

Properties of Winning Iterated Prisoner’s Dilemma Strategies.

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Abstract

Researchers have explored the performance of Iterated Prisoner’s Dilemma strategies for decades, from the celebrated performance of Tit for Tat to the introduction of the zero-determinant strategies and the use of sophisticated learning structures such as neural networks. Many new strategies have been introduced and tested in a variety of tournaments and population dynamics. Typical results in the literature, however, rely on performance against a small number of somewhat arbitrarily selected strategies, casting doubt on the generalizability of conclusions. In this work, we analyze a large collection of 195 strategies in thousands of computer tournaments, present the top performing strategies across multiple tournament types, and distill their salient features. The results show that there is not yet a single strategy that performs well in diverse Iterated Prisoner’s Dilemma scenarios, nevertheless there are several properties that heavily influence the best performing strategies. This refines the properties described by Axelrod in light of recent and more diverse opponent populations to: be nice, be provokable and generous, be a little envious, be clever, and adapt to the environment. More precisely, we find that strategies perform best when their probability of cooperation matches the total tournament population’s aggregate cooperation probabilities. The features of high performing strategies help cast some light on why strategies such as Tit For Tat performed historically well in tournaments and why zero-determinant strategies typically do not fare well in tournament settings.

1 Introduction

The Iterated Prisoner’s Dilemma (IPD) is a repeated two-player game that models behavioral interactions, specifically interactions where self-interest clashes with collective interest. It encompasses a wide range of social and biological phenomena. In each turn of the game, both players simultaneously and independently decide between cooperation (C) and defection (D). This decision is made with the memory of all prior interactions. The payoffs for each player at each turn are influenced by their own choice and the choice of the other player. To this end, the payoffs of the game are defined by

	Cooperate (C)	Defect (D)
Cooperate (C)	R, R	S, T
Defect (D)	T, S	P, P

where typically $T > R > P > S$ and $2R > T + S$. The most common values used in the literature [1] are $R = 3$, $P = 1$, $T = 5$, $S = 0$, and these are the values also used in this work.

Conceptualizing strategies and understanding the best way to play the game have been of interest to the scientific community since the formulation of the game [2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13]. This extends to both tournament settings and population dynamics. Computer tournaments became a common evaluation technique for newly designed strategies following Axelrod’s computer tournaments in the 1980s [2, 14]. The winner of both of Axelrod’s tournaments [2, 14] was the simple strategy Tit For Tat (TFT). TFT cooperates on the first turn and thereafter copies the previous action of its opponent, retaliating against defections with a defection and forgiving a defection if followed by cooperation. Axelrod concluded that the strategy’s robustness was due to four properties, which he adapted into four suggestions for success in an IPD tournament:

- (i) Do not be envious by striving for a payoff larger than the opponent’s payoff.
- (ii) Be “nice” by not being the first to defect.
- (iii) Reciprocate both cooperation and defection; Be provokable to retaliation and forgiveness.
- (iv) Do not be too clever by scheming to exploit the opponent.

Forgiveness, in this context, is a strategy’s ability to cooperate after a *DC* outcome to achieve mutual cooperation again. In environments without noise, TFT would end up in *DC* only if it had received a defection and then retaliated. Subsequently, TFT would forgive an opponent that apologizes (in a *DC* round) by returning to cooperation, as mutual cooperation is deemed better than mutual defection.

Due to the strategy’s strong performance in both tournaments and a series of evolutionary experiments [1], TFT was often claimed to be a highly robust (and sometimes the most robust) strategy for the IPD. There are strategies that have built upon TFT and the reciprocity-based approach. In [5], a strategy called Gradual was introduced, constructed to have the same qualities as those of TFT with one addition. Gradual has a memory of the previous rounds of play in the game, recording the number of defections by the opponent and punishing them with a growing number of defections. It then enters a calming state in which it cooperates for two rounds. A strategy with the same intuition as Gradual is Adaptive Tit for Tat [15]. Adaptive Tit for Tat maintains a continually updated estimate of the opponent’s behavior and uses this estimate to condition its future actions. Other research has built upon the limitations of TFT. For example, in [16, 17, 18, 19], it was shown that TFT suffered in environments with noise. This was mainly due to the strategy being too provokable and its lack of generosity and contrition. Since TFT immediately punishes a defection, in a noisy environment, it can get stuck in a repeated cycle of defections and cooperations. Some new strategies, more robust in tournaments with noise, were soon introduced, including Nice and Forgiving [16], Generous Tit For Tat [3], and Pavlov (aka Win Stay Lose Shift) [4], as well as later variants such as OmegaTFT [20].

Finally, others introduced strategies deviating completely from the originally suggested properties of success. For example, a set of “envious” Iterated Prisoner’s Dilemma (IPD) strategies were introduced, called zero-determinant strategies (ZDs), in [6]. These strategies attempt to force a linear relationship between stationary payoffs against their opponents, potentially ensuring that they receive a higher average payout. While ZDs were introduced with a small tournament in which some were reportedly successful [21], this result has not generally held in future work [22]. Furthermore, in [23], a series of “clever” strategies trained using reinforcement learning were introduced. These strategies were trained using lookup tables [24], hidden Markov models [23], and finite-state automata [25], on a set of 170 strategies.

One thing that has remained the same is that the introduction of a new strategy is often accompanied by a claim that the new strategy is the best performing strategy for the IPD, often without extensive testing

against a broad spectrum of opponents or representative classes of opponents. The lack of testing against formally defined strategies and tournament winners is understandable given the effort required to implement the hundreds of published IPD strategies. Implementing prior strategies faithfully is often extremely difficult or impossible due to insufficient descriptions and lack of published implementations or code. Despite these challenges, the absence of thorough testing raises concerns about claims regarding the superiority or robustness of newly introduced strategies.

Beyond these difficulties, we believe that limited comprehensive analyses are rooted in field conventions. Tournaments or evolutionary dynamics often rely on a select list of hand-picked strategies chosen by modelers, typically based on specific properties they wish to examine. This practice may stem from misconceptions, such as the assumption that because TFT performs relatively well, it is sufficient to test only against TFT variants. Another misconception may stem from the Press & Dyson result [6], which implies that lower memory strategies will always dominate pairwise interactions, leading some to consider only memory-one strategies. However, this result holds strictly only in pairwise interactions.

It is not only the set of strategies or the tournament parameters such as noise that may impact results but also the design of the round-robin tournament itself. To address this [26] separately examined the effects of changes in format, objective criteria, and payoff values on tournament outcomes. They demonstrated that TFT’s performance declined under certain conditions. To our knowledge, this is the only study that has reanalyzed the tournament structure and critically evaluated it. However, in this work, we employ an extensive list of strategies made possible by the **Axelrod-Python** package, an approach that would have been difficult to achieve previously. Unlike the authors of that study, we do not consider a new tournament design.

In this paper, we evaluate the performance of a significant number of IPD strategies across a diverse array of tournaments. Many of the strategies used in our analysis are drawn from well-known and named strategies in IPD literature, including previous tournament winners. This contrasts with other work that may have randomly generated essentially arbitrary strategies, often constrained to specific classes such as memory-one strategies or those of a certain structural form like finite state machines or deterministic memory-two strategies. Furthermore, our tournaments encompass variations, including standard tournaments resembling Axelrod’s original ones, tournaments with noise, probabilistic match length, and both noise and probabilistic match length. This diversity in strategies and tournament types provides new insights and tests earlier claims in alternative settings against known powerful strategies. More specifically, we show that the previous tournament winners are lacking against large enough opponent pools; they do not appear among the top-performing strategies anymore. This could be due to likely suffering from a lack of diversity in the strategies they were trained/tested against, finding it hard to adapt to the new strategies.

