Contents

1	Trial of number theory		1
	1.1 Constants and basic functions		
	1.2 Discrete Fourier transform		 1
	1.3 Number-theoretic transform		 2
	1.4 Chinese remainder theorem		 2
	1.5 Miller Rabin primality test		 3
	1.6 Pollard's Rho algorithm		 3
	1.7 Adaptive Simpson's method		 4
2	Trial of geometry		5
	2.1 Constants and basic functions		 5
	2.2 Point class		 5
	2.3 Line class		 6
	2.4 Interactions between points and lines		 6
	2.5 Centers of a triangle		
	2.6 Fermat point		
	2.7 Circle class		
	2.8 Interactions of circles		
	2.9 Convex hull		
	2.10 Minimum circle		
	2.11 Half plane intersection		
	2.12 Intersection of a polygon and a circle		
	2.12 Union of circles		
	2.13 Union of circles	• •	 9
3	Trial of graph		11
Ü	3.1 Constants and edge lists		
	3.2 SPFA improved		
	3.3 Dijkstra's shortest path algorithm		
	3.4 Tarjan		
	3.5 Vertex biconnected component		
	3.6 Edge biconnected component		
	-		
	3.8 Kuhn-Munkres		
	3.9 Stochastic weighted maximum matching		
	3.10 Weighted blossom (vfleaking ver.)		
	3.11 Maximum flow		
	3.12 Minimum cost flow		
	3.13 Dominator tree		 20
4	Th.:-1 - f - t.:		0.1
4	Trial of string		21
	4.1 KMP		21
	4.2 Suffix automaton		
	4.3 Palindromic tree		 22
5	Reference		23
IJ			
			23 23
	5.3 Operator precedence		25
	5.4 Hacks		26
	5.4.1 Formating long long in scanf & printf		
	5.4.2 Optimizing		
	5.4.3 Larger stack		26

5.5	Math	reference
	5.5.1	Catalan number
	5.5.2	Dynamic programming optimization
	5.5.3	Integration table
	5.5.4	Prefix sum of multiplicative functions
	5.5.5	Prufer sequence
	5.5.6	Spanning tree counting
5.6	Regula	ar expression
		Special pattern characters
		Quantifiers
	5.6.3	Groups
	5.6.4	Assertions
	5.6.5	Alternative
	5.6.6	Character classes

The Grimoire of Programming

Suzune Nisiyama

Chapter 1

Trial of number theory

1.1 Constants and basic functions

```
const double PI = acos (-1);
long long abs (const long long &x) { return x > 0 ? x : -x; }
long long inverse (const long long &x, const long long &mod) {
   if (x == 1) return 1;
   return (mod - mod / x) * inverse (mod % x, mod) % mod;
int fpm (int x, int n, int mod) {
   register int ans = 1, mul = x;
      register int ans - 1, ..... -,
while (n) {
   if (n & 1) ans = int (111 * ans * mul * mod);
   mul = int (111 * mul * mul * mod);
   n >>= 1;
      }
return ans;
ilong long gcd (const long long &a, const long long &b) {
   if (!b) return a;
   long long x = a, y = b;
   while (x > y ? (x = x % y) : (y = y % x));
   return x + y;
long long mul_mod (const long long &a, const long long &b, const long long &mod) {
  long long ans = 0, add = a, k = b;
  while (k) {
    if (k & 1) ans = (ans + add) % mod;
    add = (add + add) % mod;
    k >>= 1;
}
      }
return ans;
long long llfpm (const long long &x, const long long &n, const long long &mod) {
  long long ans = 1, mul = x, k = n;
  while (k) {
    if (k & 1) ans = mul_mod (ans, mul, mod);
    mul = mul_mod (mul, mul, mod);
                  = mul_mod (mul, mul, mod);
      return ans;
```

1.2 Discrete Fourier transform

```
/* Discrete Fourier transform :
    int dft::init (int n) :
        initializes the transformation with dimension n.
        Returns the recommended size.
    void dft::solve (complex *a, int n, int f) :
        transforms array a with dimension n to its image representation.
        Transforms back when f = 1. (n should be 2^k)

*/
template <int MAXN = 1000000>
struct dft {
```

```
typedef std::complex <double> complex;
complex e[2][MAXN];
int init (int n) {
    int len = 1;
    for (; len <= 2 * n; len <<= 1);
    for (int i = 0; i < len; i++) {
        e[0][i] = complex (cos (2 * PI * i / len), sin (2 * PI * i / len));
        e[1][i] = complex (cos (2 * PI * i / len), -sin (2 * PI * i / len));
    }
    return len;
}

void solve (complex *a, int n, int f) {
    for (int i = 0, j = 0; i < n; i++) {
        if (i > j) std::swap (a[i], a[j]);
        for (int t = n >> 1; (j^= t) < t; t >>= 1);
}

for (int i = 2; i <= n; i <<= 1)
    for (int j = 0; j < n; j += i)
        for (int k = 0; k < (i >> 1); k++) {
            complex A = a[j + k];
            complex B = e[f][n / i * k] * a[j + k + (i >> 1)];
        a[j + k] = A + B;
        a[j + k + (i >> 1)] = A - B;

if (f == 1) {
    for (int i = 0; i < n; i++)
        a[i] = complex (a[i].real () / n, a[i].imag ());
}
};</pre>
```

1.3 Number-theoretic transform

```
/* Number-theoretic transform :
    void ntt::solve (int *a, int n, int f, int mod, int prt) :
        transforms a[n] to its image representation.
        Converts back if f = 1. (n should be 2^k)
        Requries specific mod and corresponding prt to work. (given in MOD and PRT)
    int ntt::crt (int *a, int mod) :
        fixes the results a from module 3 primes to a certain module mod.
template <int MAXN = 1000000>
struct ntt {
            void solve (int *a, int n, int f, int mod, int prt) {
    for (register int i = 0, j = 0; i < n; i++) {
        if (i > j) std::swap (a[i], a[j]);
        for (register int t = n >> 1; (j = t) < t; t >>= 1);
                         for (register int i = 2; i <= n; i <<= 1) {
    static int exp[MAXN];
    exp[0] = 1;
    exp[1] = fpm (prt, (mod - 1) / i, mod);
    if (f == 1) exp[1] = fpm (exp[1], mod - 2, mod);
    for (register int k = 2; k < (i >> 1); k++) {
        exp[k] = int (111 * exp[k - 1] * exp[1] % mod);
}
                                      for (register int j = 0; j < n; j += i) {
    for (register int k = 0; k < (i >> 1); k++) {
        register int &pA = a[j + k], &pB = a[j + k + (i >> 1)];
        register int A = pA, B = int (lll * pB * exp[k] % mod);
        pA = (A + B) % mod;
        pB = (A - B + mod) % mod;
}
                                     }
                         }
if (f == 1) {
    register int rev = fpm (n, mod - 2, mod);
    for (register int i = 0; i < n; i++) {
        a[i] = int (111 * a[i] * rev % mod);
}</pre>
              int MOD[3] = {1045430273, 1051721729, 1053818881}, PRT[3] = {3, 6, 7};
            int sum = 1, ret = x[0] % mod;
for (int i = 1; i < 3; i ++) {
    sum = int (1LL * sum * MOD[i - 1] % mod);
    ret += int (1LL * x[i] * sum % mod);
    if (ret >= mod) ret -= mod;
}
                         return ret;
             }
1;
```

1.4 Chinese remainder theorem

1.5 Miller Rabin primality test

1.6 Pollard's Rho algorithm

1.7 Adaptive Simpson's method

Chapter 2

Trial of geometry

2.1 Constants and basic functions

```
/* Constants & basic functions :
    EPS : fixes the possible error of data.
        i.e. x == y iff |x - y| < EPS.
    PI : the value of PI.
    int sgn (const double &x) : returns the sign of x.
    int cmp (const double &x, const double &y) : returns the sign of x - y.
    double sqr (const double &x) : returns x * x.

*/
const double EPS = 1E-8;
const double PI = acos (-1);
int sgn (const double &x) { return x < -EPS ? -1 : x > EPS; }
int cmp (const double &x, const double &y) { return sgn (x - y); }
double sqr (const double &x) { return x * x; }
```

