NTH FIBONACCI NUMBER

1. DESCRIPTION

The Fibonacci numbers are a sequence of integers defined by the recurrence relation $F_n = F_{n-1} + F_{n-2}$, where F_n is the n^{th} number in the series and $F_0 = 0$ and $F_1 = 1$. This sequence appears frequently in mathematics, computer science, and even biology, and has been described in mathematical texts for centuries.

2. APPLICATIONS

- A. Golden Ratio F_n/F_{n-1} approaches φ as n approaches ∞
- B. Fibonacci heap data structure for priority queues
- C. Hilbert's Tenth Problem Fibonacci numbers used to show unsolvability
- D. Bee ancestry bee reproduction creates an unusual number of ancestors
- E. Brock-Mirman model a generalized sequence is used in an optimal control function
- F. Fibonacci Quarterly & the Fibonacci Association publishing scholarly work since 1963

3. COMPETING ALGORITHMS

A. Recursion

- i. Directly implement recurrence relation $F_n = F_{n-1} + F_{n-2}$, base cases $F_0 = 0$ and $F_1 = 1$
- ii. Creates a recursion tree of height n where each level, L, has at most 2^L sub problems
- iii. $T(n) = T(n-1) + T(n-2) \rightarrow O(2^n)$

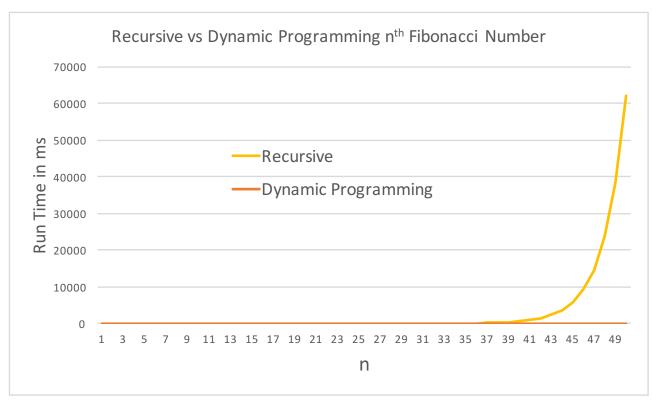
B. Dynamic Programming

- i. Store the previously calculated F_{n-1} , F_{n-2} in an array, starting with $F_0 = 0$ and $F_1 = 1$
- ii. Add F_n to the array by summing the top 2 elements only
- iii. 1 for loop of n 1 elements \rightarrow 0(n)

4. EXPERIMENTS

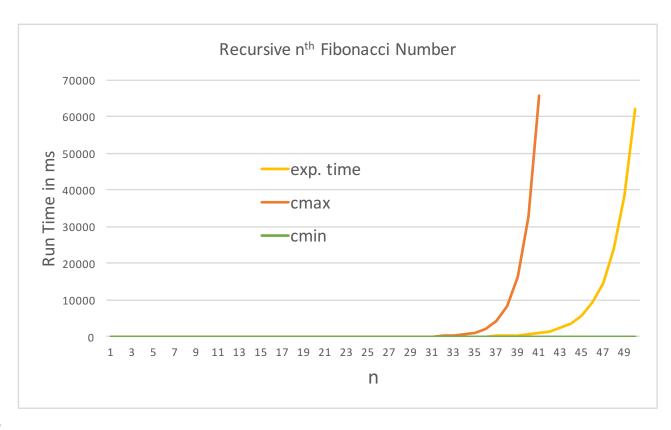
A. Direct comparison

Because the runtimes of the recursive method grow exponentially, the two algorithms can only be directly compared for relatively small values of n. The graph below shows the two algorithms' runtimes in milliseconds from n=1 up to n=50. The recursive calculation has a theoretical runtime of $O(2^n)$, while dynamic programming has a theoretical runtime of O(n). Experimental runtimes are shown on the graph below.



B. Recursive calculation

Fibonacci numbers up to F_{50} (n = 50) were calculated and the runtime for each calculation was recorded in milliseconds. For theoretical runtimes, F_0 and F_1 are assumed to take 2^n = 1 and 2ms, respectively. By comparing the experimental runtime to the theoretical runtime, the hidden constant c = (experimental RT / theoretical RT) was calculated. This was then used to bound the experimental runtime, such that $c_{min}2^n$ <= experimental runtime <= $c_{max}2^n$. This function was then graphed for n = 1 to 50.