It is important to note that we do not assert the existence of a single best-performing strategy across all tournaments or tournament types. On the contrary, our work demonstrates that such a strategy does not exist (notwithstanding a few strategies with broadly high performance). The primary objective of this paper, presented in the latter parts of the paper, is to continue the discussion on the properties of successful strategies, a conversation started by Axelrod. The results of our analysis conclude that the properties of a successful strategy in the Iterated Prisoner’s Dilemma (IPD) are:

- (i) Be a little bit envious
- (ii) Be “nice” in non-noisy environments or when game lengths are longer
- (iii) Reciprocate both cooperation and defection appropriately; Be provokable in tournaments with short matches, and generous in tournaments with noise
- (iv) It’s ok to be clever
- (v) Adapt to the environment; Adjust to the mean population cooperation

We believe that the discussion on the properties of winning strategies holds significant importance. It aims to provide guidance to researchers designing new strategies and those training strategies. Specifically, much like the recognized value of diversity in training datasets, such as variations in image perspective, skin color, etc., are critical in training accurate and generalizable machine learning models, we show that diversity in the population of opponent strategies is of paramount importance in the construction and evaluation of game theory strategies. Moreover, conducting a similar analysis can shed light on already trained strategies, aiding in understanding the key features they have autonomously developed during their training processes.

The rest of the paper is organized as follows. In section 2, we describe the data collection process. In subsection 3.1, we present the best performing strategies for each type of tournament. Subsection 3.2 explores the traits that contribute to good performance, and in subsection 3.3, we focus on the features of the winners of the tournaments. Finally, the results are summarized in section 4. This manuscript introduces several parameters, which are discussed in the following sections; the full set of parameters and their definitions are also provided in the Supplementary Material.

2 Model

The data collection of various types of tournaments and the use of different strategies are made possible due to an open-source library called **Axelrod-Python** [27] (version 3.0.0). **Axelrod-Python** enables the simulation of IPD tournaments and contains an extensive list of strategies. Most of these strategies are described in the literature, with a few exceptions contributed specifically to the package. In this paper, we use a total of 195 strategies, which can be found in the Supplementary Material. The package supports several tournament types, and this work considers standard, noisy, probabilistic ending, and noisy probabilistic ending tournaments.

Standard tournaments are similar to Axelrod’s well-known tournaments [2]. In these tournaments, there are N strategies, and each strategy plays an iterated game with n turns against all other strategies, not including self-interactions. *Noisy tournaments* also involve N strategies and n turns, but in each turn, there is a probability p_n that a player’s action is flipped. Compared to these two tournaments, in *probabilistic ending tournaments* the number of turns is not fixed. Instead, a match between strategies ends with a given probability p_e . Finally, *noisy probabilistic ending tournaments* incorporate both a noise probability p_n and an ending probability p_e . For smoother results, each tournament is repeated k times, and this repetition factor was allowed to vary to assess the impact of smoothing. The winner of each tournament is determined based on the average score achieved by a strategy from the entire set of repetitions, not by the number of wins.

To run a tournament, only a few lines of code are required (Figure 1). Specifically, one needs to define the list of strategies that the players use when participating in the tournament, the number of repetitions, the number of turns or the probability of each match ending, and the probability of noise. We demonstrate two examples in Figure 1: one with a fixed number of turns and one with a probabilistic ending. A tournament is an instance of the ‘Tournament’ class. To execute the tournament, users need to run the ‘play()’ method, which returns an instance of the ‘Results’ class. This instance includes many details of the tournament, such as the winner and the average score. Additionally, a more detailed summary of the tournament is provided, which we use in this analysis and describe in detail below.

The process of collecting tournament results is outlined in Algorithm 1. For each trial, a random size N is selected, and a random list of N strategies from the 195 available. Subsequently, one standard, one noisy, one probabilistic ending, and one noisy probabilistic ending tournament are conducted for the selected list of strategies. The parameters for the tournaments, as well as the number of repetitions, are chosen once for each trial. We have run a total of 11400 trials of Algorithm 1. For each trial, we collect the results for four

```

>>> import axelrod as axl

>>> players = [axl.TitForTat(), axl.GTFT(), axl.Gradual(), axl.EvolvedLookerUp2_2_2()]

>>> standard_tournament = axl.Tournament(players=players, turns=50, repetitions=10, noise=0)

>>> results = standard_tournament.play()

>>> prob_noisy_end_tournament = axl.Tournament(players=players, prob_end=0.5,
...                                           repetitions=10, noise=0.001)

>>> results = prob_noisy_end_tournament.play()

```

Figure 1: **Example usage of the Axelrod-Python package for running tournaments.** The strategies that players use are saved in a list, which is then passed to the Tournament class. In this example, strategies we previously discussed, such as TFT, Generous Tit-for-Tat, Gradual, and a stochastic strategy evolved through reinforcement learning, are included. We create a standard tournament where noise is set to 0, as well as a noisy tournament with probabilistic ending. Once an instance of the tournament class is defined, executing the tournament is straightforward. The play() method generates all possible pairs from the list of strategies, including pairs where each strategy plays against itself, and then iterates through each match to play it. Each match is repeated according to the specified number of repetitions, with results aggregated to a tournament summary.

different tournaments, resulting in a total of 45606 (11400×4) tournament results. Each tournament outputs a result summary in the form of Table 1.

The summary contains statistics regarding each strategy that participated in the tournament, such as its rank, cooperation rate, time spent in each state, and the probability of cooperating after each of the four possible outcomes of the previous round. In our analysis, we will use the measures for each strategy that are provided in the summary, as well as additional measures we calculated. Namely, these include the SSE error, the average cooperation in each tournament, the maximum, and the minimum values. The SSE (introduced in [28]) shows how close a strategy is to behaving as a ZDs, and subsequently, in an extortionate way. We also consider how each strategy’s cooperation rate C_r compares to those of the tournament as a whole, for example, by comparing C_r to C_{\max} .

During the data collection process, the probabilities of noise (p_n) and tournament ending (p_e) were allowed to take values between 0 and 1. However, commonly used values for these probabilities are $p_n \leq 0.1$ and $p_e \leq 0.1$. This is to make the results more interpretable. For example, consider a strategy competing in an environment with $p_n > 0.1$. In cases with a high value of noise, most of the actions the strategy takes are the complete opposite of what the strategy is designed to do. Therefore, we will focus on the tournaments for which $p_n \leq 0.1$ and $p_e \leq 0.1$. Thus, the results presented here pertain to subsets of the noisy and probabilistic ending tournaments. Specifically, the results rely on 1150 tournaments with noise, 1134 tournaments with a probabilistic ending, and 117 tournaments with both noise and a probabilistic ending. We also provide an analysis of the paper considering the entire datasets, and these results are presented in the Supplementary Material. The general results of the analysis are not affected by the restriction of the noise and probabilistic ending probabilities.

