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## Trial of data structure

### 1.1 KD tree

```
Queries the k-th closest point in O (k * n ^ (1 - 1 / k)). Stores the data in p[]. Call init (n, k). Call min_kth (d, k) / max_kth (d, k).
template <int MAXN = 200000, int MAXK = 2>
struct kd_tree {
      int k, size;
       struct point {
    int data[MAXK], id;
} p[MAXN];
       struct kd_node {
             ict kd_node {
int l, r;
point p, dmin, dmax;
kd_node() {}
kd_node() {}
kd_node (const point &rhs) : l (0), r (0), p (rhs), dmin (rhs), dmax (rhs) {}
inline void merge (const kd_node &rhs, int k) {
    for (register int i = 0; i < k; i++) {
        dmin.data[i] = std::min (dmin.data[i], rhs.dmin.data[i]);
        dmax.data[i] = std::max (dmax.data[i], rhs.dmax.data[i]);
}
</pre>
              return ret;
              ret += 111 * tmp * tmp;
                      }
return ret:
       } tree[MAXN * 4];
      struct result {
  long long dist;
  point d;
  result() {}

              result (const long long &dist, const point &d) : dist (dist), d (d) {} bool operator > (const result &rhs) const {
    return dist > rhs.dist || (dist == rhs.dist && d.id < rhs.d.id);
              bool operator < (const result &rhs) const {
   return dist < rhs.dist || (dist == rhs.dist && d.id > rhs.d.id);
       inline long long sqrdist (const point &a, const point &b) {
   register long long ret = 0;
   for (register int i = 0; i < k; i++) {
      ret += 111 * (a.data[i] - b.data[i]) * (a.data[i] - b.data[i]);
}</pre>
              return ret;
       inline int alloc() {
    tree[size].1 = tree[size].r = 0;
    return size++;
       void build (const int &depth, int &rt, const int &l, const int &r) {
              if (1 > r) return;
register int middle = (1 + r) >> 1;
std::nth_element (p + 1, p + middle, p + r + 1,
[ = ] (const point & a, const point & b) {
    return a.data[depth] < b.data[depth];</pre>
              tree[rt = alloc()] = kd_node (p[middle]);
              if (1 == r) return;
build ((depth + 1) % k, tree[rt].1, 1, middle - 1);
```

```
build ((depth + 1) % k, tree[rt].r, middle + 1, r);
if (tree[rt].l) tree[rt].merge (tree[tree[rt].l], k);
if (tree[rt].r) tree[rt].merge (tree[tree[rt].r], k);
std::priority_queue<result, std::vector<result>, std::greater<result> > heap_1;
std::priority_queue<result, std::vector<result>, std::less<result> > heap_r;
void _min_kth (const int &depth, const int &rt, const int &m, const point &d) {
   result tmp = result (sqrdist (tree[rt].p, d), tree[rt].p);
   if ((int)heap_1.size() < m) {
       heap_1.push (tmp);
   } else if (tmp < heap_1.top()) {
       heap_1.pop();
       heap_1.push (tmp);
   }
}</pre>
         int x = tree[rt].1, y = tree[rt].r;
if (x != 0 && y != 0 && sqrdist (d, tree[x].p) > sqrdist (d, tree[y].p)) std::swap (x, y);
if (x != 0 && ((int)heap_1.size() < m || tree[x].min_dist (d, k) < heap_1.top().dist)) {
    _min_kth ((depth + 1) % k, x, m, d);
}</pre>
          if (y != 0 && ((int)heap_1.size() < m || tree[y].min_dist (d, k) < heap_1.top().dist)) {
                   _min_kth ((depth + 1) % k, y, m, d);
void _max_kth (const int &depth, const int &rt, const int &m, const point &d) {
   result tmp = result (sqrdist (tree[rt].p, d), tree[rt].p);
   if ((int)heap_r.size() < m) {</pre>
          heap_r.push (tmp);
} else if (tmp > heap_r.top()) {
  heap_r.pop();
  heap_r.push (tmp);
}
         int x = tree[rt].1, y = tree[rt].r;
if (x != 0 && y != 0 && sqrdist (d, tree[x].p) < sqrdist (d, tree[y].p)) std::swap (x, y);
if (x != 0 && ((int)heap_r.size() < m || tree[x].max_dist (d, k) >= heap_r.top().dist)) {
    _max_kth ((depth + 1) % k, x, m, d);
          if (y != 0 && ((int)heap_r.size() < m || tree[y].max_dist (d, k) >= heap_r.top().dist)) {
    _max_kth ((depth + 1) % k, y, m, d);
void init (int n, int k) {
    this -> k = k; size = 0;
    int rt = 0;
    build (0, rt, 0, n - 1);
point min_kth (const point &d, const int &m) {
   heap_1 = decltype (heap_1) ();
   _min_kth (0, 0, m, d);
   return heap_1.top ().d;
point max_kth (const point &d, const int &m) {
   heap_r = decltype (heap_r) ();
   _max_kth (0, 0, m, d);
   return heap_r.top ().d;
```

### 1.2 Link cut tree

```
/* Link Cut Tree :
    Needs formatting.
struct MsgNode {
    int leftColor, rightColor, answer;
    MsgNode() {
        leftColor = -1;
        rightColor = -1;
        rightColor = -1;
        answer = 0;
}
MsgNode (int c) {
        leftColor = rightColor = c;
        answer = 1;
}
MsgNode operator + (const MsgNode &p)const {
        if (answer == 0) return p;
        if (p.answer == 0) return *this;
        MsgNode ret;
        ret.leftColor = leftColor;
        ret.rightColor = p.rightColor;
        ret.answer = answer + p.answer - (rightColor == p.leftColor);
        return ret;
}
d[MaXN], g[MaXN];
int n, m, c[MaXN][2], f[MaXN], p[MaXN], s[MaXN], flag[MaXN];
bool r[MaXN];
void init (int x, int value) {
        d[x] = g[x] = MsgNode (value);
        c[x][0] = c[x][1] = 0;
        f[x] = p[x] = flag[x] = -1;
        s[x] = s[c[x][0]] + s[c[x][1]] + 1;
        g[x] = MsgNode();
        if (c[x][0 ^ r[x]]) g[x] = g[x] + g[c[x][0 ^ r[x]]];
        g[x] = g[x] + d[x];
        if (c[x][1 ^ r[x]]) g[x] = g[x] + g[c[x][1 ^ r[x]]];
}
void makesame (int x, int c) {
        flag[x] = c;
        d[x] = MsgNode (c);
}
```

```
g[x] = MsgNode (c);
}
void pushdown (int x) {
    if (r[x]) {
        std::swap (c[x][0], c[x][1]);
        r[c[x][0]] ^= 1;
        r[c[x][1]] ^= 1;
        std::swap (g[c[x][0]].leftColor, g[c[x][0]].rightColor);
        std::swap (g[c[x][1]].leftColor, g[c[x][1]].rightColor);
        r[x] = false;
}
              fif (flag[x] != -1) {
    if (c[x][0]) makesame (c[x][0], flag[x]);
    if (c[x][1]) makesame (c[x][1], flag[x]);
    flag[x] = -1;
              }
 food rotate (int x, int k) {
   pushdown (x); pushdown (c[x][k]);
   int y = c[x][k]; c[x][k] = c[y][k ^ 1]; c[y][k ^ 1] = x;
   if (f[x]]!= -1) c[f[x]][c[f[x]][1] == x] = y;
   f[y] = f[x]; f[x] = y; f[c[x][k]] = x; std::swap (p[x], p[y]);
   update (x); update (y);
}

}
void splay (int x, int s = -1) {
   pushdown (x);
   while (f[x] != s) {
      if (f[f[x]] != s) rotate (f[f[x]], (c[f[f[x]]][1] == f[x]) ^ r[f[f[x]]]);
      rotate (f[x], (c[f[x]][1] == x) ^ r[f[x]]);
}

}
void access (int x) {
   int y = 0;
   while (x != -1) {
      splay (x); pushdown (x);
      f[c[x][1]] = -1; p[c[x][1]] = x;
      c[x][1] = y; f[y] = x; p[y] = -1;
      update (x); x = p[y = x];
}
              }

}
void setroot (int x) {
    access (x);
    splay (x);
    r[x] ^= 1;
    std::swap (g[x].leftColor, g[x].rightColor);
}

 void link (int x, int y) {
    setroot (x);
    p[x] = y;
}
void cut (int x, int y) {
    access (x); splay (y, -1);
    if (p[y] == x) p[y] = -1;
    else {
        access (y);
        splay (x, -1);
        p[x] = -1;
    }
}
```

