

Q1 Team Name

0 Points

flash

Q2 Commands

10 Points

List the commands used in the game to reach the ciphertext

go
enter
pick
c
back
give
back
back
thrnxtzy
read

Q3 Analysis

50 Points

Give a detailed analysis of how you figured out the password? (Explain in less than 500 words)

We were presented with a screen that had some hints and equations related to multiplicative groups. It is written that "password for this level is an element of the multiplicative group Z_p^* " where p is a prime. g is also an element of this group. This gave us the idea to use Modular arithmetic for the analysis.

Given, the prime $p = 455470209427676832372575348833$. Considering the multiplicative group under modulo p .

Using the equations of the form $(a, \text{password} * g^a)$ given in the hint:

$$\begin{aligned} \text{password} * g^{429} &= 431955503618234519808008749742 (\text{say} = x_1) \bmod p \quad \text{---} > (1) \\ \text{password} * g^{1973} &= 176325509039323911968355873643 (\text{say} = x_2) \bmod p \quad \text{---} > (2) \\ \text{password} * g^{7596} &= 98486971404861992487294722613 (\text{say} = x_3) \bmod p \quad \text{---} > (3) \end{aligned}$$

Dividing (2) by (1), we get,

$$g^{1544} = x_2 * x_1^{-1} \bmod p \quad \text{---} > (4)$$

Dividing (3) by (1), we get,

$$g^{7167} = x_3 * x_1^{-1} \bmod p \quad \text{---} > (5)$$

Dividing (3) by (2), we get,

$$g^{5623} = x_3 * x_2^{-1} \bmod p \quad \text{---} > (6)$$

In order to perform modular division, we need to find the modular inverse of the denominators $(x_1 \text{ and } x_2)$, if it exists.

From properties of the multiplicative group of integers modulo n , we know that $\gcd(x_1, p) = 1$ and $\gcd(x_2, p) = 1$ that means x_1 and x_2 are co-prime to p . Therefore, the modular inverse of x_1 and x_2 exists under modulo p .

Modular inverse of x_1 is a number y_1 such that $(x_1 * y_1) \bmod p = 1$ (identity of the multiplicative group).

Since it is given that p is a prime number, therefore we can use Fermat's little theorem to compute the modular inverse.

Therefore using Fermat's little theorem which states that,

$$a^{p-1} \equiv 1 \pmod{p}$$

Multiplying both sides by a^{-1} we get,

$$a^{p-2} \equiv a^{-1} \pmod{p}$$

After rearranging the above equation we get,

$$a^{-1} \equiv a^{p-2} \pmod{p}$$

Using the above equation to find inverse of x_1 and x_2 we get

,

$$x_1^{-1} = y_1 = x_1^{p-2} \pmod{p} = 70749996790223471732904681640 \pmod{p}$$

$$x_2^{-1} = y_2 = x_2^{p-2} \pmod{p} = 228947149478752602606353685125 \pmod{p}$$

Substituting the value of x_2 and x_1^{-1} in (4) we get,

$$g^{1544} = 176325509039323911968355873643 * 70749996790223471732904681640 \pmod{p}$$

$$g^{1544} = 111590994894663139264552154672 \pmod{p} - \text{---} > (7)$$

In the similar way solving (5) and (6) we get,

$$g^{7167} = 110411376670918912626907526185 \pmod{p} - \text{---} > (8)$$

$$g^{5623} = 420413074251022028027270785553 \pmod{p} - \text{---} > (9)$$

Multiplying both sides of (9) by inverse of $(g^{1544})^3$ gives,

$$g^{5623} * ((g^{1544})^3)^{-1} = 420413074251022028027270785553 * ((g^{1544})^3)^{-1} \pmod{p}$$

$$g^{991} = 161798558270556961732424822635 \pmod{p}$$

In the similar way, the following calculations can be carried out

$$g^{230} = g^{7167} * ((g^{991})^7)^{-1} \pmod{p} =$$

$$263509268584013168241508095725(modp)$$

$$g^{164} = g^{1544} * ((g^{230})^6)^{-1}(modp) = 365054265331735818534782577890(modp)$$

$$g^7 = g^{991} * ((g^{164})^6)^{-1}(modp) = 296995893405083794209776051694(modp)$$

$$g^6 = g^{991} * ((g^7)^{32})^{-1}(modp) = 313908172363422710658462184069(modp)$$

$$g = g^7 * (g^6)^{-1}(modp) = 52565085417963311027694339(modp)$$

In the hint we were given that g looks like

5__50__4____31____94__9. The g we computed looks similar to it. Thus it gives us confirmation that we have computed g correctly.

