# JSR 292 on Android

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### Us

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# Why?

- We like dynamic language runtime devs
  - No need to keep two versions of a runtime (pre-invokedynamic/post-invokedynamic)
- We like dynamic language users
  - Some languages only run on the JVM
- We like Java
  - Java 8 (lambda) uses invokedynamic

# but first a demo!

# BTW, we're lying

We don't implement JSR 292 on Android because Dalvik is not a Java VM

We first need to adapt the JSR 292 to Android

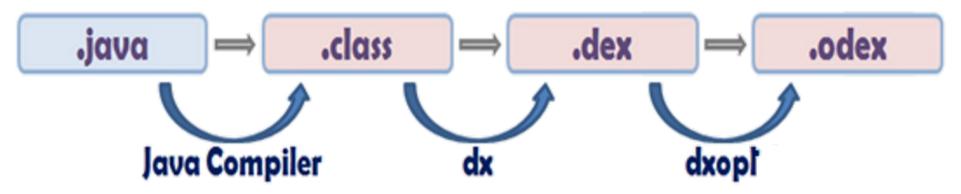
- Dalvik opcodes are register based
- Better spec (obviously!)

and then implement it

Smartphones/Tablets are constraint environments

# A Glimpse to Android

#### **Toolchain**



#### The DEX file

- is read-only (mapped)
- contains several classes (more like a jar)
- one constant pool by type (String, Prototype, ...)
- optimized ODEX
   verified offline (at install time)

### Dalvik

- Register based opcodes
- Several interpreters
  - Portable, armv\*, x86
- A 'java' dev kit (libcore) written from scratch
- An incremental non-generational GC
- A very small JIT (for benchmark)

# JSR 292 implementation for Android

#### A new DEX format

- A new header
- 5 new instructions
  - invoke-exact/invoke-generic/invoke-dynamic
  - const-methodtype/const-methodhandle
- 4 new constant pools
  - methodType/methodHandle/invokeDynamic/bsmArgs

# new invoke opcodes

- invoke-exact/range: 40/42 invoke-generic/range: 41/43
  - opcode {vC, vD, vE, vF, vG} mt@BBBB
  - opcode {vCCCC...vNNNN} mt@BBBB
- invoke-dynamic/range: 73/79

  - opcode {vCCCC...vNNNN} indy@BBBB #ZZZZZZZZZ

callsite number

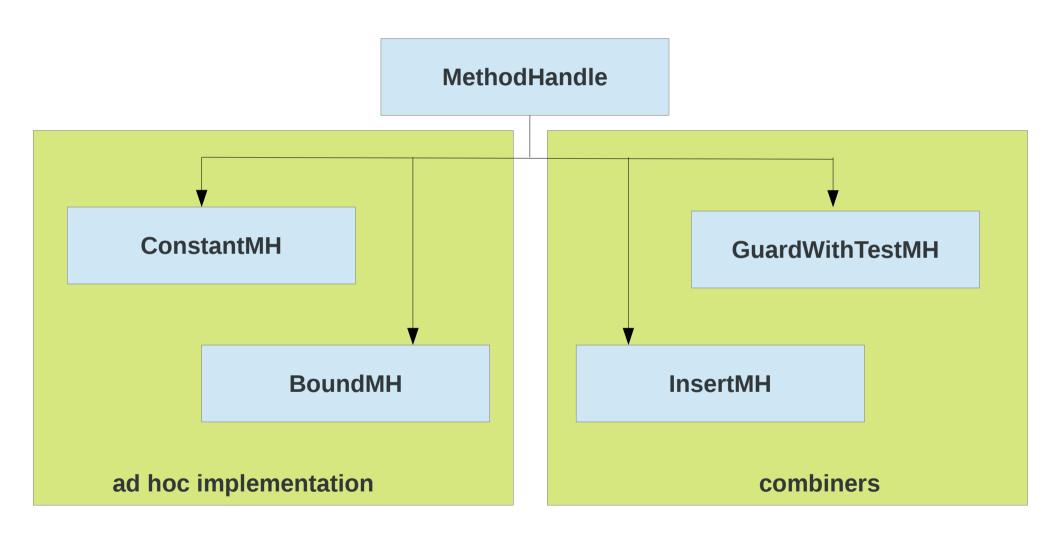
# MethodType

- Interned in a concurrent weak hashmap in Java
- Opaque pointer in C
- invoke-exact/invoke-generic have an offset in a table of resolved MethodType
- Invoke-dynamic and const-methodType doesn't intern it