## Trial of number theory

### 2.1 Constants and basic functions

### 2.2 Discrete Fourier transform

```
/* Discrete Fourier transform :
    int dft::init (int n) :
        initializes the transformation with dimension n.
        Returns the recommended size.
    void dft::solve (complex *a, int n, int f) :
        transforms array a with dimension n to its image representation.
        Transforms back when f = 1. (n should be 2^k)

*/
template <int MAXN = 1000000>
struct dft {
    typedef std::complex <double> complex;
    complex e[2][MAXN];
```

```
int init (int n) {
    int len = 1;
    for (; len <= 2 * n; len <<= 1);
    for (int i = 0; i < len; i++) {
        e[0][i] = complex (cos (2 * PI * i / len), sin (2 * PI * i / len));
        e[1][i] = complex (cos (2 * PI * i / len), -sin (2 * PI * i / len));
    }
    return len;
}

void solve (complex *a, int n, int f) {
    for (int i = 0, j = 0; i < n; i++) {
        if (i > j) std::swap (a[i], a[j]);
        for (int t = n >> 1; (j ^= t) < t; t >>= 1);
}

for (int i = 2; i <= n; i <<= 1)
    for (int j = 0; j < n; j += i)
        for (int k = 0; k < (i >> 1); k++) {
            complex A = a[j + k];
            complex B = e[f][n / i * k] * a[j + k + (i >> 1)];
        a[j + k] = A + B;
        a[j + k] = A + B;
        a[j + k] = (int i = 0; i < n; i++)
        a[i] = complex (a[i].real () / n, a[i].imag ());
}

/* Number-theoretic transform :
    void ntt::solve (int *a, int n, int f, int mod, int prt) :</pre>
```

### 2.3 Number-theoretic transform

```
/* Number-theoretic transform :
    void ntt::solve (int *a, int n, int f, int mod, int prt) :
        transforms a[n] to its image representation.
        Converts back if f = 1. (n should be 2^k)
        Requries specific mod and corresponding prt to work. (given in MOD and PRT)
    int ntt::crt (int *a, int mod) :
        fixes the results a from module 3 primes to a certain module mod.
template <int MAXN = 1000000>
struct ntt [
           void solve (int *a, int n, int f, int mod, int prt) {
    for (register int i = 0, j = 0; i < n; i++) {
        if (i > j) std::swap (a[i], a[j]);
        for (register int t = n >> 1; (j ^= t) < t; t >>= 1);
                        for (register int i = 2; i <= n; i <<= 1) {</pre>
                                   (register int 1 = 2; 1 <= n; 1 <<= 1) {
    static int exp[MAXN];
    exp[0] = 1;
    exp[1] = fpm (prt, (mod - 1) / i, mod);
    if (f == 1) exp[1] = fpm (exp[1], mod - 2, mod);
    for (register int k = 2; k < (i >> 1); k++) {
        exp[k] = int (111 * exp[k - 1] * exp[1] % mod);
    }
}
                                    for (register int j = 0; j < n; j += i) {
    for (register int k = 0; k < (i >> 1); k++) {
        register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
        register int A = pA, B = int (111 * pB * exp[k] % mod);
        pA = (A + B) % mod;
        pB = (A - B + mod) % mod;
}
                                   }
                       }
if (f == 1) {
    register int rev = fpm (n, mod - 2, mod);
    for (register int i = 0; i < n; i++) {
        a[i] = int (111 * a[i] * rev % mod);
}</pre>
            int MOD[3] = {1045430273, 1051721729, 1053818881}, PRT[3] = {3, 6, 7};
           }
int sum = 1, ret = x[0] % mod;
for (int i = 1; i < 3; i ++) {
    sum = int (1LL * sum * MOD[i - 1] % mod);
    ret += int (1LL * x[i] * sum % mod);
    if (ret >= mod) ret -= mod;
}
                        return ret;
            }
};
```

#### 2.4 Chinese remainder theorem

### 2.5 Baby step giant step algorithm

### 2.6 Miller Rabin primality test

## 2.7 Pollard's Rho algorithm

### 2.8 Adaptive Simpson's method

**}**;

## Trial of geometry

### 3.1 Constants and basic functions

```
/* Constants & basic functions :
    EPS : fixes the possible error of data.
        i.e. x == y iff |x - y| < EPS.
    PI : the value of PI.
    int sgn (const double &x) : returns the sign of x.
    int cmp (const double &x, const double &y) : returns the sign of x - y.
    double sqr (const double &x) : returns x * x.

*/
const double EPS = 1E-8;
const double PI = acos (-1);
int sgn (const double &x) { return x < -EPS ? -1 : x > EPS; }
int cmp (const double &x, const double &y) { return sgn (x - y); }
double sqr (const double &x) { return x * x; }
```

### 3.2 Point class

```
/* struct point : defines a point and its various utility.
                    point : defines a point and its various utility.
point (const double &x, const double &y) gives a point at (x, y).
It also represents a vector on a 2D plane.
point unit () const : returns the unit vector of (x, y).
point rot90 () const :
    returns a point rotated 90 degrees counter-clockwise with respect to the origin.
point _rot () const : same as above except clockwise.
point rotate (const double &t) const : returns a point rotated t radian(s) counter-clockwise.
Operators are mostly vector operations. i.e. vector +, -, *, / and dot/det product.
*/
struct point {
    double x, y;
    explicit point (const double &x = 0, const double &y = 0) : x (x), y (y) {}
    double norm () const { return sqrt (x * x + y * y); }
    double norm2 () const { return x * x + y * y; }
    point unit () const {
        double l = norm ();
        return point (x / 1, y / 1);
    }
}
          point rot90 () const {return point (-y, x); }
point _rot90 () const {return point (y, -x); }
point rotate (const double &t) const {
    double c = cos (t), s = sin (t);
    return point (x * c - y * s, x * s + y * c);
}
bool operator == (const point &a, const point &b) {
   return cmp (a.x, b.x) == 0 && cmp (a.y, b.y) == 0;
bool operator != (const point &a, const point &b) {
   return ! (a == b);
bool operator < (const point &a, const point &b) {
   if (cmp (a.x, b.x) == 0) return cmp (a.y, b.y) < 0;
   return cmp (a.x, b.x) < 0;</pre>
point operator - (const point &a) { return point (-a.x, -a.y); }
point operator + (const point &a, const point &b) {
   return point (a.x + b.x, a.y + b.y);
point operator - (const point &a, const point &b) {
   return point (a.x - b.x, a.y - b.y);
point operator * (const point &a, const double &b) {
   return point (a.x * b, a.y * b);
point operator / (const point &a, const double &b) {
   return point (a.x / b, a.y / b);
double dot (const point &a, const point &b) {
   return a.x * b.x + a.y * b.y;
double det (const point &a, const point &b) {
    return a.x * b.y - a.y * b.x;
```

```
}
double dis (const point &a, const point &b) {
   return sqrt (sqr (a.x - b.x) + sqr (a.y - b.y));
}
```

### 3.3 Line class

### 3.4 Interactions between points and lines

```
bool point_on_line (const point &a, const line &b) {
   return sgn (det (a - b.s, b.t - b.s)) == 0 && sgn (dot (b.s - a, b.t - a)) <= 0;</pre>
bool two_side (const point &a, const point &b, const line &c) {
   return sgn (det (a - c.s, c.t - c.s)) * sgn (det (b - c.s, c.t - c.s)) < 0;</pre>
bool intersect judgement (const line &a, const line &b) {
        if (point_on_line (b.s, a) || point_on_line (b.t, a)) return true;
if (point_on_line (a.s, b) || point_on_line (a.t, b)) return true;
return two_side (a.s, a.t, b) && two_side (b.s, b.t, a);
point line_intersect (const line &a, const line &b) {
    double s1 = det (a.t - a.s, b.s - a.s);
    double s2 = det (a.t - a.s, b.t - a.s);
    return (b.s * s2 - b.t * s1) / (s2 - s1);
 double point_to_line (const point &a, const line &b) {
   return fabs (det (b.t - b.s, a - b.s)) / dis (b.s, b.t);
double point_to_segment (const point &a, const line &b) {
    if (sgn (dot (b.s - a, b.t - b.s) * dot (b.t - a, b.t - b.s)) <= 0)
        return fabs (det (b.t - b.s, a - b.s)) / dis (b.s, b.t);
    return std::min (dis (a, b.s), dis (a, b.t));</pre>
*/
if (point_on_line (p, line (a, b))) return true;
int x = sgn (det (p - a, b - a)), y = sgn (a.y - p.y), z = sgn (b.y - p.y);
if (x > 0 && y <= 0 && z > 0) counter++;
if (x < 0 && z <= 0 && y > 0) counter--;
        return counter != 0;
double polygon_area (const std::vector <point> &a) {
    double ans = 0.0;
    for (int i = 0; i < (int) a.size (); ++i)
        ans += det (a[i], a[ (i + 1) % a.size ()]) / 2.0;
    return ans;</pre>
point project_to_line (const point &a, const line &b) {
    return b.s + (b.t - b.s) * (dot (a - b.s, b.t - b.s) / (b.t - b.s).norm2 ());
```

## 3.5 Centers of a triangle

```
point circumcenter (const point &a, const point &b, const point &c) {
   point p = b - a, q = c - a, s (dot (p, p) / 2, dot (q, q) / 2);
   double d = det (p, q);
   return a + point (det (s, point (p.y, q.y)), det (point (p.x, q.x), s)) / d;
}
point orthocenter (const point &a, const point &b, const point &c) {
   return a + b + c - circumcenter (a, b, c) * 2.0;
}
```

