# Parallel Computing Lab Nilay Ganvit - 200001053 30th November 2022

# **Assignment**

Q1.

```
//including headers
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <mpi.h>
#include <time.h>
int main(int argc, char **argv){
  MPI Init(&argc, &argv);
  MPI Comm size (MPI COMM WORLD, &proc num);
  MPI Comm rank (MPI COMM WORLD, &rank usr);
  MPI Request req[proc num + 1];
  MPI Status status;
  int Send Dat[proc num + 1][3];
  int filled[1000];
  int weights[1000];
  int mat graph[1000][1000];
  int num max colour = 0;
  int vertices, edges;
       filled[l] = -1;
      scanf("%d", &edges);
```

```
weights[i] = rand() % 42 + rand() % 42 + rand() % 42 + rand() %
42;
           filled[i] = -1;
       int start, end;
       for (i = 0; i < edges; i++)
           scanf("%d %d", &start, &end);
          mat graph[start][end] = 1;
          mat graph[end][start] = 1;
       clock t begin = clock();
      int max = -1;
      for (i = 0; i < vertices; i++)
               if (mat graph[i][j] != 0)
                   max++;
              num max colour = max;
       int num vetices per proc = vertices / (proc num - 1);
      int num vetices remaining = vertices / (proc num - 1) + (vertices %
(proc_num - 1));
       int startVertex = 0, endVertex;
      for (i = 1; i < proc num - 1; i++)
           MPI Send(&vertices, 1, MPI INT, i, i, MPI COMM WORLD);
```

```
for (k = 0; k < vertices; k++)
              MPI Send(&mat graph[k], vertices, MPI INT, i, i,
MPI COMM WORLD);
           endVertex = startVertex + num vetices per proc - 1;
           Send Dat[i][0] = startVertex;
          Send Dat[i][1] = endVertex;
          Send Dat[i][2] = num max colour;
          MPI Send(Send Dat[i], 3, MPI INT, i, i, MPI COMM WORLD);
          MPI Send(filled, vertices, MPI INT, i, i, MPI COMM WORLD);
          MPI Send(weights, vertices, MPI INT, i, i, MPI COMM WORLD);
          startVertex = endVertex + 1;
      MPI Send(&vertices, 1, MPI INT, i, i, MPI COMM WORLD);
       for (k = 0; k < vertices; k++)
          MPI Send(&mat graph[k], vertices, MPI INT, i, i,
MPI COMM WORLD);
       endVertex = startVertex + num vetices remaining - 1;
      Send Dat[i][0] = startVertex;
      Send Dat[i][1] = endVertex;
      Send Dat[i][2] = num max colour;
      MPI Send(Send Dat[i], 3, MPI INT, i, i, MPI COMM WORLD);
      MPI Send(filled, vertices, MPI INT, i, i, MPI COMM WORLD);
      MPI Send (weights, vertices, MPI INT, i, i, MPI COMM WORLD);
MPI STATUS IGNORE);
      max = -1;
           if (filled[i] > max)
              max = filled[i];
```

```
printf("%d\n", max + 1);
           printf("%d\n", filled[i]);
       clock t final = clock();
       printf("Time taken in seconds is: %lf\n", ((double)(final/8 -
begin/8) /CLOCKS PER SEC));
      int vertices;
      MPI Recv(&vertices, 3, MPI INT, 0, rank usr, MPI COMM WORLD,
MPI STATUS IGNORE);
       int mat graph[vertices][vertices];
       for (i = 0; i < vertices; i++)
           MPI Recv(&mat graph[i], vertices, MPI INT, 0, rank usr,
MPI COMM WORLD, MPI STATUS IGNORE);
      MPI Request receiver req;
      MPI Status receiver status;
      MPI Recv(Rec Dat, 3, MPI INT, 0, rank usr, MPI COMM WORLD,
MPI STATUS IGNORE);
       int filled[vertices], weights[vertices];
       MPI Recv(filled, vertices, MPI INT, 0, rank usr, MPI COMM WORLD,
MPI STATUS IGNORE);
       MPI Recv(weights, vertices, MPI INT, 0, rank usr, MPI COMM WORLD,
MPI STATUS IGNORE);
       int round = 0;
      while (round < vertices) //outer loop</pre>
           int all filled = 1;
```

```
if (filled[vertex curr] == -1)
                   int weight_usr = weights[vertex_curr];
                   int adjacent;
                   for (adjacent = 0; adjacent < vertices; adjacent++)</pre>
                        if (adjacent != vertex curr)
                            if (mat graph[vertex curr][adjacent] == 1 &&
filled[adjacent] == -1 && weights[adjacent] > weights[vertex curr])
                        int colour filled = -1;
                        for (colour filled = 0; colour filled < Rec Dat[2];</pre>
colour filled++)
                            int avail col = 1;
                            for (adjacent = 0; adjacent < vertices;</pre>
adjacent++)
                                if (mat graph[vertex curr][adjacent] == 1
&& filled[adjacent] != -1 && colour_filled == filled[adjacent])
                                    avail col = 0;
```

