Database Normalization – Exercise Practices on 3NF

Q.1 Suppose a relational schema R (A B C D E) and set of functional dependencies

F:
$$\{A \rightarrow B \\ B \rightarrow E \\ C \rightarrow D \}$$

Check out that relation is in 3NF or not? If not decompose it in 3NF.

Solution1:

Firstly find the candidate key in the relation:

$$(AC)^+ = ABCDE$$

AC is the candidate key, because closure of AC has all the attributes of R.

Prime attributes: AC

Non prime attributes: BDE

A relation is said to be 3NF, if it holds at least one of the following for every non trivial functional dependency $\alpha \rightarrow \beta$:

- α is super key.
- β is prime attribute.

 $A \rightarrow B$ -- Neither A is super key, nor B is prime attribute.

 $B \rightarrow E$ -- Neither B is not super key, nor E is prime attribute.

 $C \rightarrow D$ -- Neither C is not super key, nor D is prime attribute.

So, the relation is not in 3NF as it is not following the rules of 3NF.

Therefore, R(ABCDE) needs to be divided into following:

$$R1(A B E)$$
 $R_{11}(A B)$ $R_{12}(B E)$ $R_{3}(A C)$

Now R_{11} , R_{12} , R_{2} , R_{3} are in 3NF.

Q.2 Suppose a relational schema R (A B C D E F G H I) and set of functional dependencies

F:
$$\{AB \rightarrow C,$$

$$AD \rightarrow GH,$$

$$BD \rightarrow EF,$$

$$A \rightarrow I,$$

$$H \rightarrow J \}$$

Check out that relation is in 3NF or not? If not decompose it in 3NF.

Solution 2:

Firstly find the candidate key in the relation:

 $(ABD)^+ = A B C D E F G H I$

ABD is the candidate key, because closure of ABD has all the attributes of R.

Prime attributes: A B D

Non prime attributes: C E F G H I

A relation is said to be 3NF, if it holds at least one of the following for every non trivial functional dependency $\alpha \rightarrow \beta$:

- α is super key.
- β is prime attribute.

 $AB \rightarrow C$ -- Neither AB is super key, nor C is prime attribute.

AD → GH -- Neither AD is not super key, nor GH is prime attribute.

BD → EF -- Neither BD is not super key, nor EF is prime attribute.

 $A \rightarrow I$ -- Neither A is not super key, nor I is prime attribute.

H → J -- Neither H is not super key, nor J is prime attribute.

So, the relation is not in 3NF as it is not following the rules of 3NF.

Therefore, R(ABCDE) needs to be divided into following:

R1(A B C I)
$$R_{11}$$
 (A B C) R_{12} (A I)

R2 (A D G H J) R_{21} (A D G H) R_{22} (H J)

R3(B D E F)

R4(A B D)

Now R₁₁, R₁₂, R₂₁, R₂₂, R3, R4 are in 3NF.