### 3.6 Fermat point

### 3.7 Circle class

```
/* struct circle defines a circle.
    circle (point c, double r) gives a circle with center c and radius r.
*/
struct circle {
    point c;
    double r;
    explicit circle (point c = point (), double r = 0) : c (c), r (r) {}
};
bool operator == (const circle &a, const circle &b) {
    return a.c == b.c && cmp (a.r, b.r) == 0;
}
bool operator != (const circle &a, const circle &b) {
    return ! (a == b);
}
```

#### 3.8 Interactions of circles

```
std::pair <line, line> tangent (const point &a, const circle &b) {
   circle p = make_circle (a, b.c);
   auto d = circle_intersect (p, b);
   return std::make_pair (line (d.first, a), line (d.second, a));
```

### 3.9 Convex hull

### 3.10 Minimum circle

### 3.11 Half plane intersection

```
#/
bool turn left (const line &1, const point &p) {
    return turn_left (l.s, l.t, p);
}

std::vector <point> half plane intersect (std::vector <line> h) {
    typedef std::pair <double, line> polar;
    std::vector <point> half plane intersect (std::vector <line> h) {
        typedef std::pair <double, line> polar;
        std::vector <point> (int) i = 0; i < (int) h.size ();
        for (int) i = 0; i < (int) h.size ();
        for (int) i = 0; i < (int) h.size ();
        for (int) i = 0; i < (int) h.size ();
        for (int) i = 0; i < (int) h.size ();
        if (const polar &a, const polar &b) {
            if (const aske pair (akan2 (v.y, v.x), h[i]);
        }
        sort (g.begin (), g.end (), [] (const polar &a, const polar &b) {
            if (const aske pair (aske pair (a
```

### 3.12 Intersection of a polygon and a circle

```
Intersection of a polygon and a circle :
             double polygon_circle_intersect::solve (const std::vector <point> &p, const circle &c) :
    returns the area of intersection of polygon p (vertices in either order) and c.
struct polygon_circle_intersect {
      // The area of the sector with center (0, 0), radius r and segment ab.
      double sector_area (const point &a, const point &b, const double &r) {
   double c = (2.0 * r * r - (a - b).norm2 ()) / (2.0 * r * r);
   double al = acos (c);
   return r * r * al / 2.0;
       . The area of triangle (a, b, (0, 0)) intersecting circle (point (), r).
      double area (const point &a, const point &b, const double &r) {
  double dA = dot (a, a), dB = dot (b, b), dC = point_to_segment (point (), li:
    if (sgn (dA - r * r) <= 0 && sgn (dB - r * r) <= 0) return det (a, b) / 2.0;
  point tA = a.unit () * r;
  point tB = b.unit () * r;
  if (sgn (dC - r) > 0) return sector_area (tA, tB, r);
  if (sgn (dC - r) > 0) return sector_area (tA, tB, r);
                                                                                                                                                line (a, b), ans = 0.0;
             ir (sgn (dC - r) > 0) return sector_area (tA, tB, r);
std::pair <point, point> ret = line_circle_intersect (line (a, b), circle (point (), r));
if (sgn (dA - r * r) > 0 && sgn (dB - r * r) > 0) {
    ans += sector_area (tA, ret.first, r);
    ans += det (ret.first, ret.second) / 2.0;
    ans += sector_area (ret.second, tB, r);
    return ans;
}
              if (sgn (dA - r * r) > 0)
   return det (ret.first, b) / 2.0 + sector_area (tA, ret.first, r);
                     return det (a, ret.second) / 2.0 + sector_area (ret.second, tB, r);
      // Main procedure.
      ret += area (p[i] - c.c, p[ (i + 1) % p.size ()] - c.c, c.r);
                     else ret -= area (p[ (i + 1) % p.size ()] - c.c, p[i] - c.c, c.r);
              return fabs (ret);
};
```

### 3.13 Union of circles

```
/* Union of circles :
```

```
std::vector <double> union_circle::solve (const std::vector <circle> &c) :
                            returns the union of circle set c. The i-th element is the area covered with at least i circles.
struct union_circle {
        struct cp {
   double x, y, angle;
                  double x, y, angle,
int d;
double r;
cp (const double &x = 0, const double &y = 0, const double &angle = 0,
    int d = 0, const double &r = 0) : x (x), y (y), angle (angle), d (d), r (r) {}
         double dis (const cp &a, const cp &b) {
    return sqrt (sqr (a.x - b.x) + sqr (a.y - b.y));
         double cross (const cp &p0, const cp &p1, const cp &p2) {
    return (p1.x - p0.x) * (p2.y - p0.y) - (p1.y - p0.y) * (p2.x - p0.x);
        int cir_cross (cp p1, double r1, cp p2, double r2, cp &cp1, cp &cp2) {
    double mx = p2.x - p1.x, sx = p2.x + p1.x, mx2 = mx * mx;
    double my = p2.y - p1.y, sy = p2.y + p1.y, my2 = my * my;
    double sq = mx2 + my2, d = - (sq - sqr (r1 - r2)) * (sq - sqr (r1 + r2));
    if (sgn (d) < 0) return 0;
    if (sgn (d) <= 0) d = 0;
    else d = sqrt (d);
    double x = mx * ((r1 + r2) * (r1 - r2) + mx * sx) + sx * my2;
    double y = my * ((r1 + r2) * (r1 - r2) + my * sy) + sy * mx2;
    double dx = mx * d, dy = my * d;
    sq *= 2:</pre>
                  double dx = mx * d, dy = m

sq *= 2;

cp1.x = (x - dy) / sq;

cp1.y = (y + dx) / sq;

cp2.x = (x + dy) / sq;

cp2.y = (y - dx) / sq;

if (sgn (d) > 0) return 2;

else return 1;
         static bool circmp (const cp &u, const cp &v) {
   return sgn (u.r - v.r) < 0;</pre>
         static bool cmp (const cp &u, const cp &v) {
   if (sgn (u.angle - v.angle)) return u.angle < v.angle;
   return u.d > v.d;
        std::vector <double> solve (const std::vector <circle> &c) {
  int n = c.size ();
  std::vector <cp> cir, tp;
  std::vector <double> area;
  cir recirc (n):
                  std::vector <double> area;
cir.resize (n);
tp.resize (2 * n);
area.resize (n + 1);
for (int i = 0; i < n; i++)
    cir[i] = cp (c[i].c.x, c[i].c.y, 0, 1, c[i].r);
cp cp1, cp2;
std::sort (cir begin () cir end () circmn);</pre>
                  cp2.angle = atan2 (cp2.y = cff[1].y, cp2.x cp1.d = 1;
tp[tn++] = cp1;
cp2.d = -1;
tp[tn++] = cp2;
if (sgn (cp1.angle - cp2.angle) > 0) cnt++;
                            fp[tn++] = cp (cir[i].x - cir[i].r, cir[i].y, PI, -cnt);
tp[tn++] = cp (cir[i].x - cir[i].r, cir[i].y, -PI, cnt);
std::sort (tp.begin (), tp.begin () + tn, cmp);
int p, s = cir[i].d + tp[0].d;
for (int j = 1; j < tn; ++j) {
    p = s;
}</pre>
                                     p = 0,
s += tp[j].d;
area[p] += calc (cir[i], tp[j - 1], tp[j]);
                  return area;
        }
1;
```

## Trial of graph

### 4.1 Constants and edge lists

```
const int INF = 1E9;
template <int MAXN = 100000, int MAXM = 100000>
struct edge_list {
       int size;
int begin[MAXN], dest[MAXM], next[MAXM];
       void clear (int n) {
    size = 0;
    std::fill (begin, begin + n, -1);
        edge_list (int n = MAXN) {
    clear (n);
       void add_edge (int u, int v) {
   dest[size] = v; next[size] = begin[u]; begin[u] = size++;
template <int MAXN = 100000, int MAXM = 100000>
template tint MARN = 1000000, int MARN = 10000000
struct cost_edge_list {
  int size;
  int begin[MAXN], dest[MAXM], next[MAXM], cost[MAXM];
  void clear (int n) {
     size = 0;
     std::fill (begin, begin + n, -1);
}
       cost_edge_list (int n = MAXN) {
    clear (n);
        void add_edge (int u, int v, int c) {
    dest[size] = v; next[size] = begin[u]; cost[size] = c; begin[u] = size++;
};
template <int MAXN = 100000, int MAXM = 100000>
struct flow_edge_list {
       int size;
int begin[MAXN], dest[MAXM], next[MAXM], flow[MAXM], inv[MAXM];
void clear (int n) {
    size = 0;
    std::fill (begin, begin + n, -1);
       flow_edge_list (int n = MAXN) {
    clear (n);
        void add_edge (int u, int v, int f) {
    dest[size] = v; next[size] = begin[u]; flow[size] = f; inv[size] = size + 1; begin[u] = size++;
    dest[size] = u; next[size] = begin[v]; flow[size] = 0; inv[size] = size - 1; begin[v] = size++;
template <int MAXN = 100000, int MAXM = 100000>
template <int maxx = 100000, Int MAXX = 100000;
struct cost_flow_edge_list {
  int size;
  int begin[MAXN], dest[MAXM], next[MAXM], cost[MAXM], flow[MAXM], inv[MAXM];
  void clear (int n) {
     size = 0;
     std::fill (begin, begin + n, -1);
}</pre>
       cost_flow_edge_list (int n = MAXN) {
    clear (n);
       void add_edge (int u, int v, int c, int f) {
    dest[size] = v; next[size] = begin[u]; cost[size] = c;
    flow[size] = f; inv[size] = size + 1; begin[u] = size++;
    dest[size] = u; next[size] = begin[v]; cost[size] = c;
    flow[size] = 0; inv[size] = size - 1; begin[v] = size++;
}
} ;
```

