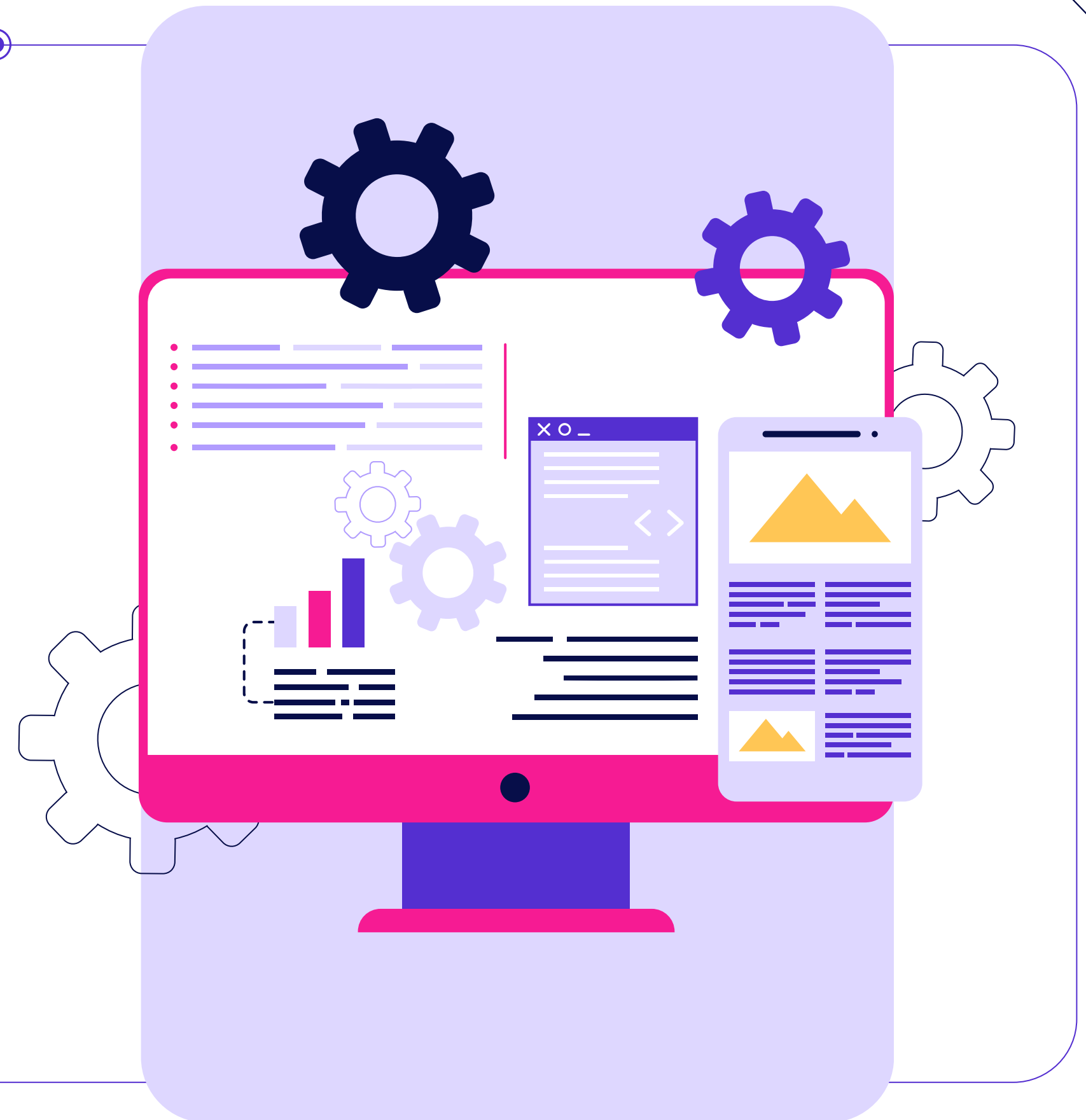


IT250 MINI PROJECT

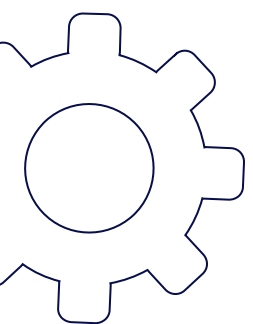
