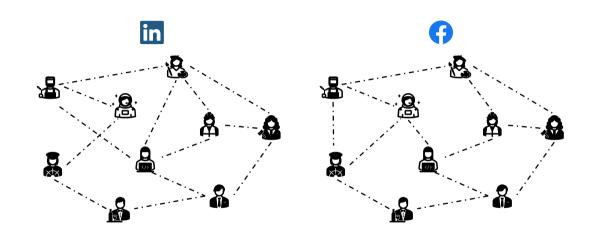
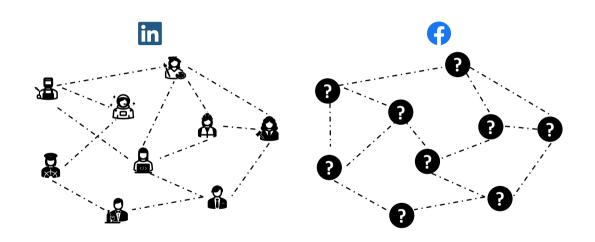
Attributed Graph Alignment

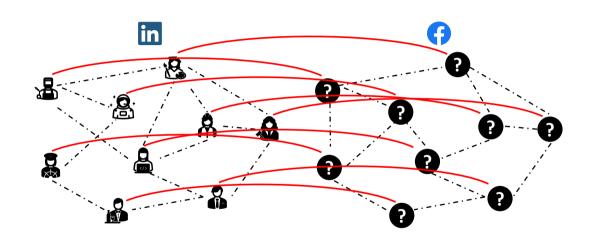
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Ning Zhang * Weina Wang † Lele Wang *
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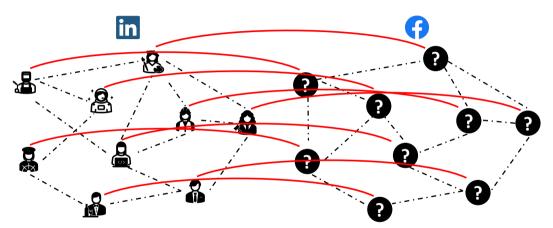
*University of British Columbia †Carnegie Mellon University

International Symposium on Information Theory (ISIT)
June 2021









Other applications: computer vision, biological network analysis, natural language processing...

Related works

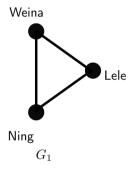
• Information-theoretic limit:

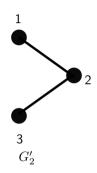
Pedarsani and Grossglauser (2011), Cullina and Kiyavash (2016), Cullina and Kiyavash (2017), Wu et al. (2021)

• Efficient algorithms:

- ► Convex optimization: Lyzinski et al. (2015);
- ► Spectral method: Nassar et al. (2018), Feizi et al. (2019), Fan et al. (2020) etc.;
- ▶ Belief propagation: Onaran and Villar (2020)
- ▶ Others: Barak et al. (2019)

What if graph structure is not enough?





We do know more



Ning Zhang

Intro

- Studies ECE at University of British Columbia
- Studied Physics Major at Nankai University, China
- Lives in Vancouver, British Columbia
- From Shijiazhuang Shi, Hebei, China
- Joined February 2017

Edit Details

A Hiking

Badminton



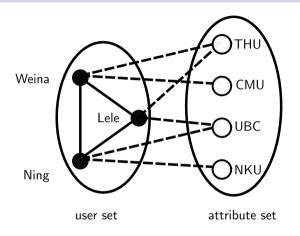
Ning Zhang

Information theory; MASc@UBC; Vancouver

We do know more ...

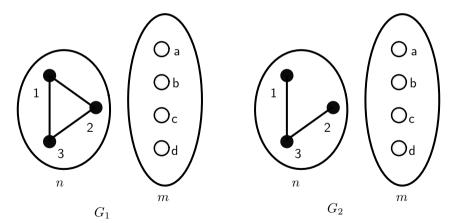


Attributes as vertices

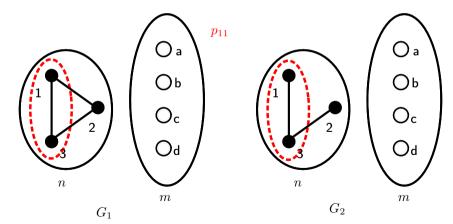


How much benefit can vertex attributes bring?

