

R in Java

FastR: an implementation of the R language



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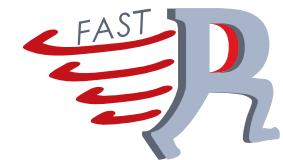
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<https://github.com/allr>





What we do...

- **TimeR** — an instrumentation-based profiler for GNU-R
- **TracR** — a trace analysis framework for GNU-R
- **CoreR** — a formal semantics for a fragment of R
- **TestR** — a testing framework for the R language
- **FastR** — a new R virtual machine written in Java

$\frac{\mathbf{e} \Gamma; H \rightarrow \mathbf{e}' H'}{C[\mathbf{e}] \Gamma * S; H \implies C[\mathbf{e}'] \Gamma * S; H'} \quad \frac{H(\delta) = \mathbf{e} \Gamma'}{C[\delta] \Gamma * S; H \implies \mathbf{e} \Gamma' * C[\delta] \Gamma * S; H}$	$\frac{}{[EXP]}$	$\frac{}{[FORCEP]}$
$\frac{\mathbf{getfun}(H, \Gamma, x) = \delta}{C[x(\bar{a})] \Gamma * S; H \implies \delta \Gamma * C[x(\bar{a})] \Gamma * S; H} \quad \frac{\mathbf{getfun}(H, \Gamma, x) = \nu}{C[x(\bar{a})] \Gamma * S; H \implies C[\nu(\bar{a})] \Gamma * S; H}$	$\frac{}{[FORCEF]}$	$\frac{}{[GETF]}$
$\frac{H(\nu) = \lambda \bar{f}. \mathbf{e}, \Gamma' \quad \mathbf{args}(\bar{f}, \bar{a}, \Gamma, \Gamma', H) = F, \Gamma'', H'}{C[\nu(\bar{a})] \Gamma * S; H \implies \mathbf{e} \Gamma'' * C[\nu(\bar{a})] \Gamma * S; H'}$	$\frac{}{[INVF]}$	$\frac{}{[RETP]}$
$H' = H[\delta / \nu]$		
$R[\nu] \Gamma' * C[\delta] \Gamma * S; H \implies C[\delta] \Gamma * S; H' \quad R[\nu] \Gamma'' * C[\nu'(\bar{a})] \Gamma * S; H \implies C[\nu] \Gamma * S; H'$		

Evaluation Contexts:

```

C ::= [] | x <- C | x[[C]] | x[[e]] <- C | x[[C]] <- ν | {C; e} | {ν; C}
    | attr(C, e) | attr(ν, C) | attr(e, e) <- C | attr(C, e) <- ν | attr(ν, C) <- ν
R ::= [] | {ν; R}

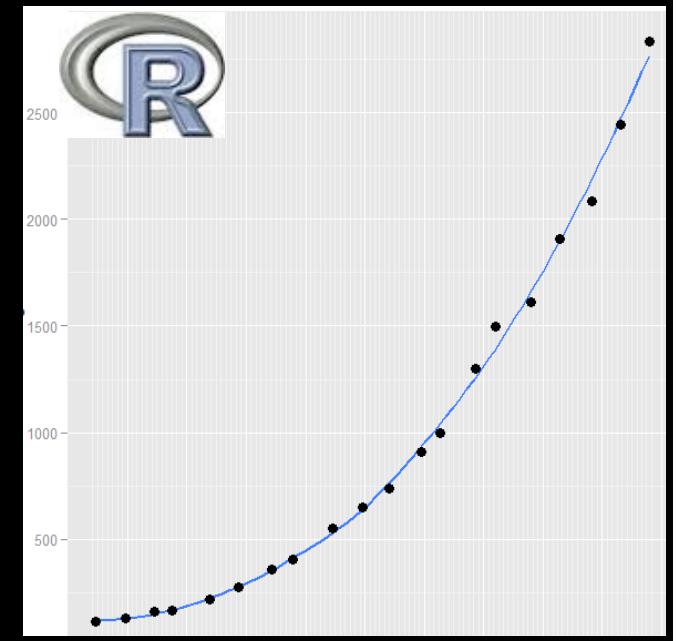
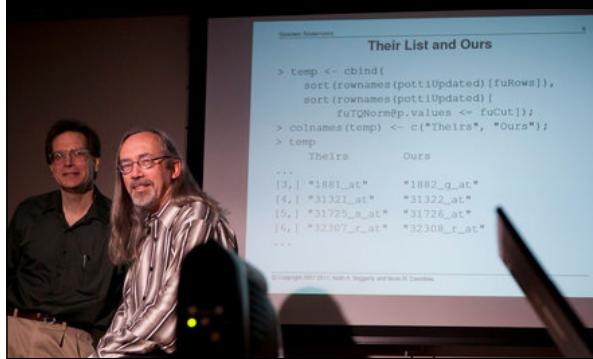
```

$\frac{\nu \text{ fresh} \quad \alpha = \perp \perp}{H' = H[\nu / \text{mem}^{\perp \perp}_\alpha]} \quad \frac{\nu \text{ fresh} \quad \alpha = \perp \perp}{H' = H[\nu / \text{str}[\alpha^\perp]]} \quad \frac{\nu \text{ fresh} \quad \alpha = \perp \perp}{H' = H[\nu / \lambda \mathbf{e}. \Gamma]}$	$\frac{\nu \text{ fresh} \quad \alpha = \perp \perp}{\pi \Gamma; H \rightarrow \nu; H'} \quad \frac{\nu \text{ fresh} \quad \alpha = \perp \perp}{\sigma \Gamma; H \rightarrow \nu; H'} \quad \frac{\nu \text{ fresh} \quad \alpha = \perp \perp}{\text{function}(\mathbf{e}; \Gamma; H \rightarrow \nu; H')}$	$\frac{}{[\MEM]}$
$\frac{\Gamma(H, x) = \nu}{x \Gamma; H \rightarrow \nu; H'} \quad \frac{H(\delta) = \nu}{\delta \Gamma; H \rightarrow \nu; H'}$		
$\frac{}{\text{cpy}(H, \nu) = H', \nu'} \quad \frac{\Gamma = \star \Gamma' \quad H(\star) = F \quad F' = F[\mathbf{g}/\nu']}{\text{cpy}(H, \nu) = H', \nu'} \quad \frac{}{[\ASS]}$	$\frac{}{\Gamma(\mathbf{g}, \nu) = H'} \quad \frac{}{[\ASS]}$	
$\frac{}{\mathbf{x}[\nu] \Gamma; H \rightarrow \nu; H'}$	$\frac{\mathbf{x}[\nu] \Gamma; H \rightarrow \nu; H'}{\text{assign}(\mathbf{x}, \nu', \Gamma', H') = H''}$	
$\frac{}{\Gamma(H, x) = \nu'} \quad \frac{}{\text{readin}(\nu, H) = \mathbf{n}}$	$\frac{}{\Gamma(H, \mathbf{x}) = \nu'} \quad \frac{\mathbf{x}[\nu] \Gamma; H \rightarrow \nu'; H'}{\text{readin}(\nu, H) = \mathbf{n}} \quad \frac{}{[\GET]}$	
$\frac{}{\text{cpy}(H, \nu') = H', \nu'} \quad \frac{\Gamma = \star \Gamma' \quad \mathbf{x}(\nu', H, \mathbf{x}) = \nu''}{\text{cpy}(H, \nu') = H', \nu'} \quad \frac{}{[\SETL]}$	$\frac{}{\text{readin}(\nu, H') = \mathbf{n}} \quad \frac{\text{set}(\nu', \mathbf{n}, \nu'', H') = H''}{\text{readin}(\nu, H') = \mathbf{n}} \quad \frac{}{[\SETL]}$	
$\frac{}{\text{cpy}(H, \nu') = H', \nu'} \quad \frac{\Gamma = \star \Gamma' \quad H(\star) = F \quad \mathbf{x} \notin F \quad \Gamma'(H', \mathbf{x}) = \nu''}{\text{cpy}(H, \nu') = H', \nu'} \quad \frac{}{[\SETG]}$	$\frac{}{\text{readin}(\nu, H') = \mathbf{n}} \quad \frac{P' = F[\mathbf{x}/\nu'''] \quad H''' = H''[\nu'/F']}{\text{readin}(\nu, H') = \mathbf{n}} \quad \frac{}{[\SETG]}$	
$\frac{}{H(\nu) = \kappa^\alpha \quad \alpha = \nu_\perp \nu'_\perp} \quad \frac{\mathbf{index}(\nu', \nu'_\perp, H) = \mathbf{n}}{\mathbf{attr}(\nu, \nu'_\perp) \Gamma; H \rightarrow \nu'^\perp; H'}$	$\frac{}{H(\nu) = \kappa^\alpha \quad \alpha = \nu_\perp \nu'_\perp} \quad \frac{\mathbf{index}(\nu', \nu'_\perp, H) = \mathbf{n}}{\mathbf{attr}(\nu, \nu'_\perp) \Gamma; H \rightarrow \nu'^\perp; H'} \quad \frac{}{[\GETA]}$	
$\frac{}{\text{cpy}(H, \nu') = H', \nu'''} \quad \frac{H'(\nu) = \kappa^{+\perp} \quad \mathbf{index}(\nu', \nu'_\perp, H) = \perp}{\text{attr}(\nu, \nu'_\perp) \Gamma; H \rightarrow \nu'^\perp; H'}$	$\frac{}{\text{attr}(\nu, \nu'_\perp) \Gamma; H \rightarrow \nu'^\perp; H'} \quad \frac{}{[\GETA]}$	
$\frac{}{H'(\nu_\perp) = \mathbf{gen}[\nu]} \quad \frac{H'(\nu') = \text{str}[\nu]}{\text{attr}(\nu, \nu') \Gamma; H \rightarrow \nu'^\perp; H'}$	$\frac{}{\text{attr}(\nu, \nu') \Gamma; H \rightarrow \nu'^\perp; H'} \quad \frac{}{[\GETB]}$	
$\frac{}{\text{cpy}(H, \nu') = H', \nu^3} \quad \frac{H'(\nu) = \kappa^{\perp \perp} \quad \nu^3 \text{ fresh} \quad \mathbf{reads}(\nu', H') = \mathbf{s}}{\text{attr}(\nu, \nu') \Gamma; H \rightarrow \nu'^\perp; H'}$	$\frac{}{\text{attr}(\nu, \nu') \Gamma; H \rightarrow \nu'^\perp; H'} \quad \frac{}{[\GETB]}$	



Why?