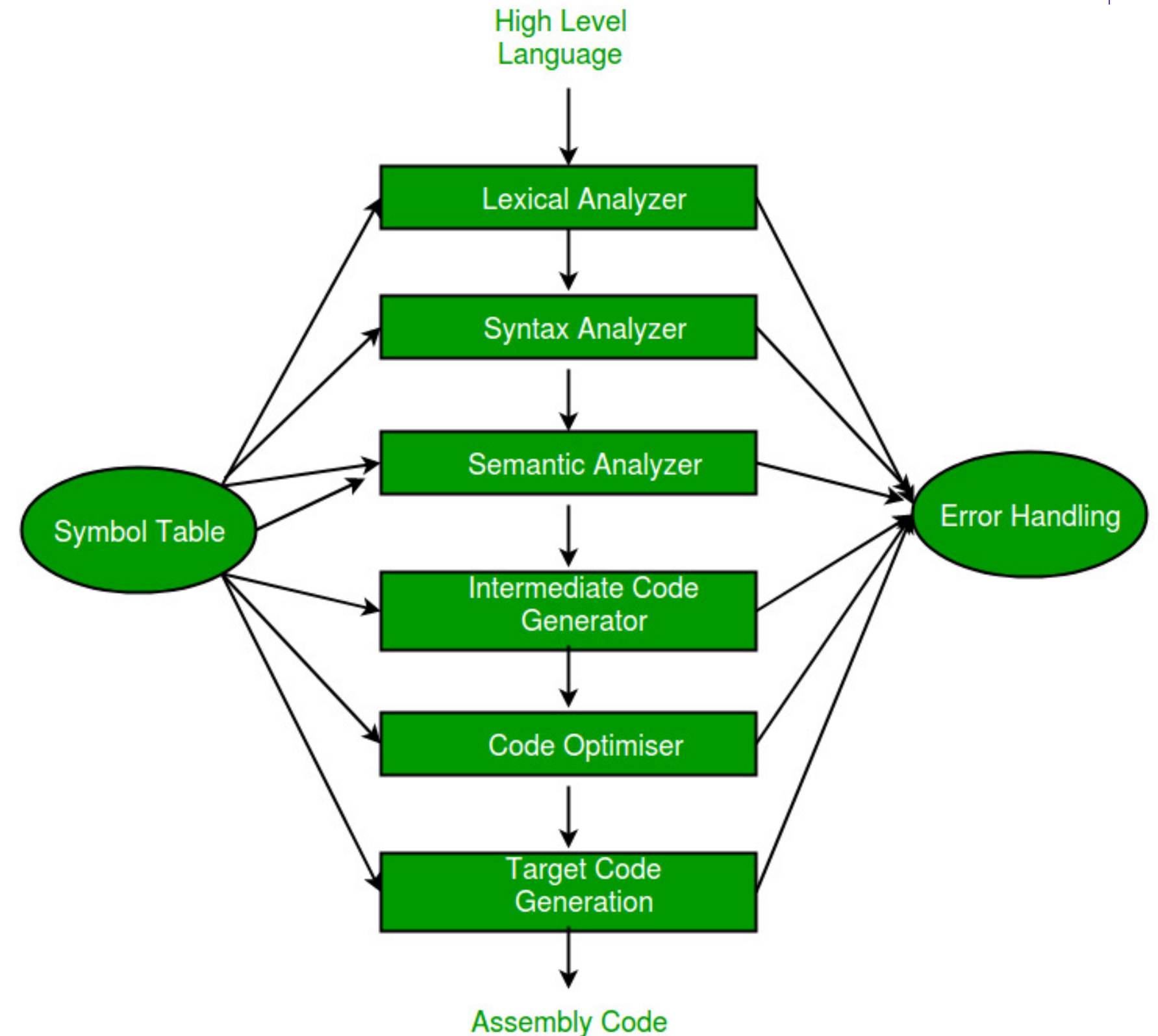
Nithin S
Ayush Kumar
Jay Chavan