## 4.2 SPFA improved

```
/* SPFA :
```

### 4.3 Dijkstra's shortest path algorithm

## 4.4 Tarjan

### 4.5 Vertex biconnected component

### 4.6 Edge biconnected component

## 4.7 Hopcoft-Carp

```
/* Hopcoft-Carp algorithm :
```

```
unweighted maximum matching for bipartition graphs with complexity O (m * n^0.5).
                 struct hopcoft_carp :
Usage : solve () for maximum matching. The matching is in matchx and matchy.
*/
template <int MAXN = 100000, int MAXM = 100000>
struct hopcoft_carp {
        int n, m;
        int matchx[MAXN], matchy[MAXN], level[MAXN];
        bool dfs (edge_list <MAXN, MAXM> &e, int x) {
   for (int i = e.begin[x]; ~i; i = e.next[i]) {
                           int y = e.dest[i];
int w = matchy[y];
                           if (w == -1 || (level[x] + 1 == level[w] && dfs (e, w))) {
   matchx[x] = y;
   matchy[y] = x;
   return true;
                 level[x] = -1;
return false;
        fint solve (edge_list <MAXN, MAXM> &e, int n, int m) {
   std::fill (matchx, matchx + n, -1);
   std::fill (matchy, matchy + m, -1);
   for (int answer = 0; ;) {
      std::vector <int> queue;
      for (int i = 0; i < n; ++i) {
        if (matchx[i] == -1) {
            level[i] = 0;
            queue.push_back (i);
      } else {</pre>
                                   } else {
   level[i] = -1;
                           for (int head = 0; head < (int) queue.size(); ++head) {
   int x = queue[head];
   for (int i = e.begin[x]; ~i; i = e.next[i]) {</pre>
                                            int y = e.desf[x],
int w = matchy[y];
if (w != -1 && level[w] < 0) {
    level[w] = level[x] + 1;
    queue.push_back (w);
}</pre>
                                   }
                          for (int i = 0;
    for (int i = 0; i < n; ++i)
        if (matchx[i] == -1 && dfs (e, i)) delta++;
if (delta == 0) return answer;
else answer += delta;</pre>
        }
1:
```

### 4.8 Kuhn-Munkres

### 4.9 Stochastic weighted maximum matching

```
Weighted matching algorithm :
    maximum matching for general weighted graphs. Not stable.
    struct weighted_match :
                            Usage: Set k to the size of vertices, w to the weight matrix. Note that k has to be even for the algorithm to work.
template <int MAXN = 500>
struct weighted_match {
        int k;
long long w[MAXN][MAXN];
int match[MAXN], path[MAXN], p[MAXN], long long d[MAXN];
bool v[MAXN];
      }
--len;
v[i] = false;
return false;
        long long solve () {
    if (k & 1) ++k;
    for (int i = 0; i < k; ++i) p[i] = i, match[i] = i ^ 1;
    int cnt = 0;
    for (;;) {
        len = 0;
        bool flag = false;
}</pre>
                           len = 0;
bool flag = false;
std::fill (d, d + k, 0);
std::fill (v, v + k, 0);
for (int i = 0; i < k; ++i) {
    if (dfs (p[i])) {
        flag = true;
        int t = match[path[len - 1]], j = len - 2;
        while (path[j] != path[len - 1]) {
            match[t] = path[j];
            std::swap (t, match[path[j]]);
            --i;</pre>
                                                match[t] = path[j];
match[path[j]] = t;
                            }
if (!flag) {
    if (++cnt >= 2) break;
    std::random_shuffle (p, p + k);
                  }
long long ans = 0;
for (int i = 0; i < k; ++i
    ans += w[i][match[i]];
return ans / 2;</pre>
         }
};
```

## 4.10 Weighted blossom (vfleaking ver.)

```
};
int n, n_x;
~[MAX]
 int n, n_x;
edge g[MAXN * 2 + 1][MAXN * 2 + 1];
int lab[MAXN * 2 + 1];
int match[MAXN * 2 + 1], slack[MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 + 1];
int flower_from[MAXN * 2 + 1][MAXN + 1], S[MAXN * 2 + 1], vis[MAXN * 2 + 1];
std::vector<int> flower[MAXN * 2 + 1];
 std::queue<int> q;
int e_delta (const edge &e) {
   return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
 void update_slack (int u, int x) {
    if (!slack[x] || e_delta (g[u][x]) < e_delta (g[slack[x]][x]))slack[x] = u;</pre>
 }
void set_slack (int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
        if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)update_slack (u, x);
 }
void set_st (int x, int b) {
    st[x] = b;
    if (x > n) for (size_t i = 0; i < flower[x].size(); ++i)
        set_st (flower[x][i], b);</pre>

}
int get_pr (int b, int xr) {
    int pr = find (flower[b].begin(), flower[b].end(), xr) - flower[b].begin();
    if (pr % 2 == 1) {
        reverse (flower[b].begin() + 1, flower[b].end());
        return (int)flower[b].size() - pr;
}
}

}
void set_match (int u, int v) {
    match[u] = g[u][v].v;
    if (u > n) {
        edge e = g[u][v];
        int xr = flower_from[u][e.u], pr = get_pr (u, xr);
        for (int i = 0; i < pr; ++i)set_match (flower[u][i], flower[u][i ^ 1]);
        set_match (xr, v);
        rotate (flower[u].begin(), flower[u].begin() + pr, flower[u].end());
}
</pre>

}
void augment (int u, int v) {
    for (;;) {
        int xnv = st[match[u]];
        set_match (u, v);
        if (!xnv)return;
        set_match (xnv, st[pa[xnv]]);
        u = st[pa[xnv]], v = xnv;
}

 if (u == 0) continue;
if (vis[u] == t) return u;
vis[u] = t;
u = st[match[u]];
if (u) u = st[pa[u]];
           return 0;
set_slack (b);
set_stack (b),

void expand_blossom (int b) {
    for (size_t i = 0; i < flower[b].size(); ++i)
        set_st (flower[b][i], flower[b][i]);
    int xr = flower_from[b][g[b][pa[b]].u], pr = get_pr (b, xr);
    for (int i = 0; i < pr; i += 2) {
        int xs = flower[b][i], xns = flower[b][i + 1];
        pa[xs] = g[xns][xs].u;
        S[xs] = 1, S[xns] = 0;
        slack[xs] = 0, set_slack (xns);
        q_push (xns);
}</pre>
           for (size_t i = pr + 1; i < flower[b].size(); ++i) {</pre>
```

```
int xs = flower[b][i];
S[xs] = -1, set_slack (xs);
                         st[b] = 0;
            bool on_found_edge (const edge &e) {
                       l on_found_edge (const edge &e) {
int u = st[e.u], v = st[e.v];
if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push (nu);
} else if (S[v] == 0) {
    int lca = get_lca (u, v);
    if (!lca)return augment (u, v), augment (v, u), true;
    else add_blossom (u, lca, v);
}
                         return false;
         le (q.size()) {
   int u = q.front();
   q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)
      if (g[u][v].w > 0 && st[u] != st[v]) {
        if (e_delta (g[u][v]) == 0) {
            if (on_found_edge (g[u][v])) return true;
        } else update_slack (u, st[v]);
   }

}
int d = INF;
for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b && S[b] == 1)d = std::min (d, lab[b] / 2);
for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x]) {
        if (S[x] == -1)d = std::min (d, e_delta (g[slack[x]][x]));
        else if (S[x] == 0)d = std::min (d, e_delta (g[slack[x]][x]) / 2);
}
</pre>
                                                 } (int u = 1; u <= n; ++u) {
if (S[st[u]] == 0) {
    if (lab[u] <= d)return 0;
    lab[u] -= d;
} else if (S[st[u]] == 1)lab[u] += d;</pre>
                                     }
for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b) {
        if (S[st[b]] == 0)lab[b] += d * 2;
        else if (S[st[b]] == 1)lab[b] -= d * 2;</pre>
                                    q = std::queue<int>();
for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x] && st[slack[x]] != x && e_delta (g[slack[x]][x]) == 0)
        if (on_found_edge (g[slack[x]][x]))return true;
for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)expand_blossom (b);</pre>
                         }
return false;
            std::pair <long long, int> solve () {
    std::fill (match + 1, match + n + 1, 0);
    n_x = n;
                         int n_matches = 0;
                        int n_matches = 0;
long long tot_weight = 0;
for (int u = 0; u <= n; ++u)st[u] = u, flower[u].clear();
int w_max = 0;
for (int u = 1; u <= n; ++u)
    for (int v = 1; v <= n; ++v) {
        flower_from[u][v] = (u == v ? u : 0);
        w_max = std::max (w_max, g[u][v].w);
}</pre>
                        for (int u = 1; u <= n; ++u)lab[u] = w_max;
while (matching())++n_matches;
for (int u = 1; u <= n; ++u)
    if (match[u] && match[u] < u)
        tot_weight += g[u][match[u]].w;
return std::make_pair (tot_weight, n_matches);</pre>
           }
void init () {
    for (int u = 1; u <= n; ++u)
        for (int v = 1; v <= n; ++v)
        g[u][v] = edge (u, v, 0);</pre>
};
                Maximum flow
       Sparse graph maximum flow :
```

