

Regular Expressions

Made of $\Sigma, \epsilon, +, \circ, *$

Regular Sets

Accepted by FA

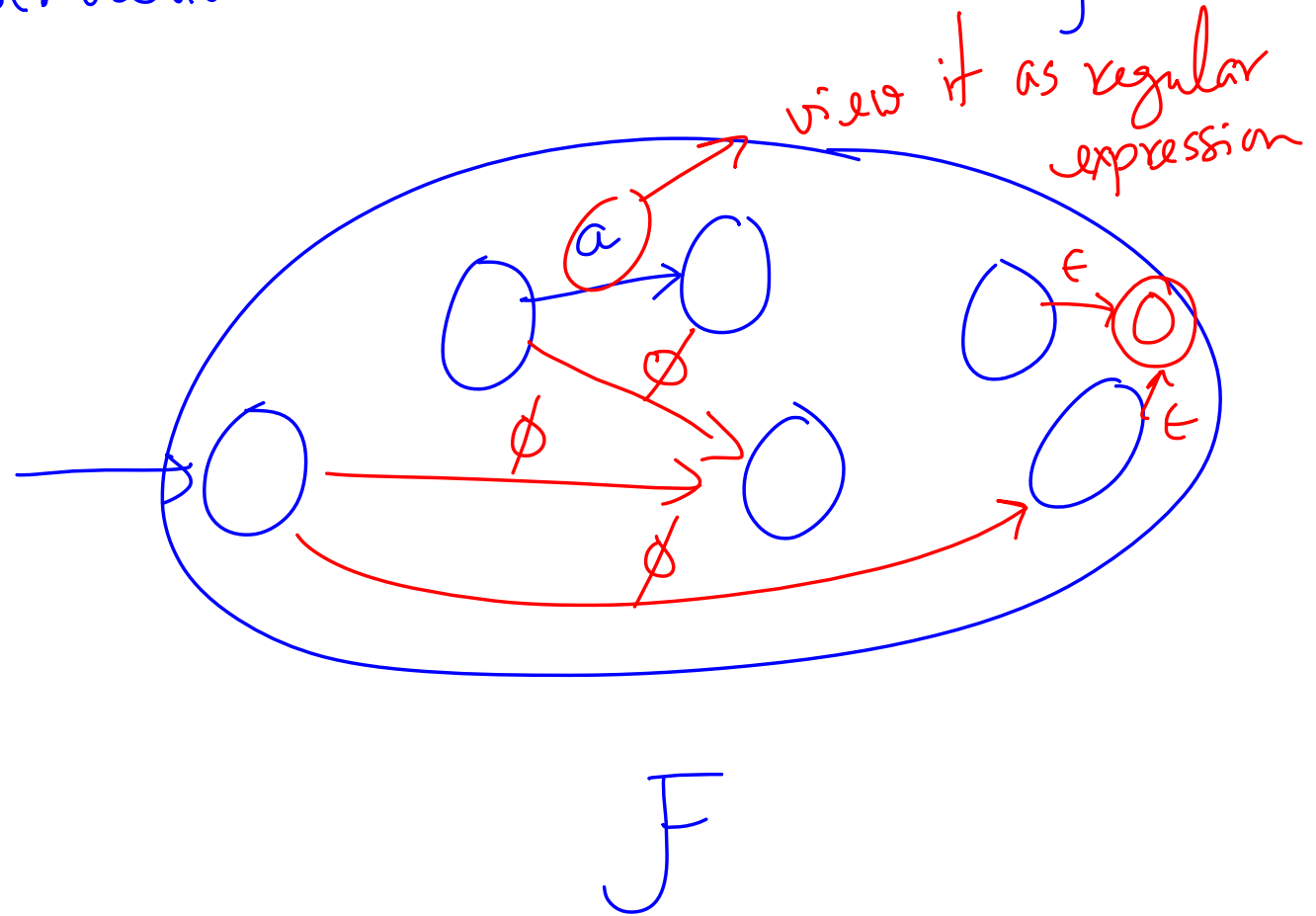
Theorem: Set A is regular if and only if
there exists a regular expression r with $A = L(r)$.