- ... language for data analysis and graphics
- ... used in statistics, biology, finance ...
- ... books, conferences, user groups
- ... 4,338 packages
- ... 3 millions users



Scripting data

read data into variables

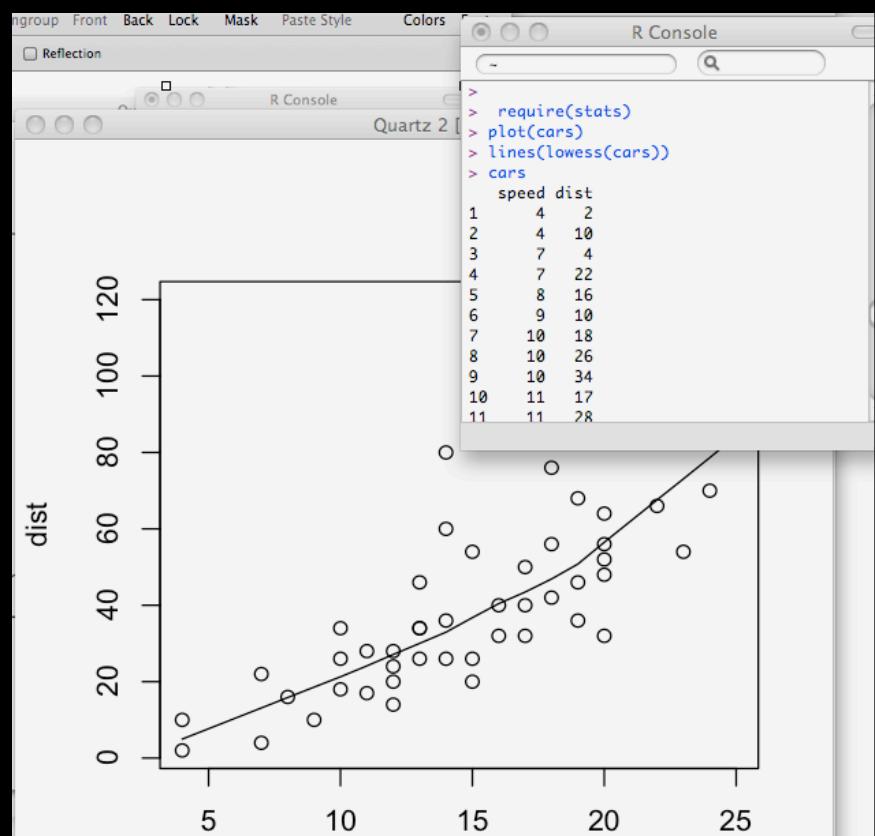
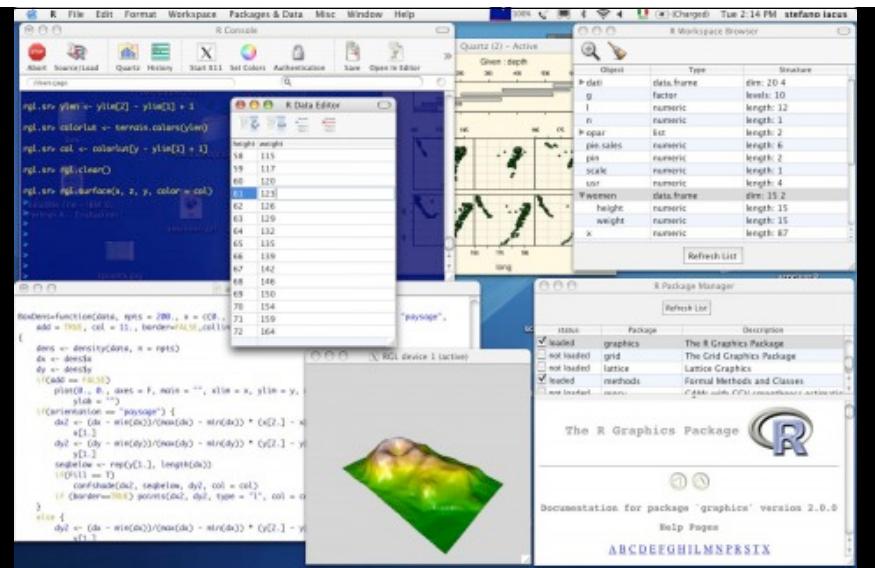
make plots

compute summaries

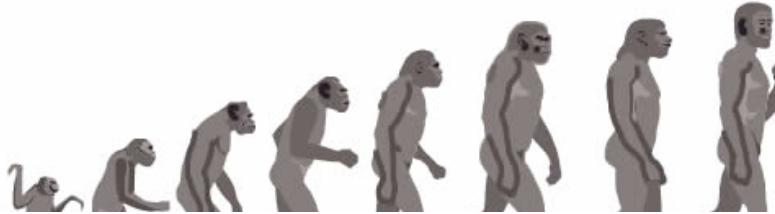
more intricate modeling

develop simple functions
to automate analysis

...



R history



- 1976 S

John Chambers @ Bell Labs, then S-Plus
(closed-source owned by Tibco)

- 1993 R

Ross Ihaka and Robert Gentleman,
started R as new language at the
University of Auckland, NZ

- Today, The R project



<http://www.r-project.org>
<http://cran.r-project.org>

Core team ~ 20 people, released under
GPL license. Continued development of
language & libraries: namespaces ('11),
bytecode ('11), indexing beyond 2GB ('13)

What R is...

- vectorized
- functional
- object-oriented
- lazy
- portable
- interactive

What R isn't...

- fast
- low-footprint
- concurrent
- distributed
- formally specified
- standardized