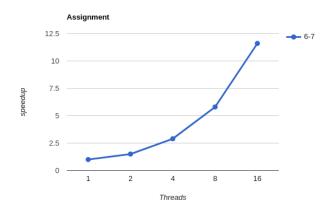
```
for (m = 1; m <= proc num - 1; m++)
                  MPI Send(filled, vertices, MPI INT, m, m,
MPI COMM WORLD);
           for (m = 1; m <= proc num - 1; m++)
               int col usr[vertices], weight curr usr[vertices];
MPI COMM WORLD, MPI STATUS IGNORE);
               int count = 0;
                   if (filled[count] == -1 && col usr[count] != -1 &&
col usr[count] < Rec Dat[2] && col usr[count] >= 0)
                       filled[count] = col usr[count];
           int count;
           round++;
           sleep(3);
```

MPI Send(filled, vertices, MPI INT, 0, rank usr, MPI COMM WORLD);

Graph for Speedup: 1 core-1, 2 core~1.51, 4 core~2.92 8 core~5.83, 16 core~11.64

Time taken in seconds is: 2.280828

nilay@Nilay-PC:~/Documents/cs359/Assignment\$



Time taken in seconds is: 4.519935

nilay@Nilay-PC:~/Documents/cs359/Assignment\$

# Description:

The algorithm used here is basically a greedy approach to the problem.

Create every conceivable combination of colors. The total number of possible color combinations is mV since each node can be coloured using any of the m-accessible colors. Verify whether or not the adjacent vertices share the same color after generating a color configuration. Print the combination and end the loop if the conditions are satisfied.

To solve the issue, adhere to the suggested instructions.

- Make a loop that accepts the output color array, the number of vertices, and the current index.
- Whether the number of vertices and the current index match. Verify that the output color configuration is secure. That is, ensure that no two neighboring vertices have the same color. Print the settings and break if the conditions are met.
- Give a vertex a color (1 to m).
- Loop for each color assigned, passing it to the following index and the number of vertices.
- Break the loop and return true if any loop returns true.

# Inference & Parallel Exploitation:

While the outermost loops can't parallelise because of communication overheads the inner loops can be run in parallel as the overheads then will weigh in as much as the other calculations. And the Program seems to give diminishing speedup which is common.

# Q2.

```
//Including Header files
#include <stdio.h>
#include <mpi.h>
#include <math.h>
#include <stdlib.h>
#include <time.h>
#define MAX 50

//DFS
int depth_search(int begin, int n, int mat_graph[MAX][MAX], int k)
{
   if (k == 0)
      return 1;
   int cnt_tot = 0, i = 1;
```

```
if (mat_graph[begin][i] == 1)
          cnt tot += depth search(i, n, mat graph, k - 1);
int main(int argc, char **argv)
  int proc num, rank;
  MPI Init(&argc, &argv);
  MPI Comm rank(MPI COMM WORLD, &rank);
  MPI Comm size (MPI COMM WORLD, &proc num);
  MPI Status status;
  int flag = 5, flag1 = 10;
  int walk tot = 0;
  int mat graph[MAX][MAX] = {0};
  if (rank == 0)
      scanf("%d %d %d", &n, &e, &k);
          s=rand()%n;
          d=rand()%n;
          mat graph[s][d] = 1;
      int num node on proc = n / (proc num - 1);
      int num node remain = n % (proc num - 1);
```