Algorithm 1: Tournament Data Collection Algorithm

for $seed \in [0, 11420]$ **do** $N \leftarrow$ randomly select integer $\in [3, 195]$;
players \leftarrow randomly select N players;
 $k \leftarrow$ randomly select integer $\in [10, 100]$;
 $n \leftarrow$ randomly select integer $\in [1, 200]$;
 $p_n \leftarrow$ randomly select float $\in [0, 1]$;
 $p_e \leftarrow$ randomly select float $\in [0, 1]$;

result standard \leftarrow Axelrod.tournament(players, n, k);
result noisy \leftarrow Axelrod.tournament(players, n, p_n, k);
result probabilistic ending \leftarrow Axelrod.tournament(players, p_e, k);
result noisy probabilistic ending \leftarrow Axelrod.tournament(players, p_n, p_e, k);**return** result standard, result noisy, result probabilistic ending, result noisy probabilistic ending;

Rank	Name	Median score	Cooperation rating (C_r)	Win	Initial C	Rates							
						CC	CD	DC	DD	CC to C	CD to C	DC to C	DD to C
0	EvolvedLookerUp2 2 2	2.97	0.705	28.0	1.0	0.639	0.066	0.189	0.106	0.836	0.481	0.568	0.8
1	Evolved FSM 16 Noise 05	2.875	0.697	21.0	1.0	0.676	0.020	0.135	0.168	0.985	0.571	0.392	0.07
2	PSO Gambler 1 1 1	2.874	0.684	23.0	1.0	0.651	0.034	0.152	0.164	1.000	0.283	0.000	0.136
3	PSO Gambler Mem1	2.861	0.706	23.0	1.0	0.663	0.042	0.145	0.150	1.000	0.510	0.000	0.122
4	Winner12	2.835	0.682	20.0	1.0	0.651	0.031	0.141	0.177	1.000	0.441	0.000	0.462
...

Table 1: **Result Summary Example of a Tournament.** A result summary consists of N rows, with each row containing information for each strategy that participated in the tournament. This information includes the strategy’s rank (R), median score, the cooperation rate (C_r), the number of match wins, and the probability that the strategy cooperated in the opening move. Additionally, it provides the probabilities of a strategy being in any of the four states (CC, CD, DC, DD) and the cooperation rate after each state.

3 Results

3.1 Top ranked strategies across tournaments

A strategy has participated in multiple tournaments of each type, and to evaluate its overall performance, we introduce a measure called the *normalized rank*. In each tournament, the strategies receive a rank (R), where 0 denotes that the strategy was the winner, and $N - 1$ indicates that the strategy came last in the tournament. The normalized rank, denoted as r , is calculated as $r = \frac{R}{N-1}$. Thus, the rank a strategy achieved over the number of players in the tournament. The performance of the strategies is assessed based on the *median of the normalized rank*, denoted as \bar{r} .

For example, let’s consider the well-known strategies TFT and Gradual. Each strategy participated in several tournaments of each type. In Figure 2 we show the distribution of the normalised ranks of these strategies in each of the four tournaments. We can observe that TFT looks to be normally distributed normalized rank. In comparison, Gradual’s performance has longer tails, indicating that there were tournaments where the strategy performed very well or very poorly. Overall, Gradual achieves a lower median rank, signifying that

it performs better than TFT except in the case of noisy and probabilistic ending tournaments (lower rank is better).

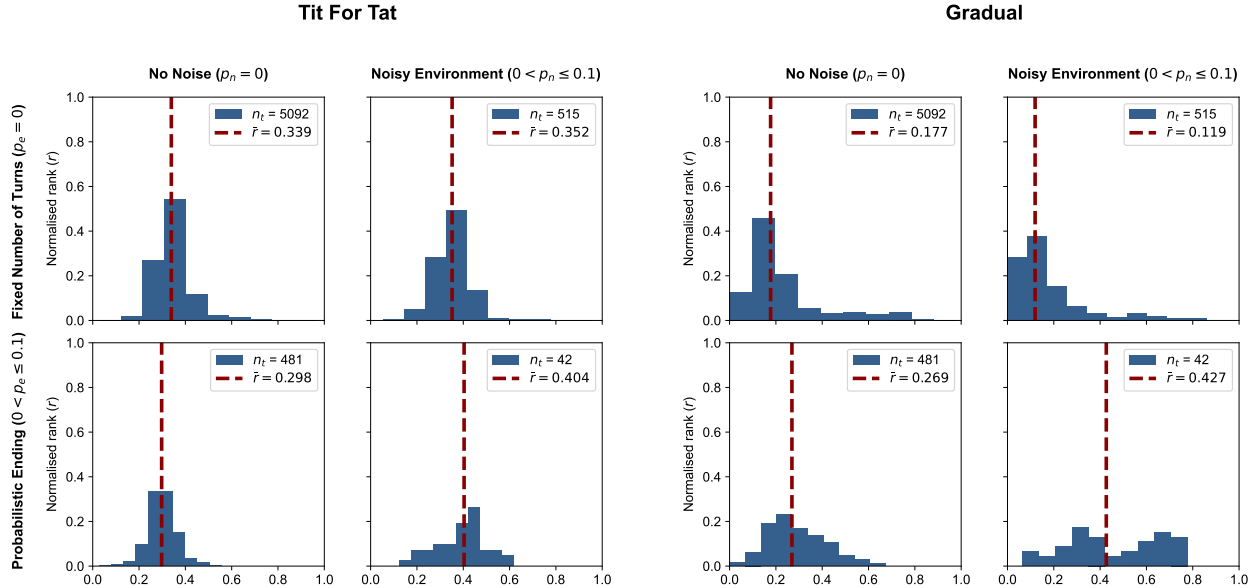


Figure 2: **Examples of normalized rank distributions for two strategies, TFT and Gradual.** We plot the distributions of r for the two strategies in the four tournament types. As a reminder, lower values of r correspond to better performances. The top left quadrant of each plot shows the distribution for standard tournaments (fixed number of turns and no noise). The top right quadrant shows the distribution for noisy tournaments (fixed number of turns and noise). The bottom left quadrant shows the distribution for probabilistic ending tournaments (no noise and probabilistic ending). Finally, the bottom right quadrant shows the distribution for noisy probabilistic ending tournaments (noise and probabilistic ending). In each quadrant, we also show the number of data points. Both strategies participated in a similar number of tournaments. Based on the median rank, which we use in this work to define overall performance, TFT performs best in probabilistic ending tournaments, whereas Gradual was in standard tournaments.

The top 15 strategies for each tournament type, based on \bar{r} , are presented in Table 2, while the r distributions for the top-ranked strategies can be found in Figure 3.

In standard tournaments dominating strategies were those trained using reinforcement learning techniques. 10 out of the 15 top strategies were introduced in [23]. These strategies are based on finite state automata (FSM), hidden Markov models (HMM), artificial neural networks (ANN), lookup tables (LookerUp), and stochastic lookup tables (Gambler). They have been trained using reinforcement learning algorithms (evolutionary and particle swarm algorithms) to perform well against a subset of the strategies in **Axelrod-Python** in a standard tournament. Thus, their performance in the specific setting was anticipated, although still noteworthy given the random sampling of tournament participants. DoubleCrosser and BackStabber, both from the **Axelrod-Python**, use the number of turns and are set to defect in the last two rounds. These strategies can be characterized as “cheaters” because their source code allows them to know the number of turns (unless the match has a probabilistic ending). These strategies were expected to not perform as well in tournaments where the number of turns is not specified. Finally, Winner 12 [22] and DBS [29] are both from the literature. DBS is a strategy specifically designed for noisy environments; however, it ranks highly in standard tournaments as well. Similarly, the fourth-ranked player, Evolved FSM 16 Noise 05, was trained for noisy tournaments yet performs well in standard tournaments.