Now for computing password

Multiplying $(g^{429})^{-1}$ on both sides of (1) we have,

$$password * g^{429} * (g^{429})^{-1} = 431955503618234519808008749742 * (g^{429})^{-1}(modp)$$

For the multiplicative group $a * a^{-1} = 1$

$$So, password = 431955503618234519808008749742 * (g^{429})^{-1}(modp)$$

Computing $(g^{429})^{-1}$ using fermat's little theorem and substituting it we get,

$$password = 431955503618234519808008749742 * 442956820316148690889301696615(modp)$$

$$Therefore, password = 134721542097659029845273957$$

We then used this password to clear the level.

We used python to implement Fermat's little theorem to find the modular inverse and do all the computation required.

Q4 Password

10 Points

What was the final command used to clear this level?

134721542097659029845273957

Q5 Codes

0 Points

Upload any code that you have used to solve this level

▼ Crypto_Assignment3.ipynb

 Download

Implementation of Fermat's little theorem

In [76]:

```
# Function to compute gcd of two
number a and b

def gcd(a,b):
    if(b==0):
        return a
    else:
        return gcd(b, a % b)
```

In [77]:

```
# Function to compute x^y under modulo
p

def power_under_modulo(x,y,p):
    if(y==0):
        return 1

    a=power_under_modulo(x,y//2,p)%p
    a=(a*a)%p
    if(y%2==0):
        return a
    else:
        return (x*a)%p
```

In [78]:

```
# Function to find modular inverse of
a under modulo p

def modInverse(a, p):
    if(gcd(a,p)!= 1):
        print("Inverse of",a,"doesn't
exist !!")
    else:
        q=power_under_modulo(a,p-2, p)
        print("Modular multiplicative
```

```
inverse of",a,"is :",q)
return q
```

In [79]:

```
x1=431955503618234519808008749742
x2=176325509039323911968355873643
x3=98486971404861992487294722613
p=455470209427676832372575348833
```

In [80]:

```
x1_inv=modInverse(x1,p)
x2_inv=modInverse(x2,p)
```

Modular multiplicative inverse of 431955503
Modular multiplicative inverse of 176325509

In [81]:

```
g_1544=(x2*x1_inv)%p
print("g_1544 :",g_1544)
```

g_1544 : 111590994894663139264552154672

In [82]:

```
g_7167=(x3*x1_inv)%p
print("g_7167 :",g_7167)
```

g_7167 : 110411376670918912626907526185

In [83]:

```
g_5623=(x3*x2_inv)%p
print("g_5623 :",g_5623)
```

g_5623 : 420413074251022028027270785553

In [84]:

```
g_991=
(g_5623*modInverse(power_under_modulo(g_1544,p)))%p
print("g_991 :",g_991)
```

Modular multiplicative inverse of 345360603
g_991 : 161798558270556961732424822635

In [85]:

```
g_230=
(g_7167*modInverse(power_under_modulo(g_991,p)))%p
print("g_230 :",g_230)
```

Modular multiplicative inverse of 451843559
g_230 : 263509268584013168241508095725

In [86]:

```
g_164=
(g_1544*modInverse(power_under_modulo(g_230,p)))%p
```

```
print("g_164 :",g_164)
```

Modular multiplicative inverse of 199249906
g_164 : 365054265331735818534782577890

In [87]:

```
g_7=  
(g_991*modInverse(power_under_modulo(g_164,  
print("g_7 :",g_7)
```

Modular multiplicative inverse of 398578713
g_7 : 296995893405083794209776051694

In [88]:

```
g_6=  
(g_230*modInverse(power_under_modulo(g_7,32  
print("g_6 :",g_6)
```

Modular multiplicative inverse of 303462035
g_6 : 313908172363422710658462184069

In [89]:

```
g=(g_7*modInverse(g_6,p))%p  
print("The value of g is :",g)
```

Modular multiplicative inverse of 313908172
The value of g is : 52565085417963311027694

In [90]:

```
password=  
(x1*modInverse(power_under_modulo(g,429,p),  
print("The password is :",password)
```

Modular multiplicative inverse of 190938394
The password is : 1347215420976590298452739

Assignment 3



GROUP

Ravi Shankar Das
Nimit Jain
Prashant Mishra
 [View or edit group](#)

TOTAL POINTS
70 / 70 pts

QUESTION 1
[Team Name](#) **0 / 0 pts**

QUESTION 2
[Commands](#) **10 / 10 pts**

QUESTION 3
[Analysis](#) **50 / 50 pts**

QUESTION 4
[Password](#) **10 / 10 pts**

QUESTION 5
[Codes](#) **0 / 0 pts**