221IT085
221IT015
221IT020

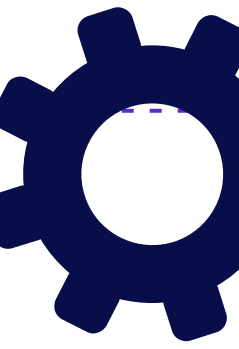


COMPILER

A compiler is a special program that processes statements written in a particular programming language and turns them into machine language or "code" that a computer's processor uses.



PHASES OF A COMPILER



01 Lexical Analyzer

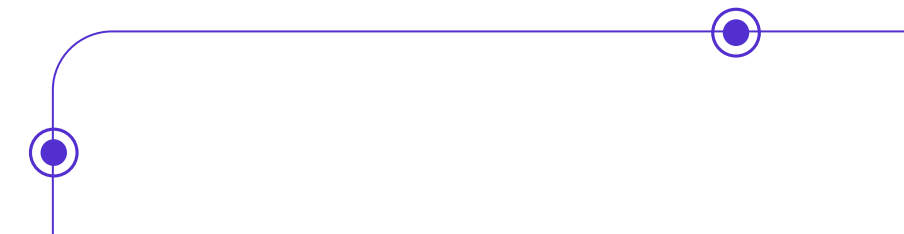
04 Intermediate Code Generator

02 Syntax Analyzer

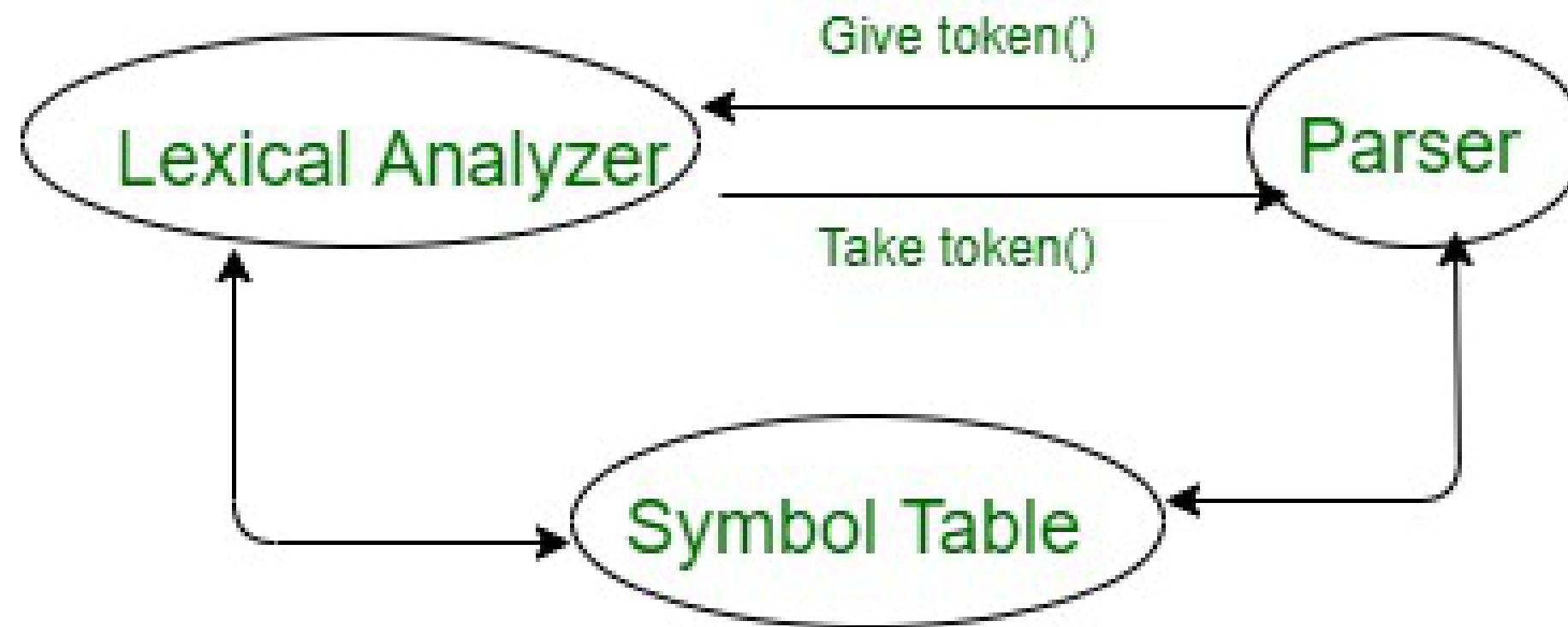
05 Code Optimizer

03 Semantic Analyzer

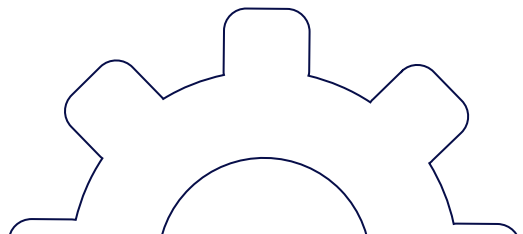
06 Target Code Generator



LEXICAL ANALYZER



- Scans the Pure High Level Language Code Line by Line
- Takes Lexemes as input and produces Tokens using DFA for pattern matching
- Removes Comments and Whitespaces from the Pure High Level Code
- Helps in macro expansion in the Pure HLL Code
- Creates a Symbol Table



Input

```
// test case to check loop statements

int main(){

    int i, a, b;
    int nume=3.45;
    for(i = 0; i < 10; i++){
        a=i;
    }
    i=1;
}
```

