

Information Theory, Part II

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This document contains lecture notes from Harker's Advanced Topics in Mathematics class in Information Theory II, taught by Dr. Anuradha Aiyer. This course is the second part of a two part offering that explores the basic concepts of Information Theory, as initially described by Claude Elwood Shannon at Bell Labs in 1948. In Part 2 of the course, we explore other applications of Information Theory to the disciplines of Gambling, Statistics, Physics, Computer science, Economics and Philosophy. These notes were taken using TeXShop and $\text{\LaTeX}2\epsilon$ and will be updated for each class. The reader is advised to note any errata at the source control repository <https://github.com/mananshah99/infotheory>.

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1 Unit 1: Gambling

We'll discuss the duality between the growth rate of investment (i.e. a horse race) and the entropy rate of the horse race and how the side information's financial value is tied to mutual information.

Definition 1 (Horse Race). We have m horses in a race in which the i th horse wins with probability p_i . If horse i wins, the payoff is o_i for 1^1 . We'll assume that the gambler invests his wealth across all horses and doesn't hold on to any of his money. Specifically, b_i is the fraction of wealth invested in horse i where $b_i \geq 0$ and $\sum b_i = 1$. If horse i wins, the gambler wins $o_i b_i$; this case occurs with probability p_i . The wealth at the end of the race is a random variable which we will attempt to maximize.

1.1 Repeated Gambling

Define S_n as the total growth in the gambler's wealth after n races. We have

$$S_n = \prod_{j=1}^n S(X_j) \quad \text{and} \quad S(X_j) = b(X_j) o(X_j)$$

with X representing the horse that wins (this changes between races). Here, $S(X_i)$ represents the factor by which the gambler's wealth grows. We can define the doubling rate of a race W as

$$E(\log S(X)) = \sum_{k=1}^m p_k \log(b_k o_k) = W(b, p)$$

Theorem 1. Let race outcomes $X_1, X_2, X_3, \dots, X_n$ be identically and independently distributed $\sim p(x)$. The wealth of a gambler using betting strategy b grows exponentially at the rate $W(b, p)$ such that $S_n = 2^{nW(b, p)}$.

Proof. Functions of independent random variables are also independent, so $\log S(X_1), \dots, \log S(X_n)$ are i.i.d. From our earlier definition of S_n we have

$$\frac{1}{n} \log S_n = \frac{1}{n} \sum \log S(X_i)$$

By the weak law of large numbers², this equates to $E(\log S(X)) = W(b, p)$. So we can conclude that $S_n = 2^{nW(b, p)}$ and the proof is complete. So if to maximize S_n , we'll need to maximize W . \square

Definition 2. The optimum doubling rate over all choices of b_i is

$$W^*(p) = \max_b W(b, p) = \max_{b: b_i \geq 0, \sum b_i = 1} \sum_i p_i \log b_i o_i$$

We must formally maximize $W(b, p)$ such that $\sum b_i = 1$. To do this, we'll apply Lagrange optimization. We have

$$J(b) = \sum p_i \log b_i o_i + \lambda \sum b_i$$

¹There are two ways to describe a bet: either a for 1 or b to 1. The first notation indicates an exchange that happens prior to the race, and the latter indicates an exchange that happens post-race (although in both cases the horses are picked before the race). More concretely, a for 1 indicates that if one places \$1 on a particular horse before the race, the payoff is \$ a iff the horse wins and \$0 if the horse loses. b to 1 indicates that one would pay \$1 after the race if a particular horse loses and win \$ b if the horse wins. The equivalency between these scenarios is $b = a - 1$.

²See <http://mathworld.wolfram.com/WeakLawofLargeNumbers.html> for more information.

Taking the partial with respect to b_i and setting it equal to 0,

$$\frac{\partial J}{\partial b_i} = \frac{p_i}{b_i} + \lambda$$

where $i \in \{1 \dots m\}$. $\sum b_i = 1$ results in $b_i = p_i$. Technically, we'd have to take the second derivative to prove that this is a maximum; this verification is left to the reader.

Theorem 2. $W^* = \sum p_i \log o_i - H(p)$

Appendix A—Quotes

- “I have a problem. It’s called gambling.” (Dr. Aiyer)
- “What’s E ? Entropy?” (David Zhu)
- “So is it the strong law of weak numbers?” (Jerry Chen)
- “It’s like a half life... but it’s a double life” (Steven Cao)
- “Isn’t this just Lagrange?”
10 minutes later..
“Wait, how do you do Lagrange again?” (Swapnil Garg)