CS 105: Department Introductory Course on Discrete Structures

Instructor : S. Akshay Guest Lecture by : R. Govind

Sep 04, 2023 Lecture 12 – Basic Mathematical Structures Chains and Antichains

Recap: Partial order relations

Last class we saw

- ▶ Partial orders: definition and examples
- ► Posets
- ▶ Graphical representation as Directed Acyclic Graphs

Definition

- ► A partial order is a relation which is reflexive, transitive and anti-symmetric.
- ▶ A total order is a partial order in which every pair of elements is comparable.
- ▶ A poset is a set S with a partial order \leq \subseteq S \times S.

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Examples

- \triangleright (\mathbb{Z} , \leq): integers with the usual less than or equal to relation.
- \triangleright $(\mathcal{P}(S), \subseteq)$: powerset of any set with the subset relation.
- \triangleright (\mathbb{Z}^+ , |): positive integers with divisibility relation.

Recap: Partial order relations

Let $S = \{1, 2, 3\}$. Recall the poset $(\mathcal{P}(S), \subseteq)$. How does the graph of $(\mathcal{P}(S), \subseteq)$ look like?

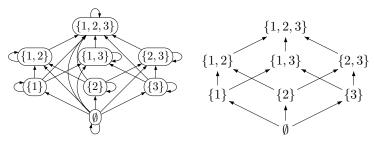
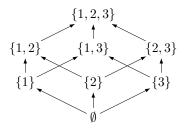


Figure: Graph of a poset and its Hasse diagram

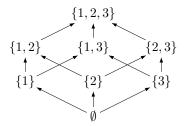
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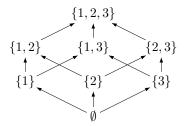
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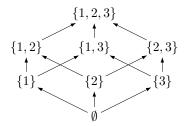
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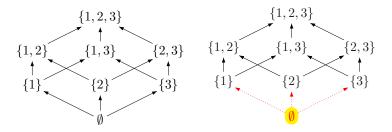
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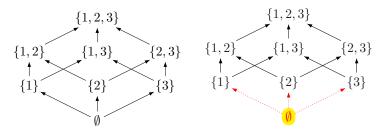
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- ▶ What are the minimal & maximal elements in $(\mathbb{Z}^+, |)$.
- ▶ Is there always a unique minimal/maximal element?
- ▶ What are the minimal element(s?) in $(\mathbb{Z}_{>1}, |)$.

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Proof by induction?(H.W)

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What about infinite posets?

Posets: Chains and Antichains

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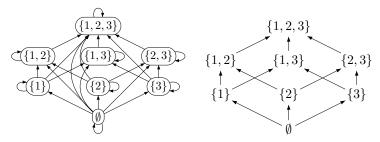


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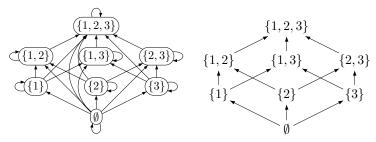


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- ► Subsets that are totally ordered?
- ▶ Subsets that are unordered?

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Let (S, \preceq) be a poset. A subset $B \subseteq S$ is called

- \triangleright a chain if every pair of elements in B is related by \preceq .
- ▶ That is, $\forall a, b \in B$, we have $a \leq b$ or $b \leq a$ (or both).

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- ▶ an anti-chain if no two distinct elements of A are related to each other under \preceq .
- ▶ That is, $\forall a, b \in A, a \neq b$, we have neither $a \leq b$ nor $b \leq a$.

Chains and Anti-chains: examples

▶ Let $S = \{1, 2, 3\}$.

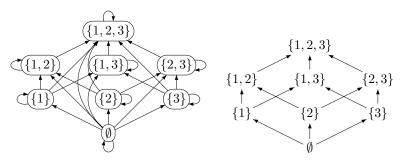


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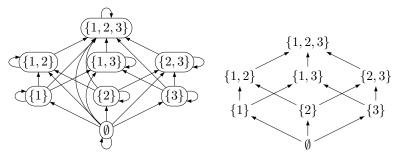


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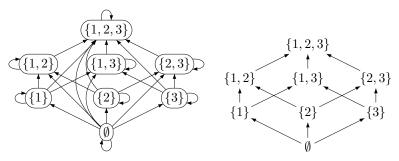


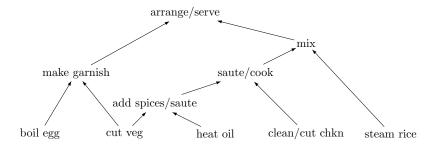
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- ▶ What are the chains in this poset?
- ▶ What are the anti-chains in this poset?
- Give an example of an infinite chain & anti-chain in $(\mathbb{Z}^+, |)$.

Examples and applications

A task scheduling example

Let us represent a recipe for making Chicken Biriyani as a poset!

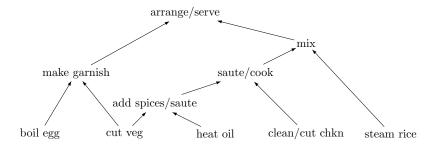


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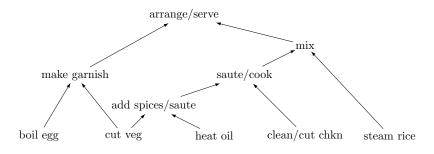


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- ► Clearly, this shows the dependencies.
- ▶ But when you cook you need a total order, right?
- ► Further, this total order must be consistent with the po.
- ► This is called a linearization or a topological sorting.

Topological sorting

Definition

A topological sort or a linearization of a poset (S, \preceq) is a poset (S, \preceq_t) with a total order \preceq_t such that $x \preceq y$ implies $x \preceq_t y$.

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Proof: (H.W)

- ► Recall the lemma:
 - Every finite non-empty poset has at least one minimal element $(x \text{ is minimal if } \not\exists y, y \leq x).$
- ▶ Then, construct a (new) chain to complete the proof.