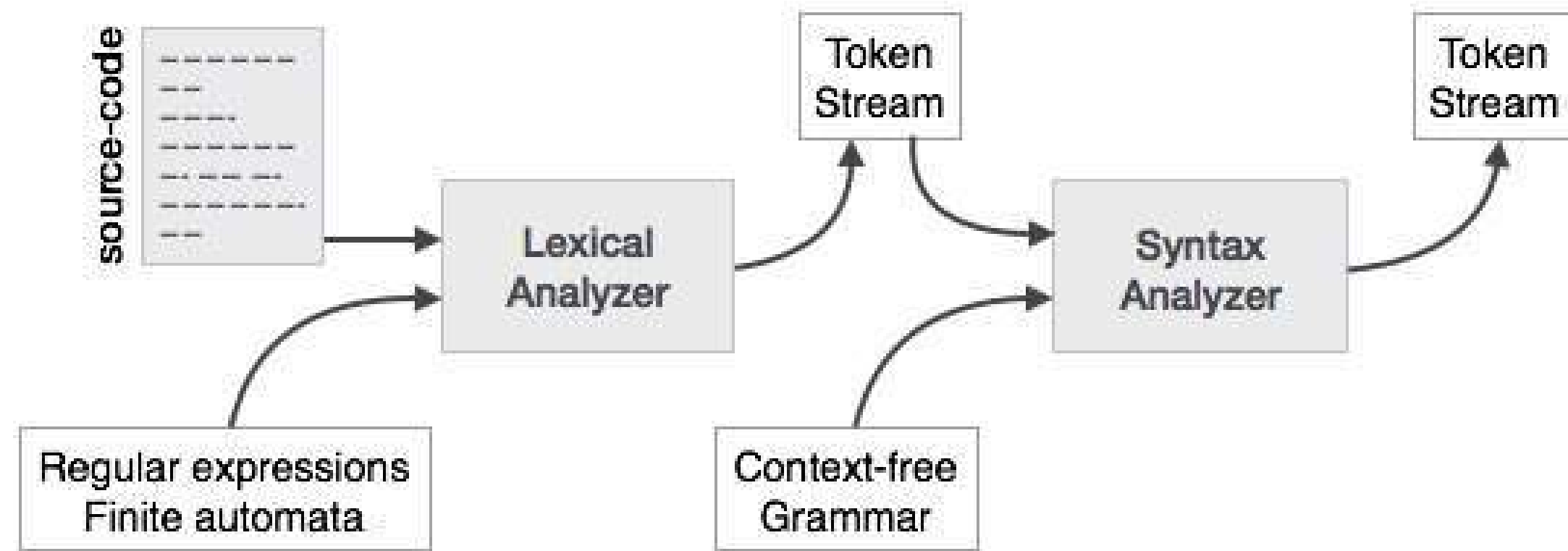
Output

Table:				
Lexeme	Token	Attribute Value	Line Number	
int	Keyword	0	3	
main	Procedure	1	3	
{	Punctuator	2	3	
int	Keyword	0	5	
i	Identifier	3	5	
,	Punctuator	4	5	
a	Identifier	5	5	
,	Punctuator	4	5	
b	Identifier	6	5	
;	Punctuator	7	5	
int	Keyword	0	6	
nume	Identifier	8	6	
=	Assignment Op	9	6	
3.45	Float Constant	10	6	
;	Punctuator	7	6	
for	Keyword	11	7	
(Punctuator	12	7	
i	Identifier	3	7	
=	Assignment Op	9	7	
0	Integer Constant	13	7	
;	Punctuator	7	7	
i	Identifier	3	7	
<	Relational Op	14	7	
10	Integer Constant	15	7	
;	Punctuator	7	7	
i	Identifier	3	7	
+	Arithmetic Op	16	7	
+	Arithmetic Op	16	7	
)	Punctuator	17	7	
{	Punctuator	2	7	
a	Identifier	5	8	
=	Assignment Op	9	8	
i	Identifier	3	8	
;	Punctuator	7	8	
}	Punctuator	18	9	
i	Identifier	3	10	
=	Assignment Op	9	10	
1	Integer Constant	19	10	
;	Punctuator	7	10	
}	Punctuator	18	11	

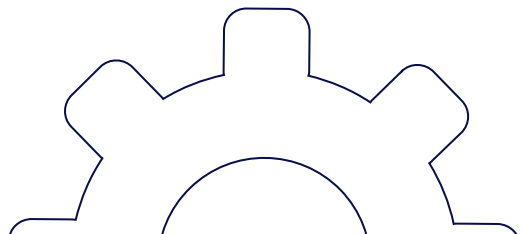
MultiLineComment (0 lines):

SingleLineComment :
test case to check loop statements

SYNTAX ANALYZER



- Checks for syntactic errors like missing semicolons, mismatched parentheses, etc.
- May involve resolving ambiguities in the grammar.
- Implements parsing algorithms such as LL, LR, or Recursive Descent.
- Constructs an Abstract Syntax Tree (AST) representing the hierarchical structure of the code.
- Handles language features like function prototypes, declarations, and definitions.



Input

```
// test case to check loop statements

int main(){

    int i, a, b;
    int nume=3.45;
    for(i = 0; i < 10; i++){
        a=i;
    }
    i=1;
}
```