2.2 Point class

```
/* struct point : defines a point and its various utility.
                    point : defines a point and its various utility.
point (const double &x, const double &y) gives a point at (x, y).
It also represents a vector on a 2D plane.
point unit () const : returns the unit vector of (x, y).
point rot90 () const :
    returns a point rotated 90 degrees counter-clockwise with respect to the origin.
point _rot () const : same as above except clockwise.
point rotate (const double &t) const : returns a point rotated t radian(s) counter-clockwise.
Operators are mostly vector operations. i.e. vector +, -, *, / and dot/det product.
*/
struct point {
    double x, y;
    point (const double &x = 0, const double &y = 0) : x (x), y (y) {}
    double norm () const { return sqrt (x * x + y * y); }
    double norm2 () const { return x * x + y * y; }
    point unit () const {
        double 1 = norm ();
        return point (x / 1, y / 1);
}
          point rot90 () const {return point (-y, x); }
point _rot90 () const {return point (y, -x); }
point rotate (const double &t) const {
    double c = cos (t), s = sin (t);
    return point (x * c - y * s, x * s + y * c);
}
bool operator == (const point &a, const point &b) {
   return cmp (a.x, b.x) == 0 && cmp (a.y, b.y) == 0;
bool operator != (const point &a, const point &b) {
   return ! (a == b);
bool operator < (const point &a, const point &b) {
   if (cmp (a.x, b.x) == 0) return cmp (a.y, b.y) < 0;
   return cmp (a.x, b.x) < 0;</pre>
point operator - (const point &a) { return point (-a.x, -a.y); }
point operator + (const point &a, const point &b) {
   return point (a.x + b.x, a.y + b.y);
point operator - (const point &a, const point &b) {
   return point (a.x - b.x, a.y - b.y);
point operator * (const point &a, const double &b) {
   return point (a.x * b, a.y * b);
point operator / (const point &a, const double &b) {
   return point (a.x / b, a.y / b);
double dot (const point &a, const point &b) {
   return a.x * b.x + a.y * b.y;
double det (const point &a, const point &b) {
    return a.x * b.y - a.y * b.x;
```

```
}
double dis (const point &a, const point &b) {
  return sqrt (sqr (a.x - b.x) + sqr (a.y - b.y));
}
```

2.3 Line class

2.4 Interactions between points and lines

```
bool point_on_line (const point &a, const line &b) {
   return sgn (det (a - b.s, b.t - b.s)) == 0 && sgn (dot (b.s - a, b.t - a)) <= 0;</pre>
bool two_side (const point &a, const point &b, const line &c) {
   return sgn (det (a - c.s, c.t - c.s)) * sgn (det (b - c.s, c.t - c.s)) < 0;</pre>
bool intersect judgement (const line &a. const line &b) {
        if (point_on_line (b.s, a) || point_on_line (b.t, a)) return true;
if (point_on_line (a.s, b) || point_on_line (a.t, b)) return true;
return two_side (a.s, a.t, b) && two_side (b.s, b.t, a);
point line_intersect (const line &a, const line &b) {
    double s1 = det (a.t - a.s, b.s - a.s);
    double s2 = det (a.t - a.s, b.t - a.s);
    return (b.s * s2 - b.t * s1) / (s2 - s1);
 double point_to_line (const point &a, const line &b) {
   return fabs (det (b.t - b.s, a - b.s)) / dis (b.s, b.t);
double point_to_segment (const point &a, const line &b) {
    if (sgn (dot (b.s - a, b.t - b.s) * dot (b.t - a, b.t - b.s)) <= 0)
        return fabs (det (b.t - b.s, a - b.s)) / dis (b.s, b.t);
    return std::min (dis (a, b.s), dis (a, b.t));</pre>
*/
if (point_on_line (p, line (a, b))) return true;
int x = sgn (det (p - a, b - a)), y = sgn (a.y - p.y), z = sgn (b.y - p.y);
if (x > 0 && y <= 0 && z > 0) counter++;
if (x < 0 && z <= 0 && y > 0) counter--;
        return counter != 0;
double polygon_area (const std::vector <point> &a) {
    double ans = 0.0;
    for (int i = 0; i < (int) a.size (); ++i)
        ans += det (a[i], a[ (i + 1) % a.size ()]) / 2.0;
    return ans;</pre>
point project_to_line (const point &a, const line &b) {
    return b.s + (b.t - b.s) * (dot (a - b.s, b.t - b.s) / (b.t - b.s).norm2 ());
```

2.5 Centers of a triangle

```
point circumcenter (const point &a, const point &b, const point &c) {
   point p = b - a, q = c - a, s (dot (p, p) / 2, dot (q, q) / 2);
   double d = det (p, q);
   return a + point (det (s, point (p.y, q.y)), det (point (p.x, q.x), s)) / d;
}
point orthocenter (const point &a, const point &b, const point &c) {
   return a + b + c - circumcenter (a, b, c) * 2.0;
}
```

2.6 Fermat point

2.7 Circle class

```
/* struct circle defines a circle.
    circle (point c, double r) gives a circle with center c and radius r.
*/
struct circle {
    point c;
    double r;
    circle (point c = point (), double r = 0) : c (c), r (r) {}
};
bool operator == (const circle &a, const circle &b) {
    return a.c == b.c && cmp (a.r, b.r) == 0;
}
bool operator != (const circle &a, const circle &b) {
    return ! (a == b);
}
```

2.8 Interactions of circles

```
std::pair <line, line> tangent (const point &a, const circle &b) {
    circle p = make_circle (a, b.c);
    return circle_intersect (p, b);
}
```

2.9 Convex hull

2.10 Minimum circle

2.11 Half plane intersection

2.12 Intersection of a polygon and a circle

2.13 Union of circles

```
/* Union of circles :
```

```
std::vector <double> union_circle::solve (const std::vector <circle> &c) :
                            returns the union of circle set c. The i-th element is the area covered with at least i circles.
struct union_circle {
        struct cp {
   double x, y, angle;
                  double x, y, angle,
int d;
double r;
cp (const double &x = 0, const double &y = 0, const double &angle = 0,
    int d = 0, const double &r = 0) : x (x), y (y), angle (angle), d (d), r (r) {}
         double dis (const cp &a, const cp &b) {
    return sqrt (sqr (a.x - b.x) + sqr (a.y - b.y));
         double cross (const cp &p0, const cp &p1, const cp &p2) {
    return (p1.x - p0.x) * (p2.y - p0.y) - (p1.y - p0.y) * (p2.x - p0.x);
        int cir_cross (cp p1, double r1, cp p2, double r2, cp &cp1, cp &cp2) {
    double mx = p2.x - p1.x, sx = p2.x + p1.x, mx2 = mx * mx;
    double my = p2.y - p1.y, sy = p2.y + p1.y, my2 = my * my;
    double sq = mx2 + my2, d = - (sq - sqr (r1 - r2)) * (sq - sqr (r1 + r2));
    if (sgn (d) < 0) return 0;
    if (sgn (d) <= 0) d = 0;
    else d = sqrt (d);
    double x = mx * ((r1 + r2) * (r1 - r2) + mx * sx) + sx * my2;
    double y = my * ((r1 + r2) * (r1 - r2) + my * sy) + sy * mx2;
    double dx = mx * d, dy = my * d;
    sq *= 2:</pre>
                  double dx = mx * d, dy = m

sq *= 2;

cp1.x = (x - dy) / sq;

cp1.y = (y + dx) / sq;

cp2.x = (x + dy) / sq;

cp2.y = (y - dx) / sq;

if (sgn (d) > 0) return 2;

else return 1;
         static bool circmp (const cp &u, const cp &v) {
   return sgn (u.r - v.r) < 0;</pre>
         static bool cmp (const cp &u, const cp &v) {
   if (sgn (u.angle - v.angle)) return u.angle < v.angle;
   return u.d > v.d;
        std::vector <double> solve (const std::vector <circle> &c) {
  int n = c.size ();
  std::vector <cp> cir, tp;
  std::vector <double> area;
  cir recirc (n):
                  std::vector <double> area;
cir.resize (n);
tp.resize (2 * n);
area.resize (n + 1);
for (int i = 0; i < n; i++)
    cir[i] = cp (c[i].c.x, c[i].c.y, 0, 1, c[i].r);
cp cp1, cp2;
std::sort (cir begin () cir end () circmn);</pre>
                  cp2.angle = atan2 (cp2.y = cff[1].y, cp2.x cp1.d = 1;
tp[tn++] = cp1;
cp2.d = -1;
tp[tn++] = cp2;
if (sgn (cp1.angle - cp2.angle) > 0) cnt++;
                            fp[tn++] = cp (cir[i].x - cir[i].r, cir[i].y, PI, -cnt);
tp[tn++] = cp (cir[i].x - cir[i].r, cir[i].y, -PI, cnt);
std::sort (tp.begin (), tp.begin () + tn, cmp);
int p, s = cir[i].d + tp[0].d;
for (int j = 1; j < tn; ++j) {
    p = s;
}</pre>
                                     p = 0,
s += tp[j].d;
area[p] += calc (cir[i], tp[j - 1], tp[j]);
                  return area;
        }
1;
```