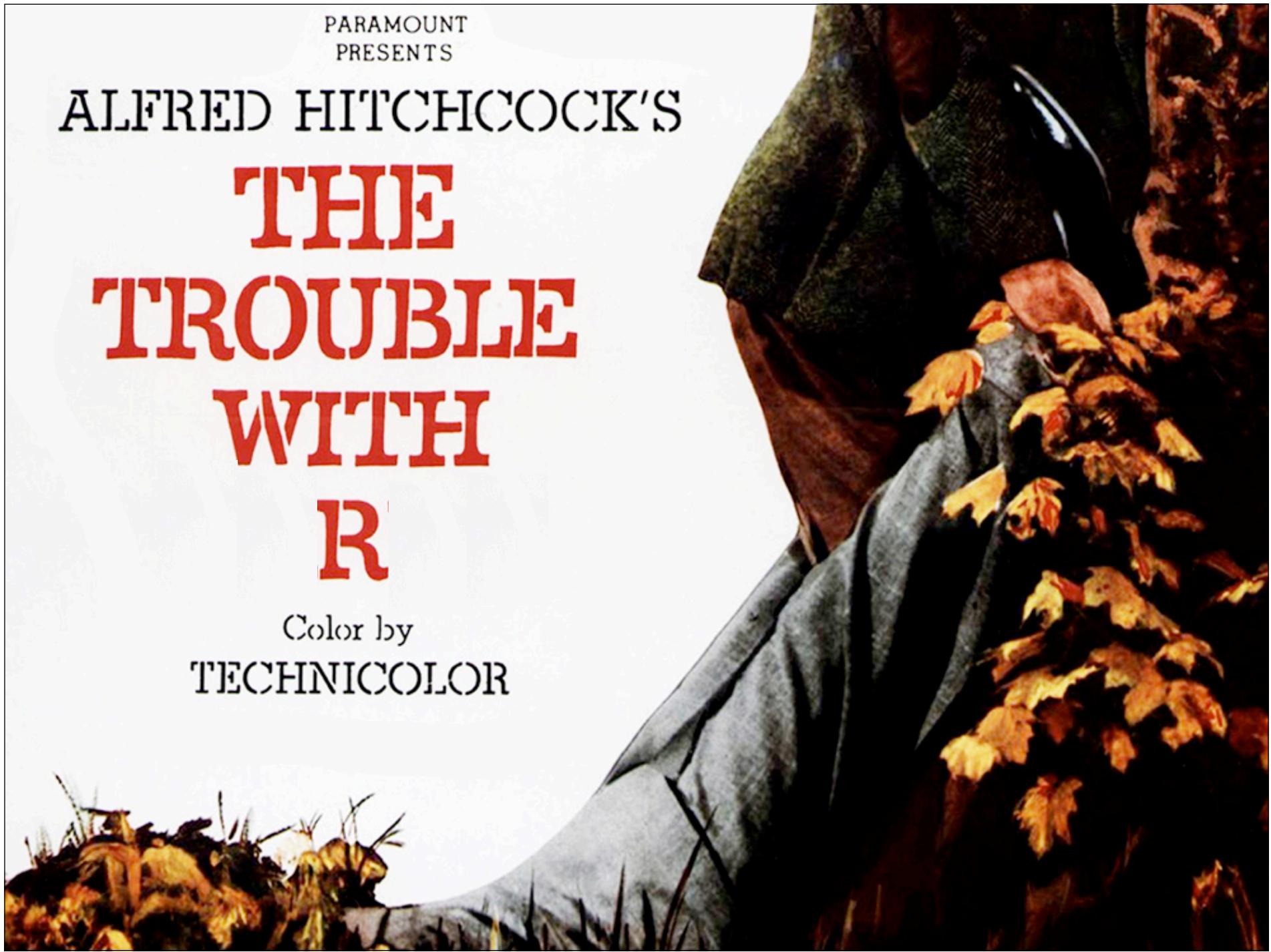


PARAMOUNT
PRESENTS

ALFRED HITCHCOCK'S

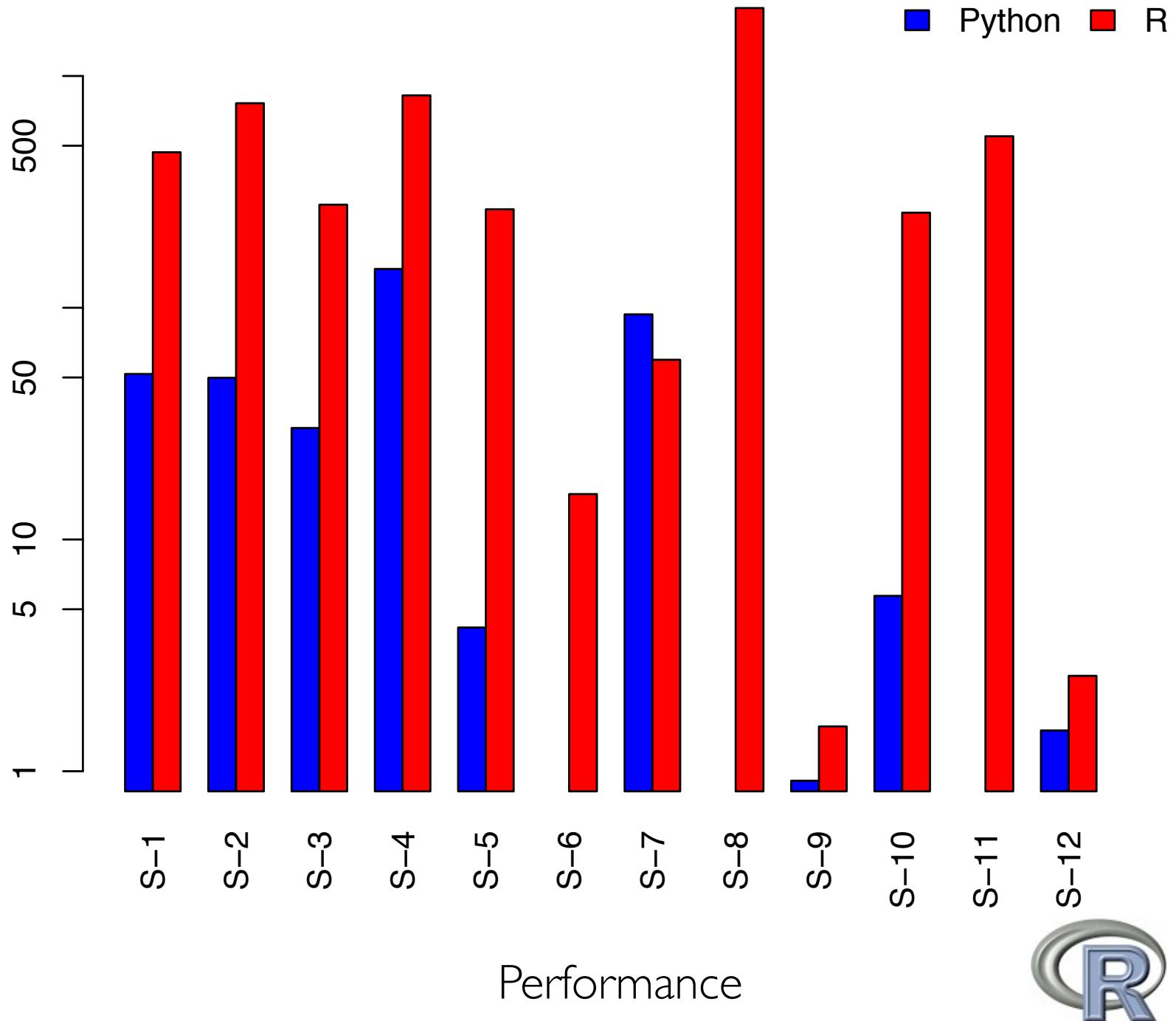
THE TROUBLE WITH R

Color by
TECHNICOLOR

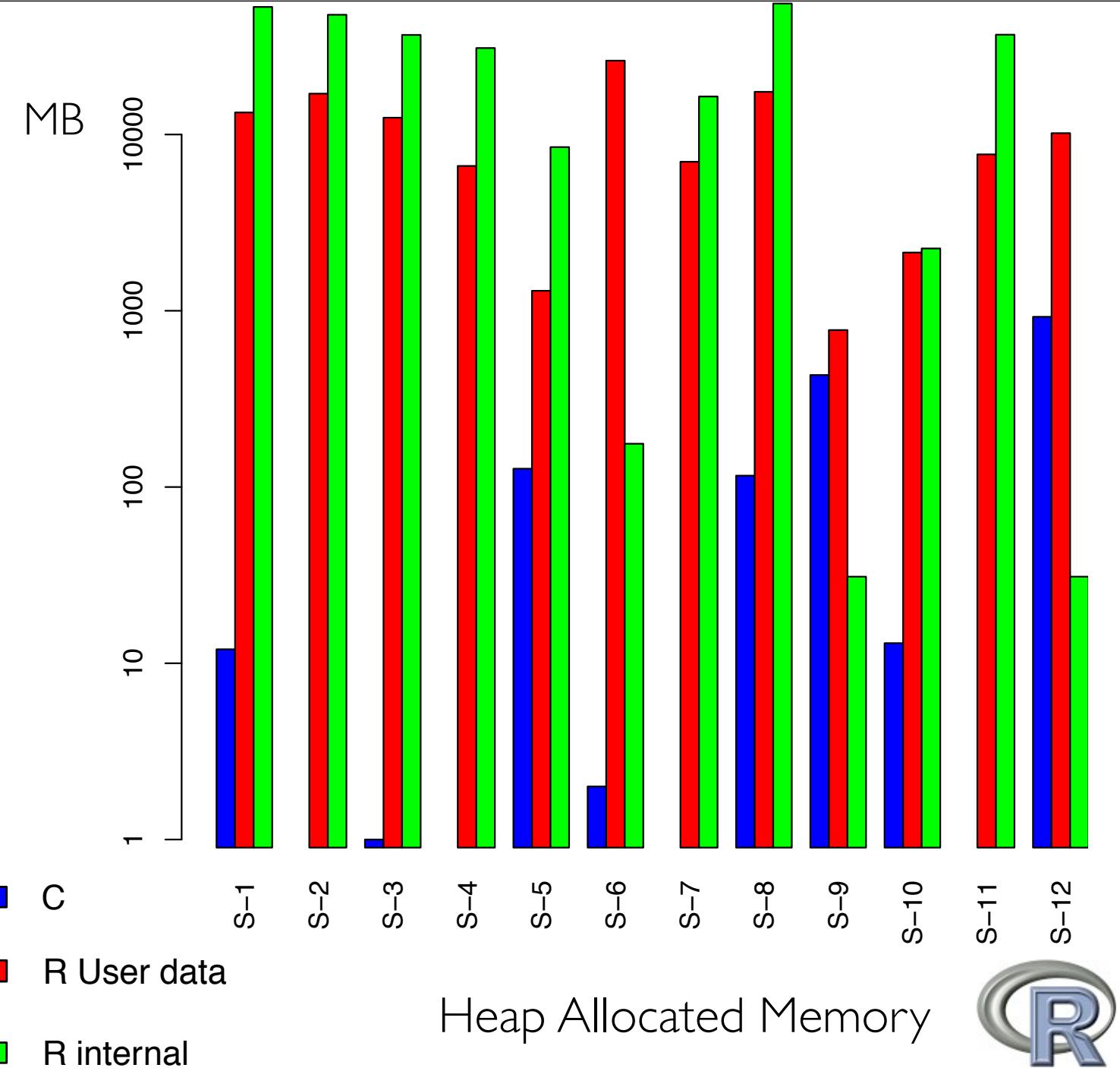


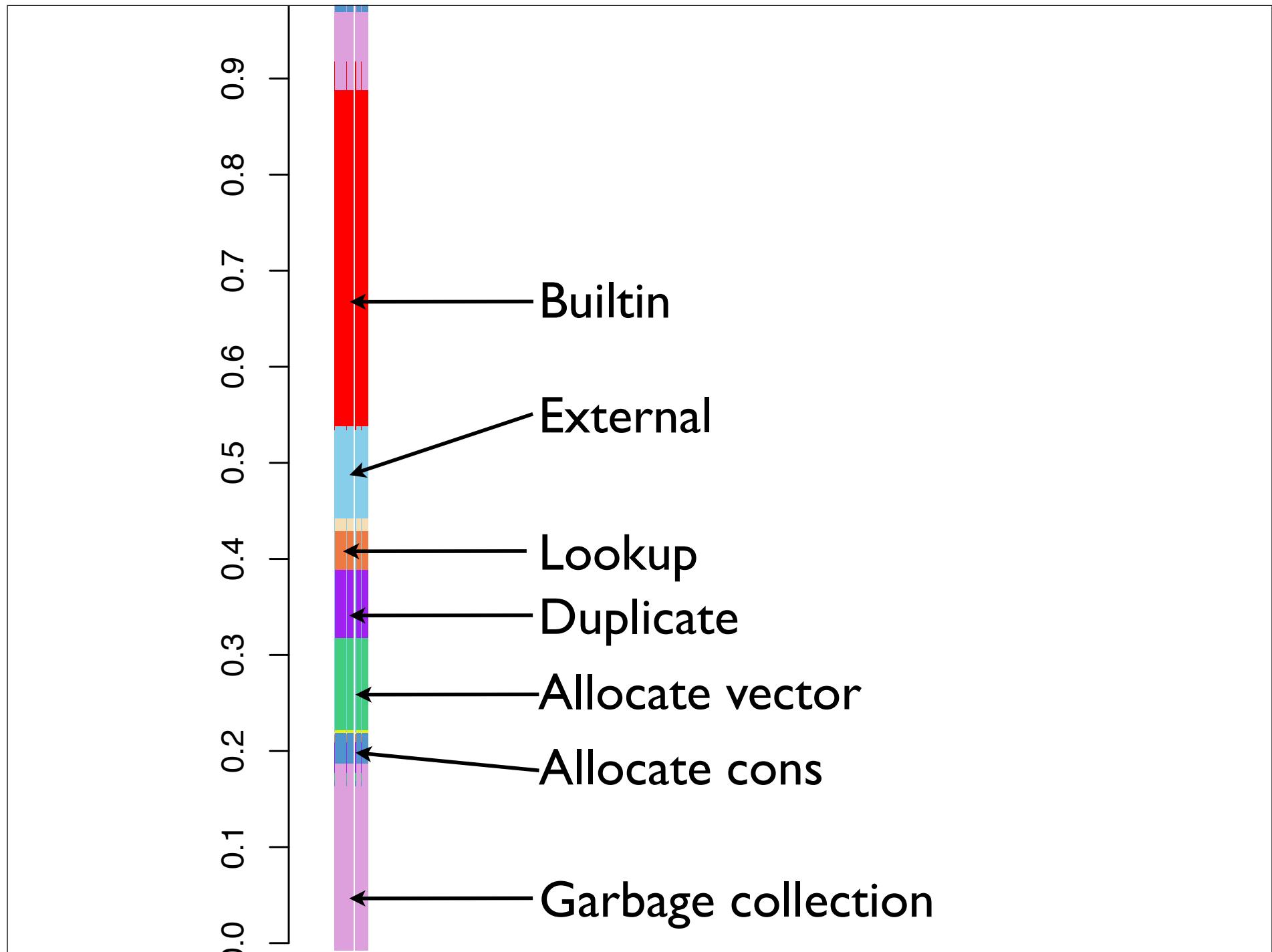
Intel X5460. 3.16GHz, Linux 2.6.34. R 2.12.1, GCC v4.4.5