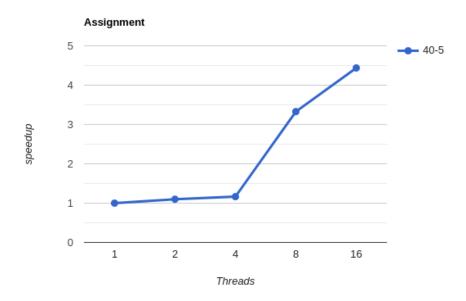
```
int boundl = 1, boundu = num node on proc;
      if (num node remain != 0)
           boundu = num node remain;
           for (int node = boundl; node <= boundu; ++node)</pre>
               walk tot += depth search(node, n, mat graph, k);
          bound1 = boundu + 1;
          boundu = bound1 + num node on proc - 1;
      for (int dest = 1; dest < proc num; dest++)</pre>
           MPI Send(&boundl, 1, MPI INT, dest, flag, MPI COMM WORLD);
          MPI Send(&boundu, 1, MPI INT, dest, flag + 1, MPI COMM WORLD);
          MPI Send(&k, 1, MPI INT, dest, flag + 2, MPI COMM WORLD);
          MPI Send(&n, 1, MPI INT, dest, flag + 3, MPI COMM WORLD);
          bound1 = boundu + 1;
          boundu += num node on proc;
  MPI Bcast(&mat graph, MAX * MAX, MPI INT, 0, MPI COMM WORLD);
  if (rank > 0)
      int boundl, boundu, k, n;
      MPI Recv(&boundl, 1, MPI INT, 0, flag, MPI COMM WORLD, &status);
      MPI Recv(&boundu, 1, MPI INT, 0, flag + 1, MPI COMM WORLD,
&status);
      MPI_Recv(&k, 1, MPI_INT, 0, flag + 2, MPI COMM WORLD, &status);
```

```
MPI Recv(&n, 1, MPI INT, 0, flag + 3, MPI COMM WORLD, &status);
      for (int node = boundl; node <= boundu; node++)</pre>
           int x = depth search(node, n, mat graph, k);
      MPI Send(&walk curr, 1, MPI INT, 0, flag1, MPI COMM WORLD);
  if (rank == 0)
      for (int src = 1; src < proc_num; ++src)</pre>
          MPI_Recv(&x, 1, MPI_INT, src, flag1, MPI_COMM_WORLD, &status);
      printf("%d\n", walk tot);
      printf("Time taken in seconds is: %lf\n", ((double)(end - start)
/CLOCKS PER SEC));
  MPI Finalize();
```

# Input/Output:

```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ mpiexec -n 2 ./mpi
 40 40 5
 27
 Time taken in seconds is: 0.000460
o nilay@Nilay-PC:~/Documents/cs359/Assignment$
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ mpiexec -n 4 ./mpi
 40 40 5
 27
 Time taken in seconds is: 0.000434
o nilay@Nilay-PC:~/Documents/cs359/Assignment$
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ mpiexec -n 8 ./mpi
 40 40 5
 27
 Time taken in seconds is: 0.000152
o nilay@Nilay-PC:~/Documents/cs359/Assignment$
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ mpiexec -n 16 ./mpi
 40 40 5
 27
 Time taken in seconds is: 0.000114
o nilay@Nilay-PC:~/Documents/cs359/Assignment$
```

Graph:1 core-1, 2 core~1.1, 4 core~1.16, 8 core~3.328, 16 core~4.43



### Description:

The algorithm is basically Depth First Search or DFS

From the source vertex, perform a depth-first search, counting the edges along the way.

- Check the node to see if you have reached the 'v'th node after traversing the K-edge; otherwise, leave that path and look for alternative options.
- Similarly, print the queue and look for any further pathways that might exist if we observe the 'v th node after the 'k' edges.

# Inference & Parallel Exploitation:

The Master Node first takes in input to randomly generate a graph. Then it allocates nodes to each process equally and remaining to the last one. After that it sets bounds and calculates the offsets. Lastly it sends all the data to the workers.

The Graph is Broadcasted to all the nodes.

The Worker Nodes receive data from the master node and the calculate walks for the given data and send the calculated result to the master

Lastly the master receives walks from the worker nodes and Calculates the final walks.

It can be seen that Better speedups aren't received till between 4 & 8 processors. This might be due to not enough parallelisation in 2 or 4 cores.