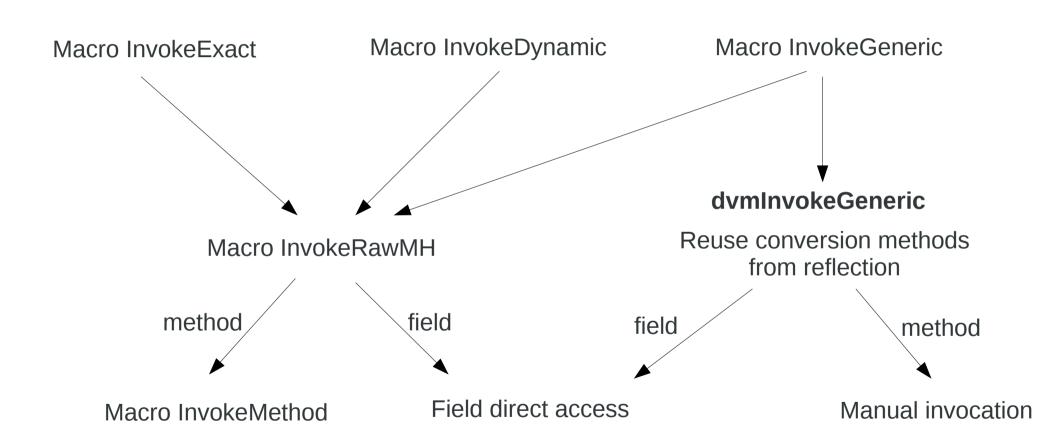
### MethodHandle

- 3 rings of MethodHandle
  - Constant method handle
  - Method Handle for lambdas (bind/constant)
  - Method Handle combiners
- In memory
  - kind (field:get/set\*instance/static\*volatile +
     method:static/direct/super/virtual/interface/newInstance +
     bind + constant +
     AST-interpreter + mini-interpreter (combiner))
  - slot (field index, directMethod index, vtable/itable index)
  - declaringClass
  - methodType

# MethodHandle hierarchy



# Macros and functions in the interpreter



### InvokeExact

method handle

- 40 {vC, vD, vE, vF, vG} mt@BBBB42 {vCCCC...vNNNN} mt@BBBB

#### Implementation scketch:

get methodHandle (vC)

rebuild arguments (drop vC)

get methodType (mt@BBBB) get methodType from methodHandle check method types

if false then throw WrongMethodTypeException!

call the macro InvokeRawMH

### InvokeGeneric

method handle

- 41 {vC, vD, vE, vF, vG} mt@BBBB43 {vCCCC...vNNNN} mt@BBBB

```
Implementation scketch:
```

```
get methodHandle (vC)
```

rebuild arguments (drop vC)

get methodType (mt@BBBB) get methodType from methodHandle check method types if true then like invokeExact! else then

> convert arguments (boxing, unboxing, cast, ...) create a new stack frame and push arguments

# InvokeDynamic

- 73 {vC, vD, vE, vF, vG} indy@BBBB #ZZZZZZZZ
- 79 {vCCCC...vNNNN} indy@BBBB #ZZZZZZZZ

```
Implementation scketch:
 get the current class
 get call site index (#ZZZZZZZ)
 get call site (currentClass → callSites[callsite idx])
   if NULL then
       get indy reference (indy@BBBB)
       get method handle, method type and bsm arguments (indy)
      java call: call_bootstrap(mh, mt, bsm_args)
          in java: create lookup and invoke-exact? the bootstrap method
      intern the call site (CAS)
call the macro InvokeRawMH (callSite → target) // a volatile read!
```

# Garbage Collector interaction

- Some opcodes use Java references
  - invoke-dynamic access to a CallSite
  - invoke-exact/generic access to a MethodType
- MethodTypes for invoke-exact/generic are interned, stored as strong reference in the global hashtable (in Java)
- CallSites are stored in ClassObject (VM) and are roots of the GC marking phase



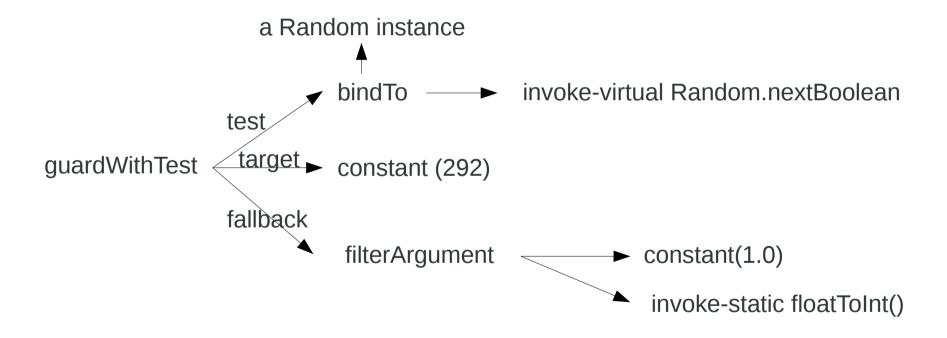
### MethodHandle combiners

- Combiners are small reusable bricks, language runtimes tend to use a bunch of them
- Some method handles are called once, some are called in hot-loop

- Two implementations
  - In Java (AST interpreter)
  - In C, with a specific stack frame (mini-interpreter)

# Example of combiners

(random.nextBoolean())? 292: (int) 1.0;