#### 4.11

```
int isap::solve (flow_edge_list &e, int n, int s, int t) :
              e : edge list.
n : vertex size.
s : source.
t : sink.
template <int MAXN = 1000, int MAXM = 100000>
struct isap {
```

```
int pre[MAXN], d[MAXN], gap[MAXN], cur[MAXN];
         int solve (flow_edge_list <MAXN, MAXN> &e, int n, int s, int t) {
   std::fill (pre, pre + n + 1, 0);
   std::fill (d, d + n + 1, 0);
   std::fill (gap, gap + n + 1, 0);
   for (int i = 0; i < n; i++) cur[i] = e.begin[i];
   gap[0] = n;
   int u = pre[s] = s, u maxflow = 0;</pre>
                    int u = pre[s] = s, v, maxflow = 0;
while (d[s] < n) {
    v = n;</pre>
                             v = n;
for (int i = cur[u]; ~i; i = e.next[i])
   if (e.flow[i] && d[u] == d[e.dest[i]] + 1) {
      v = e.dest[i];
      cur[u] = i;
      the cur [u] = i;
                              if (v < n) {
    pre[v] = u;
    u = v;
    - +)</pre>
                                       maxflow += dflow;
                                                  e.flow[e.inv[cur[p]]] += dflow;
                             } else {
   int mindist = n + 1;
   int mindist = n + 2;
}
                                       for (int i = e.begin[u]; ~i; i = e.next[i])
    if (e.flow[i] && mindist > d[e.dest[i]]) {
        mindist = d[e.dest[i]];
        cur[u] = i.
                                                           cur[u] = i;
                                       }
if (!--gap[d[u]]) return maxflow;
gap[d[u] = mindist + 1]++;
u = pre[u];
                   return maxflow;
         }
};
         Dense graph maximum flow :
    int dinic::solve (flow_edge_list &e, int n, int s, int t) :
                              e : edge list
                             n : vertex size.
s : source.
t : sink.
template <int MAXN = 1000, int MAXM = 100000>
struct dinic {
         int n, s, t;
          int d[MAXN], w[MAXN], q[MAXN];
         int d[MAXN], w[MAXN], q[MAXN];
int bfs (flow_edge_list <MAXN, MAXM> &e) {
    for (int i = 0; i < n; i ++) d[i] = -1;
    int l, r;
    q[l = r = 0] = s, d[s] = 0;
    for (; l <= r; l ++)
        for (int k = e.begin[q[l]]; k > -1; k = e.next[k])
        if (d[e.dest[k]] == -1 && e.flow[k] > 0) d[e.dest[k]] = d[q[l]] + 1, q[++r] = e.dest[k];
    return d[t] > -1 ? 1 : 0;
}
         }
int dfs (flow_edge_list <MAXN, MAXM> &e, int u, int ext) {
    if (u == t) return ext;
    int k = w[u], ret = 0;
    for (; k > -1; k = e.next[k], w[u] = k) {
        if (ext == 0) break;
        if (d[e.dest[k]] == d[u] + 1 && e.flow[k] > 0) {
            int flow = dfs (e, e.dest[k], std::min (e.flow[k], ext));
            if (flow > 0) {
                  e.flow[k] -= flow, e.flow[e.inv[k]] += flow;
                  ret += flow, ext -= flow;
        }
}
                   }
if (k == -1) d[u] = -1;
return ret;
          void solve (flow_edge_list <MAXN, MAXM> &e, int n, int s, int t) {
    dinic::n = n; dinic::s = s; dinic::t = t;
    while (bfs (e)) {
        for (int i = 0; i < n; i ++) w[i] = e.begin[i];
        dfs (e, s, INF);</pre>
};
```

### 4.12 Minimum cost flow

```
returns the flow and the cost respectively.
*/
template <int MAXN = 1000, int MAXM = 100000>
struct minimum_cost_flow {
       int n, source, target;
int prev[MAXN];
int dist[MAXN], occur[MAXN];
             laugment (cost_flow_edge_list <MAXN, MAXM> &e) {
  std::vector <int> queue;
  std::fill (dist, dist + n, INF);
  std::fill (occur, occur + n, 0);
  dist[source] = 0;
  occur[source] = true;
  queue.push_back (source);
  for (int head = 0; head < (int)queue.size(); ++head) {
    int x = queue[head];
    for (int i = e.begin[x]; ~i; i = e.next[i]) {
        int y = e.dest[i];
        if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
            dist[y] = dist[x] + e.cost[i];
            prev[y] = i;
            if (!occur[y]) {
                  occur[y] = true;
                 queue.push_back (y);
        }
}
       bool augment (cost_flow_edge_list <MAXN, MAXM> &e) {
                              }
                      occur[x] = false;
               return dist[target] < INF;</pre>
       std::pair <int, int> solve (cost_flow_edge_list <MAXN, MAXM> &e, int n, int s, int t) {
              ::pair <int, int> solve (cost_ilow_eage_ilst value, int)
minimum_cost_flow::n = n;
source = s; target = t;
std::pair <int, int> answer = std::make_pair (0, 0);
while (augment (e)) {
   int number = INF;
   for (int i = target; i != source; i = e.dest[e.inv[prev[i]]]) {
        number = std::min (number, e.flow[prev[i]]);
   }
}
                      answer.first += number;
for (int i = target; i != source; i = e.dest[e.inv[prev[i]]]) {
    e.flow[prev[i]] -= number;
    e.flow[e.inv[prev[i]]] += number;
    answer.second += number * e.cost[prev[i]];
                       }
               return answer;
       }
      edge list
                      n : vertex size.
s : source.
                              sink
                       returns the flow and the cost respectively.
template <int MAXN = 1000, int MAXM = 100000>
struct zkw_flow {
       int n, s, t, totFlow, totCost;
int dis[MAXN], slack[MAXN], visit[MAXN];
       if (delta == INF) return 1;
for (int i = 0; i < n; i++) if (visit[i]) dis[i] += delta;
return 0;</pre>
       int dfs (cost_flow_edge_list <MAXN, MAXM> &e, int x, int flow) {
               if (x == t) {
   totFlow += flow;
   totCost += flow * (dis[s] - dis[t]);
                       return flow;
             slack[y] = std::min (slack[y], dis[y] + e.cost[i] - dis[x]);
               return flow - left;
       std::pair <int, int> solve (cost_flow_edge_list <MAXN, MAXM> &e, int n, int s, int t) {
   zkw_flow::n = n; zkw_flow::s = s; zkw_flow::t = t;
   totFlow = 0; totCost = 0;
   std::fill (dis + 1, dis + t + 1, 0);
}
                       do {
                       std::fill (visit + 1, visit + t + 1, 0);
} while (dfs (e, s, INF));
```

```
} while (!modlable ());
    return std::make_pair (totFlow, totCost);
};
```

### 4.13 Dominator tree

```
void dominator_tree::solve (int s, int n, const edge_list <MAXN, MAXM> &succ) :
    solves for the immediate dominator (idom[]) of each node,
    idom[x] will be x if x does not have a dominator,
    and will be -1 if x is not reachable from s.
template <int MAXN = 100000, int MAXM = 100000>
struct dominator_tree {
        int dfn[MAXN], sdom[MAXN], idom[MAXN], id[MAXN], fa[MAXN], smin[MAXN], stamp;
        void predfs (int x, const edge_list <MAXN, MAXM> &succ) {
  id[dfn[x] = stamp++] = x;
  for (int i = succ.begin[x]; ~i; i = succ.next[i]) {
                         int y = succ.dest[i];
if (dfn[y] < 0) {
    f[y] = x;
    predfs (y, succ);</pre>
                         }
                 }
        }
        int getfa (int x) {
    if (fa[x] == x) return x;
    int ret = getfa (fa[x]);
    if (dfn[sdom[smin[fa[x]]]] < dfn[sdom[smin[x]]]) {
        smin[x] = smin[fa[x]];
    }
}</pre>
                 return fa[x] = ret;
       }
if (dfn[sdom[x]] > dfn[p]) {
    sdom[x] = p;
                                           1
                                  }
tmp.add_edge (sdom[x], x);
                         while (~tmp.begin[x]) {
    int y = tmp.dest[tmp.begin[x]];
    tmp.begin[x] = tmp.next[tmp.begin[x]];
    getfa (y);
    if (x != sdom[smin[y]]) {
        idom[v] = smin[y];
    }
}
                                  } else {
  idom[y] = x;
                         for (int i = succ.begin[x]; ~i; i = succ.next[i]) {
    if (f[succ.dest[i]] == x) {
        fa[succ.dest[i]] = x;
    }
}
                         }
                } idom[s] = s;
for (int i = 1; i < stamp; ++i) {
   int x = id[i];
   if (idom[x] != sdom[x]) {
      idom[x] = idom[idom[x]];
   }
}</pre>
                }
        }
};
```