#### Q3.

```
#include <bits/stdc++.h>
#include <time.h>
#include <pthread.h>
using namespace std;
#define max 512
pthread mutex t mutex g;
int ans final = 0, n, thread num, n on thread;
int array g[10000][2] = \{0\};
void *thread func(void *s)
  int *index = (int *)s;
  int index usr = *(index);
  for (int i = 0; i < n on thread; ++i)
```

```
pthread mutex lock(&mutex g); //lock
          ans final += ((array g[index_usr + i][0] * array_g[index_usr +
i + 1][1]) - (array_g[index_usr + i][1] * array_g[index_usr + i + 1][0]));
          pthread mutex unlock(&mutex g);
  pthread exit(0);
int main(int argc, char **argv)
  pthread t p threads[max];
  clock t begin, end;
  pthread attr init(&attr);
  pthread attr setscope(&attr, PTHREAD SCOPE SYSTEM);
  pthread mutex init(&mutex g, NULL);
  n = 10000;
      array g[i][0] = rand() % 42;
      array g[i][1] = rand() % 42;
  array g[n][0] = array g[0][0];
  array g[n][1] = array g[0][1];
  begin = clock();
  double area sequential = 0;
```

```
area sequential += (array g[i][0] * array g[i + 1][1] -
array g[i][1] * array g[i + 1][0]);
   end = clock();
  double time sequential = (double(end - begin)) / CLOCKS PER SEC;
  cout << "Sequential Area = " << area sequential / 2 << endl;</pre>
   cout << "Time taken for Serial = " << time sequential << endl;</pre>
   cout << "Threads = ";</pre>
  cin >> thread num;
   n on thread = n / thread num;
  int *index usr = new int(0);
   int num index[thread num] = {0};
  begin = clock();
       *(index usr) += n on thread;
   for (int i = 0; i < thread_num; ++i)</pre>
       pthread create(&p threads[i], &attr, thread func, (void
*)(&num index[i]));
   for (int i = 0; i < thread num; ++i)</pre>
       pthread join(p threads[i], NULL);
  end = clock();
  double parallel time = (double(end - begin)) / CLOCKS PER SEC;
  cout << "Time taken by parallel = " << parallel time << endl;</pre>
  cout << "Speedup = " << double(time sequential / parallel time) <<</pre>
endl;
```

# Input/Output

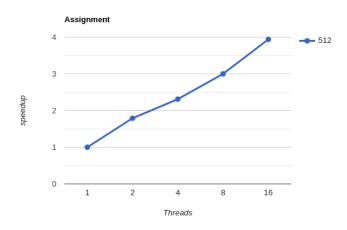
```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
Sequential Area = -1328.5
  Time taken for Serial = 0.0089
  Threads = 2
  Area Parallel = -1328
  Time taken by parallel = 0.004956
  Speedup = 1.7958
```

```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
Sequential Area = -1328.5
  Time taken for Serial = 0.0088
  Threads = 4
  Area Parallel = -1328
  Time taken by parallel = 0.003801
  Speedup = 2.31518
• nilay@Nilay-PC:~/Documents/cs359/Assignment$
```

```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
Sequential Area = -1328.5
  Time taken for Serial = 0.0088
  Threads = 8
  Area Parallel = -1328
  Time taken by parallel = 0.002925
  Speedup = 3.00855
• nilay@Nilay-PC:~/Documents/cs359/Assignment$
```

```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
Sequential Area = -1328.5
  Time taken for Serial = 0.0103
  Threads = 16
  Area Parallel = -1328
  Time taken by parallel = 0.002613
  Speedup = 3.94183
• nilay@Nilay-PC:~/Documents/cs359/Assignment$
```

Speedup Graph: 1 core-1, 2 core~1.79, 4 core~2.31, 8 core~3.01, 16 core~3.94



### Description:

A polygon can always be divided into triangles. The cross-product, which yields the area of a parallelogram, is taken, multiplied by 2, and used to calculate the (signed) area of the triangle with a vertex at the origin. This creates the area formula. These triangles with +ve and -ve areas will overlap as one loop around the polygon, canceling out and adding to zero the areas between the origin and the polygon, leaving only the area inside the reference triangle.

Area = 
$$|1/2[(x_1y_2 + x_2y_3 + ... + x_{n-1}y_n + x_ny_1) - (x_2y_1 + x_3y_2 + ... + x_ny_{n-1} + x_1y_n)]|$$

# Inference & Parallel Exploitation:

Polygon Clipping is used as parallel exploitation includes division of the whole polygon in parts and calculating area of the parts and then finally adding them up.

An Almost Linear Graph is obtained for Speedup This shows the nature of the algorithm in mathematical form and how the ratio of work increment and core increment (in this case exponentially power of 2) from 2 to 4 to 8 to 16 results in a linear fashion of datapoints.

