for all self-reactive strategies $\tilde{\mathbf{p}} \in \mathcal{S}_n$. The following result simplifies our task even further.

Theorem 1 (To any reactive strategy, there is a best response among the pure self-reactive strategies)

For any reactive strategy $\mathbf{p} \in \mathcal{R}_n$ there is some pure self-reactive strategy $\tilde{\mathbf{p}} \in \mathcal{S}_n$ such that

$$\pi^1(\tilde{\mathbf{p}}, \mathbf{p}) \geq \pi^1(\sigma', \mathbf{p}) \text{ for all other strategies } \sigma' \text{ for which the limit (14) exists.} \quad (25)$$

This result implies that we only need to check finitely many other strategies if we are to verify that some given reactive- n strategy is a Nash equilibrium.

Corollary 1 (An algorithm to check whether a reactive- n strategy is a Nash equilibrium)

A reactive strategy $\mathbf{p} \in \mathcal{R}_n$ is a Nash equilibrium if and only if $\pi^1(\mathbf{p}, \mathbf{p}) \geq \pi^1(\tilde{\mathbf{p}}, \mathbf{p})$ for all pure self-reactive strategies $\tilde{\mathbf{p}} \in \mathcal{S}_n$.

Corollary 1 gives us a straightforward procedure to check whether a given reactive strategy \mathbf{p} is a Nash equilibrium (for a depiction, see Algorithm 1). To verify that \mathbf{p} is a Nash equilibrium, we merely need to compare its payoff against itself to the payoff of a deviation towards one of the 2^n pure self-reactive strategies.

A more efficient way to calculate payoffs. For the remainder of this section, we thus assume that player 1 uses a self-reactive- n strategy $\tilde{\mathbf{p}} = (\tilde{p}_{\mathbf{h}^i})_{\mathbf{h}^i \in H^i}$, whereas player 2 uses a reactive- n strategy $\mathbf{p} = (p_{\mathbf{h}^{-i}})_{\mathbf{h}^{-i} \in H^{-i}}$. Our algorithm to compute payoffs for the two players in Section 1.3 would require us to interpret the two strategies as memory- n strategies. We would thus compute a left eigenvector of a $2^{2n} \times 2^{2n}$ transition matrix. In the following, however, we show that for games between reactive and self-reactive players, it suffices to consider a $2^n \times 2^n$ transition matrix. This efficiency gain is possible because both players only consider

Algorithm 1: An algorithm to verify whether a given reactive strategy \mathbf{p} is a Nash equilibrium.

input: \mathbf{p}, n
 $\text{pure_self_reactive_strategies} \leftarrow \{\tilde{\mathbf{p}} \mid \tilde{\mathbf{p}} \in \{0, 1\}^{2^n}\};$
 $\text{isNash} \leftarrow \text{True};$
for $\tilde{\mathbf{p}} \in \text{pure_self_reactive_strategies}$ **do**
 if $\pi^1(\mathbf{p}, \mathbf{p}) < \pi^1(\tilde{\mathbf{p}}, \mathbf{p})$ **then**
 $\text{isNash} \leftarrow \text{False};$
return $(\mathbf{p}, \text{isNash})$;

player 1's past actions. Instead of taking the space of all of the game's n -histories $H = H^1 \times H^2$ as the state space, we can thus take the space H^1 . Let $h^1 = (a_{-n}^1, \dots, a_{-1}^1)$ be the state in the current round. Then we obtain the following probability that the state after one round is $\tilde{h}^1 = (\tilde{a}_{-n}^1, \dots, \tilde{a}_{-1}^1)$,

$$\tilde{M}_{h^1, \tilde{h}^1} = \begin{cases} \tilde{p}_{h^1} & \text{if } \tilde{a}_{-1}^1 = C, \text{ and } \tilde{a}_{-t}^1 = a_{-t+1}^1 \text{ for all } t \in \{2, \dots, n\} \\ 1 - \tilde{p}_{h^1} & \text{if } \tilde{a}_{-1}^1 = D, \text{ and } \tilde{a}_{-t}^1 = a_{-t+1}^1 \text{ for all } t \in \{2, \dots, n\} \\ 0 & \text{if } \tilde{a}_{-t}^1 \neq a_{-t+1}^1 \text{ for some } t \in \{2, \dots, n\}. \end{cases} \quad (26)$$

Similar to the vector \mathbf{v} for matrix M , let $\tilde{\mathbf{v}} = (\tilde{v}_{h^1})_{h^1 \in H^1}$ be the limiting distribution of the dynamics defined by \tilde{M} (which only in exceptional cases depends on player 1's behavior during the first n rounds). Then the players' payoffs are given by

$$\begin{aligned} \pi^1(\tilde{\mathbf{p}}, \mathbf{p}) &= \sum_{h^1 \in H^1} \tilde{v}_{h^1} \left(\tilde{\mathbf{p}}_{h^1} \mathbf{p}_{h^1} \cdot R + \tilde{\mathbf{p}}_{h^1} (1 - \mathbf{p}_{h^1}) \cdot S + (1 - \tilde{\mathbf{p}}_{h^1}) \mathbf{p}_{h^1} \cdot T + (1 - \tilde{\mathbf{p}}_{h^1}) (1 - \mathbf{p}_{h^1}) \cdot P \right), \\ \pi^2(\tilde{\mathbf{p}}, \mathbf{p}) &= \sum_{h^1 \in H^1} \tilde{v}_{h^1} \left(\tilde{\mathbf{p}}_{h^1} \mathbf{p}_{h^1} \cdot R + \tilde{\mathbf{p}}_{h^1} (1 - \mathbf{p}_{h^1}) \cdot T + (1 - \tilde{\mathbf{p}}_{h^1}) \mathbf{p}_{h^1} \cdot S + (1 - \tilde{\mathbf{p}}_{h^1}) (1 - \mathbf{p}_{h^1}) \cdot P \right). \end{aligned} \quad (27)$$

Example: Payoffs and best responses with one-round memory. To illustrate the above results, we consider the case $n = 1$. Assume player 1's self-reactive strategy is $\tilde{\mathbf{p}}^1 = (\tilde{p}_C^1, \tilde{p}_D^1)$ and player 2's reactive strategy is $\mathbf{p}^2 = (p_C^2, p_D^2)$. If we use the algorithm in Section 1.3, we first formally represent these strategies as memory-1 strategies, $\mathbf{m}^1 = (\tilde{p}_C^1, \tilde{p}_C^1, \tilde{p}_D^1, \tilde{p}_D^1)$ and $\mathbf{m}^2 = (p_C^2, p_D^2, p_C^2, p_D^2)$. The respective transition matrix according to Eq. (12) is

$$M = \begin{pmatrix} \tilde{p}_C^1 p_C^2 & \tilde{p}_C^1 (1 - p_C^2) & (1 - \tilde{p}_C^1) p_C^2 & (1 - \tilde{p}_C^1) (1 - p_C^2) \\ \tilde{p}_C^1 p_D^2 & \tilde{p}_C^1 (1 - p_D^2) & (1 - \tilde{p}_C^1) p_D^2 & (1 - \tilde{p}_C^1) (1 - p_D^2) \\ \tilde{p}_D^1 p_C^2 & \tilde{p}_D^1 (1 - p_C^2) & (1 - \tilde{p}_D^1) p_C^2 & (1 - \tilde{p}_D^1) (1 - p_C^2) \\ \tilde{p}_D^1 p_D^2 & \tilde{p}_D^1 (1 - p_D^2) & (1 - \tilde{p}_D^1) p_D^2 & (1 - \tilde{p}_D^1) (1 - p_D^2) \end{pmatrix}. \quad (28)$$

Assuming player 1's strategy is different from the one-round repeat strategy, $\tilde{\mathbf{p}}^1 \neq (1, 0)$, this transition matrix has a unique invariant distribution,

$$\mathbf{v} = \left(\frac{\tilde{p}_D^1 (\tilde{p}_C^1 (p_C^2 - p_D^2) + p_D^2)}{1 - (\tilde{p}_C^1 - \tilde{p}_D^1)}, \frac{\tilde{p}_D^1 (1 - \tilde{p}_C^1 (p_C^2 - p_D^2) - p_D^2)}{1 - (\tilde{p}_C^1 - \tilde{p}_D^1)}, \frac{(1 - \tilde{p}_C^1) (\tilde{p}_D^1 (p_C^2 - p_D^2) + p_D^2)}{1 - (\tilde{p}_C^1 - \tilde{p}_D^1)}, \frac{(1 - \tilde{p}_C^1) (1 - \tilde{p}_D^1 (p_C^2 - p_D^2) - p_D^2)}{1 - (\tilde{p}_C^1 - \tilde{p}_D^1)} \right).$$