### 4.14 Stoer Wagner

```
/* Stoer Wagner algorithm :
    Finds the minimum cut of an undirected graph.
    Usage :
        input : n, edge[][] : the size and weights of the graph.
        int stoer_wagner::solve () : returns the minimum cut.
*/
template <int MAXN = 500>
struct stoer_wagner {
    int n, edge[MAXN][MAXN];
```

## Trial of string

### 5.1 KMP

```
/* KMP algorithm :
    void kmp::build (const std::string & str) :
        initializes and builds the failure array. Complexity 0 (n).
    int kmp::find (const std::string & str) :
            finds the first occurence of match in str. Complexity 0 (n).
    Note : match is cylic when L % (L - 1 - fail[L - 1]) == 0 &&
            L / (L - 1 - fail[L - 1]) > 1, where L = match.size ().

*/

template <int MAXN = 1000000>
struct kmp {
    std::string match;
    int fail[MAXN];
    void build (const std::string & str) {
        match = str; fail[0] = -1;
        for (int i = 1; i < (int) str.size (); ++i) {
            int j = fail[i - 1];
            while (~j && str[i] != str[j + 1]) j = fail[j];
            fail[i] = str[i] == str[j + 1] ? j + 1 : -1;
        }
    int find (const std::string & str) {
        for (int i = 0, j = -1; i < (int) str.size (); j += str[i] == match[j + 1], ++i) {
            if (j == match.size () - 1) return i - match.size ();
            while (~j && str[i] != match[j + 1]) j = fail[j];
        }
        return str.size ();
    }
}</pre>
```

### 5.2 Suffix array

```
/* Suffix Array:
    Suffix array implementation.
    void suffix_array::solve (int *a, int n):
        sa[i]: the beginning position of the ith smallest suffix.
        rk[i]: the rank of the suffix beginning at position i.
        height[i]: the longest common prefix of sa[i] and sa[i - 1].

*/
template <int MAXN = 1000000, int MAXC = 26>
struct suffix_array {
    int rk[MAXN], height[MAXN], sa[MAXN];
    int cmp (int *x, int a, int b, int d) {
        return x[a] == x[b] && x[a + d] == x[b + d];
    }

void doubling (int *a, int n) {
        static int sRank[MAXN], tmpA[MAXN], tmpB[MAXN];
        int m = MAXC;
        int *x = tmpA, *y = tmpB;
        for (int i = 0; i < m; ++i) sRank[i] = 0;
        for (int i = 0; i < m; ++i) sRank[i] += sRank[x = 1];
        for (int i = 1; i < m; ++i) sRank[i] += sRank[x = 1];
        for (int d = 1, p = 0; p < n; m = p, d <<= 1) {
            p = 0; for (int i = n - d; i < n; ++i) y[p++] = i;
            for (int i = 0; i < m; ++i) if (sa[i] >= d) y[p++] = sa[i] - d;
            for (int i = 0; i < m; ++i) sRank[i] = 0;
            for (int i = 0; i < m; ++i) sRank[x = 0; i < int i = 0; i < n; ++i) sRank[x = 0; i < int i = 0;
```

```
if (rk[i]) {
    if (cur) cur--;
    for (; a[i + cur] == a[sa[rk[i] - 1] + cur]; ++cur);
    height[rk[i]] = cur;
}
};
```

### 5.3 Suffix automaton

```
initializes the automaton with an empty string.
void suffix_automaton::extend (int token):
    extends the string with token. Complexity 0 (1).
                 head : the first state.
tail : the last state.
Terminating states can be reached via visiting the ancestors of tail.
                 state::len : the longest length of the string in the state.
state::parent : the parent link.
state::dest : the automaton link.
template <int MAXN = 1000000, int MAXC = 26>
struct suffix_automaton {
        struct state {
                int len;
state *parent, *dest[MAXC];
state (int len = 0) : len (len), parent (NULL) {
    memset (dest, 0, sizeof (dest));

/ node_pool[MAXN * 2], *tot_node, *null = new state();
state *head, *tail;

void extend (int token) {
    state *p = tail;
    state *np = tail -> dest[token] ? null : new (tot_node++) state (tail -> len + 1);
    while (p && !p -> dest[token])
        p -> dest[token] = np, p = p -> parent;
    if (!p) np -> parent = head;
    else {
        state *g = p -> dest[token];
    }
}

        } node_pool[MAXN * 2], *tot_node, *null = new state();
                         state *q = p -> dest[token];
if (p -> len + 1 == q -> len) {
    np -> parent = q;
} else {
                                  se {
    state *nq = new (tot_node++) state (*q);
    nq -> len = p -> len + 1;
    np -> parent = q -> parent = nq;
    while (p && p -> dest[token] == q) {
        p -> dest[token] = nq, p = p -> parent;
}
                         }
                 tail = np == null ? np -> parent : np;
        void init () {
   tot_node = node_pool;
   head = tail = new (tot_node++) state();
        suffix_automaton () {
   init ();
};
```

### 5.4 Palindromic tree

```
last -> fail = now -> child[token];
}
return true;
}
void init() {
    text[size = 0] = -1;
    tot_node = node pool;
    last = even = new (tot_node++) node (0); odd = new (tot_node++) node (-1);
    even -> fail = odd;
}
palindromic_tree () {
    init ();
}
};
```

## Reference

### 6.1 Vimrc

```
set ru nu ts=4 sts=4 sw=4 si sm hls is ar bs=2 mouse=a
syntax on
nm <F3> :vsplit %<.in <CR>
nm <F4> :!gedit % <CR>
au Bufenter *.cpp set cin
au Bufenter *.cpp nm <F5> :!time ./%< <CR>|nm <F7> :!gdb ./%< <CR>|nm <F8> :!time ./%< < %<.in <CR>|
\nm <F9> :!g++ % -o %< -g -std=c++11 -O2 && size %< <CR>
au Bufenter *.java nm <F5> :!time java %< <CR>|nm <F8> :!time java %< < %<.in <CR>|nm <F9> :!g++ %
```

### 6.2 Java reference

```
/* Java reference
                  References on Java IO, structures, etc.
*/
import java.io.*;
import java.lang.*;
import java.math.*;
import java.util.*;
/* Regular usage:
Slower IO
                  Slower IO:
    Scanner in = new Scanner (System.in);
    Scanner in = new Scanner (new BufferedInputStream (System.in));
                            in.nextInt () / in.nextBigInteger () / in.nextBigDecimal () / in.nextDouble ()
in.nextLine () / in.hasNext ()
Output :
                                     System.out.print (...);
System.out.println (...
                  System.out.printf (...);
Faster IO:
Shown below.
                  BigInteger :
                            BigInteger.valueOf (int) : convert to BigInteger.
                            consists of a BigInteger value and a scale.
The scale is the number of digits to the right of the decimal point.
                            divide (BigDecimal) : exact divide.
                            divide (BigDecimal, int scale, RoundingMode roundingMode):
    divide with roundingMode, which may be:
        CEILING / DOWN / FLOOR / HALF_DOWN / HALF_EVEN / HALF_UP / UNNECESSARY / UP.
BigDecimal setScale (int newScale, RoundingMode roundingMode):
    returns a BigDecimal with newScale.
doubleValue () / toPlainString (): converts to other types.
avs:
                  addfirst <E>:
  addLast (E) / getFirst / getLast / removeFirst / removeLast () :
        deque implementation.
  clear () / add (int, E) / remove (int) : clear, add & remove.
  size () / contains / removeFirstOccurrence / removeLastOccurrence (E) :
                  size () / contains / removeFirstOccurrence / removeLastOccurrence (E) :
    deque methods.
ListIterator <E> listIterator (int index) : returns an iterator :
    E next / previous () : accesses and iterates.
    hasNext / hasPrevious () : checks availablity.
    nextIndex / previousIndex () : returns the index of a subsequent call.
    add / set (E) / remove () : changes element.

PriorityQueue <E> (int initcap, Comparator <? super E> comparator) :
    add (E) / clear () / iterator () / peek () / poll () / size () :
        priority queue implementations.

TreeMap <K, V> (Comparator <? super K> comparator) :
    Map.Entry <K, V> ceilingEntry / floorEntry / higherEntry / lowerEntry (K):
        getKey / getValue () / setValue (V) : entries.
```

```
clear () / put (K, V) / get (K) / remove (K) : basic operation.
    size () : size.
StringBuilder : .
                      Mutable string
                      Mutable string.

StringBuilder (string) : generates a builder.

append (int, string, ...) / insert (int offset, ...) : adds objects.

charAt (int) / setCharAt (int, char) : accesses a char.

delete (int, int) : removes a substring.

reverse () : reverses itself.

length () : returns the length.
                       toString () : converts to string.
              String:
Immutable string.
String.format (String, ...): formats a string. i.e. sprintf.
toLowerCase / toUpperCase (): changes the case of letters.
/* Examples on Comparator :
public class Main {
       public static class Point {
               public int x;
public int y;
public Point () {

\begin{array}{l}
\mathbf{x} = 0; \\
\mathbf{y} = 0;
\end{array}

               public Point (int xx, int yy) {
    x = xx;
    y = yy;
        };
      public static class Cmp implements Comparator <Point> {
   public int compare (Point a, Point b) {
     if (a.x < b.x) return -1;
     if (a.x = b.x) {
        if (a.y < b.y) return -1;
        if (a.y == b.y) return 0;
}</pre>
                       return 1;
               }
       public static void main (String [] args) {
               Cmp c = new Cmp ();
TreeMap <Point, Point> t = new TreeMap <Point, Point> (c);
               return:
1;
       Another way to implement is to use Comparable. However, equalTo and hashCode must be rewritten.
        Otherwise, containers may fail.
        Example :
       public static class Point implements Comparable <Point> {
              public int x;
public int y;
               public Point () {
    x = 0;
    y = 0;
               public Point (int xx, int yy) {
    x = xx;
    y = yy;
              public int compareTo (Point p) {
    if (x < p.x) return -1;
    if (x == p.x) {
        if (y < p.y) return -1;
        if (y == p.y) return 0;
}</pre>
                       return 1;
               public boolean equalTo (Point p) {
    return (x == p.x && y == p.y);
               public int hashCode () {
    return x + y;
       1;
//Faster IO :
public class Main {
       static class InputReader {
              public BufferedReader reader;
public StringTokenizer tokenizer;
public InputReader (InputStream stream) {
    reader = new BufferedReader (new InputStreamReader (stream), 32768);
    tokenizer = null;
               public String next() {
    while (tokenizer == null || !tokenizer.hasMoreTokens()) {
                                     String line = reader.readLine();
tokenizer = new StringTokenizer (line);
                              } catch (IOException e) {
   throw new RuntimeException (e);
                       return tokenizer.nextToken();
               public BigInteger nextBigInteger() {
                      return new BigInteger (next (), 10);
                                                                                                 // customize the radix here.
               public int nextInt() {
    return Integer.parseInt (next());
```

```
}
public double nextDouble() {
    return Double.parseDouble (next());
}

public static void main (String[] args) {
    InputReader in = new InputReader (System.in);
    // Put your code here.
}
```