In the case of noisy tournaments, the top-performing strategies include strategies specifically designed for

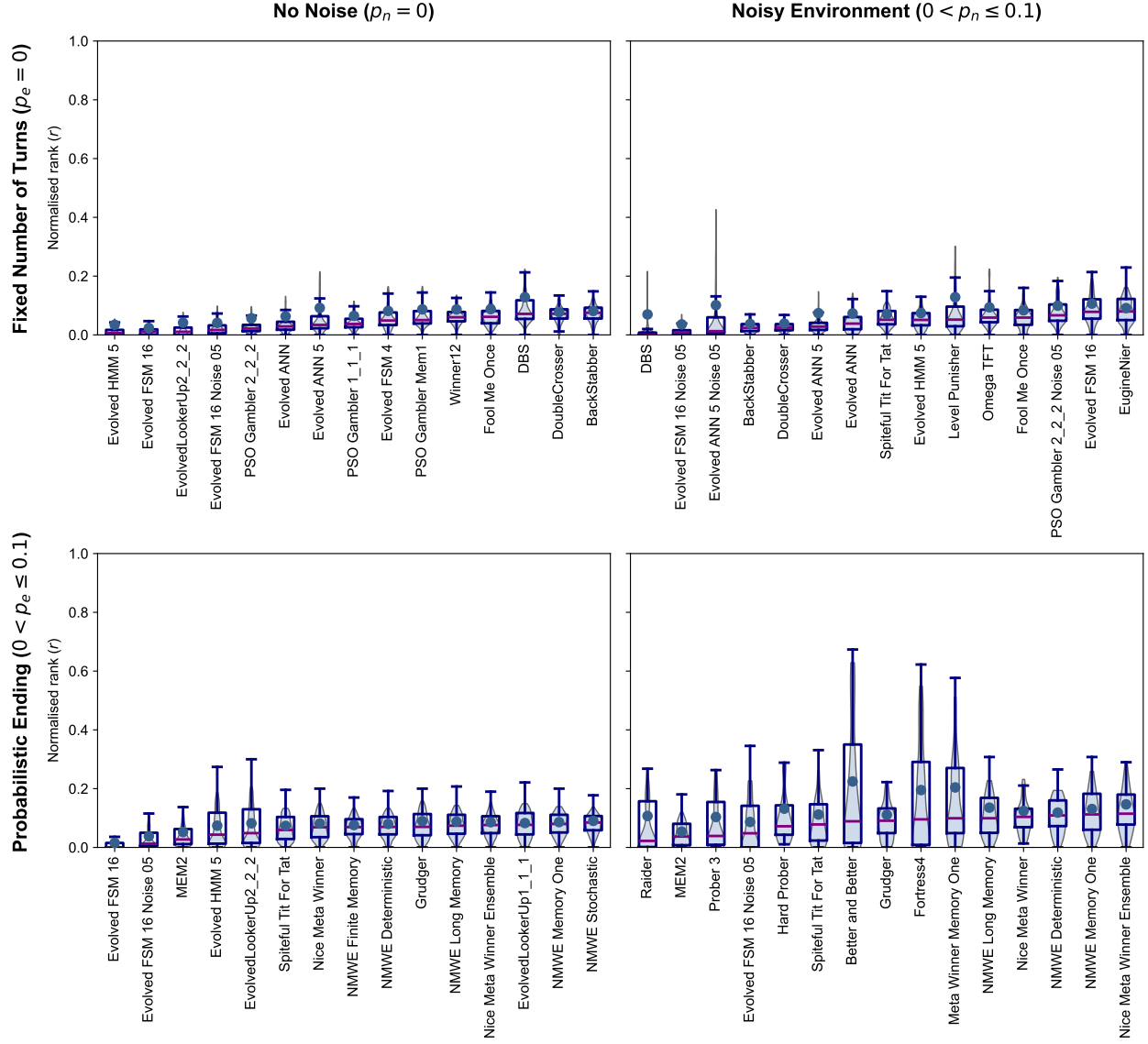


Figure 3: r distributions of the top 15 strategies in different environments. A lower value of \bar{r} corresponds to a more successful performance. A strategy's r distribution skewed towards zero indicates that the strategy ranked highly in most tournaments it participated in. Most distributions are skewed towards zero.

Standard		Noisy ($p_n \leq 0.1$)		Probabilistic ending ($p_e \leq 0.1$)		Noisy probabilistic ending	
Name	\bar{r}	Name	\bar{r}	Name	\bar{r}	Name	\bar{r}
0 Evolved HMM 5	0.007	DBS	0.0	Evolved FSM 16	0.0	Raider	0.022
1 Evolved FSM 16	0.01	Evolved FSM 16 Noise 05	0.008	Evolved FSM 16 Noise 05	0.013	MEM2	0.037
2 EvolvedLookerUp2.2.2	0.011	Evolved ANN 5 Noise 05	0.013	MEM2	0.027	Prober 3	0.039
3 Evolved FSM 16 Noise 05	0.017	BackStabber	0.024	Evolved HMM 5	0.043	Evolved FSM 16 Noise 05	0.048
4 PSO Gambler 2.2.2	0.022	DoubleCrosser	0.025	EvolvedLookerUp2.2.2	0.049	Hard Prober	0.072
5 Evolved ANN	0.029	Evolved ANN 5	0.028	Spiteful Tit For Tat	0.059	Spiteful Tit For Tat	0.078
6 Evolved ANN 5	0.034	Evolved ANN	0.038	Nice Meta Winner	0.069	Better and Better	0.089
7 PSO Gambler 1.1.1	0.037	Spiteful Tit For Tat	0.051	NMWE Finite Memory	0.069	Grudger	0.091
8 Evolved FSM 4	0.049	Evolved HMM 5	0.051	NMWE Deterministic	0.07	Fortress4	0.096
9 PSO Gambler Mem1	0.05	Level Punisher	0.052	Grudger	0.07	Meta Winner Memory One	0.099
10 Winner12	0.06	Omega TFT	0.059	NMWE Long Memory	0.074	NMWE Long Memory	0.099
11 Fool Me Once	0.061	Fool Me Once	0.059	Nice Meta Winner Ensemble	0.076	Nice Meta Winner	0.104
12 DBS	0.071	PSO Gambler 2.2.2 Noise 05	0.067	EvolvedLookerUp1.1.1	0.077	NMWE Deterministic	0.109
13 DoubleCrosser	0.072	Evolved FSM 16	0.078	NMWE Memory One	0.08	NMWE Memory One	0.112
14 BackStabber	0.075	EugeneNier	0.08	NMWE Stochastic	0.085	Nice Meta Winner Ensemble	0.115

Table 2: Top performances for each tournament type based on \bar{r} . The results of each type are based on 11420 unique tournaments. The results for noisy tournaments with $p_n < 0.1$ are based on 1151 tournaments, and for probabilistic ending tournaments with $p_e < 0.1$ on 1139. The top ranks indicate that trained strategies perform well in a variety of environments, but so do simple deterministic strategies. For noisy tournaments DBS is the top ranked strategy with $\bar{r} = 0$, thus DBS won every tournament it participated in. The same for Evolved FSM 16 Noise 05 in probabilistic ending.

noisy tournaments. These are DBS, Evolved FSM 16 Noise 05, Evolved ANN 5 Noise 05, PSO Gambler 2 2 2 Noise 05, and Omega Tit For Tat [20]. Omega TFT, a strategy designed to break the deadlocking cycles of *CD* and *DC* that TFT can fall into in noisy environments, places 10th. The rest of the top ranks are occupied by strategies that performed well in standard tournaments and deterministic strategies such as Spiteful Tit For Tat [30], Level Punisher [31], Eugene Nier [32].

Furthermore, in tournaments with probabilistic endings, the highly ranked strategies leaned towards defecting strategies and trained finite state automata, as demonstrated by the works of Ashlock et al.[33, 34]. The most effective strategies in probabilistic ending tournaments are also a series of ensemble Meta strategies, trained strategies that performed well in standard tournaments, and Grudger[27] and Spiteful Tit for Tat [30]. The Meta strategies [27] utilize a team of strategies and aggregate the potential actions of the team members into a single action in various ways.

