# **Menezes-Vanstone EC Cryptosystem**

#### System design

1. A and B agreed to select an elliptic curve *EC*:

$$y^2 = x^3 + ax + b \pmod{n}$$
 {p:=n};

- 2. They also agreed to select a point G on EC
- 3. A and B select integers  $n_A$  and  $n_B$  {these are private keys of A and B respectively}:  $n_A < n$ ; and  $n_B < n$ ;
- 4. Users compute  $P_i := n_i G$

{ public key of *i*-th user: i=A, B, C,...}

## **Encryption**

{Suppose *A* wants to send message *m* to *B* via an open channel}: m=[m(1), m(2)];m(1) < n; m(2) < n;  $\{m \text{ is a integer, } not \text{ a point on } E\}$ ; Comment: m is "split" onto two parts m(1) and m(2); 5. A selects a secret key k < n and computes  $y_0 = kG$ ; {"Hint/Clue"-point on EC}; 6. A computes a "Mask/Screen/Veil"  $[c(1), c(2)] = kP_B$ ; { c(1), c(2) are coordinates of EC point};

## 7. *A* computes

$$y_1 = c_1 m_1 \mod n;$$
  
 $y_2 = c_2 m_2 \mod n;$  {Masking/Hiding/Veiling of  $m$ };

8. A sends to B ciphertext four integers

y=[y(0), y(1), y(2)] via open channel.

# **Decryption**

9. *B* computes  $n_B * y(0) = [c(1), c(2)]$ 

$$n_B*y(0)=[c(1), c(2)]$$

- 10. B finds inverses of c(1) and c(2) using FISH algorithm
- 11. *B* computes

$$m_1 = y_1 c_1^{-1} \mod n$$

$$m_2 = y_2 c_2^{-1} \bmod n$$

and recovers m(1) and m(2).

### **Numerical Example**

- 1. Both A and B select *EC*:  $y^2 = x^3 + x + 6 \pmod{11}$ ;
- 2. Both A and B agree to select generator G=(2,7);
- 3. B selects n(B)=8 and pre-computes P(B)=(3,5);
- 4. A selects a secret number k=6<11;
- 5. A pre-computes "hint/clue" kG=6\*(2,7)=(7,9);
- 6. To hide *m*, *A* pre-computes "mask/screen/veil":

$$kP(B)=6*(3,5)=(10, 9); \{point on EC\}$$

7. 
$$c(1)=10; c(2)=9;$$

- 8. Let plaintext m=91;
- 9. Represent m as m=(9,1)

 $\{m \text{ is NOT a point on elliptic curve } E\};$ 

10. A computes 
$$y(1)=(10*9) \mod 11=2$$
;

$$y(2)=(9*1)\mod 11=9;$$

- 11. A sends ciphertext  $\{(7,9); 2;9\}$  to B via open channel;
- 12. {Decryption by B}: [c(1), c(2)] = 8\*(7,9) = (10,9);
- 13.  $m_1 = 2 \times 10^{-1} \mod 11 = (2*10) \mod 11 = 9;$

$$m_2 = 9 \times 9^{-1} \mod 11 = (9*5) \mod 11 = 1.$$

#### **Homework:**

Example 1: Send plaintext m=2410=(24,10) from A to B;

Example 2: Select t=9 as a secret number for A and show how to send a plaintext m=2319 from B to A.