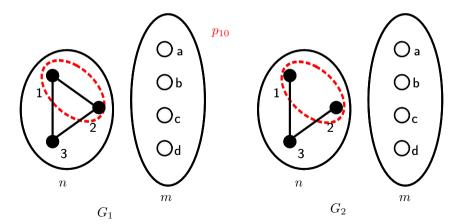
- $(G_1, G_2) \sim \mathcal{G}(n, p; m, q)$
 - ▶ For $e \in \mathcal{E}_u$, $(\mathbb{1}\{e \in E_1\}, \mathbb{1}\{e \in E_2\}) \stackrel{\mathsf{i.i.d}}{\sim} p$, where $p = (p_{11}, p_{10}, p_{01}, p_{00})$.



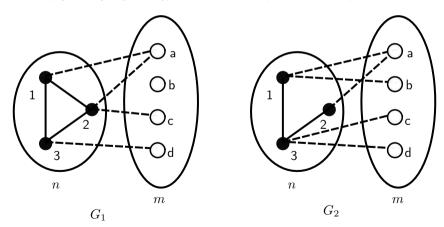
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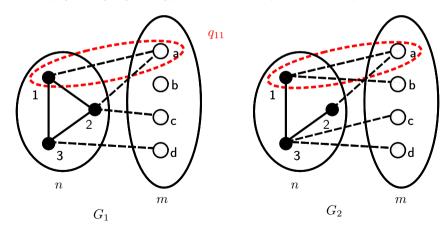
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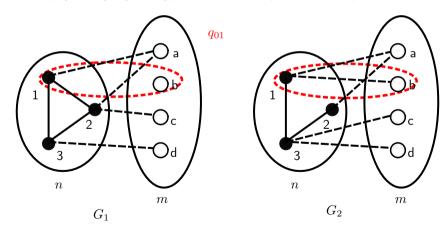
- $(G_1, G_2) \sim \mathcal{G}(n, \boldsymbol{p}; m, \boldsymbol{q})$
 - ▶ For $e \in \mathcal{E}_u$, $(\mathbb{1}\{e \in E_1\}, \mathbb{1}\{e \in E_2\}) \stackrel{\mathsf{i.i.d}}{\sim} p$, where $p = (p_{11}, p_{10}, p_{01}, p_{00})$.
 - ▶ For $e \in \mathcal{E}_a$, $(\mathbb{1}\{e \in E_1\}, \mathbb{1}\{e \in E_2\}) \stackrel{\text{i.i.d}}{\sim} \mathbf{q}$, where $\mathbf{q} = (q_{11}, q_{10}, q_{01}, q_{00})$.



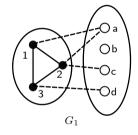
- $(G_1, G_2) \sim \mathcal{G}(n, \boldsymbol{p}; m, \boldsymbol{q})$
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 - ▶ For $e \in \mathcal{E}_a$, $(\mathbb{1}\{e \in E_1\}, \mathbb{1}\{e \in E_2\}) \stackrel{\text{i.i.d}}{\sim} q$, where $q = (q_{11}, q_{10}, q_{01}, q_{00})$.

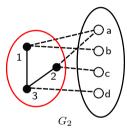


- $(G_1, G_2) \sim \mathcal{G}(n, \boldsymbol{p}; m, \boldsymbol{q})$
 - ▶ For $e \in \mathcal{E}_u$, $(\mathbb{1}\{e \in E_1\}, \mathbb{1}\{e \in E_2\}) \stackrel{\text{i.i.d}}{\sim} \boldsymbol{p}$, where $\boldsymbol{p} = (p_{11}, p_{10}, p_{01}, p_{00})$.
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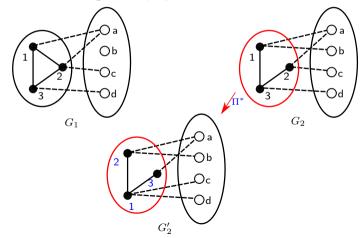


- $(G_1, G_2) \sim \mathcal{G}(n, p; m, q)$
- Anonymization: Obtain G_2' by applying random Π^* on user vertex set

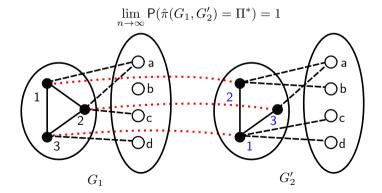




- $(G_1, G_2) \sim \mathcal{G}(n, p; m, q)$
- Anonymization: Obtain G_2' by applying random Π^* on user vertex set



- $(G_1, G_2) \sim \mathcal{G}(n, \boldsymbol{p}; m, \boldsymbol{q})$
- Anonymization: Obtain G_2' by applying random Π^* on user vertex set
- Exact alignment: Recover $\hat{\pi}(G_1, G_2')$ s.t.



