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**DWARKADAS J. SANGHVI COLLEGE OF ENGINEERING**

(Autonomous College Affiliated to the University of Mumbai)

NAAC Accredited with "A" Grade (CGPA : 3.18)



## DEPARTMENT OF INFORMATION TECHNOLOGY

**COURSE CODE: DJ19ITL502**

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**COURSE NAME: Advanced Data Structures Laboratory**

**CLASS/Div: A**

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### EXPERIMENT NO. 7

**CO/LO: LO1**

**AIM:** To Implement Perfect hashing

#### DESCRIPTION OF EXPERIMENT:

We say a hash function is perfect for S if all lookups involve  $O(1)$  work

This is used when the keys stored in the hash table are expected to be static. In this case perfect hashing guarantees excellent average as well as worst-case performance. This is how real dictionaries work. You most of the time just need to read from a dictionary.

This hashing technique is called perfect hashing because it takes constant time,  $O(1)$  to search for an item in the table.

**Perfect hashing** is implemented using two hash tables, one at each level. Each of the table uses universal hashing. The first level is the same a hashing with chaining such that n elements is hashed into m slots in the hash table. This is done using a has function selected from a universal family of hash functions.

The second level of uses a second hash table (instead of a linked list used in chaining). Elements that hash to the same slot j in the first hash table are stored in a second hash table. This second hash table is known as secondary hash table . The hash function  $h_j$  is carefully chosen such that there are no collision in the secondary table.

To ensure there are no collisions in the secondary hash table  $S_j$ , we need to make the size  $m_j$  of the secondary table equal to the square of the number of keys hashing into slot j in the first table. That is:

$$m_j = n_j^2$$

The hash functions for the primary hash table is carefully chosen so that we limit the expected total amount of space used to be  $O(n)$



**Technology stack used:** Python

**Program:**

```
import random

m=5
p=101
keys=[]
for i in range(0,m):
    keys.append(random.randint(1,2000))
collision_frequency={}
d={}
A=random.randint(1,2000)
B=random.randint(1,2000)

def gethash(a,b,p,e,m):
    #print(a,b,p,e,m)
    #print(((a*e+b)%p)%m)
    return ((a*e+b)%p)%m

for i in keys:
    h=gethash(A,B,p,i,m)
    try:
        collision_frequency[h]=collision_frequency[h]+1
    except:
        collision_frequency[h]=1

print(collision_frequency)
for i in collision_frequency.keys():
    M=collision_frequency[i]

    size=M*M
    l=[]
    for k in range(0,size+3):
        l.append(None)
    if size==1:
        a=0
        b=0
    else:
        a=random.randint(1,2000)
        b=random.randint(1,2000)
    l[0]=a
    l[1]=b
    l[2]=M
```



```
d[i]=1
#print(d)
for i in keys:
    h=gethash(A,B,p,i,m)
    a=d[h]
    h1=gethash(a[0],a[1],p,i,a[2]*a[2])
    print(i,h,h1+3)
    #print(h1)
    d[h][h1+3]=i

print(d)
```

### Output:

```
PS C:\Users\SHREE RAM\Desktop\ads> python -u "c:\Users\SHREE RAM\Desktop\ads\try.py"
{4: 2, 2: 2, 3: 1}
28 4 4
1490 4 6
650 2 3
1354 3 3
602 2 4
{4: [1208, 371, 2, None, 28, None, 1490], 2: [449, 294, 2, 650, 602, None, None], 3: [0, 0, 1, 1354]}
PS C:\Users\SHREE RAM\Desktop\ads>
```

### APPLICATIONS:

Some application of perfect hashing includes:

- data storage on a CD ROM
- set of reserved words in a programming language

### Conclusion:

Thus we have implemented perfect hashing.