Chapter 3

Trial of graph

3.1 Constants and edge lists

```
const int INF = 1E9;
template <int MAXN = 100000, int MAXM = 100000>
struct edge_list {
       int size;
int begin[MAXN], dest[MAXM], next[MAXM];
       void clear (int n) {
    size = 0;
    std::fill (begin, begin + n, -1);
       edge_list (int n = MAXN) {
    clear (n);
       void add_edge (int u, int v) {
   dest[size] = v; next[size] = begin[u]; begin[u] = size++;
template <int MAXN = 100000, int MAXM = 100000>
template tint MARN = 1000000, int MARN = 10000000
struct cost_edge_list {
  int size;
  int begin[MAXN], dest[MAXM], next[MAXM], cost[MAXM];
  void clear (int n) {
     size = 0;
     std::fill (begin, begin + n, -1);
}
       cost_edge_list (int n = MAXN) {
    clear (n);
       void add_edge (int u, int v, int c) {
    dest[size] = v; next[size] = begin[u]; cost[size] = c; begin[u] = size++;
};
template <int MAXN = 100000, int MAXM = 100000>
struct flow_edge_list {
       int size;
int begin[MAXN], dest[MAXM], next[MAXM], flow[MAXM], inv[MAXM];
void clear (int n) {
    size = 0;
    std::fill (begin, begin + n, -1);
       flow_edge_list (int n = MAXN) {
    clear (n);
       void add_edge (int u, int v, int f) {
    dest[size] = v; next[size] = begin[u]; flow[size] = f; inv[size] = size + 1; begin[u] = size++;
    dest[size] = u; next[size] = begin[v]; flow[size] = 0; inv[size] = size - 1; begin[v] = size++;
template <int MAXN = 100000, int MAXM = 100000>
struct cost_flow_edge_list {
      int size;
int begin[MAXN], dest[MAXM], next[MAXM], cost[MAXM], flow[MAXM], inv[MAXM];
void clear (int n) {
    size = 0;
    std::fill (begin, begin + n, -1);
}
       cost_flow_edge_list (int n = MAXN) {
    clear (n);
       void add_edge (int u, int v, int c, int f) {
    dest[size] = v; next[size] = begin[u]; cost[size] = c;
    flow[size] = f; inv[size] = size + 1; begin[u] = size++;
    dest[size] = u; next[size] = begin[v]; cost[size] = c;
    flow[size] = 0; inv[size] = size - 1; begin[v] = size++;
}
} ;
```

3.2 SPFA improved

```
/* SPFA :
```

3.3 Dijkstra's shortest path algorithm

3.4 Tarjan

3.5 Vertex biconnected component

3.6 Edge biconnected component

3.7 Hopcoft-Carp

```
unweighted maximum matching for bipartition graphs with complexity O (m \star n^0.5). struct hopcoft_carp :

Usage : solve () for maximum matching. The matching is in matchx and matchy.
        template <int MAXN = 100000, int MAXM = 100000>
struct hopcoft_carp {
                 int n, m;
                 int matchx[MAXN], matchy[MAXN], level[MAXN];
                 bool dfs (edge_list <MAXN, MAXM> &e, int x) {
   for (int i = e.begin[x]; ~i; i = e.next[i]) {
                                    int y = e.dest[i];
int w = matchy[y];
                                   if (w == -1 || (level[x] + 1 == level[w] && dfs (e, w))) {
   matchx[x] = y;
   matchy[y] = x;
   return true;
                                   }
                          level[x] = -1;
return false;
                }
for (int head = 0; head < (int) queue.size(); ++head) {
   int x = queue[head];
   for (int i = e.begin[x]; ~i; i = e.next[i]) {
      int y = e.dest[i];
      int w = matchy[y];
      if (w != -1 && level[w] < 0) {
            level[w] = level[x] + 1;
            queue.push_back (w);
      }
}</pre>
                                            }
                                   int delta = 0;
for (int i = 0; i < n; ++i)
    if (matchx[i] == -1 && dfs (e, i)) delta++;
if (delta == 0) return answer;
else answer += delta;</pre>
                 }
        };
                Kuhn-Munkres
3.8
              Kuhn Munkres algorithm :
                          weighted maximum matching algorithm for bipartition graphs. Complexity O (N^3). struct kuhn_munkres:

Initialize: pass nx, ny as the size of both sets, w as the weight matrix.

Usage: solve() for the minimum matching. The exact matching is in link[].
         template <int MAXN = 500>
         struct kuhn_munkres {
                 int nx, ny;
int w[MAXN][MAXN];
                 int lx[MAXN], ly[MAXN], visx[MAXN], visy[MAXN], slack[MAXN], link[MAXN];
                 if (t == 0) {
    visy[y] = 1;
                                            if (link[y] == -1 || dfs (link[y])) {
   link[y] = x;
   return 1;
                                    } else slack[y] = std::max (slack[y], t);
                           return 0;
                         solve () {
int i, j;
std::fill (link, link + ny, -1);
std::fill (ly, ly + ny, 0);
for (i = 0; i < nx; i++)
    for (j = 0, lx[i] = INF; j < ny; j++)
        lx[i] = std::min (lx[i], w[i][j]);
for (int x = 0; x < nx; x++) {
    for (i = 0; i < ny; i++) slack[i] = -INF;
    while (true) {
        std::fill (visx, visx + nx, 0);
        std::fill (visy, visy + ny, 0);
        if (dfs (x)) break;
        int d = INF;
        for (i = 0; i < ny; i++)
              if (!visy[i] && clack[i]) d = slack[i];
        for (i = 0; i < nx; i++)</pre>
                  int solve () {
```

/* Hopcoft-Carp algorithm

3.9 Stochastic weighted maximum matching

```
Weighted matching algorithm : maximum matching for general weighted graphs. Not stable.
                Usage: Set k to the size of vertices, w to the weight matrix.

Note that k has to be even for the algorithm to work.
template <int MAXN = 500>
struct weighted_match {
       int k;
int k;
long long w[MAXN][MAXN];
int match[MAXN], path[MAXN], p[MAXN], len;
long long d[MAXN];
bool v[MAXN];
       pact.
if (v[i]) return erac,
v[i] = true;
for (int j = 0; j < k; ++j) {
    if (i != j && match[i] != j && !v[j]) {
        int kok = match[j];
        if (d[kok] < d[i] + w[i][j] - w[j][kok]) {
            d[kok] = d[i] + w[i][j] - w[j][kok];
            if (dfs (kok)) return true;
}</pre>
               }
--len;
v[i] = false;
return false;
     match[t] = path[j];
match[path[j]] = t;
                                         break;
                         }
if (!flag) {
    if (++cnt >= 2) break;
    std::random_shuffle (p, p + k);
                forg long ans = 0;
for (int i = 0; i < k; ++i
    ans += w[i][match[i]];
return ans / 2;</pre>
        }
1:
```