While no single strategy consistently outperforms all others in any of the distinct tournament types or across various tournament types, certain types of strategies consistently achieve top rankings. These include strategies that have undergone training, those that retaliate, and those that adapt their behavior based on preassigned rules to optimize outcomes. These findings challenge some of Axelrod’s suggestions, particularly the advice to “not be clever” and “not be envious”.

3.2 The effect of strategy features on performance

For each strategy, we have a variety of features as described in Table 3. These features capture measures related to a strategy’s behavior in the tournaments it competed in, as well as intrinsic properties, such as whether a strategy is deterministic or stochastic. The correlation coefficients between the features for performance evaluation, the median score and the median normalised rank are given by Table 4. The correlation coefficients between all features have also been calculated and a graphical representation can be found in the Supplementary Material.

feature	feature explanation	source	value type	min value	max value
stochastic	If a strategy is stochastic	strategy classifier from APL	boolean	Na	Na
makes use of game	If a strategy makes used of the game information	strategy classifier from APL	boolean	Na	Na
makes use of length	If a strategy makes used of the number of turns	strategy classifier from APL	boolean	Na	Na
memory usage	The memory size of a strategy divided by the number of turns	memory size from APL	float	0	1
SSE	A measure of how far a strategy is from ZD behaviour	method described in [28]	float	0	1
max cooperating rate (C_{max})	The biggest cooperating rate in a given tournament	result summary	float	0	1
min cooperating rate (C_{min})	The smallest cooperating rate in a given tournament	result summary	float	0	1
median cooperating rate (C_{median})	The median cooperating rate in a given tournament	result summary	float	0	1
mean cooperating rate (C_{mean})	The mean cooperating rate in a given tournament	result summary	float	0	1
C_r / C_{max}	A strategy's cooperating rate divided by the maximum	result summary	float	0	1
C_{min} / C_r	The minimum divided by a strategy's cooperating rate	result summary	float	0	1
C_r / C_{median}	A strategy's cooperating rate divided by the median	result summary	float	0	1
C_r / C_{mean}	A strategy's cooperating rate divided by the mean	result summary	float	0	1
C_r	The cooperating ratio of a strategy	result summary	float	0	1
CC to C rate	The probability a strategy will cooperate after a mutual cooperation	result summary	float	0	1
CD to C rate	The probability a strategy will cooperate after being betrayed by the opponent	result summary	float	0	1
DC to C rate	The probability a strategy will cooperate after betraying the opponent	result summary	float	0	1
DD to C rate	The probability a strategy will cooperate after a mutual defection	result summary	float	0	1
p_n	The probability of a player's action being flip at each interaction	trial summary	float	0	1
n	The number of turns	trial summary	integer	1	200
p_e	The probability of a match ending in the next turn	trial summary	float	0	1
N	The number of strategies in the tournament	trial summary	integer	3	195
k	The number of repetitions of a given tournament	trial summary	integer	10	100

Table 3: **Included features for performance evaluation analysis.** Stochastic, makes use of length and makes use of game are APL classifiers that determine whether a strategy is stochastic or deterministic, whether it makes use of the number of turns or the game's payoffs. The memory usage is calculated as the number of turns the strategy considers to make an action (which is specified in the APL) divided by the number of turns. The SSE (introduced in [28]) shows how close a strategy is to behaving as a ZDs, and subsequently, in an extortionate way. The method identifies the ZDs closest to a given strategy and calculates the algebraic distance between them as the sum of squared error (SSE). A SSE value of 1 indicates no extortionate behaviour at all whereas a value of 0 indicates that a strategy is behaving as a ZDs. The memory usage of strategies is the number of rounds of play used by the strategy when deciding on an action, divided by the number of turns in each match. For example, Winner12 uses the previous two rounds of play, and if participating in a match with 100 turns its memory usage would be 2/100. For strategies with an infinite memory size, for example Evolved FSM 16 Noise 05, memory usage is equal to 1. Note that for tournaments with a probabilistic ending the number of turns was not collected, so the memory usage feature is not used for probabilistic ending tournaments. The rest of the features considered are the CC to C , CD to C , DC to C , and DD to C rates as well as cooperating ratio of a strategy, the minimum (C_{min}), maximum (C_{max}), mean (C_{mean}) and median (C_{median}) cooperating ratios of each tournament.

	Standard		Noisy $p_n \leq 0.1$		Probabilistic ending $p_e \leq 0.1$		Noisy probabilistic ending	
	r	median score	r	median score	r	median score	r	median score
CC to C rate	-0.501	0.501	-0.210	0.194	-0.336	0.348	0.087	0.015
CD to C rate	0.226	-0.199	0.337	-0.235	0.458	-0.352	0.609	-0.372
DC to C rate	0.127	-0.100	0.227	-0.111	0.164	-0.105	0.410	-0.203
DD to C rate	0.412	-0.396	0.549	-0.391	0.433	-0.378	0.615	-0.407
C_r	-0.323	0.383	0.298	-0.051	-0.060	0.160	0.595	-0.213
C_{max}	0.000	0.050	-0.000	0.244	-0.000	0.079	-0.000	0.296
C_{min}	0.000	0.085	0.000	-0.070	0.000	0.128	0.000	0.000
C_{median}	0.000	0.209	0.000	0.572	-0.000	0.324	0.000	0.667
C_{mean}	0.000	0.229	-0.000	0.583	-0.000	0.354	-0.000	0.689
C_r / C_{max}	-0.323	0.381	0.307	-0.076	-0.060	0.156	0.608	-0.246
C_{min} / C_r	0.109	-0.080	-0.141	-0.011	0.024	0.029	-0.335	0.092
C_r / C_{median}	-0.330	0.353	0.326	-0.258	-0.065	0.111	0.614	-0.464
C_r / C_{mean}	-0.331	0.357	0.325	-0.228	-0.066	0.114	0.617	-0.431
N	-0.000	-0.009	-0.000	-0.017	-0.000	0.011	0.000	0.139
k	-0.000	-0.002	-0.000	-0.003	-0.000	0.010	-0.000	0.035
n	-0.000	-0.125	-0.000	-0.392	-	-	-	-
p_n	-	-	0.000	-0.244	-	-	0.000	-0.272
p_e	-	-	-	-	0.000	0.257	0.000	0.568
Make use of game	-0.003	-0.022	-0.047	0.014	-0.046	0.022	-0.110	0.057
Make use of length	-0.158	0.124	-0.224	0.139	-0.173	0.128	-0.206	0.115
SSE	0.473	-0.452	0.589	-0.412	0.458	-0.418	0.571	-0.383
stochastic	0.006	-0.024	0.010	-0.007	-0.001	0.001	-0.001	0.002
memory usage	-0.098	0.108	-0.080	0.114	-	-	-	-

Table 4: **Correlations between the features of Table 3 and the normalised rank and the median score.** For each type of tournament, standard, noisy, probabilistic ending, and noisy probabilistic ending, we run a correlation analysis. For each tournament we check the correlation between each feature used in our analysis and the normalised random and median score. Note that the correlation coefficients are calculated using the Spearman’s rank correlation coefficient. A negative value indicates a negative correlation and in the case of the normalised rank a smaller rank translates to better position in the tournament.

