CS 105: DIC on Discrete Structures

Instructor: S. Akshay

Sept 25, 2023 Lecture 18 – Counting and Combinatorics Recurrence Relations (Contd.)

Summary and what's next

Part 1: Proofs and basic mathematical structures

Part 2: Counting and Combinatorics

- ▶ Basics of counting
 - ▶ Product principle
 - Sum principle
 - Bijection principle
 - Double counting
- Subsets, partitions, Permutations and combinations
 - 1. Binomial coefficients and Binomial theorem
 - 2. Pascal's triangle
 - 3. Permutations and combinations with repetitions
 - 4. Estimating n!

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 - 4. Estimating n!
- ▶ Recurrence relations and generating functions
- ▶ Pigeonhole Principle and its extensions

Next: Recurrence relations and generating functions

Definition

- A recurrence relation for a sequence is an equation that expresses its n^{th} term using one or more of the previous terms of the sequence.
- ▶ A linear recurrence relation is of the form

$$u_n = a_{k-1}u_{n-1} + \ldots + a_1u_{n-k+1} + a_0u_{n-k}$$

where $a_0, \ldots, a_{k-1} \in \mathbb{R}, k \in \mathbb{N}$ are constants.

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where $a_0, \ldots, a_{k-1} \in \mathbb{R}, k \in \mathbb{N}$ are constants.

- \triangleright k is called the degree/depth of the sequence.
- ▶ The first few (e.g., k elements u_0, \ldots, u_{k-1}) are initial conditions and they determine the whole sequence.

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- $\qquad \qquad n = 4: (((a+b)+c)+d), ((a+b)+(c+d)), ((a+(b+c))+d), \ldots$

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- Example: n = 3 : ((a + b) + c), (a + (b + c))
- ▶ n = 4 : (((a+b)+c)+d), ((a+b)+(c+d)), ((a+(b+c))+d), ...In general, let C(n) be the number of ways of doing this.

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▶ Thus,
$$C(n) = \sum_{i=1}^{n} C(i)C(n-i)$$
 for $n > 1$

- ▶ Initial conditions are C(0) = C(1) = 1.
- ▶ This sequence are called Catalan numbers...

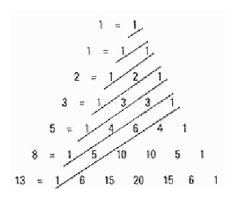
How do we solve such recurrences? We start with the Fibonacci sequence.

An aside: find the Fibonacci sequence!

```
1 5 10 10 5 1
          15 20 15 6 1
   1 7 21 35 35 21 7 1
 1 8 28 56 70 56 28 8 1
1 9 36 84 126 126 84
   45 120 200 252 200 120
```

- F(n) = F(n-1) + F(n-2).
- ► 1, 1, 2, 3, 5, 8, 13,
- ➤ Can you observe the sum of which terms in the Pascal's triangle gives rise to the terms of the Fibonacci sequence?

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By solving, we mean give a closed-form expression for n^{th} term.

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For
$$n \ge 2$$
, $F_n = F_{n-1} + F_{n-2}$, $F_0 = F_1 = 1$.

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- 2. i.e., if $F_n = F_{n-1} + F_{n-2}$, $G_n = G_{n-1} + G_{n-2}$ and $H_n = aF_n + bG_n$, then what is a recurrence for H_n ?

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- 3. So if $\alpha^2 \alpha 1 = 0$, the recurrence holds for all n.
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- 4. Solve it! $\alpha = \frac{1+\sqrt{5}}{2}$, $\beta = \frac{1-\sqrt{5}}{2}$

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- 6. How do we get a and b?

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- ▶ Recall the recurrence for Catalan Numbers:

$$C(n) = \sum_{i=1}^{n} C(i)C(n-i)$$
 for $n > 1$, $C(0) = C(1) = 1$.

No. of ways to bracket a sum of n terms s.t. it can be computed by adding two numbers at a time?

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This method does not work if we have repeated roots (this can be fixed!) and non-linear recurrences.