3.10 Weighted blossom (vfleaking ver.)

```
edge (int u = 0, int v = 0, int w = 0): u(u), v(v), w(w) {}
eage \land \};
int n, n_x;
edge g[MAXN * 2 + 1][MAXN * 2 + 1];
int lab[MAXN * 2 + 1], slack[MAXN * 2 + 1], st[MAXN * 2 + 1], pa[MAXN * 2 + 1];
int match[MAXN * 2 + 1], slack[MAXN * 2 + 1], st[MAXN * 2 + 1], vis[MAXN * 2 + 1];
int flower_from[MAXN * 2 + 1][MAXN + 1], S[MAXN * 2 + 1], vis[MAXN * 2 + 1];
std::vector<int> flower[MAXN * 2 + 1];
std::queue<int> q;
  void update_slack (int u, int x) {
   if (!slack[x] || e_delta (g[u][x]) < e_delta (g[slack[x]][x]))slack[x] = u;</pre>
  void set_slack (int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
        if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)update_slack (u, x);
  void q_push (int x) {
   if (x <= n)q.push (x);
   else for (size_t i = 0; i < flower[x].size(); i++)q_push (flower[x][i]);
.</pre>
  void set_st (int x, int b) {
    st[x] = b;
    if (x > n) for (size_t i = 0; i < flower[x].size(); ++i)
        set_st (flower[x][i], b);</pre>
 }
int get_pr (int b, int xr) {
    int pr = find (flower[b].begin(), flower[b].end(), xr) - flower[b].begin();
    if (pr % 2 == 1) {
        reverse (flower[b].begin() + 1, flower[b].end());
        return (int)flower[b].size() - pr;
    } else return pr;
}
void set_match (int u, int v) {
    match[u] = g[u][v].v;
    if (u > n) {
        edge e = g[u][v];
        int xr = flower_from[u][e.u], pr = get_pr (u, xr);
        for (int i = 0; i < pr; ++i)set_match (flower[u][i], flower[u][i ^ 1]);
        set_match (xr, v);
        rotate (flower[u].begin(), flower[u].begin() + pr, flower[u].end());
}</pre>
 }
void augment (int u, int v) {
    for (;;) {
        int xnv = st[match[u]];
        set_match (u, v);
        if (!xnv)return;
        set_match (xnv, st[pa[xnv]]);
        u = st[pa[xnv]], v = xnv;
}
              }
 }
int get_lca (int u, int v) {
    static int t = 0;
    for (++t; u || v; std::swap (u, v)) {
        if (u == 0) continue;
        if (vis[u] == t) return u;
        vis[u] = t;
        u = st[match[u]];
        if (u)u = st[pa[u]];
              return 0;
void add_blossom (int u, int lca, int v) {
   int b = n + 1;
   while (b <= n_x && st[b])++b;
   if (b > n_x)++n_x;
   lab[b] = 0, S[b] = 0;
   match[b] = match[lca];
   flower[b].clear();
   flower[b].push_back (lca);
   for (int x = u, y; x != lca; x = st[pa]);
}
            set slack (b);
set_stack (b),

void expand_blossom (int b) {
    for (size_t i = 0; i < flower[b].size(); ++i)
        set_st (flower[b][i]);
    int xr = flower_from[b][g[b][pa[b]].u], pr = get_pr (b, xr);
    for (int i = 0; i < pr; i += 2) {
        int xs = flower[b][i], xns = flower[b][i + 1];
        pa[xs] = g[xns][xs].u;
        S[xs] = 1, S[xns] = 0;
        slack[xs] = 0, set_slack (xns);
        q_push (xns);
}</pre>
              ;
S[xr] = 1, pa[xr] = pa[b];
for (size_t i = pr + 1; i < flower[b].size(); ++i) {
```

```
int xs = flower[b][i];
S[xs] = -1, set_slack (xs);
                         st[b] = 0;
            bool on_found_edge (const edge &e) {
                       l on_found_edge (const edge &e) {
int u = st[e.u], v = st[e.v];
if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q push (nu);
} else if (S[v] == 0) {
    int lca = get_lca (u, v);
    if (!lca)return augment (u, v), augment (v, u), true;
    else add_blossom (u, lca, v);
}
                         return false;
         le (q.size()) {
   int u = q.front();
   q.pop();
   if (S[st[u]] == 1) continue;
   for (int v = 1; v <= n; ++v)
      if (g[u][v].w > 0 && st[u] != st[v]) {
        if (e_delta (g[u][v]) == 0) {
            if (on_found_edge (g[u][v])) return true;
        } else update_slack (u, st[v]);
   }

}
int d = INF;
for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b && S[b] == 1)d = std::min (d, lab[b] / 2);
for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x]) {
        if (S[x] == -1)d = std::min (d, e_delta (g[slack[x]][x]));
        else if (S[x] == 0)d = std::min (d, e_delta (g[slack[x]][x]) / 2);
}
</pre>
                                                 } (int u = 1; u <= n; ++u) {
if (S[st[u]] == 0) {
    if (lab[u] <= d)return 0;
    lab[u] -= d;
} else if (S[st[u]] == 1)lab[u] += d;</pre>
                                     }
for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b) {
        if (S[st[b]] == 0)lab[b] += d * 2;
        else if (S[st[b]] == 1)lab[b] -= d * 2;</pre>
                                    q = std::queue<int>();
for (int x = 1; x <= n_x; ++x)
    if (st[x] == x && slack[x] && st[slack[x]] != x && e_delta (g[slack[x]][x]) == 0)
        if (on_found_edge (g[slack[x]][x]))return true;
for (int b = n + 1; b <= n_x; ++b)
    if (st[b] == b && S[b] == 1 && lab[b] == 0)expand_blossom (b);</pre>
                         }
return false;
            std::pair <long long, int> solve () {
    std::fill (match + 1, match + n + 1, 0);
    n_x = n;
                         int n_matches = 0;
                        int n_matches = 0;
long long tot_weight = 0;
for (int u = 0; u <= n; ++u)st[u] = u, flower[u].clear();
int w_max = 0;
for (int u = 1; u <= n; ++u)
    for (int v = 1; v <= n; ++v) {
        flower_from[u][v] = (u == v ? u : 0);
        w_max = std::max (w_max, g[u][v].w);
}</pre>
                        for (int u = 1; u <= n; ++u)lab[u] = w_max;
while (matching())++n_matches;
for (int u = 1; u <= n; ++u)
    if (match[u] && match[u] < u)
        tot_weight += g[u][match[u]].w;
return std::make_pair (tot_weight, n_matches);</pre>
           }
void init () {
    for (int u = 1; u <= n; ++u)
        for (int v = 1; v <= n; ++v)
        g[u][v] = edge (u, v, 0);</pre>
};
                Maximum flow
       Sparse graph maximum flow :
```