In standard tournaments, the features CC to C , C_r , C_r/C_{\max} , and the cooperating ratio compared to C_{median} and C_{mean} have a moderately negative effect on the normalized rank (a smaller rank is better) and a moderate positive effect on the median score. The SSE error and the DD to C rate have the opposite effects. Thus, in standard tournaments, behaving cooperatively corresponds to a more successful performance. Even though being nice generally pays off, that does not hold against defective strategies. Being more cooperative after a mutual defection, that is not retaliating, is associated with lesser overall success in terms of normalized rank. Compared to standard tournaments, in both noisy and noisy probabilistic ending tournaments, the higher the rates of cooperation, the lower a strategy's success and median score. A strategy would not want to cooperate more than both the mean and median cooperator in such settings. In probabilistic ending tournaments, the cooperation rate of the winners and its relative comparison to the cooperation rates of the tournament have no effect. The only features that have an effect are the CD to C rate, which is the tendency of a strategy to forgive, and the SSE rate, which has a positive effect on the normalized rank.

A multivariate linear regression has been fitted to model the relationship between the features and the normalized rank. Based on the graphical representation of the correlation matrices given in the Supplementary Information, several features are highly correlated and have been removed before fitting the linear regression model. The features included are given in Table 5 alongside their corresponding p values in distinct tournaments and their regression coefficients. The CD to C rate has a positively statistically significant effect on the normalized rank across all tournament types. This suggests that being generous tends to lower one's performance. In the case of probabilistic ending tournaments, the coefficient of the CD to C rate is the highest, indicating that one should be more provocative in this setting. Similarly, the SEE error rate has a positive effect on the normalized rank, suggesting that being extortionate pays off, especially in noisy tournaments. The measures of cooperation, C_r and C_r/C_{\max} , also exhibit a significant effect. In noisy probabilistic ending tournaments, this effect is positive; however, the coefficient is very close to zero. In other tournament types, the effect is negative, indicating that one should aim to be less cooperative than the mean cooperator of the tournament. However, we cannot interpret the result as suggesting that a strategy should be as uncooperative as possible.

	Standard		Noisy $p_n \leq 0.1$		Probabilistic ending $p_e \leq 0.1$		Noisy probabilistic ending	
	R adjusted: 0.541		R adjusted: 0.373		R adjusted: 0.457		R adjusted: 0.537	
	Coefficient	p -value	Coefficient	p -value	Coefficient	p -value	Coefficient	p -value
constant	0.695	0.000	0.560	0.000	0.627	0.000	0.345	0.005
CC to C rate	-0.042	0.000	-0.163	0.000	-0.046	0.000	0.032	0.029
CD to C rate	0.297	0.000	0.064	0.000	0.412	0.000	0.292	0.000
DC to C rate	0.198	0.000	0.142	0.000	0.193	0.000	0.193	0.000
SSE	0.258	0.000	0.328	0.000	0.190	0.000	0.228	0.000
C_{\max}	-0.068	0.000	-0.048	0.214	-0.040	0.347	-0.011	0.936
C_{\min}	-0.161	0.000	-0.029	0.367	-0.049	0.017	0.008	0.912
C_{mean}	0.117	0.000	-0.133	0.000	-0.159	0.000	-0.468	0.000
C_{\min} / C_r	0.057	0.000	-0.006	0.322	0.054	0.000	0.034	0.099
C_r / C_{mean}	-0.468	0.000	-0.073	0.000	-0.150	0.000	0.094	0.000
k	0.000	0.325	0.000	0.965	0.000	0.079	0.000	0.065
n	0.000	0.000	-	-	-	-	-	-
memory usage	-0.010	0.000	-0.008	0.000	-	-	-	-
C_r / C_{median}	-	-	0.069	0.001	-0.142	0.000	-	-
p_n	-	-	-0.131	0.010	-	-	-0.278	0.048
p_e	-	-	-	-	-0.071	0.016	0.320	0.024

Table 5: **Results of multivariate linear regressions with r as the dependent variable.** R squared is reported for each model. The R scores of the fitted models indicate their capability to explain some of the variation in the median rank. Most of the features have a statistically significant effect on the normalized rank. A multivariate linear regression has also be fitted on the median score. The coefficients and p values of the features can be found in Supplementary Information. Both approaches lead to similar conclusions.

The results presented here suggest that generosity/provocation and a strategy's cooperation rate, particu-

larly in comparison to the tournament averages, are significant features. The analysis suggests that strategies should be more generous in noisy tournaments and less generous in probabilistic ending tournaments. Moreover, strategies should aim to not cooperate more than the mean cooperator in their tournaments. We note the analysis is limited as we only consider a linear relationship between these parameters and the rank. To further investigate the effects of the parameters discussed in this section, we have conducted a more detailed analysis in the next section, focusing on the performances of the winners of the tournaments.

3.3 Features of top performing strategies

In Figure 4, we present the distributions of the cooperation ratio and C_r/C_{mean} for the winners of tournaments. A value of $C_r/C_{\text{mean}} = 1$ implies that the cooperation ratio of the winner was the same as the mean cooperating ratio of the tournament, and we observe that this occurs for most tournament types, apart from the case of noisy and probabilistically ending tournaments. In the case of probabilistic ending tournaments, there are several winners that cooperated much less than that, confirming the results of the previous section that defecting strategies can be winners in probabilistic ending tournaments. The distribution of the cooperation rates showcases a high cooperation rate in standard tournaments and probabilistic ending tournaments. In tournaments with noise, we observe a much less cooperative behavior, which could result from strategies being cautious of potential flip actions by the co-player or strategies not suited for noise holding grudges against defections.

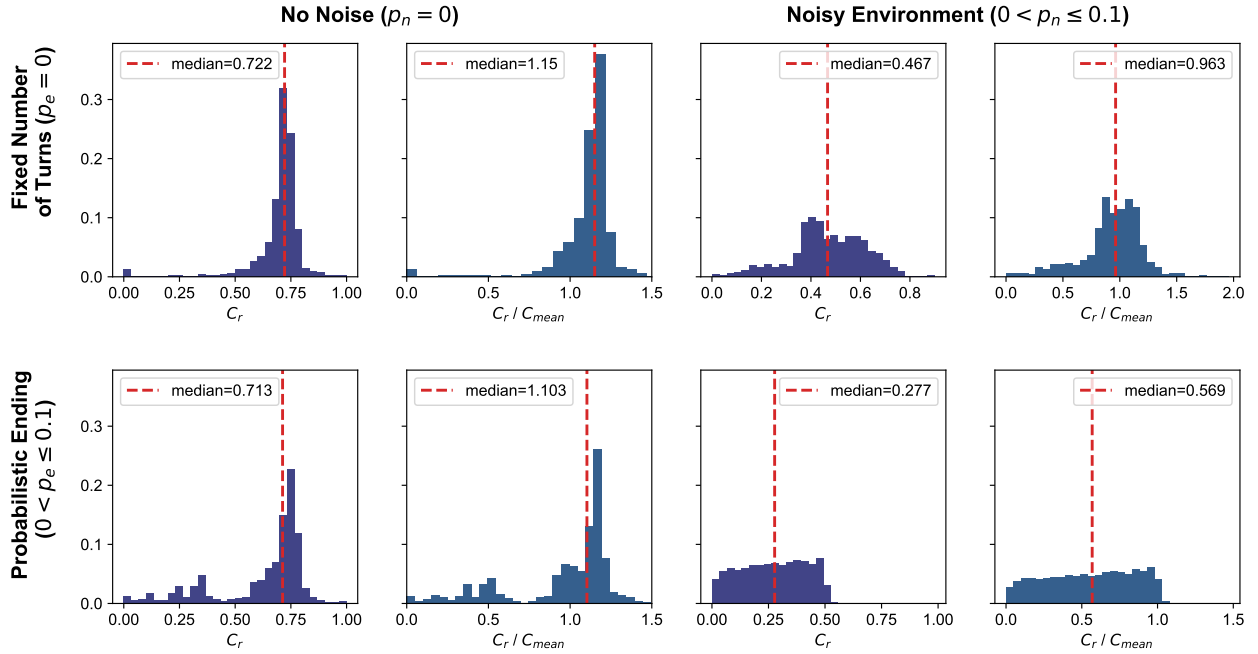


Figure 4: **Distributions of C_r and C_r/C_{mean} for the winners of tournaments.** In this distribution, we consider the winners of the tournaments, specifically the strategies that ranked first in each tournament. For each type of tournament, we plot the cooperation rate of the winner in the tournament they won, as well as the ratio of the winner’s cooperation rate to that of the entire tournament. A value of $C_r/C_{\text{mean}} = 1$ implies that the cooperation ratio of the winner was the same as the mean cooperation ratio of the tournament.