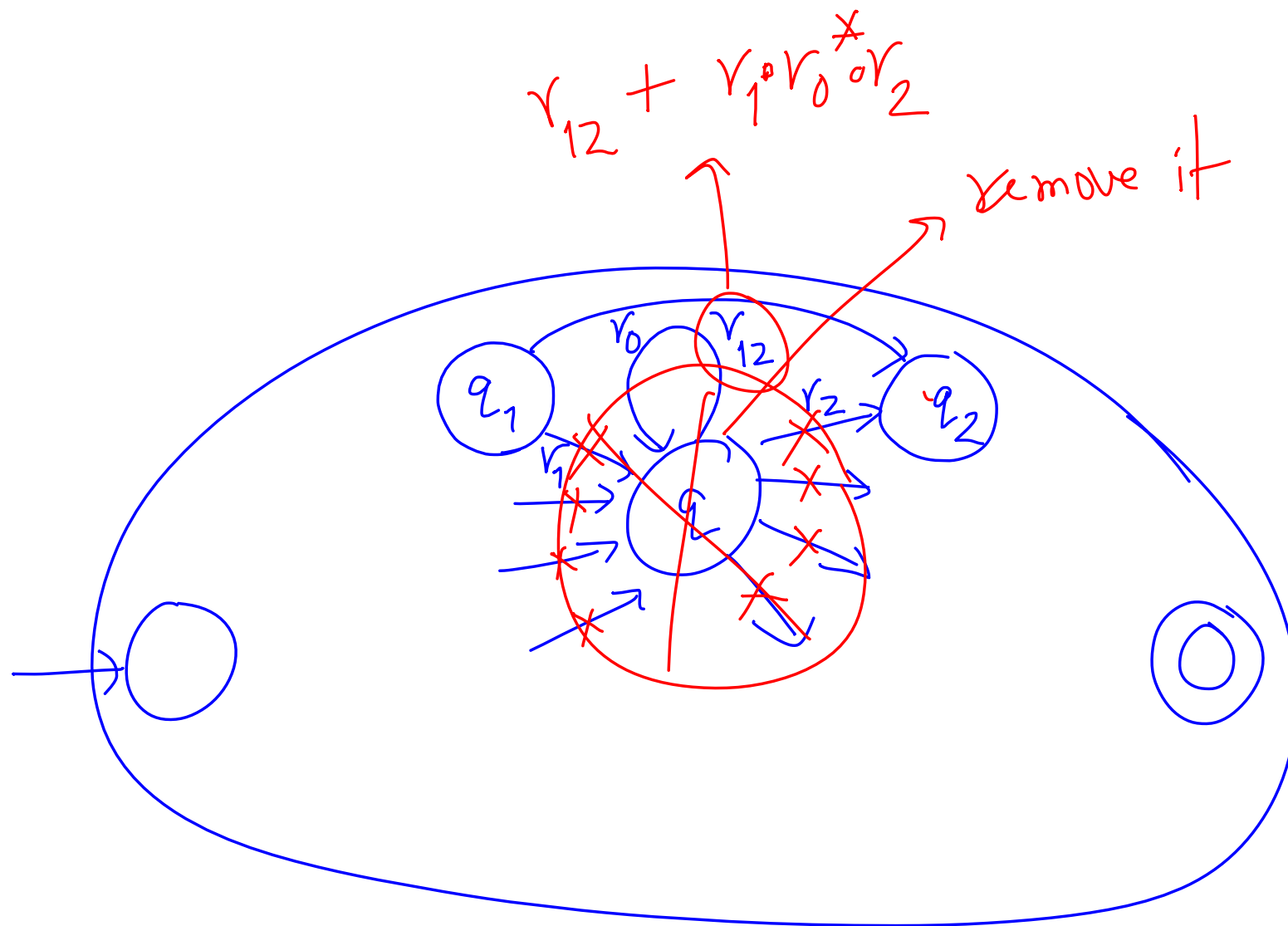
proof: Suppose A is accepted by DFA F .

A generalized NFA (GNFA) is a finite automata where transitions are done on regular expressions.