3.11

```
int isap::solve (flow_edge_list &e, int n, int s, int t) :
              e : edge list.
n : vertex size.
s : source.
t : sink.
template <int MAXN = 1000, int MAXM = 100000>
struct isap {
```

```
int pre[MAXN], d[MAXN], gap[MAXN], cur[MAXN];
          int solve (flow_edge_list <MAXN, MAXN> &e, int n, int s, int t) {
   std::fill (pre, pre + n + 1, 0);
   std::fill (d, d + n + 1, 0);
   std::fill (gap, gap + n + 1, 0);
   for (int i = 0; i < n; i++) cur[i] = e.begin[i];
   gap[0] = n;
   int u = pre[s] = s, u maxflow = 0;</pre>
                     int u = pre[s] = s, v, maxflow = 0;
while (d[s] < n) {
    v = n;</pre>
                               v = n;
for (int i = cur[u]; ~i; i = e.next[i])
    if (e.flow[i] && d[u] == d[e.dest[i]] + 1) {
       v = e.dest[i];
       cur[u] = i;
       broat
                                if (v < n) {
    pre[v] = u;
    u = v;
    - +)</pre>
                                         full v == ty;
if (v == t) {
   int dflow = INF, p = t;
   u = s;
   while (p != s) {
      p = pre[p];
      dflow = std::min (dflow, e.flow[cur[p]]);
}
                                                     maxflow += dflow;
                                                     e.flow[e.inv[cur[p]]] += dflow;
                               } else {
   int mindist = n + 1;
   int mindist = n + 1;
                                          for (int i = e.begin[u]; ~i; i = e.next[i])
    if (e.flow[i] && mindist > d[e.dest[i]]) {
        mindist = d[e.dest[i]];
        cur[u] = i.
                                                               cur[u] = i;
                                         }
if (!--gap[d[u]]) return maxflow;
gap[d[u] = mindist + 1]++;
u = pre[u];
                     return maxflow;
          }
};
         Dense graph maximum flow :
    int dinic::solve (flow_edge_list &e, int n, int s, int t) :
                                e : edge list
                               n : vertex size.
s : source.
t : sink.
template <int MAXN = 1000, int MAXM = 100000>
struct dinic {
          int n, s, t;
          int d[MAXN], w[MAXN], q[MAXN];
         int d[MAXN], w[MAXN], q[MAXN];
int bfs (flow_edge_list <MAXN, MAXM> &e) {
    for (int i = 0; i < n; i ++) d[i] = -1;
    int l, r;
    q[l = r = 0] = s, d[s] = 0;
    for (; l <= r; l ++)
        for (int k = e.begin[q[l]]; k > -1; k = e.next[k])
        if (d[e.dest[k]] == -1 && e.flow[k] > 0) d[e.dest[k]] = d[q[l]] + 1, q[++r] = e.dest[k];
    return d[t] > -1 ? 1 : 0;
}
         }
int dfs (flow_edge_list <MAXN, MAXM> &e, int u, int ext) {
    if (u == t) return ext;
    int k = w[u], ret = 0;
    for (; k > -1; k = e.next[k], w[u] = k) {
        if (ext == 0) break;
        if (d[e.dest[k]] == d[u] + 1 && e.flow[k] > 0) {
            int flow = dfs (e, e.dest[k], std::min (e.flow[k], ext));
            if (flow > 0) {
                  e.flow[k] -= flow, e.flow[e.inv[k]] += flow;
                  ret += flow, ext -= flow;
        }
}
                     }
if (k == -1) d[u] = -1;
return ret;
          void solve (flow_edge_list <MAXN, MAXM> &e, int n, int s, int t) {
    dinic::n = n; dinic::s = s; dinic::t = t;
    while (bfs (e)) {
        for (int i = 0; i < n; i ++) w[i] = e.begin[i];
        dfs (e, s, INF);</pre>
};
```

3.12 Minimum cost flow

```
returns the flow and the cost respectively.
*/
template <int MAXN = 1000, int MAXM = 100000>
struct minimum_cost_flow {
       int n, source, target;
int prev[MAXN];
int dist[MAXN], occur[MAXN];
              laugment (cost_flow_edge_list <MAXN, MAXM> &e) {
  std::vector <int> queue;
  std::fill (dist, dist + n, INF);
  std::fill (occur, occur + n, 0);
  dist[source] = 0;
  occur[source] = true;
  queue.push_back (source);
  for (int head = 0; head < (int)queue.size(); ++head) {
    int x = queue[head];
    for (int i = e.begin[x]; ~i; i = e.next[i]) {
        int y = e.dest[i];
        if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
            dist[y] = dist[x] + e.cost[i];
            prev[y] = i;
            if (!occur[y]) {
                  occur[y] = true;
                 queue.push_back (y);
        }
}
       bool augment (cost_flow_edge_list <MAXN, MAXM> &e) {
                               }
                       occur[x] = false;
                return dist[target] < INF;</pre>
        std::pair <int, int> solve (cost_flow_edge_list <MAXN, MAXM> &e, int n, int s, int t) {
               ::pair <int, int> solve (cost_ilow_eage_ilst value, int)
minimum_cost_flow::n = n;
source = s; target = t;
std::pair <int, int> answer = std::make_pair (0, 0);
while (augment (e)) {
   int number = INF;
   for (int i = target; i != source; i = e.dest[e.inv[prev[i]]]) {
        number = std::min (number, e.flow[prev[i]]);
   }
}
                       answer.first += number;
for (int i = target; i != source; i = e.dest[e.inv[prev[i]]]) {
    e.flow[prev[i]] -= number;
    e.flow[e.inv[prev[i]]] += number;
    answer.second += number * e.cost[prev[i]];
                       }
                return answer;
       }
       Dense graph minimum cost flow
               std::pair <int, int> zkw_flow::solve (cost_flow_edge_list &e, int n, int s, int t) :
                               edge list
                       n : vertex size.
s : source.
                               sink
                        returns the flow and the cost respectively.
template <int MAXN = 1000, int MAXM = 100000>
struct zkw_flow {
       int n, s, t, totFlow, totCost;
int dis[MAXN], slack[MAXN], visit[MAXN];
       if (delta == INF) return 1;
for (int i = 0; i < n; i++) if (visit[i]) dis[i] += delta;
return 0;</pre>
        int dfs (cost_flow_edge_list <MAXN, MAXM> &e, int x, int flow) {
               if (x == t) {
   totFlow += flow;
   totCost += flow * (dis[s] - dis[t]);
                       return flow;
              slack[y] = std::min (slack[y], dis[y] + e.cost[i] - dis[x]);
                return flow - left;
        std::pair <int, int> solve (cost_flow_edge_list <MAXN, MAXM> &e, int n, int s, int t) {
   zkw_flow::n = n; zkw_flow::s = s; zkw_flow::t = t;
   totFlow = 0; totCost = 0;
   std::fill (dis + 1, dis + t + 1, 0);
}
                       do {
                       std::fill (visit + 1, visit + t + 1, 0);
} while (dfs (e, s, INF));
```

```
} while (!modlable ());
    return std::make_pair (totFlow, totCost);
}
```

3.13 Dominator tree

```
/* Dominator tree
                 nator tree :
void dominator_tree::solve (int s, int n, const edge_list <MAXN, MAXM> &succ) :
    solves for the immediate dominator (idom[]) of each node,
        idom[x] will be x if x does not have a dominator,
        and will be -1 if x is not reachable from s.
template <int MAXN = 100000, int MAXM = 100000>
struct dominator_tree {
        int dfn[MAXN], sdom[MAXN], idom[MAXN], id[MAXN], fa[MAXN], smin[MAXN], stamp;
       void predfs (int x, const edge_list <MAXN, MAXM> &succ) {
  id[dfn[x] = stamp++] = x;
  for (int i = succ.begin[x]; ~i; i = succ.next[i]) {
    int y = succ.dest[i];
    if (dfn[y] < 0) {
        f[y] = x;
        predfs (y, succ);
    }
}</pre>
                }
        }
       int getfa (int x) {
    if (fa[x] == x) return x;
    int ret = getfa (fa[x]);
    if (dfn[sdom[smin[fa[x]]]] < dfn[sdom[smin[x]]]) {
        smin[x] = smin[fa[x]];
    }
}</pre>
                 }
return fa[x] = ret;
       t x = id[o];
(o) {
    sdom[x] = f[x];
    for (int i = pred.begin[x]; ~i; i = pred.next[i]) {
        int p = pred.dest[i];
        if (dfn[p] < 0) continue;
        if (dfn[p] > dfn[x]) {
            getfa (p);
            p = sdom[smin[p]];
        }
}
                                                    sdom[x] = p;
                                   tmp.add_edge (sdom[x], x);
                          while (~tmp.begin[x]) {
   int y = tmp.dest[tmp.begin[x]];
   tmp.begin[x] = tmp.next[tmp.begin[x]];
                                   getfa (y);
if (x != sdom[smin[y]]) {
                                           idom[y] = smin[y];
                                  } else { idom[y] = x;
                          for (int i = succ.begin[x]; ~i; i = succ.next[i]) {
    if (f[succ.dest[i]] == x) {
        fa[succ.dest[i]] = x;
    }
                }
idom[s] = s;
for (int i = 1; i < stamp; ++i) {
   int x = id[i];
   if (idom[x] != sdom[x]) {
      idom[x] = idom[idom[x]];
   }
}</pre>
       }
} :
```