According to Eq. (16), Player 1's payoff is the scalar product

$$\pi^1(\tilde{\mathbf{p}}^1, \mathbf{p}^2) = \langle \mathbf{v}, (R, S, T, P) \rangle. \quad (29)$$

Following Corollary 1, we can use these observations to characterize under which conditions a nice reactive strategy $\mathbf{p}^2 = (1, p_D^2)$ is a partner. To this end, we compute player 1's payoff for all pure self-reactive strategies $\tilde{\mathbf{p}}^1 = (\tilde{p}_C^1, \tilde{p}_D^1)$. These are ALLC = (1, 1), ALLD = (0, 0), and Alternator = (0, 1); we can ignore the one-round repeat strategy (1, 0), because depending on the strategy's first round-behavior it is either equivalent to ALLC or to ALLD. The payoffs of these three strategies are

$$\begin{aligned} \pi^1(\text{ALLC}, \mathbf{p}^2) &= R, \\ \pi^1(\text{ALLD}, \mathbf{p}^2) &= p_D^2 \cdot T + (1 - p_D^2) \cdot P \\ \pi^1(\text{Alternator}, \mathbf{p}^2) &= p_D^2 / 2 \cdot R + (1 - p_D^2) / 2 \cdot S + 1/2 \cdot T. \end{aligned} \quad (30)$$

We conclude that player 2's reactive strategy \mathbf{p}^2 is a Nash equilibrium (and hence a partner) if none of these three payoffs exceeds the mutual cooperation payoff R . This requirement yields the condition

$$p_D^2 \leq \min \{ 1 - (T - R) / (R - S), (R - P) / (T - P) \}. \quad (31)$$

As one may expect, \mathbf{p}^2 is a partner if and only if its generosity p_D^2 does not exceed the generosity of GTFT, as defined by Eq. (8).

Instead of computing the 4×4 matrix M in (28), we could also consider the simplified 2×2 transition matrix (26). Here, the two possible states are $\mathbf{h}^1 \in \{C, D\}$, and hence the matrix is

$$\tilde{M} = \begin{pmatrix} \tilde{p}_C^1 & 1 - \tilde{p}_C^1 \\ \tilde{p}_D^1 & 1 - \tilde{p}_D^1 \end{pmatrix}. \quad (32)$$

Again, for $\tilde{\mathbf{p}}^1 \neq (1, 0)$, this transition matrix has a unique invariant distribution,

$$\tilde{\mathbf{v}} = (\tilde{v}_C, \tilde{v}_D) = \left(\frac{\tilde{p}_D^1}{1 - (\tilde{p}_C^1 - \tilde{p}_D^1)}, \frac{1 - \tilde{p}_C^1}{1 - (\tilde{p}_C^1 - \tilde{p}_D^1)} \right). \quad (33)$$

If we take this invariant distribution and compute player 1's payoff according to Eq. (27), we recover the same expression as in Eq. (29), as expected.

2.4 Reactive partner strategies in the donation game

Just as in the previous example with $n = 1$, we can use the results of the previous section to characterize the partner strategies for reactive-2 and reactive 3-strategies. For simplicity, we first consider the case of the donation game. Results for the general prisoner's dilemma follow in the next section.

Reactive-2 partner strategies. We first consider the case $n = 2$. The resulting reactive-2 strategies can be represented as a vector $\mathbf{p} = (p_{CC}, p_{CD}, p_{DC}, p_{DD})$. The entries $p_{\mathbf{h}^{-i}}$ are the player's cooperation probability, depending on the co-player's actions in the previous two rounds, $\mathbf{h}^{-i} = (a_{-2}^{-i}, a_{-1}^{-i})$. For the strategy to be nice, we require $p_{CC} = 1$. Based on Corollary 1, we obtain the following characterization of partners.

Theorem 2 (Reactive-2 partner strategies in the donation game)

A nice reactive-2 strategy \mathbf{p} , is a partner strategy if and only if its entries satisfy the conditions

$$p_{CC} = 1, \quad \frac{p_{CD} + p_{DC}}{2} \leq 1 - \frac{1}{2} \cdot \frac{c}{b}, \quad p_{DD} \leq 1 - \frac{c}{b}. \quad (34)$$

The resulting conditions can be interpreted as follows: For each time a co-player has defected during the past two rounds, the reactive player's cooperation probability needs to decrease by $c/(2b)$. This reduced cooperation probability is sufficient to incentivize the co-player to cooperate. Interestingly, for the strategy to be a partner, the middle condition in (34) suggests that the exact timing of a co-player's defection is irrelevant. As long as *on average*, the respective cooperation probabilities p_{CD} and p_{DC} are below the required threshold $1 - c/(2b)$, the strategy is a Nash equilibrium.

The conditions for a partner become even simpler for *reactive- n counting strategies*. To define these strategies, let $|\mathbf{h}^{-i}|$ denote the number of C 's in a given n -history of the co-player. We say a reactive- n strategy $\mathbf{p} = (p_{\mathbf{h}^{-i}})_{\mathbf{h}^{-i} \in \mathbf{H}^{-i}}$ is a counting strategy if

$$|\mathbf{h}^{-i}| = |\tilde{\mathbf{h}}^{-i}| \Rightarrow p_{\mathbf{h}^{-i}} = p_{\tilde{\mathbf{h}}^{-i}}. \quad (35)$$

That is, the reactive player's cooperation probability only depends on the number of cooperative acts during the past n rounds and not on their timing. Such reactive- n counting strategies can be written as $n + 1$ -dimensional vectors $\mathbf{r} = (r_k)_{k \in \{n, \dots, 1\}}$, where r_i is the player's cooperation probability if the co-player cooperated k times during the past n rounds. In particular, for reactive-2 counting strategies, we associate $r_2 = p_{CC}$, $r_1 = p_{CD} = p_{DC}$, and $r_0 = p_{DD}$. The following characterization of partners among the reactive-2 counting strategies then follows immediately from Theorem 2.

Corollary 2 (Partners among the reactive-2 counting strategies)

A nice reactive-2 counting strategy $\mathbf{r} = (r_2, r_1, r_0)$ is a partner strategy if and only if

$$r_2 = 1, \quad r_1 \leq 1 - \frac{1}{2} \cdot \frac{c}{b}, \quad r_0 \leq 1 - \frac{c}{b}. \quad (36)$$

Reactive-3 Partner Strategies. Next, we focus on the case $n = 3$. Reactive-3 strategies can be represented as a vector $\mathbf{p} = (p_{CCC}, p_{CCD}, p_{CDC}, p_{CDD}, p_{DCC}, p_{DCD}, p_{DDC}, p_{DDD})$. Again, each entry $p_{\mathbf{h}^{-i}}$ refers to the player's cooperation probability, depending on the co-player's previous three actions, $\mathbf{h}^{-i} = (a_{-3}^{-i}, a_{-2}^{-i}, a_{-1}^{-i})$. For the respective partner strategies, we obtain the following characterization.

Theorem 3 (Reactive-3 partner strategies in the donation game)

A nice reactive-3 strategy \mathbf{p} is a partner strategy if and only if its entries satisfy the conditions

$$\begin{aligned}
 p_{CCC} &= 1 \\
 \frac{p_{CDC} + p_{DCD}}{2} &\leq 1 - \frac{1}{2} \cdot \frac{c}{b} \\
 \frac{p_{CCD} + p_{CDC} + p_{DCC}}{3} &\leq 1 - \frac{1}{3} \cdot \frac{c}{b} \\
 \frac{p_{CDD} + p_{DCD} + p_{DDC}}{3} &\leq 1 - \frac{2}{3} \cdot \frac{c}{b} \\
 \frac{p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}}{4} &\leq 1 - \frac{1}{2} \cdot \frac{c}{b} \\
 p_{DDD} &\leq 1 - \frac{c}{b}
 \end{aligned} \tag{37}$$

As before, the average of certain cooperation probabilities need to be below specific thresholds. However, compared to the case $n=2$, the respective conditions are now somewhat more difficult to interpret. The conditions again become more straightforward if we further restrict attention to reactive-3 counting strategies.

Corollary 3 (Partners among the reactive-3 counting strategies)

A nice reactive-3 counting strategy $\mathbf{r} = (r_3, r_2, r_1, r_0)$ is a partner strategy if and only if

$$r_3 = 1 \quad r_2 \leq 1 - \frac{1}{3} \cdot \frac{c}{b}, \quad r_1 \leq 1 - \frac{2}{3} \cdot \frac{c}{b}, \quad r_0 \leq 1 - \frac{c}{b}. \tag{38}$$

As in the case of $n=2$ we observe here that with each additional defection of the opponent in memory, the focal player reduces its conditional cooperation probability by a constant, in this case $c/(3b)$.