The programming language shootout C / Python / R

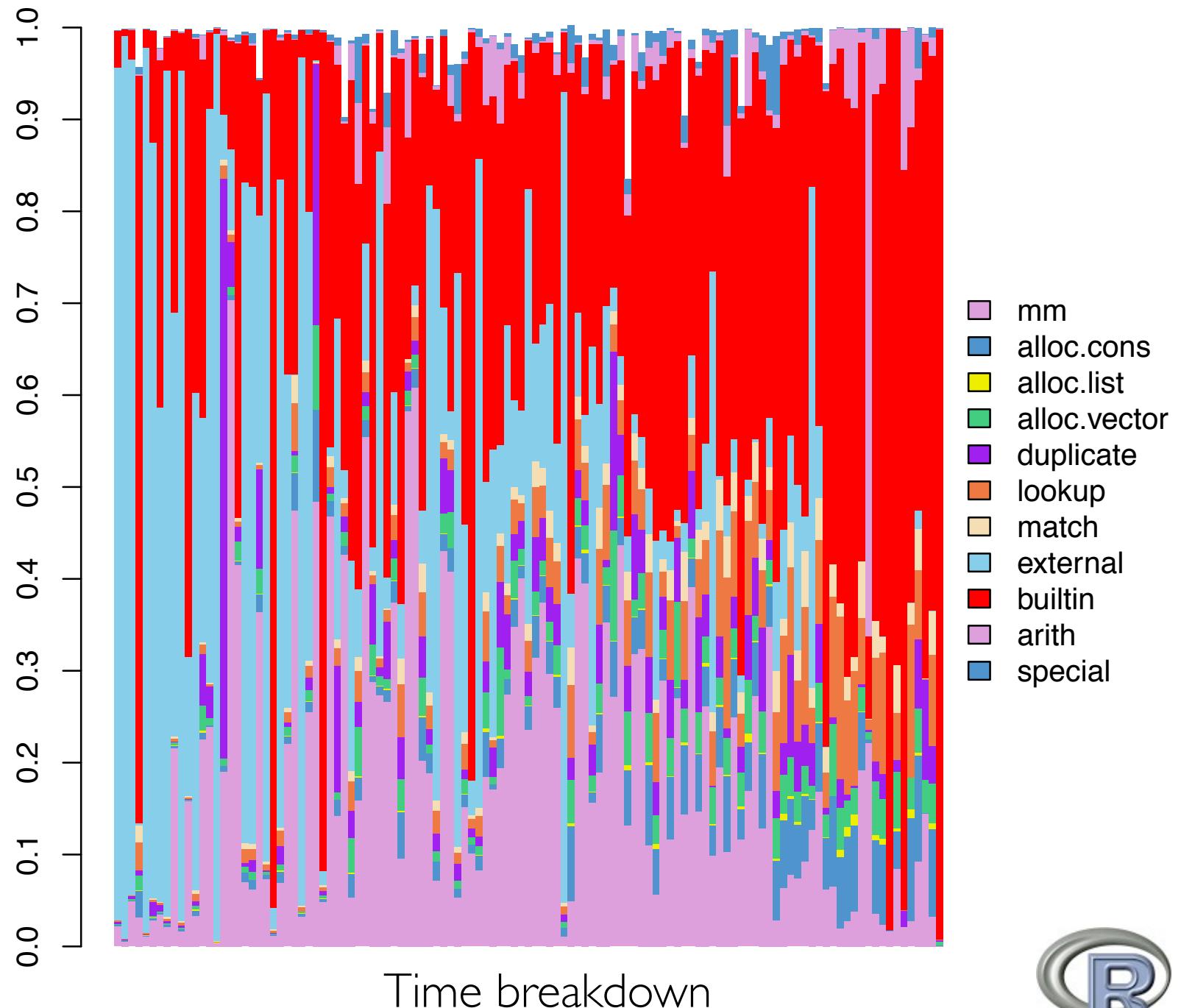


The programming language shootout C / R





Bioconductor vignettes



How is R used?

- Extract core semantics by testing
 - *R has no official semantics*
 - *A single reference implementation*
- Observational study based on a large corpus
 - *Many open source programs come with “vignettes”*
 - *Dynamic analysis gives under-approximated behaviors*
 - *Static analysis gives over-approximation*

POPULATION

Name	Bioc.	Shoot.	Misc.	CRAN	Base
# Package	630	11	7	1238	27
# Vignettes	100	11	4	–	–
R LOC	1.4M	973	1.3K	2.3M	91K
C LOC	2M	0	0	2.9M	50K



Vectors

```
x <- c(2, 7, 9, NA, 5)
```

```
c(1, 2, 3) + x[1:3]
```

```
x[is.na(x)] <- 0
```

Functions

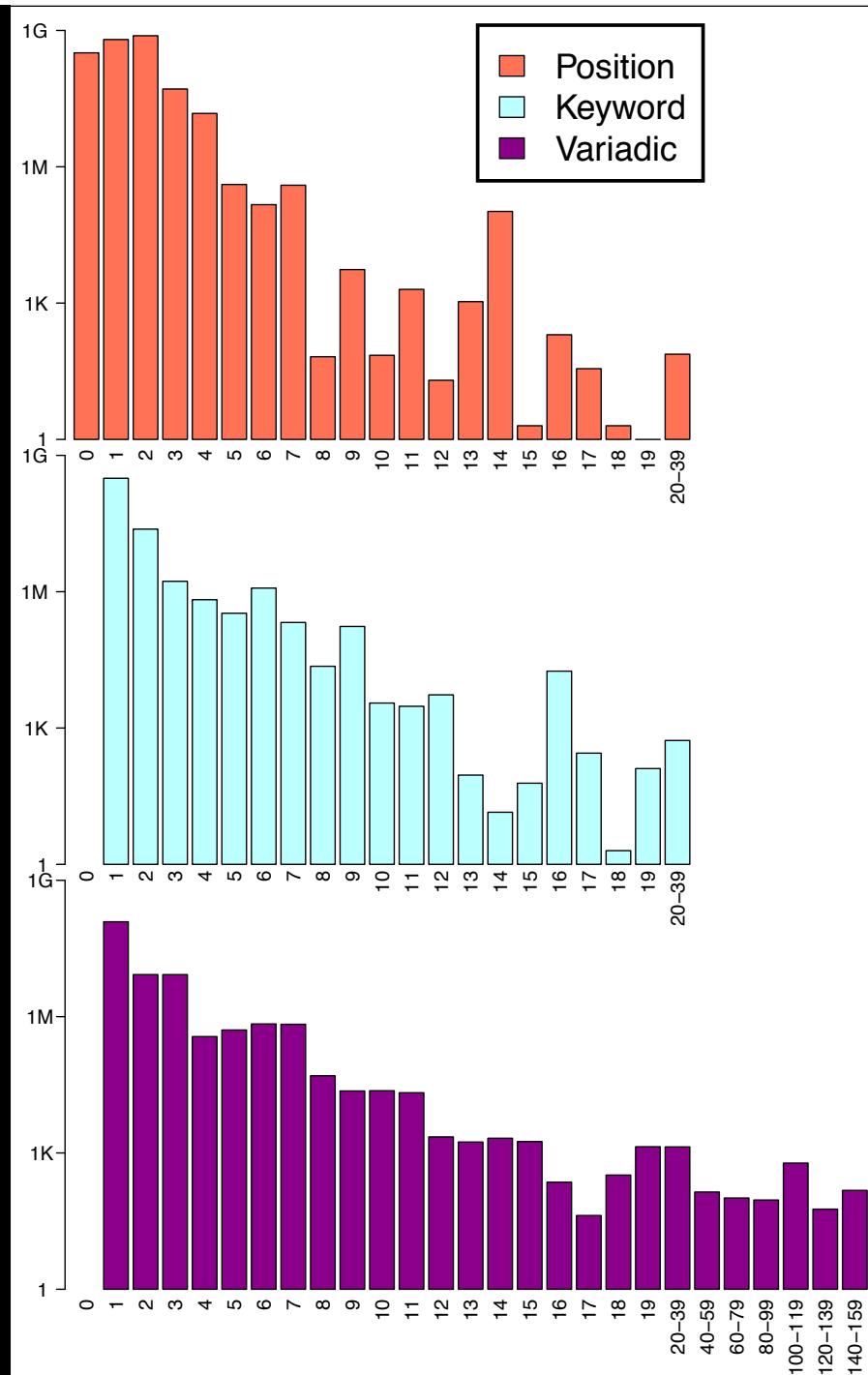
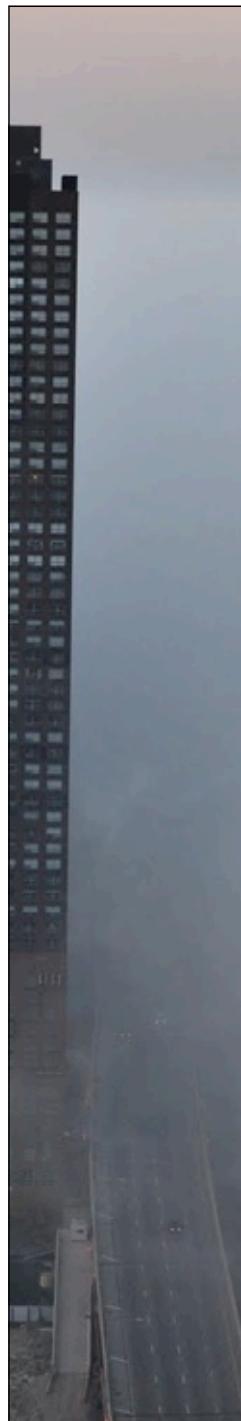
```
q<-function (x=5) x*x*x
```

```
q()
```

```
q(2)
```

```
q(x=4)
```

```
p<-function (x=5, ..., y=x+1)
```