Chapter 4

Trial of string

4.1 KMP

```
/* KMP algorithm :
    void kmp::build (const std::string & str) :
        initializes and builds the failure array. Complexity O (n).
    int kmp::find (const std::string & str) :
        finds the first occurence of match in str. Complexity O (n).
    Note : match is cylic when L % (L - 1 - fail[L - 1]) == 0 &&
        L / (L - 1 - fail[L - 1]) > 1, where L = match.size ().

*/

template <int MAXN = 1E6>
struct kmp {
    std::string match;
    int fail[MAXN];
    void build (const std::string & str) {
        match = str; fail[0] = -1;
        for (int i = 1; i < (int) str.size (); ++i) {
            int j = fail[i - 1];
            while ("j && str[i] != str[j + 1]) j = fail[j];
            fail[i] = str[i] == str[j + 1] ? j + 1 : -1;
        }
    int find (const std::string & str) {
        for (int i = 0, j = 0; i < (int) str.size (); ++i, ++j) {
            if (j == match.size ()) return i - match.size ();
            while ("j && str[i] != match[j]) j = fail[j];
        }
    return str.size ();
}
</pre>
```

4.2 Suffix automaton

```
/* Suffix automaton:
    void suffix automaton::init ():
        initializes the automaton with an empty string.
    void suffix_automaton:extend (int token):
        extends the string with token. Complexity O (1).
    head: the first state.
    tail: the last state.
        Terminating states can be reached via visiting the ancestors of tail.
    state::len: the longest length of the string in the state.
    state::parent: the parent link.
    state::parent: the parent link.

*/

template <int MAXN = 1E6, int MAXC = 26>
    struct suffix_automaton {
    state *head, *tail;
    struct suffix_automaton {
        state *parent, *dest[MAXC];
        state (int len = 0): len (len), parent (NULL) {
            memset (dest, 0, sizeof (dest));
        } state *parent, *dest[MAXC];
        state *parent, *dest[M
```

```
}
    return np == null ? np -> parent : np;
}
void init () {
    tot_node = node_pool;
    head = tail = new (tot_node++) state();
}
suffix_automaton () {
    init ();
}
void extend (int token) {
    tail = tail -> extend (head, token);
}
};
```

4.3 Palindromic tree

```
/* Palindromic tree :
    void palindromic_tree::init () : initializes the tree.
    bool palindromic_tree::extend (int) : extends the string with token.
    returns whether the tree has generated a new node.
                         Complexity O (log MAXC).
                 odd, even : the root of two trees.
                last: the node representing the last char.
node::len: the palindromic string length of the node.
template <int MAXN = 1E6, int MAXC = 26>
struct palindromic_tree {
        struct node {
  node *child[MAXC], *fail;
  int len;
  node (int len) : fail (NULL), len (len) {
      memset (child, NULL, sizeof (child));
}
        } node_pool[MAXN * 2], *tot_node;
        int size, text[MAXN];
        node *odd, *even, *last;
node *match (node *now) {
    for (; text[size - now -> len - 1] != text[size]; now = now -> fail);
    return now;
       bool extend (int token) {
    text[++size] = token;
    node *now = match (last);
    if (now -> child[token])
        return last = now -> child[token], false;
    last = now -> child[token] = new (tot_node++) node (now -> len + 2);
    if (now == odd) last -> fail = even;
else {
                 else {
  now = match (now -> fail);
  last -> fail = now -> child[token];
                 return true;
        }
        void init() {
   text[size = 0] = -1;
   tot_node = node_pool;
   last = even = new (tot_node++) node (0); odd = new (tot_node++) node (-1);
   even -> fail = odd;
        palindromic_tree () {
   init ();
        }
1;
```

Chapter 5

Reference

5.1 Vimrc

```
set ruler
set number
set tabstop=4
set softtabstop=4
set siftwidth=4
set sinftwidth=4
set smartindent
set showmatch
set locsearch
set incsearch
set incsearch
set autoread
set backspace=2
set mouse=a
syntax on
nmap <C-A> ggVG
vmap <C-C> "+y
nmap <C-P> "+p
autocmd FileType cpp set cindent
autocmd FileType cpp map <F3> :vsplit %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time ./%<.exe < %<.in <CR>
autocmd FileType cpp map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < R>
size %<.exe <CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5> :!time java %< < CR>
autocmd FileType java map <F5
```

5.2 Java reference

```
LinkedList <E> :
    addFirst / addLast (E) / getFirst / getLast / removeFirst / removeLast () :
                           deque implementation.

clear () / add (int, E) / remove (int) : clear, add & remove.

size () / contains / removeFirstOccurrence / removeLastOccurrence (E) :
    deque methods.
                           ListIterator <E> listIterator (int index) : return E next / previous () : accesses and iterates. hasNext / hasPrevious () : checks availablity
                                                                                                                          : returns an iterator :
                 hasNext / hasPrevious () : checks availablity.
nextIndex / previousIndex () : returns the index of a subsequent call.
add / set (E) / remove () : changes element.

PriorityQueue <E> (int initcap, Comparator <? super E> comparator) :
add (E) / clear () / iterator () / peek () / poll () / size () :
priority queue implementations.

TreeMap <K, V> (Comparator <? super K> comparator) :
Map.Entry <K, V> ceilingEntry / floorEntry / higherEntry / lowerEntry (K):
    getKey / getValue () / setValue (V) : entries.
clear () / put (K, V) / get (K) / remove (K) : basic operation.
size () : size.

StringBuilder :
                  StringBuilder
                          ingBuilder :
Mutable string.
StringBuilder (string) : generates a builder.
append (int, string, ...) / insert (int offset, ...) : adds objects.
charAt (int) / setCharAt (int, char) : accesses a char.
delete (int, int) : removes a substring.
reverse () : reverses itself.
length () : returns the length.
toString () : converts to string.
ing :
                  String :
                           Immutable string.
                           String.format (String, ...) : formats a string. i.e. sprintf. toLowerCase / toUpperCase () : changes the case of letters.
/* Examples on Comparator :
public class Main {
        public static class Point {
                 public int x;
public int y;
public Point () {
    x = 0;
    y = 0;
}
                 public Point (int xx, int yy) {
    x = xx;
    y = yy;

         };
        public static class Cmp implements Comparator <Point> {
   public int compare (Point a, Point b) {
     if (a.x < b.x) return -1;
     if (a.x == b.x) {
        if (a.y < b.y) return -1;
        if (a.y == b.y) return 0;
}</pre>
                           return 1:
                 }
         };
        public static void main (String [] args) {
   Cmp c = new Cmp ();
   TreeMap <Point, Point> t = new TreeMap <Point, Point> (c);
                 return:
         1
         Another way to implement is to use Comparable.
         However, equalTo and hashCode must be rewritten.
         Otherwise, containers may fail.
         Example :
        public static class Point implements Comparable <Point> {
                 public int x;
                 public int y;
public Point () {

\mathbf{x} = 0; \\
\mathbf{y} = 0;

                 public Point (int xx, int yy) {
                          y = xx;

y = yy;
                 public int compareTo (Point p) {
                          if (x < p.x) return -1;
if (x == p.x) {
   if (y < p.y) return -1;
   if (y == p.y) return 0;</pre>
                           }
return 1;
                 public boolean equalTo (Point p) {
   return (x == p.x && y == p.y);
                 public int hashCode () {
    return x + y;
                 }
       };
//Faster IO :
public class Main {
         static class InputReader {
                 public BufferedReader reader;
public StringTokenizer tokenizer;
                 public InputReader (InputStream stream) {
   reader = new BufferedReader (new InputStreamReader (stream), 32768);
   tokenizer = null;
```

5.3 Operator precedence

Precedence Operator		Description	Associativity	
1	::	Scope resolution		
	a++ a	Suffix/postfix increment and decrement		
	type()	Functional cast	Left-to-right	
2	type{}		12010 00 118110	
	a()	Function call		
	a[]	Subscript		
	>	Member access		
	++aa	Prefix increment and decrement		
	+a -a	Unary plus and minus		
	! ~	Logical NOT and bitwise NOT		
	(type)	C-style cast		
3	*a	Indirection (dereference)	Right-to-left	
	&a	Address-of		
	sizeof	Size-of		
	new new[]	Dynamic memory allocation		
	delete	Dynamic memory deallocation		
	delete[]	, , , , , , , , , , , , , , , , , , ,		
4	.* ->*	Pointer-to-member		
5	a*b a/b	Multiplication, division, and remainder		
	a%b	, , ,		
6	a+b a-b	Addition and subtraction		
7	<< >>	Bitwise left shift and right shift		
8	< <=	For relational operators $<$ and \le respectively	Left-to-right	
	> >=	For relational operators $>$ and \ge respectively	2010 00 118110	
9	== !=	For relational operators = and \neq respectively		
10	a&b	Bitwise AND		
11	^	Bitwise XOR (exclusive or)		
12		Bitwise OR (inclusive or)		
13	& &	Logical AND		
14		Logical OR		
	a?b:c	Ternary conditional		
	throw	throw operator		
15	=	Direct assignment	Right-to-left	
	+= -= *=	Compound assignment by arithmetic operation	Tenglie to fore	
	/= %=			
	<<= >>=	Compound assignment by bitwise shift		
	&= ^= =	Compound assignment by bitwise AND, XOR, and		
		OR		
16	,	Comma	Left-to-right	