Partners among the reactive- n counting strategies. Using the same methods as before, one can in principle also characterize the partners among the reactive-4 or the reactive-5 strategies. However, the respective conditions quickly become unwieldy. In case of the counting strategies, however, the simple pattern in Corollaries 2 and 3 does generalize to arbitrary memory lengths.

Theorem 4 (Partners among the reactive- n counting strategies)

A nice reactive- n counting strategy $\mathbf{r} = (r_k)_{k \in \{n, n-1, \dots, 0\}}$, is a partner strategy if and only if

$$r_n = 1 \quad \text{and} \quad r_{n-k} \leq 1 - \frac{k}{n} \cdot \frac{c}{b} \quad \text{for } k \in \{1, 2, \dots, n\}. \tag{39}$$

2.5 Reactive partner strategies in the general prisoner's dilemma

In the previous section, we have characterized the reactive partner strategies for a special case of the prisoner's dilemma, the donation game. In the following, we apply the same methods based on Section 2.3

to analyze the general prisoner's dilemma. For the case of reactive-2 strategies, we obtain the following characterization.

Theorem 5 (Reactive-2 partner strategies in the prisoner's dilemma)

A nice reactive-2 strategy \mathbf{p} is a partner strategy if and only if its entries satisfy the conditions

$$\begin{aligned}
p_{CC} &= 1, \\
(T - P) p_{DD} &\leq R - P, \\
(R - S) (p_{CD} + p_{DC}) &\leq 3R - 2S - T, \\
(T - P) p_{DC} + (R - S) p_{CD} &\leq 2R - S - P, \\
(T - P) (p_{CD} + p_{DC}) + (R - S) p_{DD} &\leq 3R - S - 2P, \\
(T - P) p_{CD} + (R - S) (p_{CD} + p_{DD}) &\leq 4R - 2S - T - P.
\end{aligned} \tag{40}$$

Compared to the donation game, there are now more conditions, and these conditions are somewhat more difficult to interpret. Reassuringly, however, the conditions simplify to the conditions (34) in the special case that the payoff values satisfy $R = b - c$, $S = -c$, $T = b$, and $P = 0$. For the case of reactive-3 strategies, the characterization is as follows.

Theorem 6 (Reactive-3 partner strategies in the prisoner's dilemma)

A nice reactive-3 strategy \mathbf{p} is a partner strategy if and only if its entries satisfy the conditions in Table 1.

Given the large number of conditions in Table 1, we do not pursue a similar characterization for $n > 3$, even though the same methods remain applicable.

$$\begin{aligned}
p_{CCC} &= 1, \\
(T - P)(p_{CDD} + p_{DCD} + p_{DDC}) + (R - S)p_{DDD} &\leq 4R - 3P - S \\
(T - P)p_{CDC} + (R - S)p_{DCD} &\leq 2R - P - S \\
(T - P)p_{DDD} &\leq R - P \\
(T - P)(p_{CCD} + p_{CDD} + p_{DDC}) + (R - S)(p_{CDC} + p_{DCC} + p_{DCD} + p_{DDD}) &\leq 8R - 3P - 4S - T \\
(T - P)p_{DCC} + (R - S)(p_{CCD} + p_{CDC}) &\leq 3R - P - 2S \\
(T - P)(p_{CCD} + p_{DCC} + p_{DDC}) + (R - S)(p_{CDC} + p_{CDD} + p_{DCD}) &\leq 6R - 3P - 3S \\
(T - P)(p_{CCD} + p_{DDC}) + (R - S)(p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD}) &\leq 7R - 2P - 4S - T \\
(T - P)(p_{CCD} + p_{CDD} + p_{DCC}) + (R - S)(p_{DDC} + p_{DDD}) &\leq 5R - 3P - 2S \\
(T - P)(p_{DCD} + p_{DDC}) + (R - S)p_{CDD} &\leq 3R - 2P - S \\
(T - P)p_{CCD} + (R - S)(p_{CDD} + p_{DCC} + p_{DDC}) &\leq 5R - P - 3S - T \\
(T - P)(p_{CCD} + p_{DCC}) + (R - S)(p_{CDD} + p_{DDC}) &\leq 4R - 2P - 2S \\
(T - P)(p_{CDC} + p_{DCD}) + (R - S)(p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}) &\leq 7R - 2P - 4S - T \\
(T - P)(p_{CDC} + p_{CDD} + p_{DCD}) + (R - S)(p_{CCD} + p_{DCC} + p_{DDC} + p_{DDD}) &\leq 8R - 3P - 4S - T \\
(T - P)(p_{CDC} + p_{DCC} + p_{DCD}) + (R - S)(p_{CCD} + p_{CDD} + p_{DDC}) &\leq 6R - 3P - 3S \\
(T - P)(p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}) + (R - S)(p_{CDC} + p_{DCD} + p_{DDD}) &\leq 7R - 4P - 3S \\
(R - S)(p_{CCD} + p_{CDC} + p_{DCC}) &\leq 4R - 3S - T \\
(T - P)(p_{CCD} + p_{CDD}) + (R - S)(p_{DCC} + p_{DDC} + p_{DDD}) &\leq 6R - 2P - 3S - T \\
(T - P)(p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD}) + (R - S)(p_{CCD} + p_{DDC} + p_{DDD}) &\leq 7R - 4P - 3S
\end{aligned}$$

Table 1: Necessary and sufficient conditions for a nice reactive-3 strategy to be a partner in the prisoner's dilemma.

2.6 Reactive defecting Nash strategies in the donation game

In the previous section, we have characterized the reactive partner strategies for a special case of the donation game and the general prisoner's dilemma. In the following, we apply the same methods based on Section 2.3 to analyze defecting Nash equilibria. For the case of reactive-1 strategies, we obtain the following characterization.

Theorem 7 (Reactive-1 defecting Nash strategies in the donation game)

A reactive-1 strategy \mathbf{p} is a defecting Nash strategy if and only if its entries satisfy the conditions

$$p_C \leq \frac{c}{b} \quad \text{and} \quad p_D = 0. \quad (41)$$

Theorem 8 (Reactive-2 defecting Nash strategies in the donation game)

A reactive-2 strategy \mathbf{p} is a defecting Nash strategy if and only if its entries satisfy the conditions

$$p_{CC} \leq \frac{c}{b}, \quad \frac{p_{CD} + p_{DC}}{2} \leq \frac{c}{2b}, \quad p_{DD} = 0. \quad (42)$$

Theorem 9 (Reactive-3 defecting Nash strategies in the donation game)

A reactive-3 strategy \mathbf{p} is a defecting Nash strategy if and only if its entries satisfy the conditions

$$\begin{aligned} p_{CCC} &\leq \frac{c}{b} \\ \frac{p_{CDC} + p_{DCC}}{2} &\leq \frac{1}{2} \cdot \frac{c}{b} \\ \frac{p_{CCD} + p_{CDC} + p_{DCC}}{3} &\leq \frac{2}{3} \cdot \frac{c}{b} \\ \frac{p_{CDD} + p_{DCD} + p_{DDC}}{3} &\leq \frac{1}{3} \cdot \frac{c}{b} \\ \frac{p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}}{4} &\leq \frac{1}{2} \cdot \frac{c}{b} \\ p_{DDD} &= 0. \end{aligned} \quad (43)$$

We can observe that for each value of n , the left-hand side of the conditions for cooperative and defective Nash are the same. Moreover, it is clear that the right-hand side of the defective Nash conditions is always strictly smaller than those of the cooperative Nash conditions. This means that within the space of feasible strategies, the volume of partner strategies is always larger than the volume of defective Nash strategies. We verified these analytical results numerically as well.