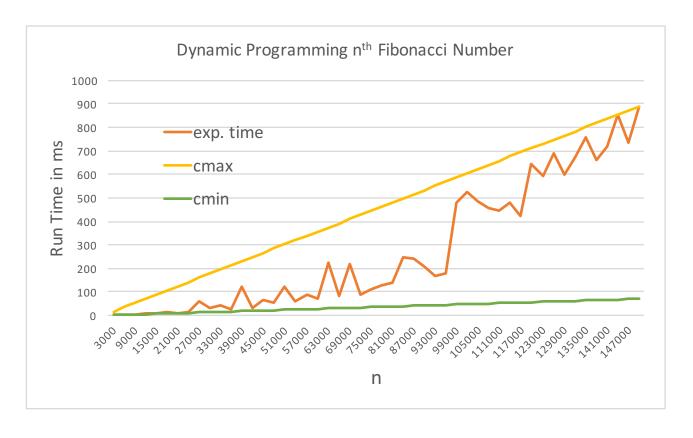


		Recursive			C
n	value	exp. time	thr. time	С	2.98
1	1	0	2	0	5.96
2	1	0	4	0	1.19
3	2	0	8	0	2.38
4	3	0	16	0	4.76
5	5	0	32	0	9.53
6	8	0	64	0	1.90
7	13	0	128	0	3.8
8	21	0	256	0	7.62
9	34	0	512	0	1.52
10	55	0	1024	0	3.05
11	89	0	2048	0	6.10
12	144	0	4096	0	0.00
13	233	0	8192	0	0.00
14	377	0	16384	0	0.000
15	610	0	32768	0	0.000
16	987	0	65536	0	0.00
17	1597	0	131072	0	0.00
18	2584	0	262144	0	0.0
19	4181	0	524288	0	0
20	6765	0	1048576	0	
21	10946	0	2097152	0	0.063
22	17711	0	4194304	0	0.124
23	28657	0	8388608	0	0.249
24	46368	0	16777216	0	0.499
25	75025	1	33554432	2.98E-08	0.999
26	121393	1	67108864	1.49E-08	1.999
27	196418	1	134217728	7.45E-09	3.999
28	317811	1	268435456	3.73E-09	7.99
29	514229	3	536870912	5.59E-09	15.99
30	832040	4	1073741824	3.73E-09	31.99
31	1346269	7	2147483648	3.26E-09	63.99
32	2178309	11	4294967296	2.56E-09	127.
33	3524578	18	8589934592	2.10E-09	255.9
34	5702887	29	1.718E+10	1.69E-09	511.9
35	9227465	47	3.436E+10	1.37E-09	1023
36	14930352	77	6.8719E+10	1.12E-09	2047
37	24157817	122	1.3744E+11	8.88E-10	4095
38	39088169	215	2.7488E+11	7.82E-10	8191
39	63245986	321	5.4976E+11	5.84E-10	1638
40	102334155	518	1.0995E+12	4.71E-10	3276
41	165580141	857	2.199E+12	3.90E-10	6553
42	267914296	1382	4.398E+12	3.14E-10	1310
43	433494437	2439		2.77E-10	2621
44	701408733	3702		2.10E-10	5242
45	1134903170	5615		1.60E-10	1048
	1836311903	9326		1.33E-10	2097
47	2971215073	14534		1.03E-10	4194
	4807526976	23867	2.8147E+14	8.48E-11	8388
	7778742049	38306		6.80E-11	1677
50		61877		5.50E-11	3355

C _{max}	C _{min}
2.98023E-08	C
5.96046E-08	C
1.19209E-07	C
2.38419E-07	C
4.76837E-07	C
9.53674E-07	C
1.90735E-06	C
3.8147E-06	C
7.62939E-06	C
1.52588E-05	
3.05176E-05	C
6.10352E-05	C
0.00012207	C
0.000244141	C
0.000488281	C
0.000976562	C
0.001953125	C
0.00390625	C
0.0078125	C
0.015625	C
0.03125	C
0.062499999	C
0.124999998	C
0.249999997	C
0.499999993	C
0.99999987	C
1.999999974	C
3.999999948	C
7.999999896	C
15.99999979	C
31.99999958	C
63.99999917	C
127.9999983	C
255.9999967	C
511.9999933	C
1023.999987	C
2047.999973	C
4095.999947	C
8191.999893	C
16383.99979	C
32767.99957	C
65535.99915	C
131071.9983	C
262143.9966	C
524287.9932	C
1048575.986	C
2097151.973	C
4194303.945	C
8388607.891	C
16777215.78	C
33554431.56	C

C. Dynamic Programming calculation

The same methods were applied to the dynamic programming algorithm, but because runtimes grew slowly, change would not be evident unless very large values of n were used. In this case, Fibonacci numbers up to $F_{150,000}$ (n = 150,000) were calculated, starting at n = 3,000 and stepping by 3,000, giving a total of 50 values. For theoretical runtimes, $F_{3,000}$ and $F_{6,000}$ are assumed to take n = 3,000 and 6,000ms, respectively. All other calculation and graphing methods are the same as those used for the recursive calculation, however Fibonacci values are thousands of digits long, so they are truncated in the table.



