## DPU

Dr. D. Y. Patil Institute of Engineering & Technology

	ASSIGNMENT NO.1
	710019111
( <del>+</del> )	TITLE: Write a Java program (using OOP features) to
(")	implement following scheduling algorithms:
	implement following scheduling algorithms:  FCFS, SJF (Preemptive), Priority (Non-Preemptive)
	and Round-Robin (Preemptive).
( <del>1</del> )	OBJECTIVES: To understand the functions of operating system
	2) To learn and understand process, resource
	and memory management
<b>(</b> ★)	SOFTWARE AND HARDWARE REQUIREMENTS:
	64 bit open source linux, 05 (ubuntu), Eclipse IDE etc.
(#)	THEORY:
Æ	SCHEDULING ALGORITHM:
	CPU Scheduling deals with the problem of deading
	which of the processes in the ready queue is to be
	allocated the CPU. There are many different CPU-scheduling
	algorithms.
(1)	FIRST-COME, FIRST-SERVED SCHEDULING.
	The FCFS is the simplest CPU scheduling algorithm.
	with this scheme the process that requests the CPU
	first is allocated the CPU first the implementation of
	FCFS policy is managed with a FIFO queue when a
	process enters the ready queue, its PCB is linked onto
	tail of queue. The running process is then removed

## PPU

Dr. D. Y. Patil Institute of Engineering & Technology

	from the queue . The FCFS scheduling algorithm is
/	non-preemptive conce the CPU has been allocated to
1 100	process the process keeps CPU until it releases CPU,
<u>A</u>	either by terminating or by requesting 110.
<u> Kani</u>	
	Example: consider, (set of processes that arrive at time o)
	Process Burst Time.
-district	P1 24
00.5	P2 3
	P3 GANTT CHART
	GANII CHAKI
	CERTAIN PLA VALUE NA PROPERTO PER PROPERTO P
15 9	
	0 24 27 30.
- 1	waiting time is 0 milliseconds for process P1, 24 millisecond
	for process p2 and 27 milliseconds for process p3.
	Avg walting time = 0+24+21 If milliseronds.
par .	Sometimes and the second secon
<u> </u>	A CONTRACTOR OF THE CONTRACTOR
(2)	SHORTEST-JOB-FIRST SCHEDULING PREEMPTIVE.
	This algorithm associates with each process the leach
la di	process next CPV burst when the and en wille
150	TO MOCPESS PERT COLL LINES !
and a	me so saveduing algorithm is amually
	set of processes. It can be either preemptive or non-preemptive
· · · · · · · ·	(Prepartive)
· 1	Example: Consider the following 4 processes with the
Lang.	length of CPII huset the processes, with the
	length of CPU burst given in milligeronds.
V. C51 1	

## DPU

Dr. D. Y. Patil Institute of Engineering & Technology

Process.	Arrival time	Burst Time	
p,	AVIIVAL FILIPE	8	
Pa	0	4	
Pa		0	
Du	2	- 9	
4	2	5.	

	GA	HTT CHA	RT.		1
P)	72	P.	1.	P	=
	12 11	14	1	13	•
0 1	- 5	10		17	₹6

Process P, is started at time o, since it is the only process in queue. Process P2 arrives at time 1. The remaining time for process P, (milliseconds) is larger that time required by process B2 (4 milliseconds), so process P, is preempted and process P3 is scheduled

Aug walting time = (i0-1)+(1-1)+(17-2)+(5-3)= 26/4 = 6.5 milliseronds

(3) PRIORITY (NON-PREEMPTIVE) SUREDULING ALGORITHM:

The GJF Algorithm is a special case of general

priority scheduling Algorithm

A priority is associated with each process and CPU Ps allocated

to process with highest priority. A non-preemptive priority scheduling algorithm will simply put the new process at head of ready queue.

A major problem with priority scheduling algorithms is indefinite blocking or starvation.

Dr. D. Y. Patil Institute of Engineering & Technology

Example: consider the following set of processes,	
assumed to have arrived at time o in	
the order P1, P2, P5 with length of CPL	2
burst given in milliseconds.	
The state of the s	

Process_	Burst Time.	Priority.	1000
P,	10	3	•
P2	1		
P3	2	4	47
P4	1	5	
PS	5	2	
V			

	1 1	$\Theta$	ANTT	CHARI	M1. 14.	T		
t				ρ,	V -, x - ( , -,	Pa	B.	
	12	15	1 . 13.	11	A	3	191	
C	) 1		6		16	t.	8 19	

The average walting time is 8.2 milliseconds.

The round-robin (RR) scheduling algorithm is designed copecially for time sharing systems. It is similar to FCFS scheduling but preemption is added to enable the system to switch blw processes.

A small unit of time called time quantum or time dice is defined, A time quantum is generally from 10 to 100 milliseconds in length.

(3)

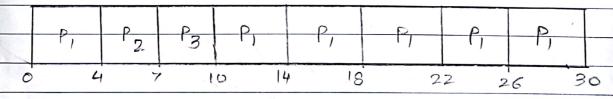
-	example: consider,
	The following set of process that arrive at time
	o, with the length of CPU burst given in
	o, with the length of CPU Dais &
	milliserande

-	Process: Buist IIIVE
	P, 24
	P <sub>2</sub> 3
	P3 3
	if me use a time quantum of 4 msec, P1 = 4 millise coods
	since it requires another so millisec, it is preempted
	after 1st quantum and CPU is given the next
	process in queue, 12. Process 12 does not need
	4 millisec so it guits before it time quantum

expires. The CPU is then given to next process, process Pa once each process has received 1 time quantum, the CPU is returned to process 81 for additional time

quantum

## GANTIT CHART



Average waiting time: P1 malts for & milliseconds

(10-4), P2 walts for 4 millisec

P3 waits 7 millisec.

: Aug walting time = 17/3 = 5:66 millisec