We selected random strategies from the feasible space of strategies and created two copies of each strategy. For one copy, we set the probability of cooperating after full cooperation of the co-player to 1 (for example, for reactive-1, $p_C = 1$). For the second copy, we set the probability of cooperating after full defections of the co-player to 0 (for example, for reactive-2, $p_{DD} = 0$). We then checked if either copy is Nash: cooperative for the first and defective for the second. We repeated this process for 10^4 randomly selected

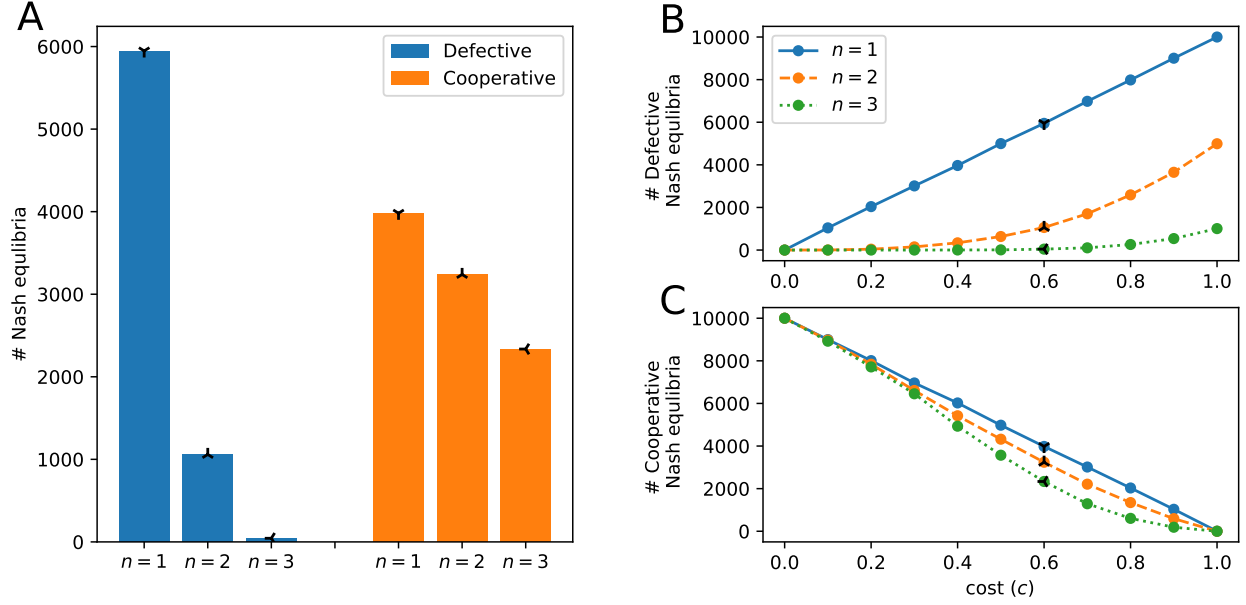


Figure S2: Volume of cooperative and defective Nash. We draw 10^4 random strategies from the feasible space of strategies and create two copies of each strategy. For one copy, we set the probability of cooperating after full cooperation of the co-player to 1. For the second copy, we set the probability of cooperating after full defections of the co-player to 0. We then checked if either copy is Nash: cooperative for the first and defective for the second. We set the benefit of cooperation to $b = 1$. **A** We plot the results for a given value of cost, $c = 0.5$. **B** The number of defective Nash strategies as a function of cost. **C** The number of cooperative Nash strategies as a function of cost.

strategies and plotted the relative volumes of cooperative and defective Nash equilibria (Figure S2). We also verified that this holds true for different values of cost.

Nash strategies are best response to themselves. This makes them robust against invasion. In the stochastic evolutionary processes considered in this work, this doesn't mean they cannot be invaded, but rather it makes it harder to do so. The volumes of cooperative Nash and defective Nash strategies can determine the evolutionary dynamics of cooperation. That is, they indicate the likelihood that a randomly selected strategy will produce sustained cooperation or defection. Therefore, since the volume of cooperative Nash increases with memory, and defective Nash decreases, we expect that as memory increases, the cooperation rate of the evolved population will also increase. This expectation is verified by the results in **Figure 4** of the paper.

3 Evolutionary Simulations

In this section, we will analyze the results of the evolutionary simulations to better understand the dynamics of the different strategies. As a reminder, in Fig. 3 of the manuscript panel, we simulated the evolutionary process twenty times, independent simulations for reactive- n strategies, and for reactive- n counting strategies. In panel A, we show for each of the three different memory values the strategies that were the most abundant from each simulation.

Now, one question that arises is how many of these strategies are actually partner strategies? And for the partner strategies, do we see all of them represented?

****Reactive-1 strategies.**** In the case of reactive-1, all the most abundant strategies are partner strategies. However, from Eq. and for a given value of c , the probability of defecting is... However, we observe only a lower value of p_{dd} , thus only some partner strategies are residents. This result is not explained by the theory. To better understand this result, we need to do an invasion analysis. For the invasion analysis, we consider a given resident strategy, and we sample 10^3 random strategies to estimate if this mutant will invade the resident. We record the number of mutants that took over. We repeat this 10^3 times. The results are shown in Fig. X, and what we observe is that a lower value of p_{dd} results in a more robust strategy to invasion, as only cooperative strategies on the GTFT can take over, in comparison with panel B where defective strategies are more likely to invade.

****Reactive-2 strategies.**** In the case of reactive-2, we observe that almost all strategies are partner strategies. However, what seems to be the case is that p_{dc} is higher on average than p_{cd} . In order to understand this, we again run an invasion analysis. This time we also consider two values of p_{dd} . The results are shown in Fig. Y.

4 Errors

So far, we have considered the case where there cannot be a mistake in the actions taken by a player; the actions of the players are realized without error. Here, we discuss what happens in the case where such an error is possible. More specifically, we consider that ϵ is the probability that a player makes a mistake in the action taken.

5 Appendix: Proofs

5.1 Proof of Lemma 1: Akin's lemma

Proof. The proof is based on a similar argument as the proof of Eq. (18), showing that different ways of calculating payoffs are equivalent. Let us first introduce some notation. Let \mathbf{m}^1 be the memory- n strategy of player 1. For $t \geq n$ and the given strategy of player 2, let $\mathbf{v}(t) = (v_{\mathbf{h}})_{\mathbf{h} \in H}$ be the probability that player 1 observes the n -history \mathbf{h} after players have made their t -th decision. By assumption, we can compute the limiting distribution

$$\mathbf{v} = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} \mathbf{v}(t). \quad (44)$$

Moreover, let $\rho^i(t)$ be player i 's cooperation probability in round t . For $t \geq n+1$, we obtain

$$\rho^1(t) = \langle \mathbf{v}(t-1), \mathbf{m}^1 \rangle = \langle \mathbf{v}(t+k-1), \mathbf{m}^{\mathbf{k-Rep}} \rangle. \quad (45)$$

That is, we either need to know how likely each n -history occurred at time $t-1$, and then we compute how likely player 1 is to cooperate in the next round, based on player 1's strategy. Or, we need to know how likely each n -history occurred after round $t+k-1$; and then we compute the correct probability by assuming player 1 cooperates in the next round if and only if the player cooperated k rounds before. Eq. (45) gives us two different ways to compute player 1's average payoff across all rounds,

$$\rho^1 := \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=1}^{\tau} \rho^1(t). \quad (46)$$

The first way is to take

$$\begin{aligned} \rho^1 &= \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=1}^{\tau} \rho^1(t) = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n+1}^{n+\tau} \rho^1(t) \\ &= \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n+1}^{n+\tau} \langle \mathbf{v}(t-1), \mathbf{m}^1 \rangle = \left\langle \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n+1}^{n+\tau} \mathbf{v}(t-1), \mathbf{m}^1 \right\rangle = \langle \mathbf{v}, \mathbf{m}^1 \rangle. \end{aligned}$$

In particular, because $\langle \mathbf{v}, \mathbf{m}^1 \rangle$ is well-defined, so is the limiting time average ρ^1 . The second way is to take

$$\begin{aligned} \rho^1 &= \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=1}^{\tau} \rho^1(t) = \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n+1}^{n+\tau} \rho^1(t) \\ &= \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n+1}^{n+\tau} \langle \mathbf{v}(t+k-1), \mathbf{m}^{\mathbf{k-Rep}} \rangle = \left\langle \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n+1}^{n+\tau} \mathbf{v}(t+k-1), \mathbf{m}^{\mathbf{k-Rep}} \right\rangle = \langle \mathbf{v}, \mathbf{m}^{\mathbf{k-Rep}} \rangle. \end{aligned}$$

We conclude $0 = \rho^1 - \rho^1 = \langle \mathbf{v}, \mathbf{m}^1 \rangle - \langle \mathbf{v}, \mathbf{m}^{\mathbf{k-Rep}} \rangle = \langle \mathbf{v}, \mathbf{m}^1 - \mathbf{m}^{\mathbf{k-Rep}} \rangle$. \square

5.2 Proof of Lemma 2: Sufficiency of testing self-reactive strategies

Proof. The proof uses similar arguments as in a study by Park on alternating games *et al* [12]. For the given game between player 1 (with arbitrary strategy σ^1) and player 2 (with reactive- n strategy \mathbf{p}^2), let $v_{\mathbf{h}}(t)$ denote the probability to observe an n -history \mathbf{h} at time $t \geq n$. By assumption, the following time averages are well-defined,

$$v_{\mathbf{h}} := \lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} v_{\mathbf{h}}(t) \quad (47)$$