$f(1, 2)$

$f(x=1, y=2)$

$f(y=1, x=2)$

$f(2, x=1)$

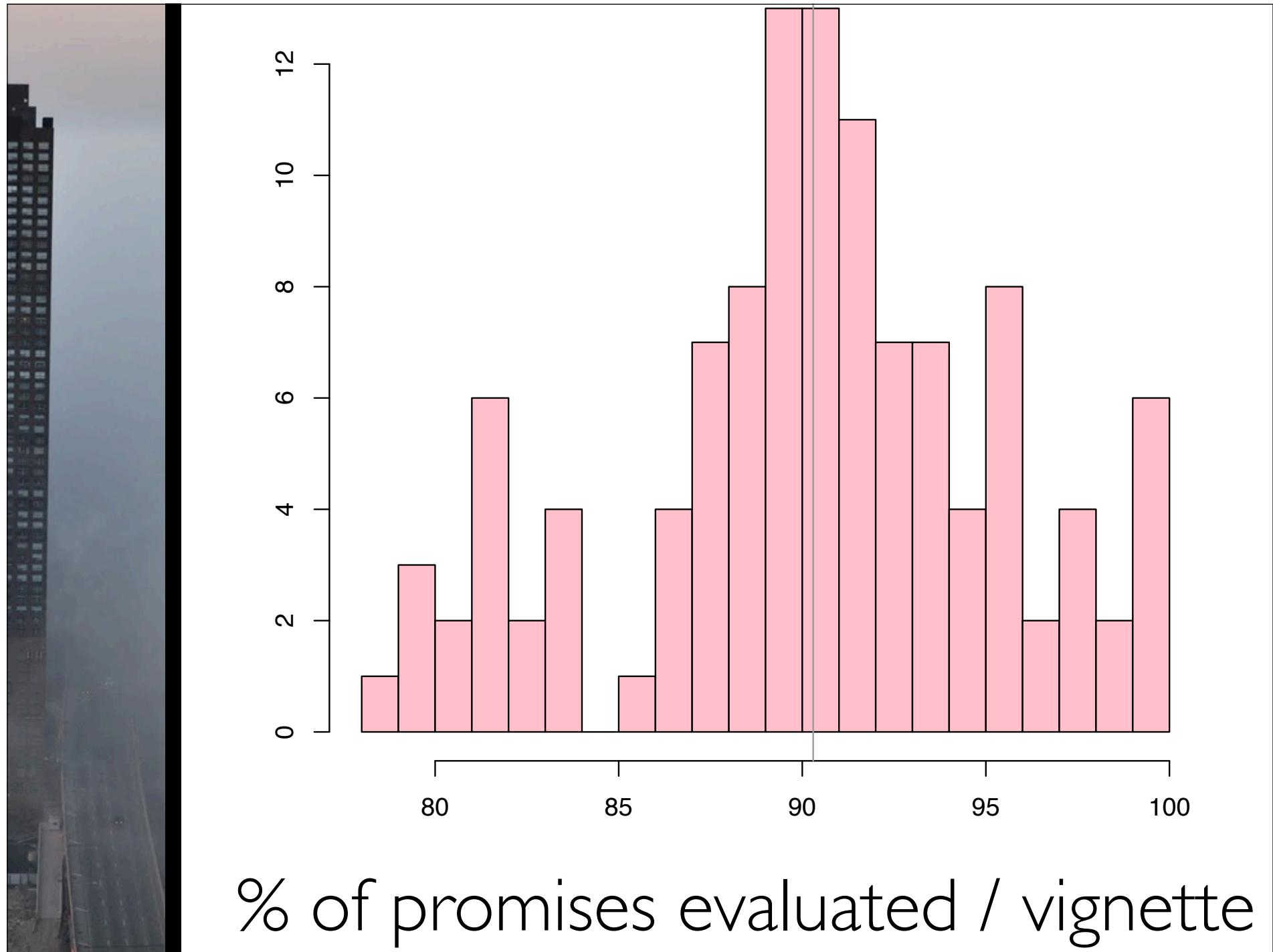
$f(x=1, 2)$

$c(1, 2, 3, 4)$

Promises

```
assert<-function (C, P)  
  if (C) print(P)
```

```
assert( x==42, print("Oops") )
```



Forcing promises

`x <- F`

`x[12] <- F`

`F ; e`

`{e ; F}`



Scoping

Lexical scoping with context sensitive name resolution

```
c <- 42
```

```
c(1,2,3)
```

```
c <- 42
```

```
d <- c
```

```
d(1,2,3)
```



less than 0.05% context sensitive
function name lookups

only symbols that rely on it are **c**
and **file**

Referential transparency

```
assert(y[[1]]==5)
```

```
f(y)
```

```
assert(y[[1]]==5)
```

```
f<-function(b) {b[[1]]<-0}
```

Assignment

x [42] <- y

$$\text{cpy}(H, \nu') = H', \nu'' \quad \Gamma = \iota * \Gamma' \quad \iota(H', x) = \nu''' \\ \text{readn}(\nu, H') = m \quad \text{set}(\nu''', m, \nu'', H') = H''$$

$$x[[\nu]] <- \nu' \Gamma; H \rightarrow \nu'; H''$$



Assignment

```
y <- c(...)  
f <- function() {  
  x [ 42 ] <- y
```

$$\begin{array}{lll} \text{cpy}(H, \nu') = H', \nu'' & \Gamma = \iota * \Gamma' & H'(\iota) = F \\ \text{cpy}(H', \nu''') = H'', \nu'''' & F' = F[x/\nu'''''] & x \notin F \\ \text{readn}(\nu, H) = m & \text{set}(\nu'''', m, \nu'', H''') = H'''' & H''' = H''[\iota/F'] \end{array}$$

$$x[[\nu]] <- \nu'; \Gamma; H \rightarrow \nu'; H''''$$

Assignment

x [42] <<- y

$$\frac{\text{cpy}(H, \nu) = H', \nu' \quad \Gamma = \iota * \Gamma' \quad \text{assign}(x, \nu', \Gamma', H') = H''}{x <<- \nu \Gamma; H \rightarrow \nu; H''}$$



45% of assignments are definitions

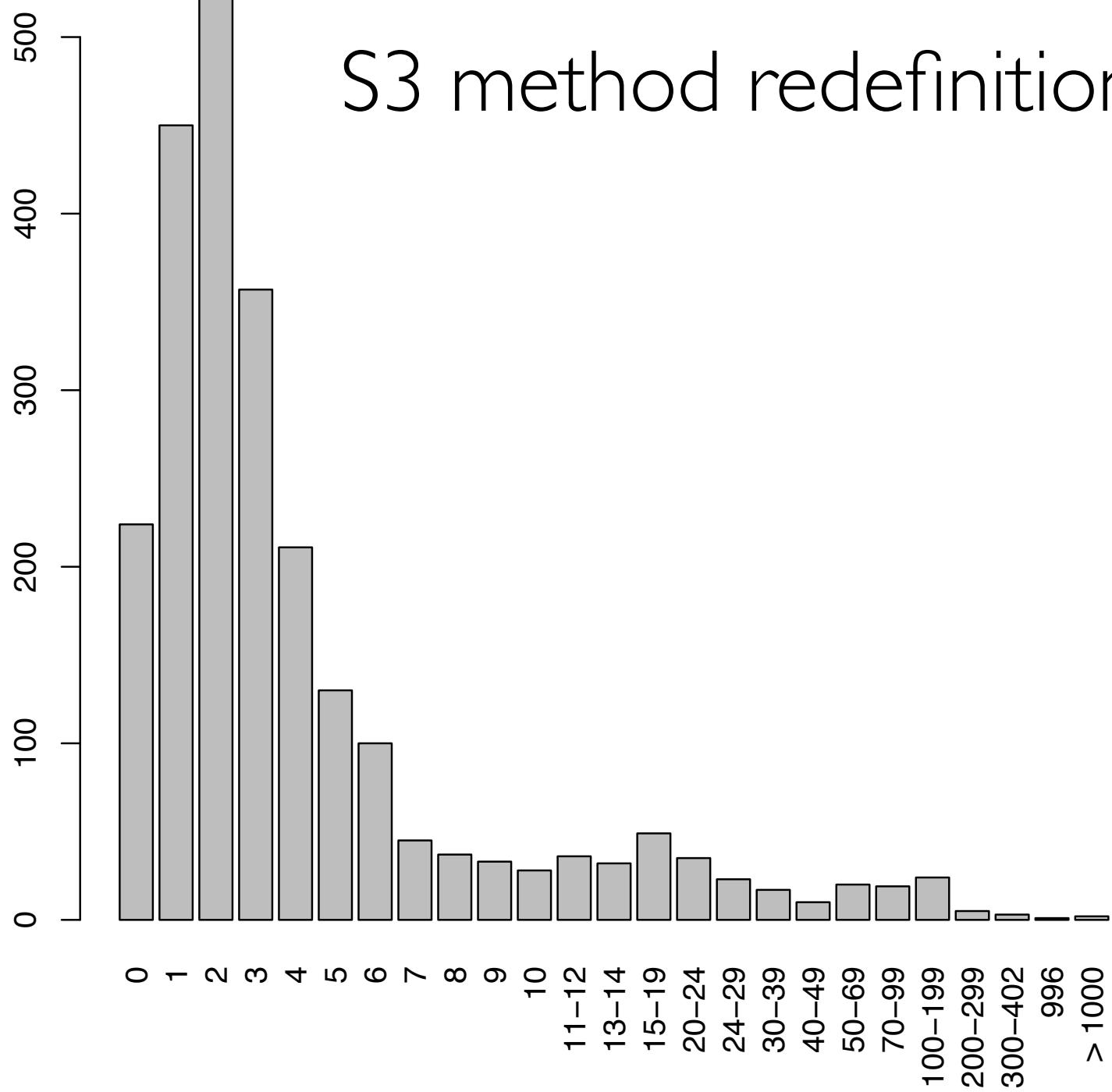
only 2 out of 217 million
assignments are definitions in a
parent frame

99.9% of side effects are local

Objects

```
who <-function(x) UseMethod("who")  
  
who.man <-function(x) print("Ceasar!")  
  
who.default <-function(x)print("??")  
  
me <- 42; who(me)  
class(me)<-`man` ; who(me)
```