- $(G_1, G_2) \sim \mathcal{G}(n, \boldsymbol{p}; m, \boldsymbol{q})$
- Anonymization: Obtain G'_2 by applying random Π^* on user vertex set
- Exact alignment: Recover $\hat{\pi}(G_1, G_2')$ s.t.

$$\lim_{n\to\infty}\mathsf{P}(\hat{\pi}(G_1,G_2')=\Pi^*)=1$$

• Our goal: characterize the information-theoretic limit

Achievability

Conditions on (n, p, m, q) to achieve exact alignment;

Converse

Conditions on (n, p, m, q) s.t. exact alignment is not possible.

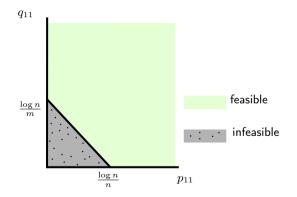
Main results

Achievability (simplified)

- $m = \Omega((\log n)^3)$: $np_{11} + mq_{11} \log n \to \infty$
- $m = o((\log n)^3)$: $np_{11} + mq_{11} - O(mq_{11}^{3/2}) - \log n \to \infty$

Converse

 $np_{11} + mq_{11} - \log n \rightarrow -\infty$



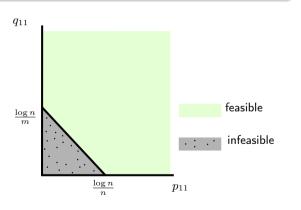
Main results

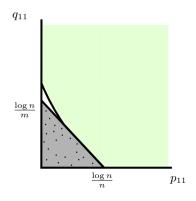
Achievability (simplified)

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- $m = o((\log n)^3)$: $np_{11} + mq_{11} - O(mq_{11}^{3/2}) - \log n \to \infty$

Converse

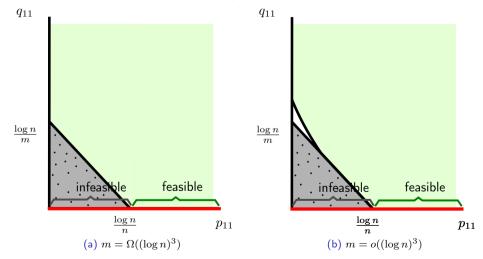
$$np_{11} + mq_{11} - \log n \to -\infty$$





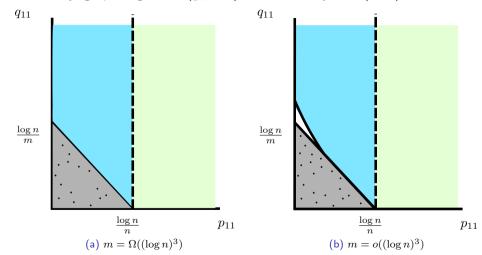
Recover the Erdős-Rényi graph alignment

1. Erdős–Rényi graph alignment $(q_{00} = 1)$ Cullina and Kiyavash (2017)



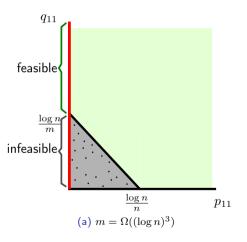
Recover the Erdős-Rényi graph alignment

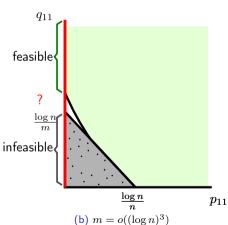
1. Erdős-Rényi graph alignment $(q_{00}=1)$ Cullina and Kiyavash (2017)