#### Q4.

```
if (iter != dist)
              mat init[iter][dist] = rand() % 42;
  double start time = omp get wtime();
  for (pivot = 0; pivot < Num; pivot++)</pre>
      int *mat ptr = mat init[pivot];
      for (iter = 0; iter < Num; iter++)</pre>
           int *dp = mat init[iter];
               dp[dist] = min(dp[dist], dp[pivot] + mat ptr[dist]);
  double time = omp get wtime() - start time;
  printf("Sequential time = %.2f sec\n", time);
  double seqt=time;
  for (thread num = 2; thread num <= 16; thread num=thread num*2)</pre>
      omp set num threads(thread num);
      double start_time = omp_get_wtime();
#pragma omp parallel shared(mat init)
       for (pivot = 0; pivot < Num; pivot++)</pre>
```

# Input/Output:

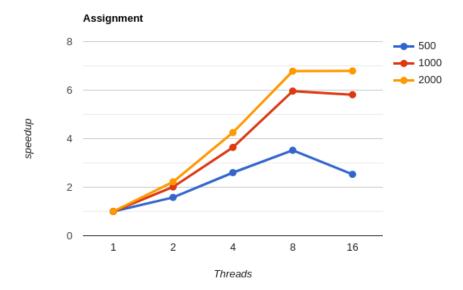
```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ gcc -fopenmp Q4.c
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
Sequential time = 0.39 sec
parallel time for 2 threads = 0.25 sec and speedup = 1.583333
parallel time for 4 threads = 0.15 sec and speedup = 2.606855
parallel time for 8 threads = 0.11 sec and speedup = 3.522010
parallel time for 16 threads = 0.16 sec and speedup = 2.533321
• nilay@Nilay-PC:~/Documents/cs359/Assignment$
```

```
    nilay@Nilay-PC:~/Documents/cs359/Assignment$ gcc -fopenmp Q4.c
    nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
    Sequential time = 5.58 sec
    parallel time for 2 threads = 2.78 sec and speedup = 2.003086
    parallel time for 4 threads = 1.53 sec and speedup = 3.643291
    parallel time for 8 threads = 0.93 sec and speedup = 5.967423
    parallel time for 16 threads = 0.96 sec and speedup = 5.806644
```

```
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ gcc -fopenmp Q4.c
• nilay@Nilay-PC:~/Documents/cs359/Assignment$ ./a.out
Sequential time = 13.41 sec
parallel time for 2 threads = 6.03 sec and speedup = 2.224819
parallel time for 4 threads = 3.15 sec and speedup = 4.252475
parallel time for 8 threads = 1.98 sec and speedup = 6.781714
parallel time for 16 threads = 1.97 sec and speedup = 6.798573
• nilay@Nilay-PC:~/Documents/cs359/Assignment$
```

# Speedup Graph:

Cores	1	2	4	8	16
n=500	1	1.58	2.60	3.52	2.53
n=1000	1	2.01	3.64	5.96	5.81
n=2000	1	2.22	4.25	6.78	6.79



### Description:

The Algorithm is Basically Floyd Warshall Algorithm.

- As a first step, initialize the solution matrix to be identical to the input graph matrix.
- The solution matrix is then updated by treating each vertex as an intermediate vertex.
- The plan is to select each vertex one at a time and update any shortest paths that use the selected vertex as an intermediate vertex.
- Vertices 0, 1, 2,.., k-1 are already considered when vertex number k is chosen as an intermediate vertex.
- There are two potential outcomes for every pair of source and destination vertices I j), respectively.
- The shortest path from I to j should not include k as an intermediary vertex. We maintain the current dist[i][j] value.

### Inference & Parallel Exploitation:

Because we are using OpenMP parallelism can be done by executing the "for" loop with the help of OpenMP directives

But in retrospect it is basically archived by running matrix multiplication in parallel as Floyd Warshal has elements of matrix multiplication

In the Speedup graph we can see that good speedup is achieved until 8 cores and then speedup either drops down or gets plateaued.

This may be due to communication overheads or maybe because the device used to execute the program has 8 Logical processors physically present as Hardware

Note: The programs may give different results depending on which hardware it is run on. It might also depend on the state of the same hardware such as if other programs are running or not or if the hardware is connected to power supply or getting DC from a battery source.