3000 4106158863 3 3000 0.001 17.770833 1 6000 3770131493 5 6000 8.33E-04 35.541666 2 9000 3461602912 5 9000 5.56E-04 53.312499 4 12000 3178322757 8 12000 6.67E-04 71.083332 5 15000 2918224824 9 15000 6.00E-04 88.854165 7 18000 2679411996 15 18000 8.33E-04 106.624998 8	2.85714282 4.28571423 5.71428564 7.14285705
3000 4106158863 3 3000 0.001 17.770833 1 6000 3770131493 5 6000 8.33E-04 35.541666 2 9000 3461602912 5 9000 5.56E-04 53.312499 4 12000 3178322757 8 12000 6.67E-04 71.083332 5 15000 2918224824 9 15000 6.00E-04 88.854165 7 18000 2679411996 15 18000 8.33E-04 106.624998 8	1.42857141 2.85714282 4.28571423 5.71428564 7.14285705 8.57142846 9.99999987 11.4285713
6000 3770131493 5 6000 8.33E-04 35.541666 2 9000 3461602912 5 9000 5.56E-04 53.312499 4 12000 3178322757 8 12000 6.67E-04 71.083332 5 15000 2918224824 9 15000 6.00E-04 88.854165 7 18000 2679411996 15 18000 8.33E-04 106.624998 8	2.85714282 4.28571423 5.71428564 7.14285705 8.57142846 9.99999987 11.4285713 12.8571427
9000 3461602912 5 9000 5.56E-04 53.312499 4 12000 3178322757 8 12000 6.67E-04 71.083332 5 15000 2918224824 9 15000 6.00E-04 88.854165 7 18000 2679411996 15 18000 8.33E-04 106.624998 8	4.28571423 5.71428564 7.14285705 8.57142846 9.99999987 11.4285713 12.8571427
12000 3178322757 8 12000 6.67E-04 71.083332 5 15000 2918224824 9 15000 6.00E-04 88.854165 7 18000 2679411996 15 18000 8.33E-04 106.624998 8	5.71428564 7.14285705 8.57142846 9.9999987 11.4285713 12.8571427
15000 2918224824 9 15000 6.00E-04 88.854165 7 18000 2679411996 15 18000 8.33E-04 106.624998 8	7.14285705 8.57142846 9.9999987 11.4285713 12.8571427
18000 2679411996 15 18000 8.33E-04 106.624998 8	8.57142846 9.99999987 11.4285713 12.8571427
	9.99999987 11.4285713 12.8571427
21000 2460142407 10 21000 4.76E-04 124.395831 9	11.4285713 12.8571427
	12.8571427
	15.7142855
36000 1605328799 25 36000 6.94E-04 213.249996 1	
	18.5714283
	19.9999997
45000 1242585401 66 45000 0.00146667 266.562495 2	
	22.8571426
51000 1047533080 124 51000 0.00243137 302.104161	24.285714
	25.7142854
	27.1428568
	28.5714282
	29.9999996
66000 6835519025	31.428571
	32.8571424
	34.2857138
	35.7142853
78000 4857966282 128 78000 0.00164103 462.041658 3	
	38.5714281
	39.9999995
87000 3760250225	
90000 3452530277 208 90000 0.00231111 533.12499 4	
93000 3169992581 170 93000 0.00182796 550.895823 4	
96000 2910576347 181 96000 0.00188542 568.666656 4	
99000 2672389432 478 99000 0.00482828 586.437489 4	
102000 2453694535 523 102000 0.00512745 604.208322 4	
105000 2252896527 483 105000 0.0046 621.979155 4	
108000 2068530817 455 108000 0.00421296 639.749988 5	
111000 1899252669 445 111000 0.00400901 657.520821 5	52.8571422
114000 1743827392 481 114000 0.0042193 675.291654 5	
	55.714285
	57.1428564
123000 1349788566 590 123000 0.00479675 728.604153 5	58.5714278
126000 1239328665 688 126000 0.00546032 746.374986 5	59.9999992
129000 1137908247 600 129000 0.00465116 764.145819 6	61.4285706
132000 1044787565 671 132000 0.00508333 781.916652	62.857142
135000 9592874106 759 135000 0.00562222 799.687485 6	64.2857135
138000 8807841583 661 138000 0.00478986 817.458318 6	65.7142849
141000 8087052171 719 141000 0.00509929 835.229151 6	67.1428563
144000 7425248535 853 144000 0.00592361 852.999984 6	68.5 <mark>71</mark> 4277
147000 6817603576 734 147000 0.0049932 870.770817 6	69.999991
150000 6259685222 887 150000 0.00591333 888.54165 7	71.4285705

5. PROGRAMMING LANGUAGE

A. Java

- i. For handling extremely large numbers, the Long and BigInteger data types must be used
- ii. Output to console is too cumbersome, results are saved in a .csv file
- iii. Runtimes are determined by taking the current time before and after the calculation, then recording the difference
- iv. Smallest runtime granularity is milliseconds

6. SOURCES

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