Improve the bipartite graph alignment

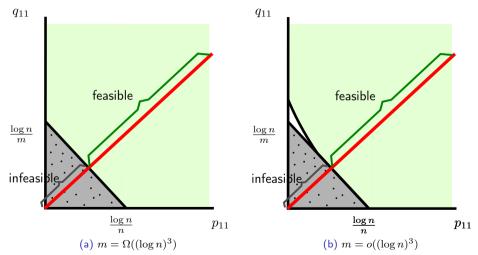
2. Bipartite graph alignment $(p_{00} = 1)$ Dai et al. (2019)





Improve the seeded graph alignment

3. Seeded graph alignment (p=q) Yartseva and Grossglauser (2013)



Proof of achievability (sketch)

Achievability

- $m = \Omega((\log n)^3)$: $np_{11} + mq_{11} \log n \to \infty$
- $m = o((\log n)^3)$: $np_{11} + mq_{11} O(mq_{11}^{3/2}) \log n \to \infty$

Key ideas:

• MAP estimator (see Lemma1)

Proof of achievability (sketch)

Achievability

- $m = \Omega((\log n)^3)$: $np_{11} + mq_{11} \log n \to \infty$
- $m = o((\log n)^3)$: $np_{11} + mq_{11} O(mq_{11}^{3/2}) \log n \to \infty$

Key ideas:

- MAP estimator (see Lemma1)
- Upper bound error probability of MAP (Cullina and Kiyavash (2017)):
 - Orbit decomposition
 - Generating functions

Converse

$$np_{11} + mq_{11} - \log n \to -\infty$$

Key ideas:

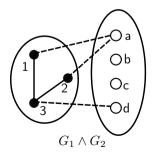
•
$$P(\hat{\pi}_{MAP}(G_1, G'_2) = \Pi^*) \le \frac{1}{|Aut(G_1 \wedge G_2)|}$$

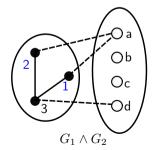
Converse

$$np_{11} + mq_{11} - \log n \to -\infty$$

Key ideas:

• $P(\hat{\pi}_{MAP}(G_1, G'_2) = \Pi^*) \le \frac{1}{|Aut(G_1 \wedge G_2)|}$



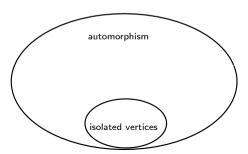


Converse

$$np_{11} + mq_{11} - \log n \to -\infty$$

Key ideas:

- $P(\hat{\pi}_{MAP}(G_1, G'_2) = \Pi^*) \le \frac{1}{|Aut(G_1 \wedge G_2)|}$
- $|Aut| \ge |Aut_{iso}|$

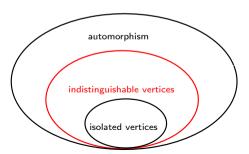


Converse

$$np_{11} + mq_{11} - \log n \to -\infty$$

Key ideas:

- $P(\hat{\pi}_{MAP}(G_1, G'_2) = \Pi^*) \le \frac{1}{|Aut(G_1 \wedge G_2)|}$
- $|Aut| \ge |Aut_{ind}| \ge |Aut_{iso}|$

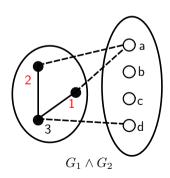


Converse

$$np_{11} + mq_{11} - \log n \to -\infty$$

Key ideas:

- $P(\hat{\pi}_{MAP}(G_1, G'_2) = \Pi^*) \le \frac{1}{|Aut(G_1 \wedge G_2)|}$
- $|Aut| \ge |Aut_{ind}| \ge |Aut_{iso}|$



 $|Aut_{iso}| = 0$ $|Aut_{ind}| = 1$

Converse

$$np_{11} + mq_{11} - \log n \to -\infty$$

Key ideas:

- $P(\hat{\pi}_{MAP}(G_1, G'_2) = \Pi^*) \le \frac{1}{|Aut(G_1 \wedge G_2)|}$
- $|Aut| \ge |Aut_{ind}| \ge |Aut_{iso}|$
- Use the second moment method to show that w.h.p. $|\operatorname{Aut}(G_1 \wedge G_2)| \geq 2$ In particular, we show # indistinguishable vertex pair >0.

Contributions

- Propose attributed Erdős–Rényi graph pair model
- Characterize the information-theoretic limits
- Unify existing models for graph alignment

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