Moreover, for any $t \geq n$ and $\mathbf{h} \in H$, let $\sigma_{\mathbf{h}}^1(t)$ denote the conditional probability that player 1 cooperates at time $t+1$, given the n -history after round t is \mathbf{h} . Depending on $(\sigma_{\mathbf{h}}^1(t))$ and \mathbf{v} , we define an associated self-reactive strategy $\tilde{\mathbf{p}}^1$ for player 1. For any given history $\mathbf{h}^1 \in H^1$, the corresponding probability $\tilde{p}_{\mathbf{h}^1}^1$ is defined as an implicit solution of the equation

$$\left(\sum_{\mathbf{h}^2 \in H^2} v_{(\mathbf{h}^1, \mathbf{h}^2)} \right) \tilde{p}_{\mathbf{h}^1}^1 = \sum_{\mathbf{h}^2 \in H^2} \left(\lim_{\tau \rightarrow \infty} \frac{1}{\tau} \sum_{t=n}^{n+\tau-1} v_{(\mathbf{h}^1, \mathbf{h}^2)}(t) \cdot \sigma_{(\mathbf{h}^1, \mathbf{h}^2)}^1(t) \right). \quad (48)$$

Note that for each $\mathbf{h} \in H$, the limit in the bracket on the right hand side exists, for otherwise the limits $v_{\mathbf{h}}$ according to Eq. (47) would not exist. Also note that if the bracket on the left hand's side is zero, the right hand side must be zero, and $\tilde{p}_{\mathbf{h}^1}^1$ can be chosen arbitrarily. Only if the bracket on the left hand side is positive, $\tilde{p}_{\mathbf{h}^1}^1$ is uniquely defined.

We are going to show: If player 1 uses $\tilde{\mathbf{p}}^1$ instead of σ^1 , then \mathbf{v} defined by Eq. (47) is an invariant distribution of the corresponding transition matrix M defined by Eq. (12) (hence it is also the limiting distribution of the resulting game if the first n moves are chosen accordingly). For simplicity, we show the required relationship $\mathbf{v} = \mathbf{v}M$ for one of the 2^{2n} equations. For the one equation we show, we consider the history according to which everyone fully cooperates, $\mathbf{h}_C = (\mathbf{h}_C^1, \mathbf{h}_C^2)$. For an arbitrary n -history $\mathbf{h}^i = (a_{-n}^i, \dots, a_{-1}^i)$, we say the n -history $\tilde{\mathbf{h}}^i = (\tilde{a}_{-n}^i, \dots, \tilde{a}_{-1}^i)$ is a *possible successor* of \mathbf{h} if $\tilde{a}_{-t}^i = a_{-t+1}^i$ for $t \in \{2, \dots, n\}$. To indicate successorship, we define a function $e_{\mathbf{h}, \tilde{\mathbf{h}}}$ that is one if $\tilde{\mathbf{h}}$ is a possible successor of \mathbf{h} , and zero otherwise. By definition of $v_{\mathbf{h}}(t)$, $\sigma_{\mathbf{h}}^1(t)$, and $p_{\mathbf{h}}^2(t)$, we obtain for $t \geq n$

$$v_{(\mathbf{h}_C^1, \mathbf{h}_C^2)}(t+1) = \sum_{\mathbf{h}^1 \in H^1} \sum_{\mathbf{h}^2 \in H^2} v_{(\mathbf{h}^1, \mathbf{h}^2)}(t) \cdot \sigma_{(\mathbf{h}^1, \mathbf{h}^2)}^1(t) \cdot p_{\mathbf{h}^1}^2 \cdot e_{\mathbf{h}^1, \mathbf{h}_C^1} \cdot e_{\mathbf{h}^2, \mathbf{h}_C^2}. \quad (49)$$

If we sum up this equation from time $t=n$ to $t=n+\tau-1$, divide by τ , and rearrange the terms, we obtain

$$\frac{1}{\tau} \sum_{t=n}^{n+\tau-1} v_{(\mathbf{h}_C^1, \mathbf{h}_C^2)}(t+1) = \sum_{\mathbf{h}^1 \in H^1} \sum_{\mathbf{h}^2 \in H^2} \left(\frac{1}{\tau} \sum_{t=n}^{n+\tau-1} v_{(\mathbf{h}^1, \mathbf{h}^2)}(t) \cdot \sigma_{(\mathbf{h}^1, \mathbf{h}^2)}^1(t) \right) \cdot p_{\mathbf{h}^1}^2 \cdot e_{\mathbf{h}^1, \mathbf{h}_C^1} \cdot e_{\mathbf{h}^2, \mathbf{h}_C^2}. \quad (50)$$

Taking the limit $\tau \rightarrow \infty$, and taking into account the relationships (47) and (48), this simplifies to

$$v_{(\mathbf{h}_C^1, \mathbf{h}_C^2)} = \sum_{\mathbf{h}^1 \in H^1} \sum_{\mathbf{h}^2 \in H^2} v_{(\mathbf{h}^1, \mathbf{h}^2)} \cdot (\tilde{p}_{\mathbf{h}^1}^1 e_{\mathbf{h}^1, \mathbf{h}_C^1}) \cdot (p_{\mathbf{h}^1}^2 e_{\mathbf{h}^2, \mathbf{h}_C^2}). \quad (51)$$

By using the definition of transition probabilities in (12), this expression further simplifies to

$$v_{\mathbf{h}_C} = \sum_{\mathbf{h}} v_{\mathbf{h}} \cdot M_{\mathbf{h}, \mathbf{h}_C} \quad (52)$$

That is, out of the 2^{2n} individual equations in the linear system $\mathbf{v} = \mathbf{v}M$, we have verified the equation for the probability to observe full cooperation \mathbf{h}_C after one round. All other equations follow analogously. \square

5.3 Proof of Theorem 1: Sufficiency of pure self-reactive strategies

By Lemma 2, there exists a best response to \mathbf{p} within the self-reactive n strategies. It remains to show that this best response $\tilde{\mathbf{p}}$ can be chosen to be pure. The proof follows from a series of auxiliary results. The first such result uses an insight by Press & Dyson [9]. They showed that given the transition matrix of a game among two memory-1 players, one can compute the players' payoffs by considering determinants of certain associated matrices. Herein, we apply their method to the transition matrix $\tilde{M} = (\tilde{M}_{\mathbf{h}, \mathbf{h}'})$ according to Eq. (26) for a given self-reactive strategy $\tilde{\mathbf{p}} \in S_n$. For some fixed n -history \mathbf{h}' , we define an associated matrix $\tilde{M}_{\mathbf{h}'}$ that one obtains from \tilde{M} with the following two steps:

1. Subtract the $2^n \times 2^n$ identity matrix I from \tilde{M} .
2. In the resulting matrix, replace the last column by a column that only contains zeros, except for the row corresponding to the history \mathbf{h}' , for which the entry is one.

These matrices $\tilde{M}_{\mathbf{h}'}$ can be used to compute the invariant distribution of the original matrix \tilde{M} as follows.

Auxiliary result 1: Let $\tilde{\mathbf{p}} \in S_n$ be such that its transition matrix \tilde{M} according to Eq. (26) has a unique invariant distribution $\tilde{\mathbf{v}} = (\tilde{v}_{\mathbf{h}^1})_{\mathbf{h}^1 \in H^1}$. Then for all $\mathbf{h}' \in H^1$ we have

$$\tilde{v}_{\mathbf{h}'} = \frac{\det(\tilde{M}_{\mathbf{h}'})}{\sum_{\mathbf{h}^1 \in H^1} \det(\tilde{M}_{\mathbf{h}^1})}. \quad (53)$$

Proof of Auxiliary result 1. The result follows from Press & Dyson's formula for the dot product of the invariant distribution $\tilde{\mathbf{v}}$ with an arbitrary vector \mathbf{f} , by taking the vector \mathbf{f} to be the unit vector with only the entry for history \mathbf{h}' being one. \square

Based on this first auxiliary result, we have an explicit representation of the payoff function $\pi^1(\tilde{\mathbf{p}}, \mathbf{p})$ that describes the payoff of a self-reactive player with strategy $\tilde{\mathbf{p}}$ against a reactive player with strategy \mathbf{p} . Specifically, by plugging Eq. (53) into (27), we obtain

$$\pi^1(\tilde{\mathbf{p}}, \mathbf{p}) = \frac{\sum_{\mathbf{h}^1 \in H^1} \det(\tilde{M}_{\mathbf{h}^1}) \left(\tilde{\mathbf{p}}_{\mathbf{h}^1} \mathbf{p}_{\mathbf{h}^1} \cdot R + \tilde{\mathbf{p}}_{\mathbf{h}^1} (1 - \mathbf{p}_{\mathbf{h}^1}) \cdot S + (1 - \tilde{\mathbf{p}}_{\mathbf{h}^1}) \mathbf{p}_{\mathbf{h}^1} \cdot T + (1 - \tilde{\mathbf{p}}_{\mathbf{h}^1}) (1 - \mathbf{p}_{\mathbf{h}^1}) \cdot P \right)}{\sum_{\mathbf{h}^1 \in H^1} \det(\tilde{M}_{\mathbf{h}^1})}. \quad (54)$$

For our purposes, the following properties of this payoff function will be important.