Analyzing the SSE distributions across different tournament types (Figure 5) suggests that successful strategies exhibit some extortionate behavior, though not consistently. ZDs are a set of strategies that are often envious, as they attempt to exploit their opponents. The winners of the tournaments considered in this work

demonstrate envious behavior, but not to the extent observed in many ZDs. While the exact interactions between matches are not recorded here, the work of [23], which introduced the trained strategies appearing in the top-ranked strategies of Section 3.1, did record such interactions. In [23], it was shown that clever strategies managed to achieve mutual cooperation with stronger strategies while exploiting weaker ones. This could explain the clever winners in our analysis and the observed SSE distributions.

This might also be the reason why ZDs fail to appear in the top ranks—they attempt to exploit all opponents and cannot actively adapt back to mutual cooperation against stronger strategies, which requires a deeper memory. It’s worth noting that ZDs tend to perform poorly in population games for a similar reason: they aim to exploit other players using ZDs, failing to form a cooperative subpopulation [35]. This makes them effective invaders but poor at resisting invasion.

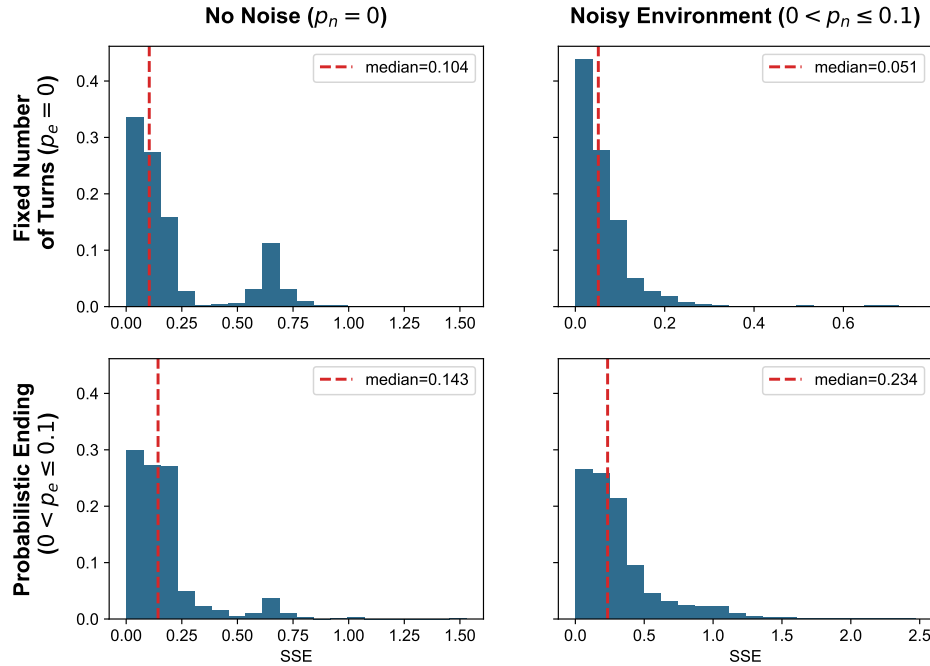


Figure 5: **Distributions of SSE error for the winners of tournaments.** Here, we again consider the winners of the tournaments, separated by type of tournament, and plot their SSE error. As a reminder, the SSE error indicates how closely a strategy behaves like a Zero-Determinant (ZD) strategy, and subsequently, in an extortionate way. An SSE value of 1 indicates no extortionate behavior at all, whereas a value of 0 indicates that a strategy is behaving as a ZD.

Finally, we examine the distributions of the cooperation rates after the outcomes CC , CD , DC , and DD , as shown in Figure 6. In the case of cooperating after mutual cooperation, the results align with expectations; the distributions skew towards 1, indicating that the winners of the tournaments are more likely to cooperate after mutual cooperation. Regarding the CD outcome and the likelihood to cooperate after such a result, capturing generosity, the distributions skew towards $1/2$, not 1, suggesting that strategies need to reduce their readiness to forgive. This aligns with the known result that Generous Tit For Tat generally outperforms TFT in most settings. In probabilistic ending tournaments, there is a peak at 0, suggesting that strategies should not be too generous in tournaments with short matches. Such a peak also appears in standard tournaments; however, not in tournaments with noise, where a strategy should be more generous.

Part of a strategy’s envious behavior can be captured by the rate of DC to C . In noisy tournaments, winners are not too envious, but in tournaments without noise, we can see that winners behave in two ways. Some

are a bit envious, whereas others are very envious. In the DD to D , we can observe that, expectedly, the results are skewed towards 0. However, there are winners that attempt to recover from a DD outcome. The remaining results are as expected, skewed towards 0.

4 Discussion

This manuscript explores the performance of 195 strategies in the IPD in thousands of computer tournaments. The collection of computer tournaments presented here is the largest and most diverse in the literature. The 195 strategies are drawn from **Axelrod-Python** library and include strategies from the IPD literature. The computer tournaments encompass four different types. So, what is the best way to play the IPD? And is there a single dominant strategy for the IPD? There was not a single strategy within the collection of 195 strategies that managed to perform well in all the tournament variations it competed in. A strategy ranking highly in a specific environment did not guarantee its success over different tournament types, with a few exceptions - strategies that generalize better. Already well-known in the AI/ML literature, adding noise to training data leads to more robust models. We see that clearly here, where the strategies trained for noise (or designed for noise) tend to be better generalists. There were instances where a few strategies trained in narrow conditions outperformed more generalist strategies, as they tend to overfit. However, the strategies trained with noise perform well in general, whilst the strategies trained specifically on no noise or small subpopulations do not.

We also examined the best-performing strategies across various tournament types and analyzed their salient features. This demonstrated that there are properties associated with the success of strategies that contradict the originally suggested properties of Axelrod [1]. We showed that complex or **clever** strategies can be effective, whether trained against a corpus of possible opponents or purposely designed to mitigate the impact of noise such as the DBS strategy. Moreover, we found some strategies designed or trained for noisy environments were also highly ranked in noise-free tournaments which reinforces the idea that strategies' complexity/cleverness is not necessarily a liability, rather it can confer adaptability to a more diverse set of environments. We also showed that while the type of exploitation attempted by ZDs is not typically effective in standard tournaments, **envious** strategies capable of both exploiting and not their opponents can be highly successful. Based on the results of [23] this could be because they are selectively exploiting weaker opponents while mutually cooperating with stronger opponents. Highly noisy or tournaments with short matches also favoured envious strategies. These environments mitigated the value of being nice. Uncertainty enables exploitation, reducing the ability of maintaining or enforcing mutual cooperation, while triggering grudging strategies to switch from typically cooperating to typically defecting.

