## CENG 280 HOMEWORK 2

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## 1 Question 1

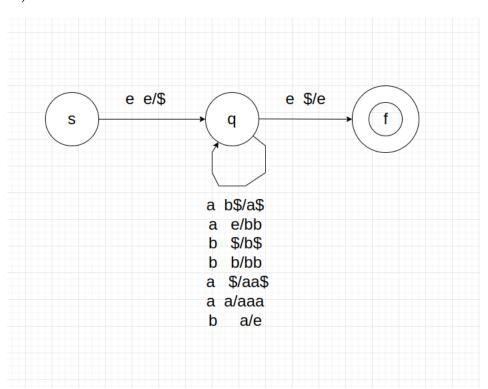
**a**)

 $S \rightarrow aSbSb|bSaSb|bSbSa|SS|e$ 

b)

 $S \to aSb|aaSb|e$ 

 $\mathbf{c})$ 



$$PDA \ P = (Q, \Sigma, \Gamma, s, \$, \{f\}, \Delta) \quad where$$

$$Q = \{s, q, f\}$$

$$\Sigma = \{a, b\}$$

$$\Gamma = \{a, b, \$\}$$

$$s = initial \ state$$

$$\{f\} \ is \ the \ set \ of \ final \ states$$

$$\Delta = \{((s, e, e), (q, \$)),$$

$$((q, b, \$), (q, b\$)),$$

$$((q, b, b), (q, b\$)),$$

$$((q, a, bb), (q, e)),$$

$$((q, a, b\$), (q, a\$)),$$

$$((q, a, b\$), (q, a\$)),$$

$$((q, a, a), (q, aa*)),$$

$$((q, a, a), (q, aaa)),$$

$$((q, b, a), (q, b)),$$

$$((q, b, a), (q, b)),$$

$$((q, b, a), (q, b)),$$

d) Let start variables for L1 nad L2 be S1 and S2 respectively. Also let start variable for L3 be S. We can define a grammar for their union as

$$S \to S_1 | S_2$$

since context free languages are closed under union. Hence context free grammar for L3 can be written as follows :

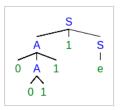
$$S \rightarrow S_1|S_2$$
 
$$S_1 \rightarrow abS_1b|bS_1aS_1b|bS_1bS_1a|S_1S_1|e$$
 
$$S_2 \rightarrow aS_2b|aaS_2b|e$$

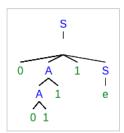
## 2 Question 2

a) For string 00111 there are two different parse trees as shown above.

$$S \rightarrow A1S \rightarrow 0A11S \rightarrow 00111S \rightarrow 00111$$

$$S \rightarrow 0A1S \rightarrow 0A11S \rightarrow 00111S \rightarrow 00111$$





Since there are two different parse trees for the same string 00111 , the grammar is ambiguous.

b)

$$S \rightarrow AS|e$$
 
$$A \rightarrow B1$$
 
$$B \rightarrow 0C1|01$$
 
$$C \rightarrow 01$$

c) 
$$S \rightarrow AS \rightarrow B1S \rightarrow 0C11S \rightarrow 00111S \rightarrow 00111$$