S3 method redefinitions



Objects

```
setClass("P", representation(x="numeric", y="numeric"))
setClass("C", representation(color="character"))
setClass("CP", contains=c("P", "C"))

r <- new("CP", x = 0, y = 0, color = "red")
r@color

setGeneric("add", function(a, b) standardGeneric("add"))
setMethod("add", signature("P", "P"),
          function(a, b) new("P", x=a@x+b@x, y=a@y+b@y))
setMethod("add", signature("CP", "CP"),
          function(a, b) new("CP", x=a@x+b@x, y=a@y+b@y, color=a@color))
```

Object usage

		Bioc	Misc	CRAN	Base	Total
S3	# classes	1 535	0	3 351	191	3 860
	# methods	1 008	0	1 924	289	2 438
	Avg. redef.	6.23	0	7.26	4.25	9.75
	Method calls	13M	58M	-	-	76M
	Super calls	697K	1.2M	-	-	2M

S4	# classes	1 915	2	1 406	63	2 893
	# singleton	608	2	370	28	884
	# leaves	819	0	621	16	1 234
	Hier. depth	9	1	8	4	9
	Direct supers	1.09	0	1.13	0.83	1.07
	# methods	4 136	22	2 151	24	5 557
	Avg. redef.	3	1	3.9	2.96	3.26
	Redef. depth	1.12	1	1.21	1.08	1.14
	# new	668K	64	-	-	668K
	Method calls	15M	266	-	-	15M
	Super calls	94K	0	-	-	94K

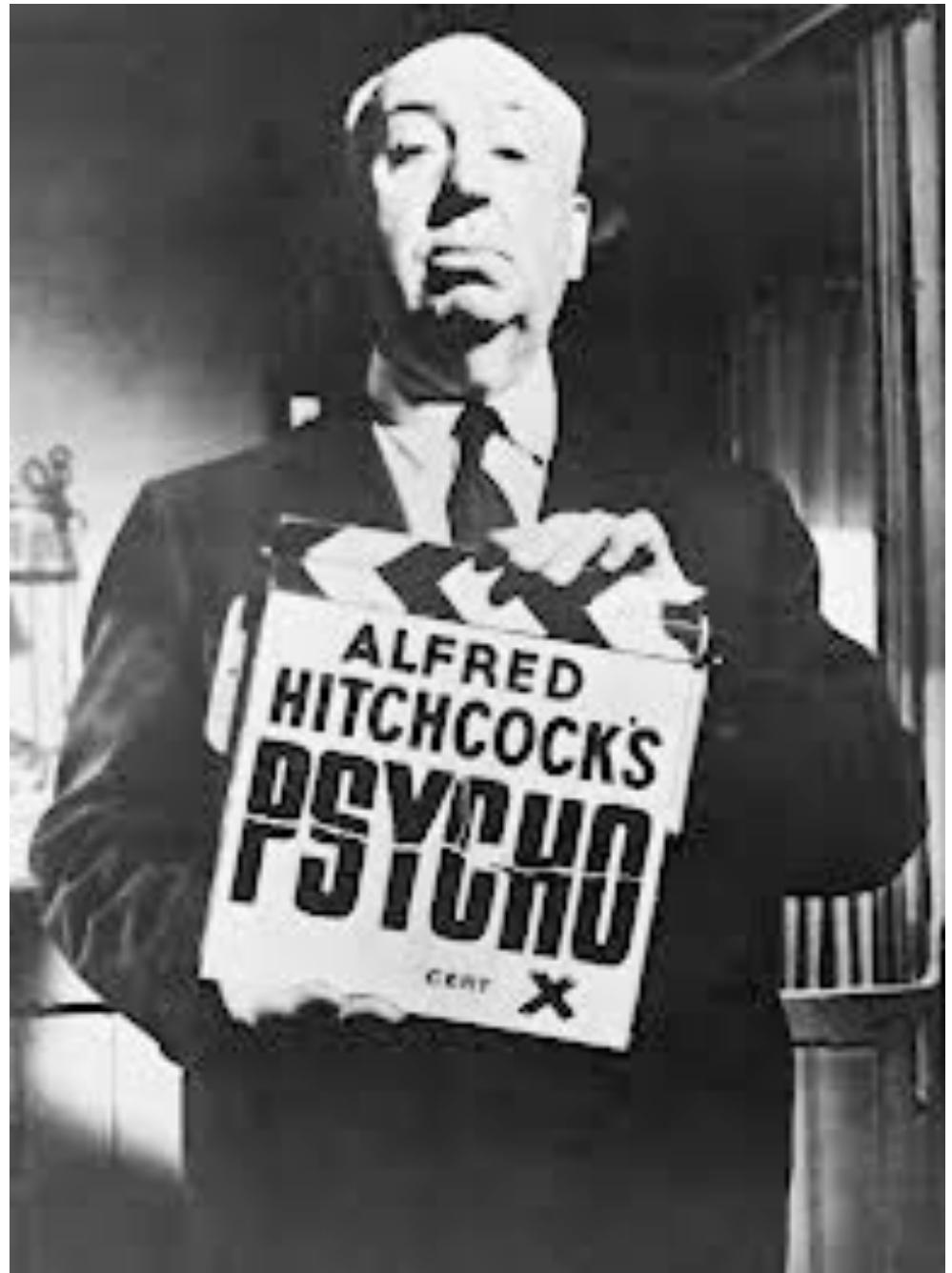
The trouble with R ?



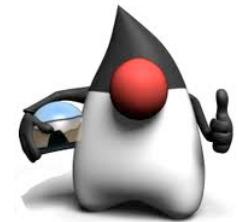
- R is slow because it lacks a JIT
- R is a memory hog because it has large objects, allocates profusely and a non-moving garbage collector
- R has tricky semantics rife with cobwebs and dark corners

R is legacy software that must be maintained, yet it must also evolve to meet new challenges

R in Java



Why Java ?



- Because JavaScript may not be fast enough
- Leverage a runtime system with zero-maintenance
- Get a good just-in-time compiler
- OpenJDK is an open source platform

Why not Java ?



- The R community dislikes Java
- Interaction with C is cumbersome
- Complete break with GNU-R

R interpreter
written in
Java

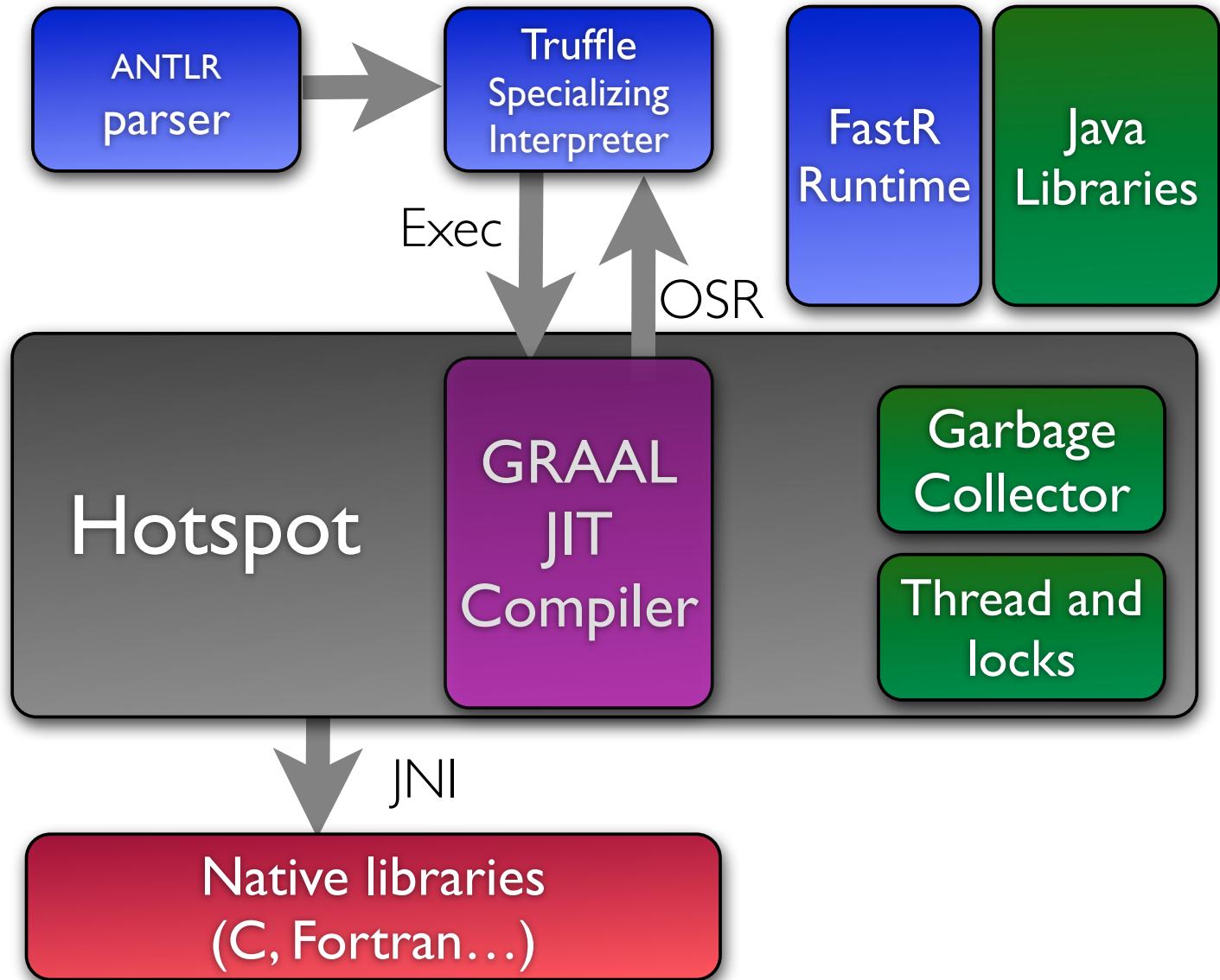
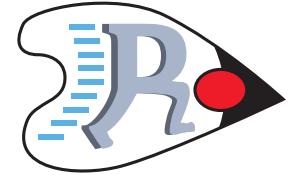
47 KLOC