Auxiliary Result 2: On its domain, the payoff function $\pi^1(\tilde{\mathbf{p}}, \mathbf{p})$ is a bounded rational function, and both its numerator and denominator are linear in each entry $\tilde{p}_{\mathbf{h}^i}$, for all $\mathbf{h}^i \in H^i$.

Proof of Auxiliary Result 2. By its definition, each $\det(\tilde{M}_{\mathbf{h}^i})$ is a polynomial. Moreover, because for each history \mathbf{h}' , the cooperation probability $\tilde{p}_{\mathbf{h}'}$ only appears in a single row of $\tilde{M}_{\mathbf{h}^i}$ (and there it appears linearly), it also appears linearly in $\det(\tilde{M}_{\mathbf{h}^i})$. Finally, we note that $\det(\tilde{M}_{\mathbf{h}^i})$ does not depend on $\tilde{p}_{\mathbf{h}^i}$. To see this, we can compute $\det(\tilde{M}_{\mathbf{h}^i})$ using Laplace expansion along the last column. As a result, we obtain that this determinant is up to its sign equal to the determinant of the matrix one obtains from $\tilde{M}_{\mathbf{h}^i}$ by deleting the last column, and the row \mathbf{h}^i (which is the only row of $\tilde{M}_{\mathbf{h}^i}$ that contains $\tilde{p}_{\mathbf{h}^i}$).

Finally, we note that the payoff function is bounded, because as an average payoff per round, payoffs need to be between T and S . Taken together, these observations imply the result for $\pi^1(\tilde{\mathbf{p}}, \mathbf{p})$. \square

The following result describes a useful property of bounded linear rational functions.

Auxiliary Result 3: Suppose $g, h : [0, 1]^k \rightarrow \mathbb{R}$ and suppose both $g(\mathbf{x})$ and $h(\mathbf{x})$ are linear in each component of $\mathbf{x} = (x_1, \dots, x_k)$. Moreover, suppose $f := g/h$ is bounded on $[0, 1]^k$. For a given \mathbf{x} and $j \in \{1, \dots, k\}$, we define an associated function $f_{\mathbf{x},j} : [-x_j, 1 - x_j] \rightarrow \mathbb{R}$ by only varying the j -th component, $f_{\mathbf{x},j}(t) = f(x_1, \dots, x_j + t, \dots, x_k)$. Then for all $\mathbf{x} \in [0, 1]^k$ and j , the function $f_{\mathbf{x},j}(t)$ is either monotonically increasing, monotonically decreasing, or constant.

Proof of Auxiliary Result 3. Let $g(\mathbf{x}) := a_0 + a_1x_1 + \dots + a_kx_k$ and $h(\mathbf{x}) := b_0 + b_1x_1 + \dots + b_kx_k$, and consider some arbitrary but fixed $\mathbf{x} \in [0, 1]^k$ and j . We compute

$$f'_{\mathbf{x},j}(t) = \frac{\partial}{\partial t} f(x_1, \dots, x_j + t, \dots, x_k) = \frac{a_j \left(\sum_{i \neq j} b_i x_i \right) - b_j \left(\sum_{i \neq j} a_i x_i \right)}{(b_0 + b_1x_1 + \dots + b_j(x_j + t) + \dots + b_kx_k)^2}. \quad (55)$$

Because f is bounded on the entire domain, the denominator in this expression for $f'_{\mathbf{x},j}(t)$ is strictly positive. Moreover, we note that the numerator is independent of t . Thus, depending on the sign of the numerator, $f'_{\mathbf{x},j}(t)$ is either monotonically increasing, monotonically decreasing, or constant. \square

After these preparations, we are ready to prove the main result.

Proof of Theorem 1. For a given reactive strategy $\mathbf{p} \in \mathcal{R}_n$, let the self-reactive $\tilde{\mathbf{p}} \in \mathcal{S}_n$ be a best response. Suppose there is some history \mathbf{h}' such that $0 < \tilde{p}_{\mathbf{h}'} < 1$. It follows from the Auxiliary Results 2 and 3 that $\pi^1(\tilde{\mathbf{p}}, \mathbf{p})$ is either monotonically increasing, monotonically decreasing, or constant in $\tilde{p}_{\mathbf{h}'}$. If it was increasing or decreasing, we end up with a contradiction, because no local improvement should be possible for a best response. Therefore, $\pi^1(\tilde{\mathbf{p}}, \mathbf{p})$ must be independent of $\tilde{p}_{\mathbf{h}'}$, and hence we can set $\tilde{p}_{\mathbf{h}'} = 0$ or $\tilde{p}_{\mathbf{h}'} = 1$ without changing $\pi^1(\tilde{\mathbf{p}}, \mathbf{p})$. By iteratively applying this reasoning to all histories \mathbf{h} for which $0 < \tilde{p}_{\mathbf{h}} < 1$, we obtain the desired result. \square

5.4 Proof of Theorem 2: Reactive-2 partner strategies in the donation game

Proof. Given that player 1 uses a nice reactive-2 strategy $\mathbf{p} = (1, p_{CD}, p_{DC}, p_{DD})$, the claim is true if and only if it is true for all deviation towards the sixteen pure self-reactive-2 strategies $\tilde{\mathbf{p}} \in \{0, 1\}^{16}$. In the following, we enumerate these sixteen strategies, $\{\tilde{\mathbf{p}}_0, \dots, \tilde{\mathbf{p}}_{15}\}$, by interpreting them as binary numbers,

$$\tilde{\mathbf{p}} = (\tilde{p}_{CC}, \tilde{p}_{CD}, \tilde{p}_{DC}, \tilde{p}_{DD}) \mapsto \tilde{p}_{CC} \cdot 2^3 + \tilde{p}_{CD} \cdot 2^2 + \tilde{p}_{DC} \cdot 2^1 + \tilde{p}_{DD} \cdot 2^0. \quad (56)$$

In particular, ALLD = (0, 0, 0, 0) is mapped to the number $j=0$, and ALLC = (1, 1, 1, 1) is mapped to $j=15$. The possible payoffs against the reactive strategy \mathbf{p} can be computed by Eq. (27), which yields

$$\begin{aligned} \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= p_{DD} \cdot b && \text{for } j \in \{0, 2, 4, 6, 8, 10, 12, 14\} \\ \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CD} + p_{DC} + p_{DD}}{3} \cdot b - \frac{1}{3} \cdot c && \text{for } j \in \{1, 9\} \\ \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CD} + p_{DC} + p_{DD}}{4} \cdot b - \frac{1}{2} \cdot c && \text{for } j \in \{3\} \\ \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CD} + p_{DC}}{2} \cdot b - \frac{1}{2} \cdot c && \text{for } j \in \{4, 5, 12, 13\} \\ \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CD} + p_{DC}}{3} \cdot b - \frac{2}{3} \cdot c && \text{for } j \in \{6, 7\} \\ \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= b - c && \text{for } j \in \{8, 9, 10, 11, 12, 13, 14, 15\} \end{aligned}$$

In this list, some strategy indices j appear multiple times. Those instances correspond to strategies that have multiple invariant distributions (such as the strategy 1-round repeat, with $j=10$). For those strategies, we have computed the payoffs for all possible initial n -histories. Requiring the payoffs in this list to be at most the mutual cooperation payoff $b-c$, we get the following unique conditions,

$$p_{DD} \leq 1 - \frac{c}{b}, \quad \frac{p_{CD} + p_{DC}}{2} \leq 1 - \frac{1}{2} \frac{c}{b}, \quad \frac{p_{CD} + p_{DC} + p_{DD}}{3} \leq 1 - \frac{2}{3} \frac{c}{b}.$$

Because the last condition is implied by the first two, we end up with the conditions in (34). \square

5.5 Proof of Theorem 3: Reactive-3 partner strategies in the donation game

Proof. The proof is similar to the previous one. Again, enumerating the 256 pure self-reactive 3 strategies $\tilde{\mathbf{p}}$ by interpreting the strategy as a binary number, we obtain the following payoffs.