# AST interpreter (Java)

- Special method handle kind(AST)
- Upcall to a non-public Java method, box all arguments, unbox return value

```
class MethodHandle {
  /*package*/ Object invokeAST(Object[] args) {
   ...
} }
```

# AST interpreter (Java)

```
class GWTMH extends MethodHandle {
 private final MethodHandle test;
 private final MethodHandle target;
 private final MethodHandle fallback;
 @Override
 /*package*/ Object invokeAST(Object[] args) {
  if ((Boolean)test.invokeAST(args)) {
   return target.invokeAST(args);
  return fallback.invokeAST(args);
 }}
                            class MethodHandle {
                             /*package*/ Object invokeAST(Object[] args) {
                               switch(args.length) {
                                case 0: return invoke();
                                case 1: return invoke(arg[0]);
```

# Mini-interpreter

 Generate a specific code corresponding to the whole method handle tree

(by walking the AST)

Do the execution in one stack frame

The whole stack frame needs to be fully typed (at generation time)

=> use index indirection

A specific interpreter inside the interpreter?

# Mini-interpreter example

```
gwt = MHs.guardWithTest(test1, target, fb1);

$\forall \test1 = \test1 = \text{MHs.insertArgument(test0, 0, new Random());}

$\forall \test0 = \text{MHs.lookup().findVirtual(Random.class, "nextBoolean", MT(boolean.class));}

$\forall \text{target} = \text{MHs.constant(int.class, 292);}

$\forall \text{fb1} = \text{MHs.filterArguments(fb0, 0, f2i);}

$\forall \text{fb0} = \text{MHs.constant(double.class, 1.0);}

$\forall \text{f2i} = \text{MHs.lookup().findStatic(Convertions.class, "floatToInt", MT(int.class, float.class);}

$\text{gwt.nb_r = 1} \text{gwt.nb_rv = 2}

$\text{gwt.mhs} = \{ \text{test1, test0, target, fb0, f2i} \}
```

r	rv1/32 or bd2	0
16) /	1 int	1
rv	true or false boolean	0
	code + 8	
	0x	
а		0
	MethodHandle	
		1

```
code {
 const r0 bd0
                           // ldc Random object
 invoke-frame-mh rv0 m1 r0 1 // call test0
 ifne rv0 :else
                                // ldc 292
 const r0 bd2
 goto :end
 :else
 const r0 bd3
                                // ldc 1.0
 invoke-frame-mh rv1 m4 r0 1 // call f2i
 const r0 rv1
 :end
 invoke-frame-mh 0 null r0 1
                                // return
```

# Mini-interpreter Stack Frame

#### Stack

Registers (32 bits integers)

Return Values (known types)

Mini Interpreter Program Counter Code Reference (Java)

MH arguments

MethodHandle

Slots size and type are known at the generation time

32/64bits

Register number representation

# Mini-Interpreter Twist

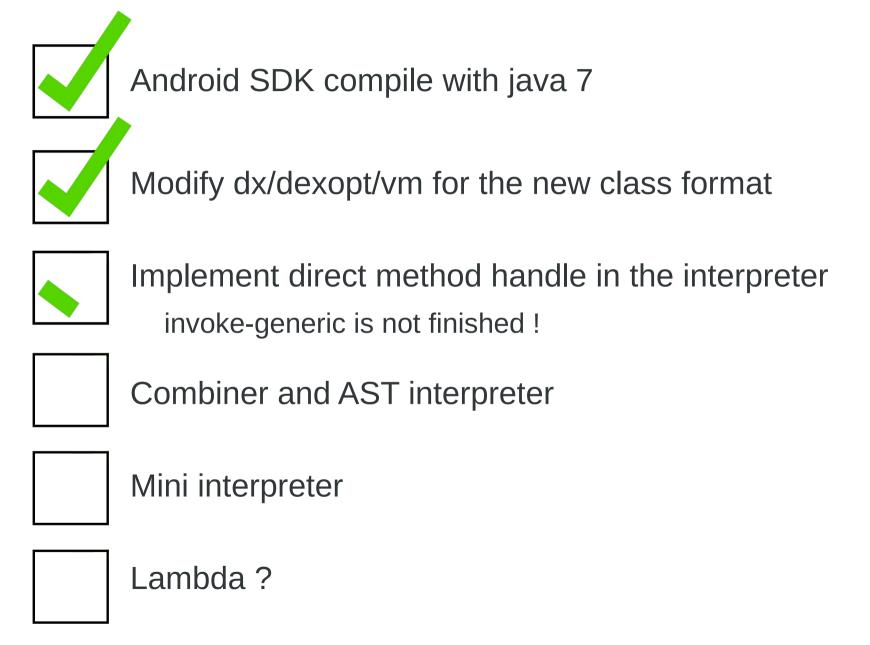
Davik is already a register based VM

So we can re-use the same opcodes

So we can re-use the same interpreter

We need one supplementary opcode
invoke-frame-mh

### Where we are?



# Question?

https://bitbucket.org/jpilliet/android-sdk-292