Up to 20x
faster on
shootout

GPL license

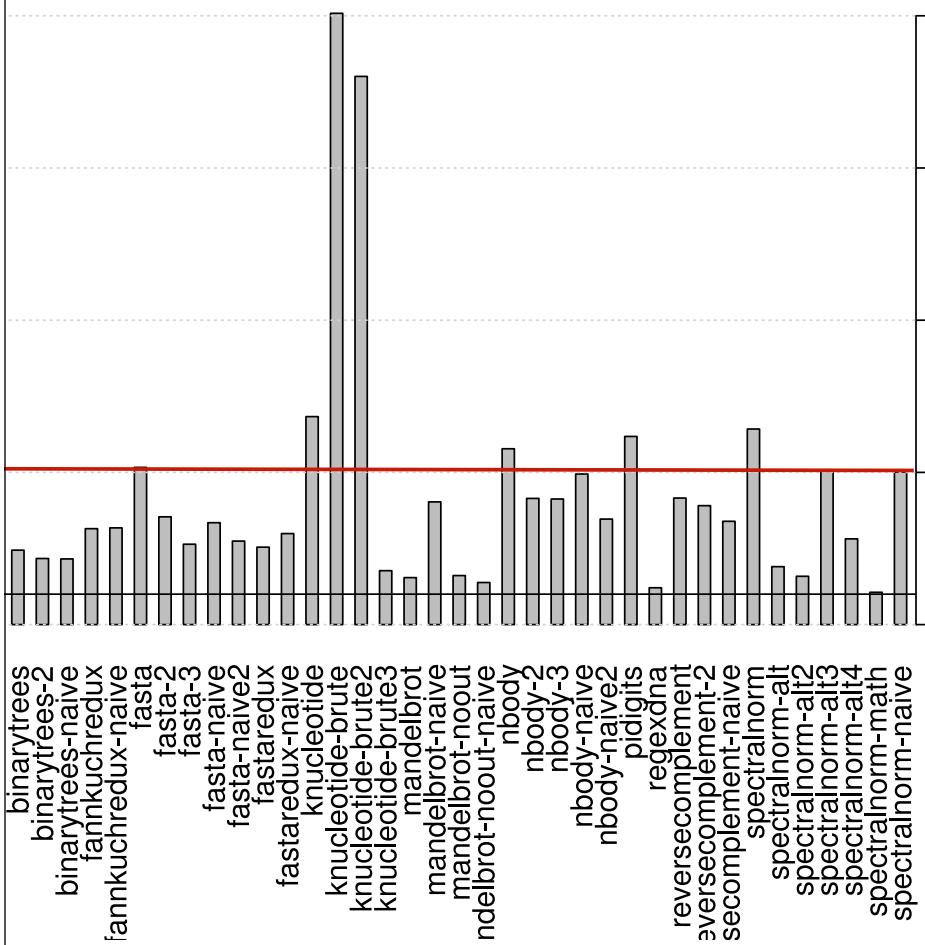
Built on
OpenJDK,
Graal, Truffle

FastR Architecture



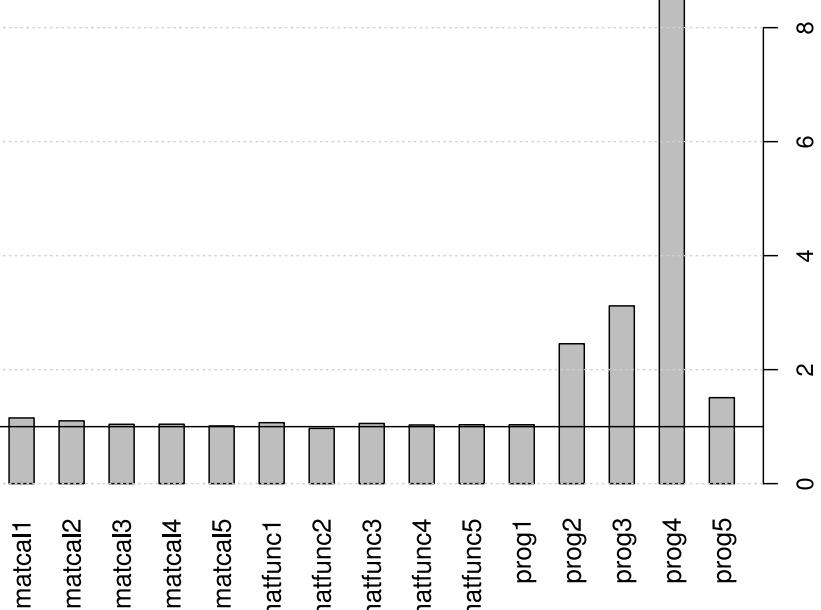
FastR Throughput (w/o JIT)

Shooutout: speedup over GNU-R (gmean: 3.34x)



Benchmark25: speedup over GNU-R (gmean: 1.41x)

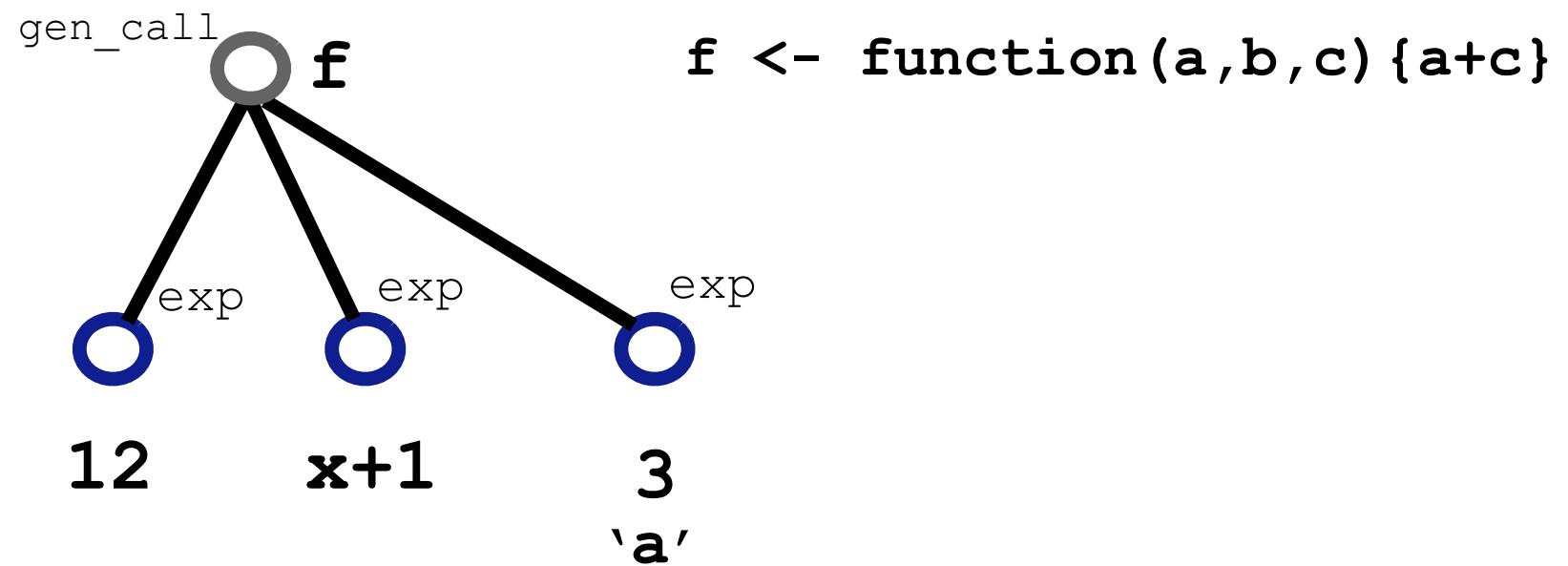
<http://r.research.att.com/benchmarks/>





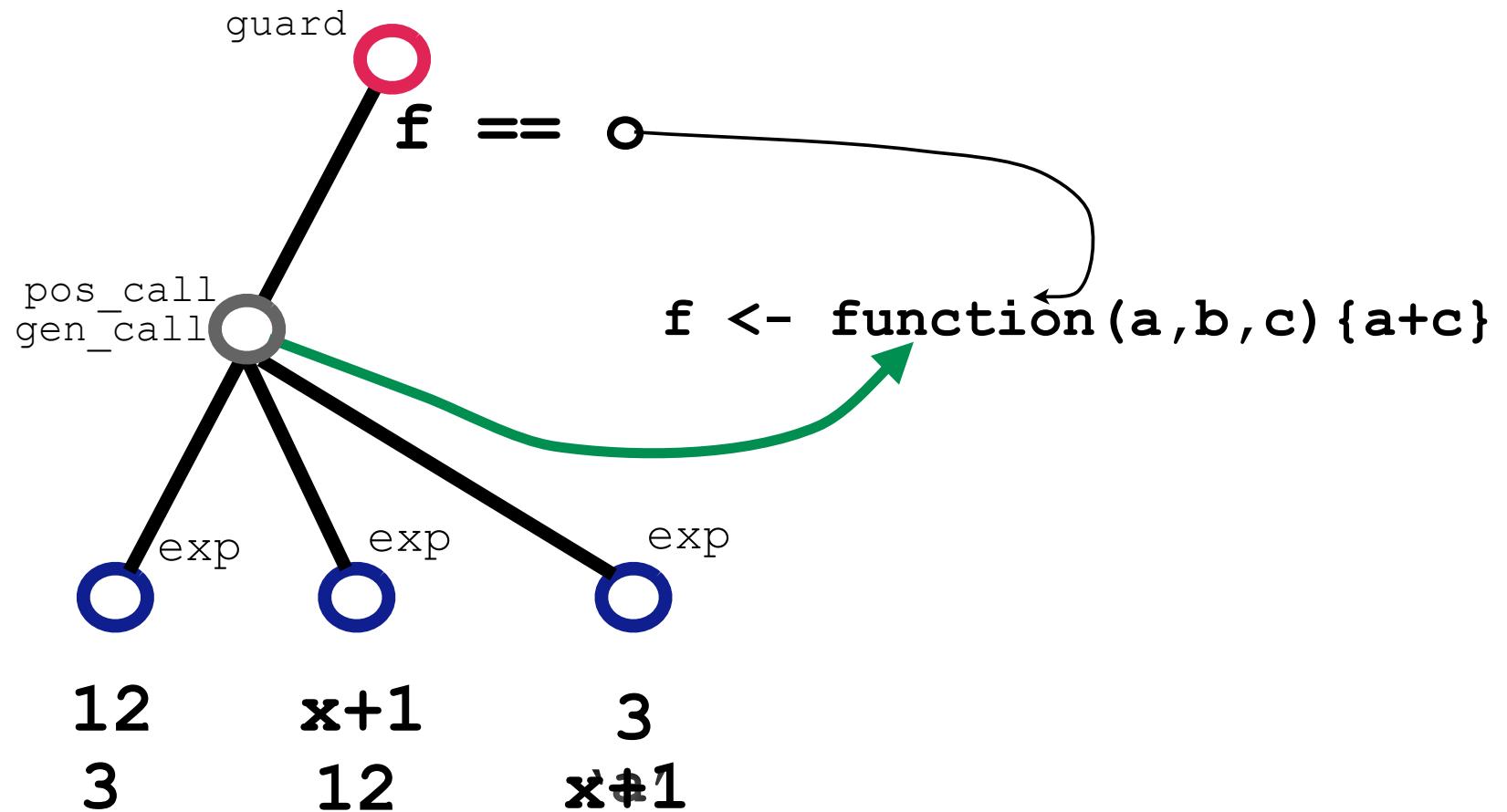
Interpretation

f (12, x+1, a=3)