## 6.3 Operator precedence

Precedence Operator		Description	Associativity
1 ::		Scope resolution	
	a++ a	Suffix/postfix increment and decrement	
2	type() type{}	Functional cast	Left-to-right
	a()	Function call	
	a[]	Subscript	
	>	Member access	
	++aa	Prefix increment and decrement	
	+a -a	Unary plus and minus	
	! ~	Logical NOT and bitwise NOT	
	(type)	C-style cast	
3	*a	Indirection (dereference)	Right-to-left
	&a	Address-of	- 0
	sizeof	Size-of	
	new new[]	Dynamic memory allocation	
	delete	, , , , , , , , , , , , , , , , , , ,	
	delete[]	Dynamic memory deallocation	
4	.* ->*	Pointer-to-member	
	a*b a/b		
5	a^D a/D a%b	Multiplication, division, and remainder	
6	a+b a-b	Addition and subtraction	
7	<< >>	Bitwise left shift and right shift	
	< <=	For relational operators $<$ and $\le$ respectively	T () 1 .
8	> >=	For relational operators $>$ and $\ge$ respectively	Left-to-right
9	== !=	For relational operators = and $\neq$ respectively	
10	a&b	Bitwise AND	
11	^	Bitwise XOR (exclusive or)	
12	1	Bitwise OR (inclusive or)	
13	& &	Logical AND	
14	11	Logical OR	
	a?b:c	Ternary conditional	
	throw	throw operator	
1.2	=	Direct assignment	D. 1 1.6
15	+= -= *=	<u> </u>	Right-to-left
	/= %=	Compound assignment by arithmetic operation	
	<<= >>=	Compound assignment by bitwise shift	
	&= ^=  =	Compound assignment by bitwise AND, XOR, and	
		OR	
16	,	Comma	Left-to-right

### 6.4 Hacks

### 6.4.1 Ultra fast functions

```
// Ultra fast functions :
    __inline void make_min (int &a, int &b) {
    asm (
        "cmpl_%2,_%0\n\t"
        "movl_%2,_%0\n\t"
        "DONE:"
        : "=r" (a)
        : "0" (a), "r" (b)
    );
}
__inline void make_max (int &a, int &b) {
    asm (
```

```
"cmpl_%2,_%0\n\t"
    "jge_DONE\n\t"
    "movl_%2,_%0\n\t"
    "DONE:"
    : "=r" (a)
    : "0" (a), "r" (b)
);
}
_inline int cmp (int a) {
    return (a >> 31) + (-a >> 31 & 1);
}
_inline int abs (int x) {
    int y = x >> 31;
    return (x + y) ^ y;
}
_inline int mul_mod (int a, int b) {
    int ret;
    asm (
        "mull_%%ebx\n\t"
        : "=d" (ret)
        : "a" (a), "b" (b), "c" (MO)
    ;
return ret;
}
```

### 6.4.2 Formating long long in scanf & printf

```
#ifdef WIN32
    #define LL "%I64d"
#else
    #define LL "%lld"
#endif
```

### 6.4.3 Optimizing

```
#pragma GCC optimize ("03")
#pragma GCC optimize ("whole-program")
```

### 6.4.4 Larger stack

### 6.4.4.1 C++

```
#pragma comment(linker, "/STACK:36777216")
```

#### 6.4.4.2 G++

### 6.5 Math reference

### 6.5.1 Catalan number

For

$$\begin{split} f(0) &= 1 \\ f(1) &= 1 \\ f(n) &= f(n-1)f(0) + f(n-2)f(1) + \ldots + f(1)f(n-2) + f(0)f(n-1) \end{split}$$

We have  $f(n) = \frac{(2n)!}{n!(n+1)!}$ .

#### 6.5.2 Dynamic programming optimization

#### 6.5.2.1 Convex hull optimization

Generally, in dynamic programming with recurrence

$$f(i) = \min_{k < i} \{a[i]b(j) + c(j)\}$$

all decisions k can be treated as a set of segments on a convex hull. By applying Graham's scanning, it is possible to maintain such hull in a monotone queue or a std::tuple <slope, intercept, x\_min>. Hence, k(i) can be obtained by performing a binary search in the hull.

#### 6.5.2.2 Divide & conquer optimization

For recurrence

$$f(i) = \min_{k < i} \{b(k) + c[k][i]\}$$

 $k(i) \le k(i+1)$  holds true if c[a][c] + c[b][d] < c[a][d] + c[b][c]. Thus, k(i) can be maintained in a monotone queue.

#### 6.5.2.3 Knuth optimization

For recurrence

$$f(i,j) = \min_{i < k < j} \{ f(i,k) + f(k,j) \} + c[i][j]$$

 $k(i, j - 1) \le k(i, j) \le k(i + 1, j)$  holds true if c[a][c] + c[b][d] < c[a][d] + c[b][c].

### 6.5.3 Integration table

### **6.5.3.1** $ax^2 + bx + c \ (a > 0)$

1. 
$$\int \frac{\mathrm{d}x}{ax^2 + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}} + C & b^2 < 4ac \\ \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right| + C & b^2 > 4ac \end{cases}$$

2. 
$$\int \frac{x}{ax^2 + bx + c} dx = \frac{1}{2a} \ln |ax^2 + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^2 + bx + c}$$

### **6.5.3.2** $\sqrt{\pm ax^2 + bx + c}$ (a > 0)

1. 
$$\int \frac{dx}{ax^2+bx+c} = \frac{1}{\sqrt{a}} \ln |2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$$

2. 
$$\int \sqrt{ax^2 + bx + c} \, dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln \left| 2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c} \right| + C$$

3. 
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln \left| 2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c} \right| + C$$

4. 
$$\int \frac{\mathrm{d}x}{\sqrt{-ax^2+bx+c}} = -\frac{1}{\sqrt{a}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5. 
$$\int \sqrt{-ax^2 + bx + c} \, dx = \frac{2ax - b}{4a} \sqrt{-ax^2 + bx + c} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

6. 
$$\int \frac{x}{\sqrt{-ax^2 + bx + c}} dx = -\frac{1}{a}\sqrt{-ax^2 + bx + c} + \frac{b}{2\sqrt{a^3}}\arcsin\frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

### 6.5.3.3 Trigonometric

1. 
$$\int \tan x \, dx = -\ln|\cos x| + C$$

2. 
$$\int \cot x \, dx = \ln|\sin x| + C$$

3. 
$$\int \sec x \, dx = \ln |\tan(\frac{\pi}{4} + \frac{x}{2})| + C = \ln |\sec x + \tan x| + C$$

4. 
$$\int \csc x \, dx = \ln\left|\tan\frac{x}{2}\right| + C = \ln\left|\csc x - \cot x\right| + C$$