5.4 Hacks

5.4.1 Formating long long in scanf & printf

```
#ifdef WIN32
     #define LL "%164d"
#else
     #define LL "%1ld"
#endif

5.4.2 Optimizing

#pragma GCC optimize ("03")
#pragma GCC optimize ("whole-program")
```

```
5.4.3 Larger stack
```

```
5.4.3.1 C++
```

```
#pragma comment(linker, "/STACK:36777216")
```

5.4.3.2 G++

5.5 Math reference

5.5.1 Catalan number

For

$$\begin{split} f(0) &= 1 \\ f(1) &= 1 \\ f(n) &= f(n-1)f(0) + f(n-2)f(1) + \ldots + f(1)f(n-2) + f(0)f(n-1) \end{split}$$

We have
$$f(n) = \frac{(2n)!}{n!(n+1)!}$$
.

5.5.2 Dynamic programming optimization

5.5.2.1 Convex hull optimization

Generally, in dynamic programming with recurrence

$$f(i) = \min_{k < i} \{a[i]b(j) + c(j)\}$$

all decisions k can be treated as a set of segments on a convex hull. By applying Graham's scanning, it is possible to maintain such hull in a monotone queue or a std::tuple <slope, intercept, x_min>. Hence, k(i) can be obtained by performing a binary search in the hull.

5.5.2.2 Divide & conquer optimization

For recurrence

$$f(i) = \min_{k < i} \{b(k) + c[k][i]\}$$

 $k(i) \le k(i+1)$ holds true if c[a][c] + c[b][d] < c[a][d] + c[b][c]. Thus, k(i) can be maintained in a monotone queue.

5.5.2.3 Knuth optimization

For recurrence

$$f(i,j) = \min_{i < k < j} \{ f(i,k) + f(k,j) \} + c[i][j]$$

$$k(i, j - 1) \le k(i, j) \le k(i + 1, j)$$
 holds true if $c[a][c] + c[b][d] < c[a][d] + c[b][c]$.

5.5.3 Integration table

5.5.3.1 $ax^2 + bx + c \ (a > 0)$

1.
$$\int \frac{\mathrm{d}x}{ax^2 + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}} + C & b^2 < 4ac \\ \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right| + C & b^2 > 4ac \end{cases}$$

2.
$$\int \frac{x}{ax^2 + bx + c} dx = \frac{1}{2a} \ln |ax^2 + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^2 + bx + c}$$

5.5.3.2 $\sqrt{\pm ax^2 + bx + c}$ (a > 0)

1.
$$\int \frac{dx}{ax^2 + bx + c} = \frac{1}{\sqrt{a}} \ln \left| 2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c} \right| + C$$

2.
$$\int \sqrt{ax^2 + bx + c} \, dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln \left| 2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c} \right| + C$$

3.
$$\int \frac{x}{\sqrt{ax^2+bx+c}} dx = \frac{1}{a} \sqrt{ax^2+bx+c} - \frac{b}{2\sqrt{a^3}} \ln \left| 2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c} \right| + C$$

4.
$$\int \frac{\mathrm{d}x}{\sqrt{-ax^2+bx+c}} = -\frac{1}{\sqrt{a}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5.
$$\int \sqrt{-ax^2 + bx + c} \, dx = \frac{2ax - b}{4a} \sqrt{-ax^2 + bx + c} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

6.
$$\int \frac{x}{\sqrt{-ax^2+bx+c}} dx = -\frac{1}{a}\sqrt{-ax^2+bx+c} + \frac{b}{2\sqrt{a^3}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5.5.3.3 Trigonometric

1.
$$\int \tan x \, \mathrm{d}x = -\ln|\cos x| + C$$

2.
$$\int \cot x \, dx = \ln|\sin x| + C$$

3.
$$\int \sec x \, dx = \ln |\tan(\frac{\pi}{4} + \frac{x}{2})| + C = \ln |\sec x + \tan x| + C$$

4.
$$\int \csc x \, dx = \ln|\tan \frac{x}{2}| + C = \ln|\csc x - \cot x| + C$$

5.
$$\int \sec^2 x \, \mathrm{d}x = \tan x + C$$

6.
$$\int \csc^2 x \, dx = -\cot x + C$$

7.
$$\int \sec x \tan x \, dx = \sec x + C$$

8.
$$\int \csc x \cot x \, dx = -\csc x + C$$

9.
$$\int \sin^2 x \, dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$$

10.
$$\int \cos^2 x \, dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$$

11.
$$\int \sin^n x \, dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x \, dx$$

12.
$$\int \cos^n x \, dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x \, dx$$

13.
$$\int \frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$$

14.
$$\int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$

15.
$$\int \cos^m x \sin^n x \, dx = \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x \, dx = -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x \, dx$$

5.5.3.4 Inverse trigonometric (a > 0)

1.
$$\int \arcsin \frac{x}{a} \, dx = x \arcsin \frac{x}{a} + \sqrt{a^2 - x^2} + C$$

2.
$$\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} - \sqrt{a^2 - x^2} + C$$

3.
$$\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2) + C$$

5.5.3.5 Exponential

$$1. \int a^x \, \mathrm{d}x = \frac{1}{\ln a} a^x + C$$

$$2. \int e^{ax} \, \mathrm{d}x = \frac{1}{a} a^{ax} + C$$

5.5.3.6 Logistic

- 1. $\int \ln x \, dx = x \ln x x + C$
- 2. $\frac{\mathrm{d}x}{x \ln x} = \ln |\ln x| + C$

Prefix sum of multiplicative functions

Define the Dirichlet convolution f * g(n) as:

$$f * g(n) = \sum_{d=1}^{n} [d|n] f(n) g(\frac{n}{d})$$

Assume we are going to calculate some function $S(n) = \sum_{i=1}^{n} f(i)$, where f(n) is a multiplicative function. Say we find some g(n) that is simple to calculate, and $\sum_{i=1}^{n} f * g(i)$ can be figured out in O(1) complexity. Then we have

$$\sum_{i=1}^{n} f * g(i) = \sum_{i=1}^{n} \sum_{d} [d|i] g(\frac{i}{d}) f(d)$$

$$= \sum_{i=1}^{n} \sum_{d=1}^{\left \lfloor \frac{n}{d} \right \rfloor} g(\frac{i}{d}) f(d)$$

$$= \sum_{i=1}^{n} \sum_{d=1}^{\left \lfloor \frac{n}{i} \right \rfloor} g(i) f(d)$$

$$= g(1) S(n) + \sum_{i=2}^{n} g(i) S(\left \lfloor \frac{n}{i} \right \rfloor)$$

$$S(n) = \frac{\sum_{i=1}^{n} f * g(i) - \sum_{i=2}^{n} g(i) S(\left \lfloor \frac{n}{i} \right \rfloor)}{g(1)}$$

It can be proven that $\left|\frac{n}{i}\right|$ has at most $O(\sqrt{n})$ possible values. Therefore, the calculation of S(n) can be reduced to $O(\sqrt{n})$ calculations of $S(\lfloor \frac{n}{i} \rfloor)$. By applying the master theorem, it can be shown that the complexity of such method is $O(n^{\frac{3}{4}})$.

Moreover, since f(n) is multiplicative, we can process the first $n^{\frac{2}{3}}$ elements via linear sieve, and for the rest of the elements, we apply the method shown above. The complexity can thus be enhaced to $O(n^{\frac{2}{3}})$.