Specialization

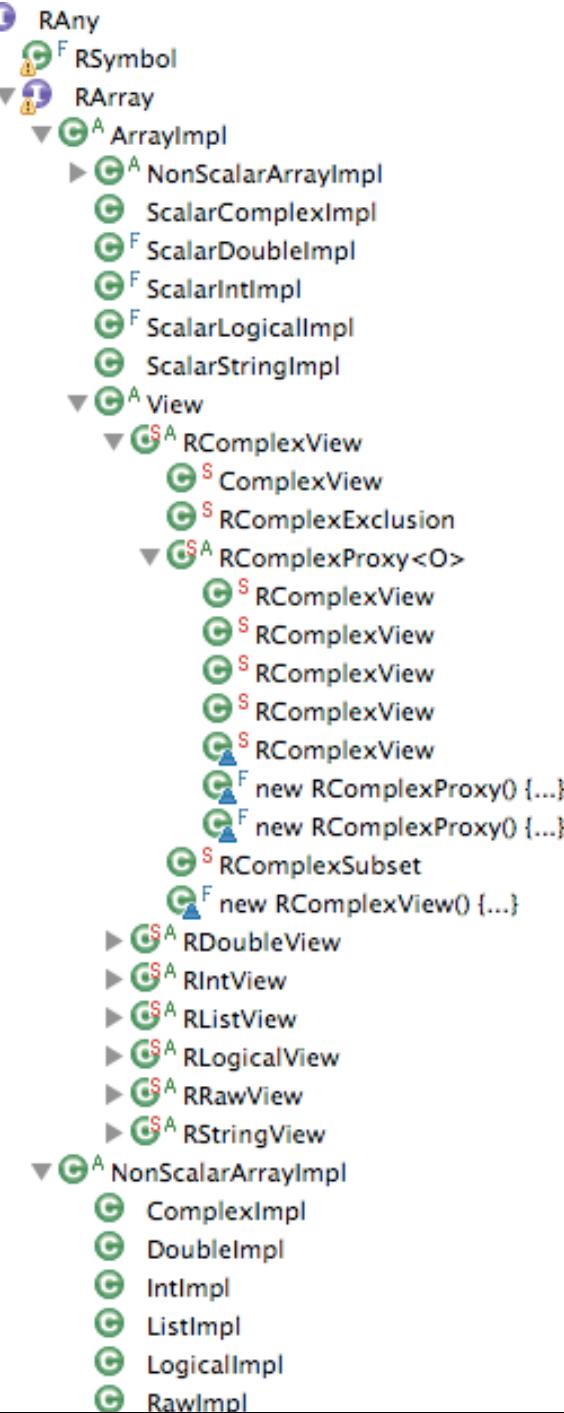




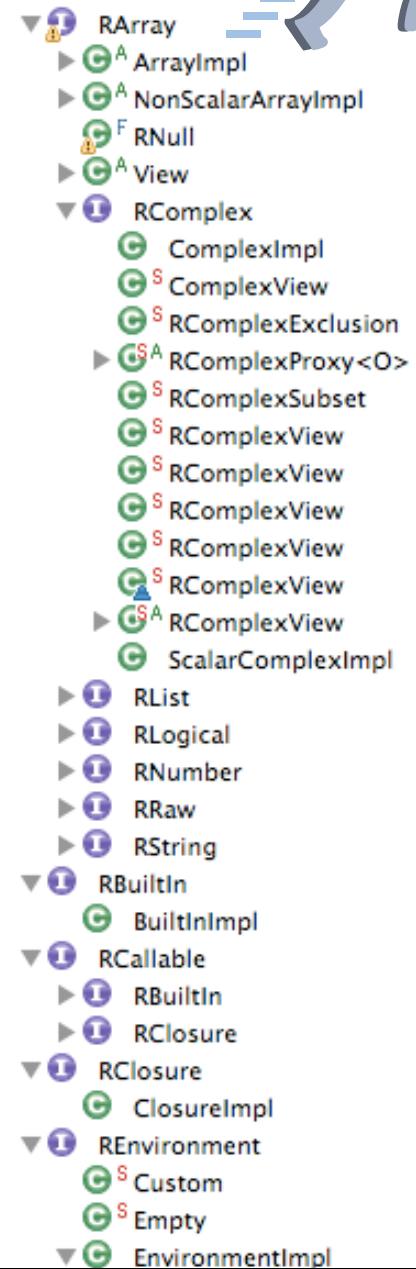
Runtime specialization

```
class If {  
    RNode condE, trueB, falseB;  
  
    Object execute(Frame f) {  
        try {  
            val = condE.executeScalarLogical(frame);  
        } catch (UnexpectedResult e) {  
            cast = ToLogical.mkNode(condE, e.result());  
            replaceChild(condE, cast);  
            return execute(frame);  
        }  
        if (val == TRUE) return trueB.execute(f);  
        if (val == FALSE) return falseB.execute(f);  
        throw unexpectedNA();  
    }  
}
```

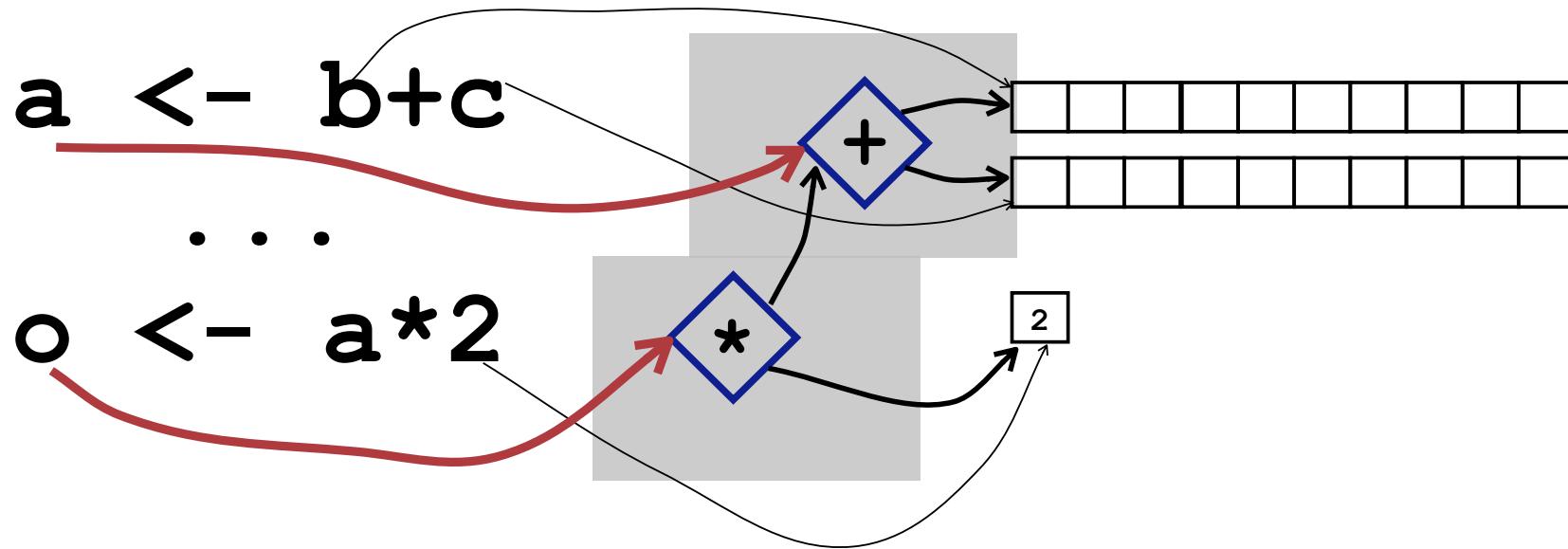
Data types



- Data types are specialized to optimize memory usage and speed
- Scalar are unboxed when possible (soon)
- Memory footprint is improved



Views



- ...delay construction of large data objects
- ...are a data-flow representation of vectors
- ...avoid unnecessary work if a subset of the data is required
- ...avoid allocation of temporary objects
- ...permit fusion of multiple data traversals into one

Experiments





Inlining

spectralnorm-naive

```
A <- function(i, j) {  
  1 / ((i + j) * (i + j + 1) / 2 + i + 1) }
```

```
B <- function(u) { ...  
  for (j in 0:n1)  
    ret[[i]] <- ret[[i]] + u[[j] +  
      * A(i - 1, j)}
```

1.8x

- Inlining is a critical optimization in modern languages
- Replace a function call with its body
- Guarded inlining leaves a slow path in the code that performs the normal function

Relite

Delite is a compiler framework and runtime for parallel embedded DSLs.

Delite provides:

- Built-in parallel execution patterns
- Optimizers for parallel code
- Code generators for Scala, C++ and CUDA
- A heterogeneous runtime for executing DSLs

Relite is a proof-of-concept R interface to Delite.

<https://github.com/TiarkRompf/Relite>

Relite

```
sapply(1:50000, function(x){sum(1:x)})  
                                4.1s  GNU-R  
                                1.5s  FastR  
Delite(sapply(1:50000,function(x) {sum(1:x)}))  
                                0.4s Relite  
sapply(1:50000,function(x) {sum((1:x)*0.1)})  
                                9.0s  GNU-R  
                                2.2s  FastR  
                                0.5s Relite  
sapply(1:50000,  
      function(x){ sum(sapply(1:x, function(y) y*0.1)))})  
                                2395s  GNU-R  
                                104s  FastR  
                                0.5s Relite
```

Conclusions

- R is an amazingly successful systems with great mindshare
- The R implementation is hard to maintain and evolve
- The FastR project aims to rethink how to implement R
- We leverage well tested technologies to build a high performance VM
- FastR can be a source of inspiration for GNU R