The features analysis of the best performing strategies demonstrated that a strategy should reciprocate, as suggested by Axelrod, but it should relax its readiness to do so and be more **generous**. For noisy environments this is inline with the results of [16, 17, 18, 19], however, we also showed that generosity pays off even in standard settings, and that in fact the only setting a strategy would want to be too provokable is when the matches are not long. Forgiveness as defined by Axelrod was not explored here. This was mainly because the two round states were not recorded during the data collection. This could be a topic of future work that examines the impact of considering more rounds of history. The features analysis also concluded that there is a significant importance in **adapting to the environment**, and more specifically, to the mean cooperator. In most tournament types, the winner of the tournament was also the average cooperator. Even in tournaments with short matches where defecting behavior could secure a win, a large number of winners were average cooperators.

This could potentially explain the early success of TFT. TFT naturally achieves a cooperation rate near C_{mean} by virtue of copying its opponent's last move while also minimizing instances where it is exploited by an opponent (cooperating while the opponent defects), at least in non-noisy tournaments. It could also explain why Tit For N Tats does not fare well for $N > 1$ - it fails to achieve the proper cooperation ratio by tolerating

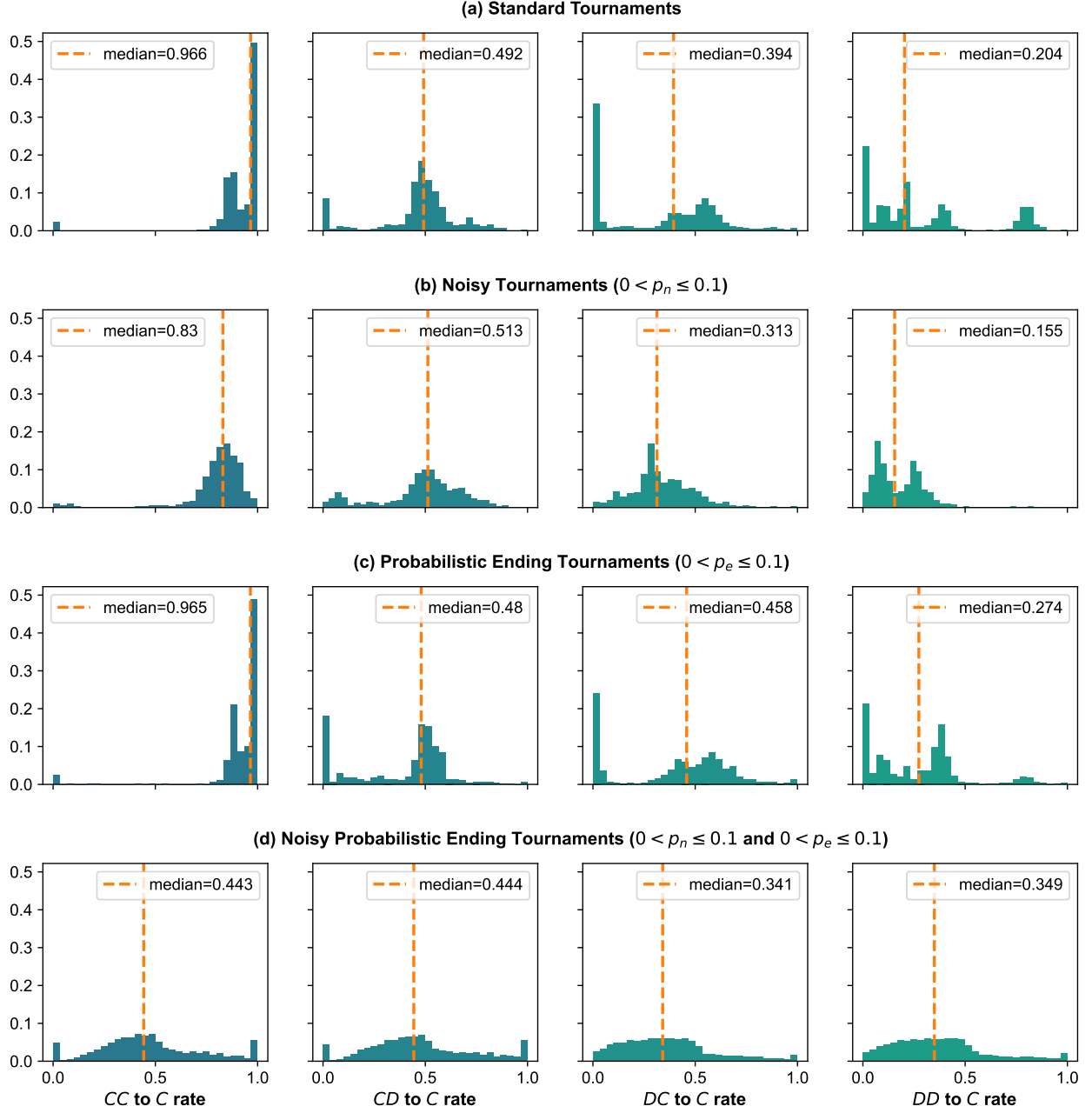


Figure 6: **Distributions of rates CC to C , CD to C , DC to C , and DD to C for the winners of tournaments.** The result summary from the tournaments records how often each strategy cooperated after each possible outcome of the previous round. Specifically, we analyze the probability with which a strategy chose C following the outcomes. Here, we plot the distributions of these probabilities for the winners of the tournaments. We separate them by type of tournament, from top to bottom: standard, noisy, probabilistic ending, and noisy probabilistic ending. From left to right, we plot the distributions of cooperation after the outcomes CC , CD , DC , and DD .

too many defections.

Similarly, our results could suggest an explanation regarding the intuitively unexpected effectiveness of memory-one strategies historically. Given that among the important features associated with success are the relative cooperation rate to the population average and the four memory-one probabilities of cooperating conditional on the previous round of play, these features can be optimized by a memory-one strategy such as TFT. Usage of more history becomes valuable when there are exploitable opponent patterns. This is indicated by the importance of SSE as a feature, showing that the first-approximation provided by a memory-one strategy is no longer sufficient. These results highlight a central idea in evolutionary game theory in this context: the fitness landscape is a function of the population (where fitness in this case is tournament performance). While that may seem obvious now, it shows why historical tournament results on small or arbitrary populations of strategies have so often failed to produce generalizable results.

To this end, many strategies, such as Win-Stay-Lose-Shift and Generous Tit For Tat, emerged due to their strong performance in evolutionary dynamics. Axelrod’s original work relied on computer tournaments, so we chose to remain consistent with this approach, as a comprehensive study like ours had not yet been undertaken. However, evolutionary settings would be an exciting direction for future study.

Overall, the five properties successful strategies need to have in a IPD competition based on the analysis that has been presented in this manuscript are:

- (i) Be “nice” in non-noisy environments or when game lengths are longer
- (ii) Be provokable in tournaments with short matches, and generous in tournaments with noise
- (iii) Be a little bit envious
- (iv) Be clever
- (v) Adapt to the environment (including the population of strategies).

The results presented here were based only on a subset of the whole data we have collected. The analysis of the full dataset is discussed in the Supplementary Information. However, we can see that the general results of our work remain the same. In the Supplementary Information, we also evaluate the importance of features using a random forest classifier and a clustering approach. The results of these analyses are also in line with the results presented here.

The data set described in this work contains the largest number of IPD tournaments, to the authors knowledge. The raw data set is available at [36] and the processed data at [37]. Further data mining could be applied and provide new insights in the field.

5 Data availability statement

The raw and processed datasets have been made publicly available [37, 36] and can be used for further analysis and insights.

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