$$\begin{aligned}
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= b \, p_{DDD} && \text{for } j \in \{0, 2, 4, 6, \dots, 250, 252, 254\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CDD} + p_{DCD} + p_{DDC} + p_{DDD}}{4} b - \frac{1}{4} c && \text{for } j \in \{1, 9, 33, 41, 65, 73, 97, 105, 129, 137, 161, \\
&&& 169, 193, 201, 225, 233\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC} + p_{DDD}}{5} b - \frac{2}{5} c && \text{for } j \in \{3, 7, 35, 39, 131, 135, 163, 167\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CDC} + p_{DCD}}{2} b - \frac{1}{2} c && \text{for } j \in \{4-7, 12-15, 20-23, 28-31, 68-71, \\
&&& 76-79, 84-87, 92-95, 132-135, \\
&&& 140-143, 148-151, 156-159, \\
&&& 196-199, 204-207, 212-215, 220-223\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC} + p_{DDD}}{6} b - \frac{1}{2} c && \text{for } j \in \{11, 15, 43, 47\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CDD} + p_{DCD} + p_{DDC}}{3} b - \frac{1}{3} c && \text{for } j \in \{16, 17, 24, 25, 48, 49, 56, 57, 80, 81, 88, \\
&&& 89, 112, 113, 120, 121, 144, 145, 152, 153, \\
&&& 176, 177, 184, 185, 208, 209, 216, 217, \\
&&& 240, 241, 248, 249\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}}{4} b - \frac{1}{2} c && \text{for } j \in \{18, 19, 22, 23, 50, 51, 54, 55, 146, 147, \\
&&& 150, 151, 178, 179, 182, 183\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}}{5} b - \frac{3}{5} c && \text{for } j \in \{26, 27, 30, 31, 58, 59, 62, 63\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CCD} + p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD} + p_{DDC} + p_{DDD}}{7} b - \frac{3}{7} c && \text{for } j \in \{37, 67, 165, 195\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CCD} + p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD} + p_{DDC} + p_{DDD}}{8} b - \frac{1}{2} c && \text{for } j \in \{45, 75\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CCD} + p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD} + p_{DDC}}{6} b - \frac{1}{2} c && \text{for } j \in \{52, 53, 82, 83, 180, 181, 210, 211\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CCD} + p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD} + p_{DDC}}{7} b - \frac{4}{7} c && \text{for } j \in \{60, 61, 90, 91\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{p_{CCD} + p_{CDC} + p_{DCC}}{3} b - \frac{2}{3} c && \text{for } j \in \{96-103, 112-119, 224-231, 240-247\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{1 + p_{CCD} + p_{CDC} + p_{DCC}}{4} b - \frac{3}{4} c && \text{for } j \in \{104-111, 120-127\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= b - c && \text{for } j \in \{128, 129, 130, \dots, 255\}
\end{aligned}$$

Requiring these payoffs to be at most equal to the mutual cooperation payoff $b - c$ gives

$$\begin{aligned}
p_{DDD} &\leq 1 - \frac{c}{b}, & \frac{p_{CDC} + p_{DCD}}{2} &\leq 1 - \frac{1}{2} \cdot \frac{c}{b}, & \frac{p_{CDD} + p_{DCD} + p_{DDC}}{3} &\leq 1 - \frac{2}{3} \cdot \frac{c}{b}, \\
\frac{p_{CCD} + p_{CDC} + p_{DCC}}{3} &\leq 1 - \frac{1}{3} \cdot \frac{c}{b}, & \frac{p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC}}{4} &\leq 1 - \frac{1}{2} \cdot \frac{c}{b}, \\
\frac{p_{CDD} + p_{DCD} + p_{DDC} + p_{DDD}}{4} &\leq 1 - \frac{3}{4} \cdot \frac{c}{b}, & \frac{p_{CCD} + p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD} + p_{DDC} + p_{DDD}}{7} &\leq 1 - \frac{4}{7} \cdot \frac{c}{b}, \\
\frac{p_{CCD} + p_{CDD} + p_{DCC} + p_{DDC} + p_{DDD}}{5} &\leq 1 - \frac{3}{5} \cdot \frac{c}{b}, & \frac{p_{CCD} + p_{CDC} + p_{CDD} + p_{DCC} + p_{DCD} + p_{DDC}}{6} &\leq 1 - \frac{1}{2} \cdot \frac{c}{b}.
\end{aligned}$$

The statement follows by noting that the five conditions in the first two rows imply the four other conditions.

□

5.6 Proof of Theorem 4: Reactive- n counting strategies in the donation game

Before we go into the details of the proof, we first start with two useful observations.

1. Assume player 1 adopts a given self-reactive strategy $\tilde{\mathbf{p}}$ and player 2 adopts the reactive- n strategy $\mathbf{r} = (r_k)_{k \in \{n, \dots, 0\}}$. For the resulting game, suppose \mathbf{v} is the limiting distribution according to Eq. (14). Then it is useful to express \mathbf{v} in terms of what the counting player can remember. To this end, let H_k^1 be the set of n -histories according to which player 1 has cooperated exactly k times,

$$H_k^1 = \left\{ \mathbf{h}^1 \in H^1 \mid |\mathbf{h}^1| = k \right\}. \quad (57)$$

Accordingly, let $\mathbf{u} = (u_k)_{k \in \{0, \dots, n\}}$ be the distribution that summarizes how often, on average, player 1 cooperates j times during n consecutive rounds,

$$u_k^1 = \sum_{\mathbf{h}^1 \in H_k^1} v_{\mathbf{h}^1}. \quad (58)$$

In particular, the entries of \mathbf{u} are normalized,

$$\sum_{k=0}^n u_k^1 = 1. \quad (59)$$

Moreover, the average cooperation rate of the two players can be written as

$$\rho^1 = \sum_{k=0}^n \frac{k}{n} u_k^1 \quad \text{and} \quad \rho^2 = \sum_{k=0}^n r_k u_k^1. \quad (60)$$

Because payoffs in the donation game only depend on the players' average cooperation rates (but not on the timing of cooperation), we conclude that player 1's payoff is

$$\pi^1(\tilde{\mathbf{p}}, \mathbf{r}) = \sum_{k=0}^n \left(r_k b - \frac{k}{n} c \right) u_k^1. \quad (61)$$

2. There is a set of strategies for which payoffs are particularly easy to compute. We refer to them as simple periodic strategies, σ_k with $k \in \{0, \dots, n\}$. A player with strategy σ_k cooperates in round t if and only if

$$t - 1 \bmod n < k. \quad (62)$$

That is, such a player cooperates in the first k rounds, then defects for $n - k$ rounds, then cooperates for another k rounds, only to defect in the $n - k$ subsequent rounds, etc. Such strategies are interesting for two reasons. First, they all can be interpreted as a round- n repeat strategy $\tilde{\mathbf{p}}^{\text{n-Rep}}$, as defined by (11). During the initial n rounds, they cooperate according to Eq. (62); thereafter, they simply

repeat whatever they have done n rounds ago. Second, players with strategy σ_k always act in such a way that according to any resulting n -history, they have cooperated exactly k times during the last n rounds. As a result, if player 1 adopts such a strategy in a donation game against a player with a reactive- n counting strategy \mathbf{r} , then player 1's average payoff is

$$\pi^1(\sigma_k, \mathbf{r}) = r_k b - \frac{k}{n} c. \quad (63)$$

After these observations, we are ready for the actual proof.

Proof of Theorem 4.

(\Rightarrow) Suppose the reactive- n counting strategy \mathbf{r} is a partner. Because it is nice, it cooperates against an unconditional cooperator, and hence $r_n = 1$. Because it is a Nash equilibrium, player 1 must not have an incentive to deviate towards any of the simple periodic strategies σ_k . By Eq. (63), this means that for all $k \in \{0, \dots, n\}$ we have

$$r_k b - \frac{k}{n} c \leq b - c. \quad (64)$$

These conditions are equivalent to $r_{n-k} \leq 1 - \frac{k}{n} \frac{c}{b}$, the inequalities in (39).