For the prefix sum of Euler's function $S(n) = \sum_{i=1}^{n} \varphi(i)$, notice that $\sum_{d|n} \varphi(d) = n$. Hence $\varphi * I(n) = id(n).(I(n)) = id(n)$ 1, id(n) = n) Now let g(n) = I(n), and we have $S(n) = \sum_{i=1}^{n} i - \sum_{i=2}^{n} S(\left\lfloor \frac{n}{i} \right\rfloor)$. For the prefix sum of Mobius function $S(n) = \sum_{i=1}^{n} \mu(i)$, notice that $\mu * I(n) = [n = 1]$. Hence $S(n) = 1 - \sum_{i=2}^{n} S(\left\lfloor \frac{n}{i} \right\rfloor)$.

5.5.5Prufer sequence

In combinatorial mathematics, the Prufer sequence of a labeled tree is a unique sequence associated with the tree. The sequence for a tree on n vertices has length n-2.

One can generate a labeled tree's Prufer sequence by iteratively removing vertices from the tree until only two vertices remain. Specifically, consider a labeled tree T with vertices 1, 2, ..., n. At step i, remove the leaf with the smallest label and set the ith element of the Prufer sequence to be the label of this leaf's neighbour.

One can generate a labeled tree from a sequence in three steps. The tree will have n+2 nodes, numbered from 1 to n+2. For each node set its degree to the number of times it appears in the sequence plus 1. Next, for each number in the sequence a[i], find the first (lowest-numbered) node, j, with degree equal to 1, add the edge (j, a[i]) to the tree, and decrement the degrees of j and a[i]. At the end of this loop two nodes with degree 1 will remain (call them u, v). Lastly, add the edge (u,v) to the tree.

The Prufer sequence of a labeled tree on n vertices is a unique sequence of length n-2 on the labels 1 to n - this much is clear. Somewhat less obvious is the fact that for a given sequence S of length n-2 on the labels 1 to n, there is a unique labeled tree whose Prufer sequence is S.

5.5.6Spanning tree counting

Kirchhoff's Theorem: the number of spanning trees in a graph G is equal to any cofactor of the Laplacian matrix of G, which is equal to the difference between the graph's degree matrix (a diagonal matrix with vertex degrees on the diagonals) and its adjacency matrix (a (0,1)-matrix with 1's at places corresponding to entries where the vertices are adjacent and 0's otherwise).

The number of edges with a certain weight in a minimum spanning tree is fixed given a graph. Moreover, the number of its arrangements can be obtained by finding a minimum spanning tree, compressing connected components of other edges in that tree into a point, and then applying Kirrchoff's theorem with only edges of the certain weight in the graph. Therefore, the number of minimum spanning trees in a graph can be solved by multiplying all numbers of arrangements of edges of different weight together.

5.6 Regular expression

```
/* C++11 supports regular expressions, see below for an example. */
char str[] = "The_thirty-three_thieves_thought_that_they_thrilled_the_throne_throughout_Thursday.";
std::regex_pattern ("(th|Th)[\\w]*", std::regex_constants::optimize | std::regex_constants::ECMAScript);
std::match_results <char *> match;
std::regex_constants::match_flag_type flag = std::regex_constants::match_default;
int begin = 0, end = strlen (str);
while (std::regex_search (str + begin, str + end, match, pattern, flag)) {
    std::cout << match[0] << "_" << match[1] << std::endl;
    begin += match.position (0) + 1;
    flag |= std::regex_constants::match_prev_avail;
}</pre>
```

5.6.1 Special pattern characters

Characters	Description	Matches
•	Not newline	Any character except line terminators (LF, CR, LS, PS).
\t	Tab (HT)	A horizontal tab character (same as \u0009).
\n	Newline (LF)	A newline (line feed) character (same as \u000A).
\v	Vertical tab (VT)	A vertical tab character (same as \u000B).
\f	Form feed (FF)	A form feed character (same as \u0000C).
\r	Carriage return (CR)	A carriage return character (same as \u0000D).
\cletter Control code of dividing the		A control code character whose code unit value is the same as the remainder of dividing the code unit value of letter by 32. For example: \ca is the same as \u00001, \cb the same as \u00002, and so on
\xhh	ASCII character	A character whose code unit value has an hex value equivalent to the two hex digits hh. For example: \x4c is the same as L, or \x23 the same as #.
\uhhhh	Unicode character	A character whose code unit value has an hex value equivalent to the four hex digits hhhh.
\0	Null	A null character (same as \u0000).
\int	Backreference	The result of the submatch whose opening parenthesis is the int-th (int shall begin by a digit other than 0). See groups below for more info.
\d	Digit	A decimal digit character (same as [[:digit:]]).
\D	Not digit	Any character that is not a decimal digit character (same as [^[:digit:]]).
\s	Whitespace	A whitespace character (same as [[:space:]]).
\S	Not whitespace	Any character that is not a whitespace character (same as [^[:space:]]).
\w	Word	An alphanumeric or underscore character (same as [_[:alnum:]]).
\W	Not word	Any character that is not an alphanumeric or underscore character (same as [^_[:alnum:]]).
\character	Character	The character character as it is, without interpreting its special meaning within a regex expression. Any character can be escaped except those which form any of the special character sequences above. Needed for: ^ \$ \ . * + ? () [] { } .
[class]	Character class	The target character is part of the class (see character classes below).
Negated abayagter		The target character is not part of the class (see character classes below).

5.6.2 Quantifiers

Characters Times		Effects	
*	0 or more	The preceding atom is matched 0 or more times.	
+	1 or more	The preceding atom is matched 1 or more times.	
?	0 or 1	The preceding atom is optional (matched either 0 times or once).	
{int}	int	The preceding atom is matched exactly int times.	
{int,}	int or more	The preceding atom is matched int or more times.	
{min,max}	Between min and	The preceding atom is matched at least min times, but not more than max	
\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\	max	The preceding atom is matched at least militatines, but not more than max.	

By default, all these quantifiers are greedy (i.e., they take as many characters that meet the condition as possible). This behavior can be overridden to ungreedy (i.e., take as few characters that meet the condition as possible) by adding a question mark (?) after the quantifier.

5.6.3 Groups

Characters Description		Effects
(subpattern)	Group	Creates a backreference.
(?:subpattern) Passive group		Does not create a backreference.

5.6.4 Assertions

Characters	Description	Condition for match
^	Beginning of line	Either it is the beginning of the target sequence, or follows a line terminator.
\$	End of line	Either it is the end of the target sequence, or precedes a line terminator.
\b	Word boundary	The previous character is a word character and the next is a non-word character (or vice-versa). Note: The beginning and the end of the target sequence are considered here as non-word characters.
\B	Not a word boundary	The previous and next characters are both word characters or both are non-word characters. Note: The beginning and the end of the target sequence are considered here as non-word characters.
(?=subpattern)	Positive lookahead	The characters following the assertion must match subpattern, but no characters are consumed.
(?!subpattern)	Negative lookahead	The characters following the assertion must not match subpattern, but no characters are consumed.

5.6.5 Alternative

A regular expression can contain multiple alternative patterns simply by separating them with the separator operator (+): The regular expression will match if any of the alternatives match, and as soon as one does.

5.6.6 Character classes

Class	Description	Equivalent (with regex_traits, default locale)
[:alnum:]	Alpha-numerical character	isalnum
[:alpha:]	Alphabetic character	isalpha
[:blank:]	Blank character	isblank
[:cntrl:]	Control character	iscntrl
[:digit:]	Decimal digit character	isdigit
[:graph:]	Character with graphical representation	isgraph
[:lower:]	Lowercase letter	islower
[:print:]	Printable character	isprint
[:punct:]	Punctuation mark character	ispunct
[:space:]	Whitespace character	isspace
[:upper:]	Uppercase letter	isupper
[:xdigit:]	Hexadecimal digit character	isxdigit
[:d:]	Decimal digit character	isdigit
[:w:]	Word character	isalnum
[:s:]	Whitespace character	isspace

Please note that the brackets in the class names are additional to those opening and closing the class definition. For example:

[[:alpha:]] is a character class that matches any alphabetic character.

[abc[:digit:]] is a character class that matches a, b, c, or a digit.

[^[:space:]] is a character class that matches any character except a whitespace.