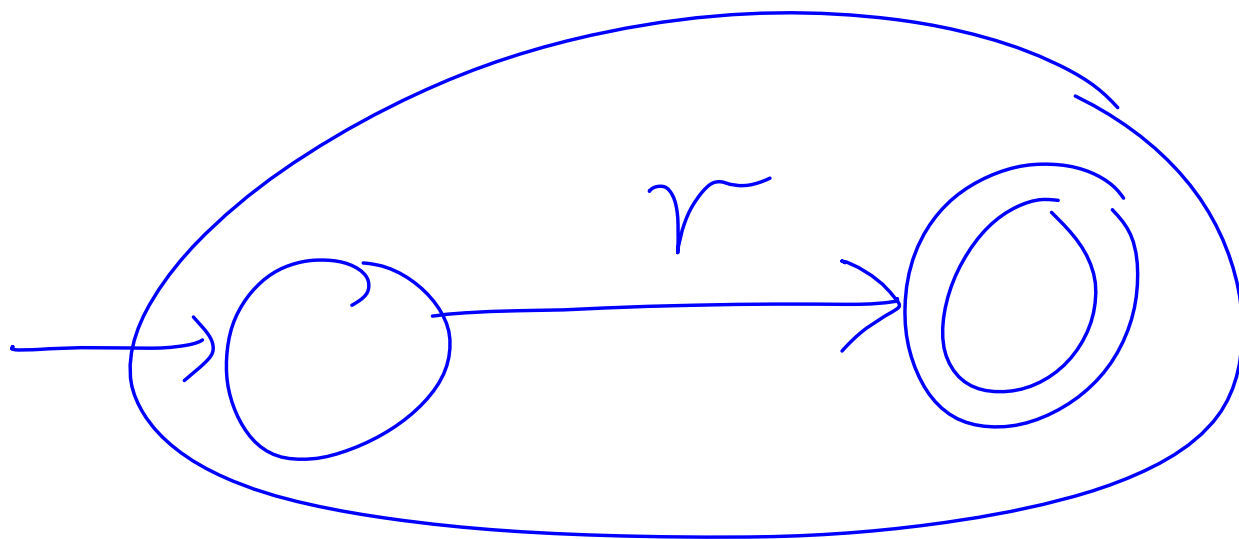
Formally, a GNFA = $(Q, \Sigma, q_0, \delta, F)$
with $\delta: Q \times R \rightarrow Q$ where R is set of all regular expressions over Σ and also includes \emptyset .

Observation: We can easily convert F to a GNFA.

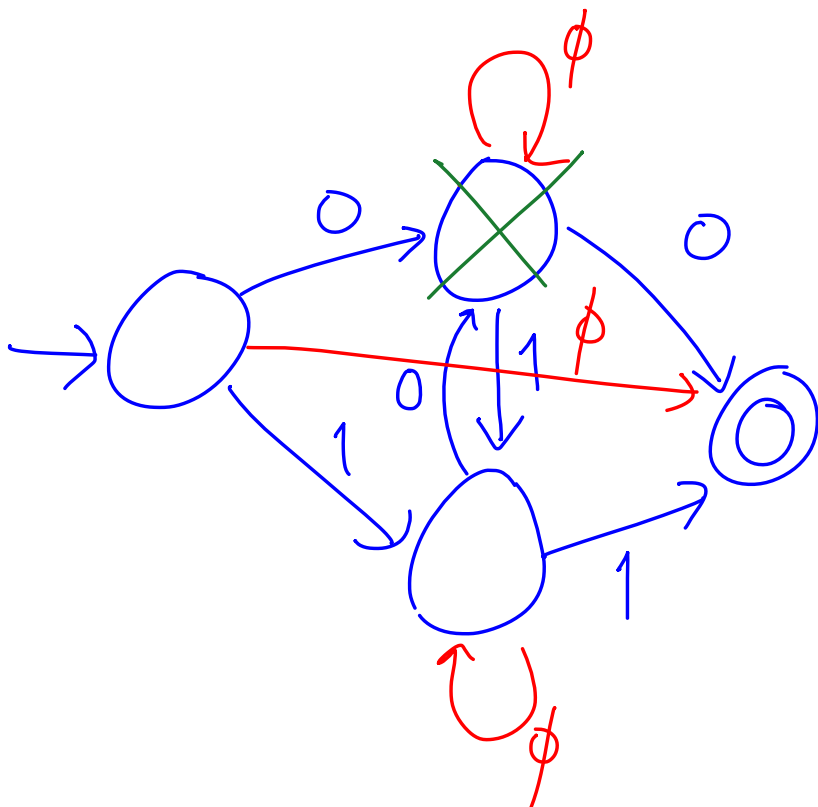




Eventually, we have:

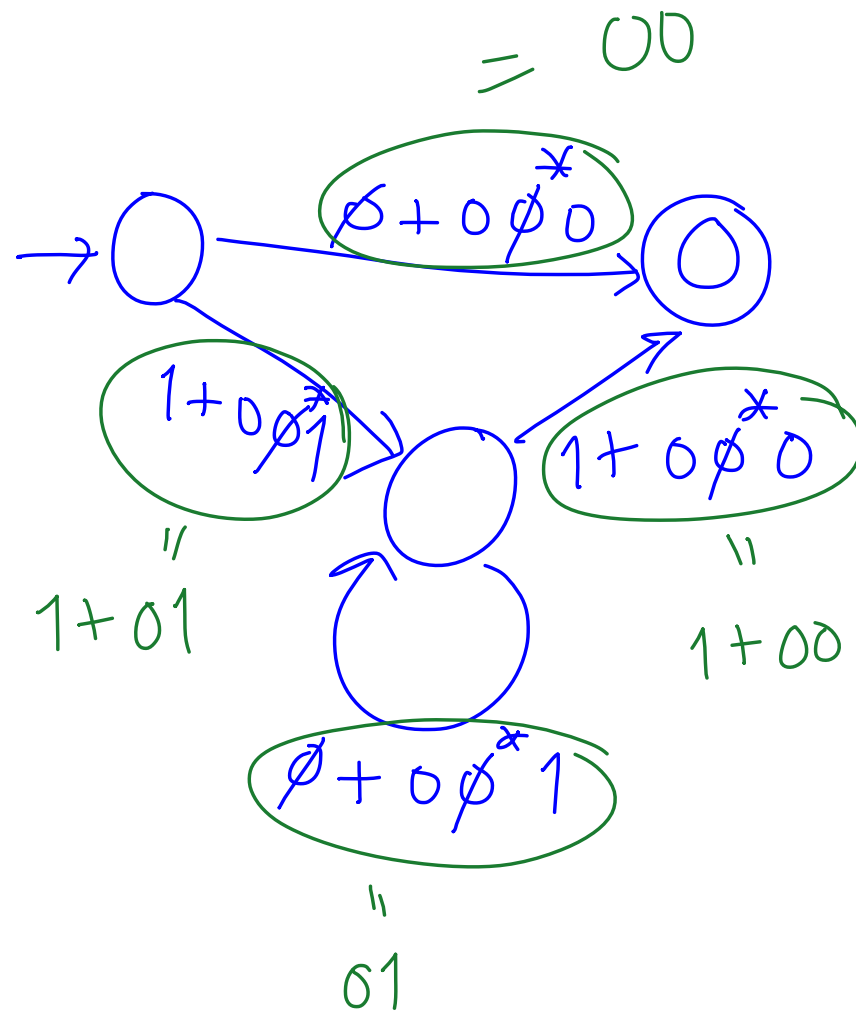


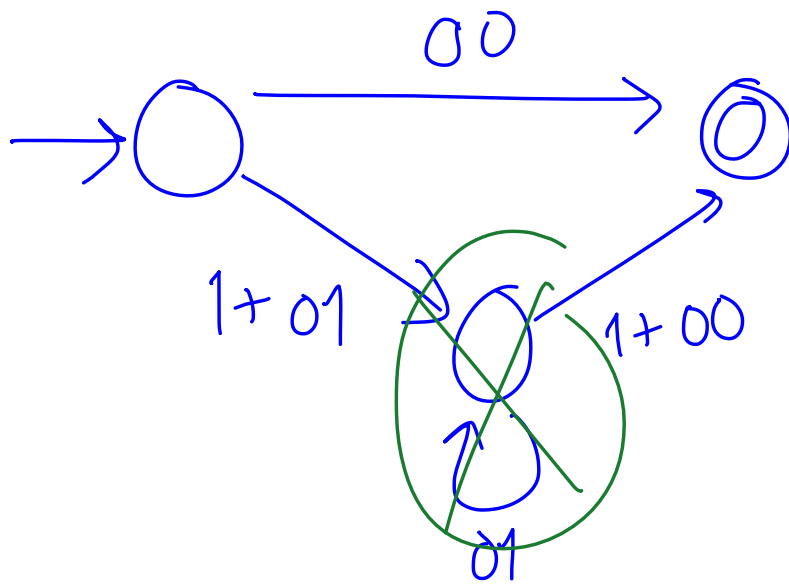
Then, $L(r) = A$.



$$\phi^* = \epsilon$$

$$\phi + s = s$$





$$\begin{array}{c}
 00 + (1+01)(01)^*(1+00) \\
 \longrightarrow \\
 \text{Start State} \longrightarrow \text{Final State}
 \end{array}$$