(\Leftarrow) Because \mathbf{r} is nice, $r_n = 1$. The proof is now by contradiction; suppose the conditions in (39) hold, but \mathbf{r} is not a Nash equilibrium. Then there needs to be some self-reactive $\tilde{\mathbf{p}}$ such that $\pi^1(\tilde{\mathbf{p}}, \mathbf{r}) > b - c$. It follows that

$$\begin{aligned} 0 &< \pi^1(\tilde{\mathbf{p}}, \mathbf{r}) - (b - c) \\ &\stackrel{(59), (61)}{=} \sum_{k=0}^n \left(r_k b - \frac{k}{n} c \right) u_k^1 - \sum_{k=0}^n (b - c) u_k^1 \\ &= (r_n - 1) b u_n + \sum_{k=0}^{n-1} \left((r_k - 1) b + \frac{n - k}{n} c \right) u_k^1 \\ &= b \cdot \sum_{k=1}^n \underbrace{\left(r_{n-k} - \left(1 - \frac{k}{n} \frac{c}{b} \right) \right)}_{\leq 0 \text{ by Eq. (39)}} u_{n-k}^1 \leq 0. \end{aligned} \quad (65)$$

We end up with $0 < 0$, a contradiction.

□

5.7 Proof of Theorem 5: Reactive-2 partner strategies in the prisoner's dilemma

Proof. The proof is analogous to the proof of Theorem 2 for the donation game. For the general prisoner's dilemma, the payoffs of the 16 pure self-reactive-2 strategies are

$$\begin{aligned}
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= P(1-p_{DD}) + Tp_{DD} && \text{for } i \in \{0, 2, 4, 6, 8, 10, 12, 14\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{Rp_{DD} + S(1-p_{DD}) + T(p_{CD} + p_{DC}) + P(2-p_{CD} - p_{DC})}{3} && \text{for } i \in \{1, 9\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{R(p_{DC} + p_{DD}) + S(2-p_{DC} - p_{DD}) + T(p_{CD} + 1) + P(1-p_{CD})}{4} && \text{for } i \in \{3\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{Rp_{CD} + S(1-p_{CD}) + Tp_{DC} + P(1-p_{DC})}{2} && \text{for } i \in \{4, 5, 12, 13\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{R(p_{CD} + p_{DC}) + S(2-p_{CD} - p_{DC}) + T}{3} && \text{for } i \in \{6, 7\} \\
\pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= R && \text{for } i \in \{8, 9, 10, 11, 12, 13, 14, 15\}
\end{aligned}$$

By requiring these expressions to be at most equal to R , we obtain

$$\begin{aligned}
(T - P)p_{DD} &\leq R - P, \\
(R - S)(p_{CD} + p_{DC}) &\leq 3R - 2S - T, \\
(T - P)p_{DC} + (R - S)p_{CD} &\leq 2R - S - P, \\
(T - P)(p_{CD} + p_{DC}) + (R - S)p_{DD} &\leq 3R - S - 2P, \\
(T - P)p_{CD} + (R - S)(p_{CD} + p_{DD}) &\leq 4R - 2S - P - T.
\end{aligned}$$

□

5.8 Proof of Theorem 6: Reactive-3 partner strategies in the prisoner's dilemma

Again, we compute payoffs for all 256 self-reactive-3 strategies. The expressions are given below,

$$\begin{aligned}
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCDD + PDCD + PDDC) + 3P + (R-S)PDDD + S}{4} & \text{for } j \in \{1, 9, 33, 41, 65, 73, 97, 105, \\
 & & 129, 137, 161, 169, 193, 201, \\
 & & 225, 233\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)PCDC + P + (R-S)PDCD + S}{2} & \text{for } j \in \{4-7, 12-15, 20-23, \\
 & & 28-31, 68-71, 76-79, \\
 & & 84-87, 92-95, 132-135, \\
 & & 140-143, 148-151, 156-159, \\
 & & 196-199, 204-207, 212-215, \\
 & & 220-223\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= -P(PDDD - 1) + TPDDD & \text{for } j \in \{0, 2, 4, \dots, 252, 254\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PCDD + PDDC) + 3P + (R-S)(PCDC + PDCC + PDCD + PDDD) + 4S + T}{8} & \text{for } j \in \{45\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)PDCC + P + (R-S)(PCDC + PCCD) + 2S}{3} & \text{for } j \in \{96-103, 112-119, \\
 & & 224-231, 240-247\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PDCC + PDDC) + 3P + (R-S)(PCDC + PCDD + PDCD) + 3S}{6} & \text{for } j \in \{52, 53, 180, 181\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PDDC) + 2P + T + (R-S)(PCDC + PCDD + PDCC + PDCD) + 4S}{7} & \text{for } j \in \{60, 61\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PCDD + PDCC) + 3P + (R-S)(PDDC + PDDD) + 2S}{5} & \text{for } j \in \{3, 7, 35, 39, 131, 135, 163, 167\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PDCD + PDDC) + 2P + (R-S)PCDD + S}{3} & \text{for } j \in \{16, 17, 24, 25, 48, 49, 56, \\
 & & 57, 80, 81, 88, 89, 112, 113, \\
 & & 120, 121, 144, 145, 152, 153, \\
 & & 176, 177, 184, 185, 208, 209, \\
 & & 216, 217, 240, 241, 248, 249\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= R & \text{for } j \in \{128, 129, \dots, 255\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)PCCD + P + T + (R-S)(PCDD + PDCC + PDDC) + 3S}{5} & \text{for } j \in \{26, 27, 30, 31, 58, 59, 62, 63\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PDCC) + 2P + (R-S)(PCDD + PDDC) + 2S}{4} & \text{for } j \in \{18, 19, 22, 23, 50, 51, 54, 55, \\
 & & 146, 147, 150, 151, 178, 179, \\
 & & 182, 183\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCDC + PDCD) + 2P + T + (R-S)(PCCD + PCDD + PDCC + PDDC) + 4S}{7} & \text{for } j \in \{90, 91\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCDC + PCDD + PDCD) + 3P + T + (R-S)(PCCD + PDCC + PDDC + PDDD) + 4S}{8} & \text{for } j \in \{75\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCDC + PDCC + PDCD) + 3P + (R-S)(PCCD + PCDD + PDDC) + 3S}{6} & \text{for } j \in \{82, 83, 210, 211\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PCDD + PDCC + PDDC) + 4P + (R-S)(PCDC + PDCD + PDDD) + 3S}{7} & \text{for } j \in \{37, 165\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{T + (R-S)(PCCD + PCDC + PDCC) + 3S}{4} & \text{for } j \in \{104-111, 120-127\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCCD + PCDD) + 2P + T + (R-S)(PDCC + PDDC + PDDD) + 3S}{6} & \text{for } j \in \{11, 15, 43, 47\} \\
 \pi^1(\tilde{\mathbf{p}}_j, \mathbf{p}) &= \frac{(T-P)(PCDC + PCDD + PDCC + PDCD) + 4P + (R-S)(PCCD + PDDC + PDDD) + 3S}{7} & \text{for } j \in \{67, 195\}
 \end{aligned}$$

By requiring the above expressions to be smaller than or equal to R , we obtain the inequalities in Table 1.

6 Supplementary References

- [1] Nowak, M. & Sigmund, K. A strategy of win-stay, lose-shift that outperforms tit-for-tat in the prisoner's dilemma game. *Nature* **364**, 56–58 (1993).
- [2] Hilbe, C., Chatterjee, K. & Nowak, M. A. Partners and rivals in direct reciprocity. *Nature human behaviour* **2**, 469–477 (2018).
- [3] Axelrod, R. & Hamilton, W. D. The evolution of cooperation. *science* **211**, 1390–1396 (1981).
- [4] Nowak, M. A. & Sigmund, K. Tit for tat in heterogeneous populations. *Nature* **355**, 250–253 (1992).
- [5] Molander, P. The optimal level of generosity in a selfish, uncertain environment. *Journal of Conflict Resolution* **29**, 611–618 (1985).
- [6] Akin, E. The iterated prisoner's dilemma: good strategies and their dynamics. *Ergodic Theory, Advances in Dynamical Systems* 77–107 (2016).
- [7] Hilbe, C., Wu, B., Traulsen, A. & Nowak, M. A. Cooperation and control in multiplayer social dilemmas. *Proceedings of the National Academy of Sciences USA* **111**, 16425–16430 (2014).
- [8] Hilbe, C., Traulsen, A. & Sigmund, K. Partners or rivals? strategies for the iterated prisoner's dilemma. *Games and economic behavior* **92**, 41–52 (2015).
- [9] Press, W. H. & Dyson, F. J. Iterated prisoner's dilemma contains strategies that dominate any evolutionary opponent. *Proceedings of the National Academy of Sciences* **109**, 10409–10413 (2012).
- [10] Stewart, A. J. & Plotkin, J. B. Small groups and long memories promote cooperation. *Scientific reports* **6**, 1–11 (2016).
- [11] Ueda, M. Memory-two zero-determinant strategies in repeated games. *Royal Society open science* **8**, 202186 (2021).
- [12] Park, P. S., Nowak, M. A. & Hilbe, C. Cooperation in alternating interactions with memory constraints. *Nature Communications* **13**, 737 (2022).