5. 
$$\int \sec^2 x \, \mathrm{d}x = \tan x + C$$

6. 
$$\int \csc^2 x \, \mathrm{d}x = -\cot x + C$$

7. 
$$\int \sec x \tan x \, dx = \sec x + C$$

8. 
$$\int \csc x \cot x \, dx = -\csc x + C$$

9. 
$$\int \sin^2 x \, dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$$

10. 
$$\int \cos^2 x \, dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$$

11. 
$$\int \sin^n x \, dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x \, dx$$

12. 
$$\int \cos^n x \, dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x \, dx$$

13. 
$$\int \frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$$

14. 
$$\int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$

15. 
$$\int \cos^m x \sin^n x \, dx = \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x \, dx = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, dx$$

#### **6.5.3.4** Inverse trigonometric (a > 0)

- 1.  $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2} + C$
- 2.  $\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} \sqrt{a^2 x^2} + C$
- 3.  $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} \frac{a}{2} \ln(a^2 + x^2) + C$

### 6.5.3.5 Exponential

- 1.  $\int a^x dx = \frac{1}{\ln a} a^x + C$
- 2.  $\int e^{ax} dx = \frac{1}{2} a^{ax} + C$

#### 6.5.3.6 Logistic

- 1.  $\int \ln x \, \mathrm{d}x = x \ln x x + C$
- 2.  $\frac{\mathrm{d}x}{x \ln x} = \ln |\ln x| + C$

### Prefix sum of multiplicative functions

Define the Dirichlet convolution f \* q(n) as:

$$f * g(n) = \sum_{d=1}^{n} [d|n]f(n)g(\frac{n}{d})$$

Assume we are going to calculate some function  $S(n) = \sum_{i=1}^{n} f(i)$ , where f(n) is a multiplicative function. Say we find some g(n) that is simple to calculate, and  $\sum_{i=1}^{n} f * g(i)$  can be figured out in O(1) complexity. Then we have

$$\begin{split} \sum_{i=1}^n f * g(i) &= \sum_{i=1}^n \sum_d [d|i] g(\frac{i}{d}) f(d) \\ &= \sum_{\frac{i}{d}=1}^n \sum_{d=1}^{\left\lfloor \frac{n}{d} \right\rfloor} g(\frac{i}{d}) f(d) \\ &= \sum_{i=1}^n \sum_{d=1}^{\left\lfloor \frac{n}{i} \right\rfloor} g(i) f(d) \\ &= g(1) S(n) + \sum_{i=2}^n g(i) S(\left\lfloor \frac{n}{i} \right\rfloor) \\ S(n) &= \frac{\sum_{i=1}^n f * g(i) - \sum_{i=2}^n g(i) S(\left\lfloor \frac{n}{i} \right\rfloor)}{g(1)} \end{split}$$

It can be proven that  $\left|\frac{n}{i}\right|$  has at most  $O(\sqrt{n})$  possible values. Therefore, the calculation of S(n) can be reduced to  $O(\sqrt{n})$ calculations of  $S(\lfloor \frac{n}{i} \rfloor)$ . By applying the master theorem, it can be shown that the complexity of such method is  $O(n^{\frac{3}{4}})$ .

Moreover, since f(n) is multiplicative, we can process the first  $n^{\frac{2}{3}}$  elements via linear sieve, and for the rest of the elements, we apply the method shown above. The complexity can thus be enhaced to  $O(n^{\frac{2}{3}})$ .

For the prefix sum of Euler's function  $S(n) = \sum_{i=1}^{n} \varphi(i)$ , notice that  $\sum_{d|n} \varphi(d) = n$ . Hence  $\varphi * I(n) = id(n).(I(n)) = id(n)$ 1, id(n) = n) Now let g(n) = I(n), and we have  $S(n) = \sum_{i=1}^{n} i - \sum_{i=2}^{n} S(\lfloor \frac{n}{i} \rfloor)$ . For the prefix sum of Mobius function  $S(n) = \sum_{i=1}^{n} \mu(i)$ , notice that  $\mu * I(n) = [n = 1]$ . Hence  $S(n) = 1 - \sum_{i=2}^{n} S(\lfloor \frac{n}{i} \rfloor)$ .

#### 6.5.5Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove the leaf with the smallest label and set the *i*th element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence a[i], find the first (lowest-numbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u, v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence is S.

### 6.5.6 Spanning tree counting

**Kirchhoff's Theorem**: the number of spanning trees in a graph G is equal to *any* cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)-matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weight together.

### 6.6 Regular expression

### 6.6.1 Special pattern characters

Characters	Description	Matches	
•	Not newline	Any character except line terminators (LF, CR, LS, PS).	
\t	Tab (HT)	A horizontal tab character (same as \u0009).	
\n	Newline (LF)	A newline (line feed) character (same as \u000A).	
\v	Vertical tab (VT)	A vertical tab character (same as \u000B).	
\f	Form feed (FF)	A form feed character (same as \u0000C).	
Carriaga raturn		A carriage return character (same as \u0000D).	
\cletter	Control code	A control code character whose code unit value is the same as the remainder of dividing the code unit value of letter by 32. For example: \ca is the same as \u00001, \cb the same as \u00002, and so on	
\xhh	ASCII character	A character whose code unit value has an hex value equivalent to the two hex digits hh. For example: \x4c is the same as L, or \x23 the same as #.	
\uhhhh	Unicode character	A character whose code unit value has an hex value equivalent to the four hex digits hhhh.	
\0	Null	A null character (same as \u0000).	
\int   Backreterence		The result of the submatch whose opening parenthesis is the int-th (int shall begin by a digit other than 0). See groups below for more info.	
\d	Digit	A decimal digit character (same as [[:digit:]]).	
\D	Any character that is not a decimal digit charac		
\s	Whitespace	A whitespace character (same as [[:space:]]).	
\S	Not whitespace	Any character that is not a whitespace character (same as [^[:space:]]).	
\w	Word	An alphanumeric or underscore character (same as [_[:alnum:]]).	
\W	Not word	Any character that is not an alphanumeric or underscore character (same a [^_[:alnum:]]).	
\character	Character	The character character as it is, without interpreting its special meaning within a regex expression. Any character can be escaped except those which form any of the special character sequences above. Needed for:  ^ \$ \ . * + ? ( ) [ ] { }  .	
[class]	Character class	The target character is part of the class (see character classes below).	
[^class] Negated character class The target character is not part of the class (see character		The target character is not part of the class (see character classes below).	

### 6.6.2 Quantifiers

Characters	Times	Effects	
*	0 or more	The preceding atom is matched 0 or more times.	
+	1 or more	The preceding atom is matched 1 or more times.  The preceding atom is optional (matched either 0 times or once).	
?	0 or 1		
{int}	int	The preceding atom is matched exactly int times.	
{int,}	int or more	The preceding atom is matched int or more times.	
(min max)	Between min and	The preceding atom is matched at least min times, but not more than max.	
{min,max}	max		

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

### **6.6.3** Groups

Characters	Characters Description Effects		
(subpattern)	Group	Creates a backreference.	
(?:subpattern) Passive group		Does not create a backreference.	

### 6.6.4 Assertions

Characters	Description	Condition for match	
^	Beginning of line	Either it is the beginning of the target sequence, or follows a line	
	Deginning of fine	terminator.	
\$	End of line	Either it is the end of the target sequence, or precedes a line	
7		terminator.	
		The previous character is a word character and the next is a	
\b	Word boundary	non-word character (or vice-versa). Note: The beginning and the	
d/		end of the target sequence are considered here as non-word	
		characters.	
	Not a word	The previous and next characters are both word characters or	
\ D		both are non-word characters. Note: The beginning and the end	
\B	boundary	of the target sequence are considered here as non-word	
		characters.	
(2-gubpa++orn)	Positive lookahead	The characters following the assertion must match subpattern,	
(?=subpattern)		but no characters are consumed.	
(2 Laubnattorn)	Negative lookahead	The characters following the assertion must not match	
(?!subpattern)	riegative lookallead	subpattern, but no characters are consumed.	

### 6.6.5 Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (+): The regular expression will match if any of the alternatives match, and as soon as one does.

### 6.6.6 Character classes

Class	Description	Equivalent (with regex_traits, default locale)
[:alnum:]	Alpha-numerical character	isalnum
[:alpha:]	Alphabetic character	isalpha
[:blank:]	Blank character	isblank
[:cntrl:]	Control character	iscntrl
[:digit:]	Decimal digit character	isdigit
[:graph:]	Character with graphical representation	isgraph
[:lower:]	Lowercase letter	islower
[:print:]	Printable character	isprint
[:punct:]	Punctuation mark character	ispunct
[:space:]	Whitespace character	isspace
[:upper:]	Uppercase letter	isupper
[:xdigit:]	Hexadecimal digit character	isxdigit
[:d:]	Decimal digit character	isdigit
[:w:]	Word character	isalnum
[:s:]	Whitespace character	isspace

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example: [[:alpha:]] is a character class that matches any alphabetic character.

[abc[:digit:]] is a character class that matches a, b, c, or a digit.

[^[:space:]] is a character class that matches any character except a whitespace.