

# **CS 466: Design and Analysis of Algorithms**

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Notes written from Mark Petrick's lectures.

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# 1 Travelling Salesman Problem

## 1.1 Definition

Given a graph  $G = (V, E)$  with weights on edges  $w : E \rightarrow \mathbb{R}^{+ve \cup \{0\}}$ , find a **TSP Tour**:

- a cycle  $C$  that visits every vertex exactly once
- has  $\min \sum_{e \in C} w(e)$

This is NP-Complete

## 1.2 Metric TSP

Create an approximation problem with TSP by defining distances in a space:

- symmetry:  $d(u, v) = d(v, u)$
- triangle inequality:  $d(u, v) \leq d(u, w) + d(w, v)$

### Approximation Algorithm (1977)

- find a minimum spanning tree and walk the tour
- take shortcuts to avoid re-visiting the same vertices
- by triangle inequality, the length is  $\leq 2 * l_{MST}$  and therefore total  $l \leq 2 * l_{TSP}$  (2-approx)
- ends up with runtime of MST,  $O(m \log n)$

## 1.3 Further Approximation

- TSP General: no constant factor approx
- Metric space: 1.5-approx, proven limit is 1.0045
- Euclidean space:  $(1 + \epsilon)$ -approx in polytime

# 2 Binomial Heap

## 2.1 Data Structures

### Priority Queue

- store  $n$  elements each w/ integer key
- insert, deleteMin, decreaseKey, build (heapify)

### Heap

- binary tree of elements
- structure (min key at root for min-heap)

- shape (almost complete binary tree, store in array)
- insert (+ bubbleUp)  $O(\log n)$
- deleteMin  $O(\log n)$
- decreaseKey  $O(\log n)$
- heapify  $O(n)$
- merge  $O(n)$

## 2.2 Binomial Tree

### 2.2.1 Definition

- binomial tree of order 0 is a single node
- binomial tree of order  $k$  has a root node and its children are the binomial trees of order  $k - 1 \dots 0$

### 2.2.2 Claims

- $B_k$  has  $2^k$  nodes  
 $2^{k-1} + 2^{k-1} = 2^k$
- the height of  $B_k$  is  $k$   
 $k - 1 + 1 = k$
- there are  $\binom{k}{i}$  nodes at depth  $i$  in  $B_k$   
 $\binom{k-1}{i} + \binom{k-1}{k-i} = \binom{k}{i}$

## 2.3 Binomial Heap

### 2.3.1 Definition

collection of binomial trees at most one of each rank (each one uses one heap order)

**Example 2.1.** Make a binomial tree of size  $n = 13$

Expressing  $n$  in binary tells us what trees to build.  $13 = 1101$  so build  $k = 3, k = 2, k = 0$