Output

```
nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/2_Syntax_Analyzer$ lex parseTree.l
nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/2_Syntax_Analyzer$ yacc -d parseTree.y
parseTree.y:787.11-18: warning: POSIX Yacc does not support string literals [-Wyacc]
787 |         | "INC_OP"      { unaryop = 5; }
    |         ^~~~~~
parseTree.y:788.11-18: warning: POSIX Yacc does not support string literals [-Wyacc]
788 |         | "DEC_OP"      { unaryop = 6; }
    |         ^~~~~~
nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/2_Syntax_Analyzer$ cc lex.yy.c y.tab.c
nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/2_Syntax_Analyzer$ ./a.out < TestCases/forloop.c

Line:6: 'float' to 'int'

ST

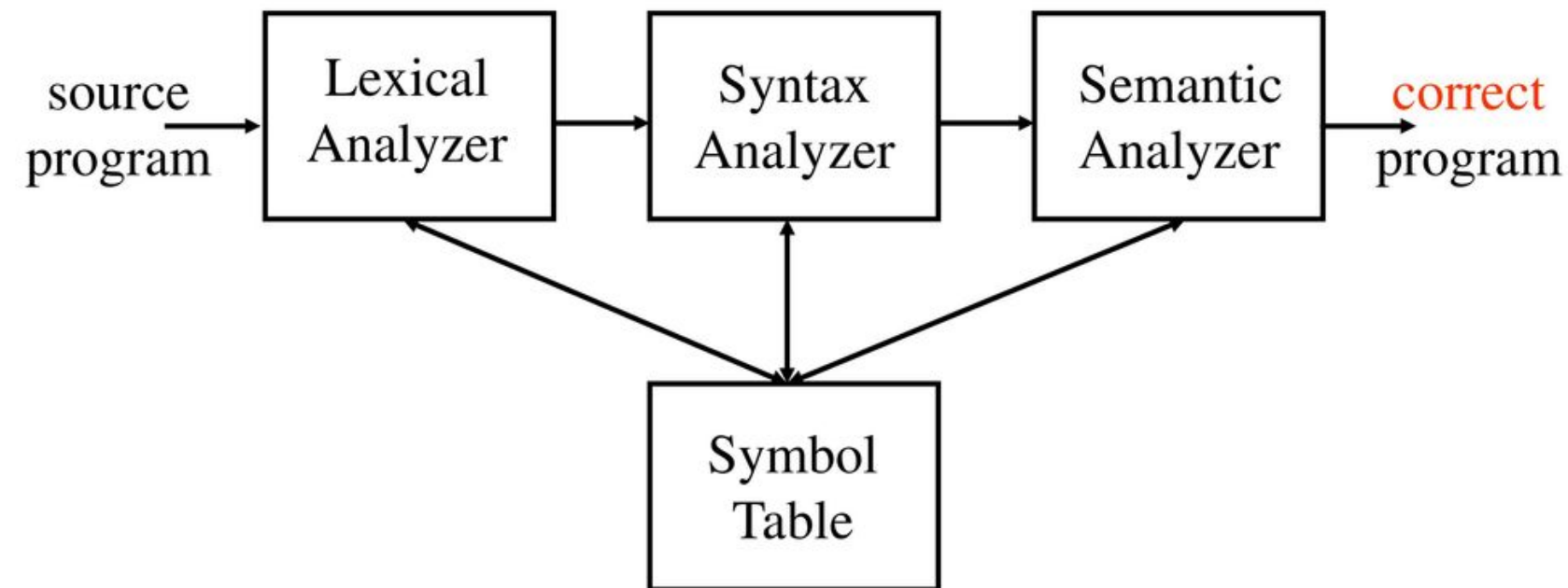
SYMBOL      NAME      TYPE      SCOPE      LINE #      VALUE
identifier  i          int       1          5          1
identifier  a          int       1          5          0
identifier  b          int       1          5          -
identifier  nume       int       1          6          3

Parse Tree

main
  stmt
    stmt =
      = for i 1
      i 0 ++ =
      < i a i
      i 10

Preorder Traversal of Parse Tree:
main ( stmt ( stmt ( = i 0 ) ( for ( ++ ( < i 10 ) i ) ( = a i ) ) ) ( = i 1 ) ) )
nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/2_Syntax_Analyzer$ |
```

SEMANTIC ANALYZER



- Performs type checking to ensure operations are performed on compatible types.
- Handles scope and namespace resolution.
- Detects and reports semantic errors such as type mismatches or undeclared variables.
- Ensures that data types are used in a way consistent with their definition.
- Keeps a check that control structures are used in a proper manner. (example: no break statement outside a loop)

Input

```
// test case to check loop statements

int main(){

    int i, a, b;
    int nume=3.45;
    for(i = 0; i < 10; i++){
        a=i;
    }
    i=1;
}
```

Output

```
nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/3_Semantic_Analyzer$ ./a.out < TestCases/forloop.c
PASSED: Semantic Phase

PRINTING SYMBOL TABLE

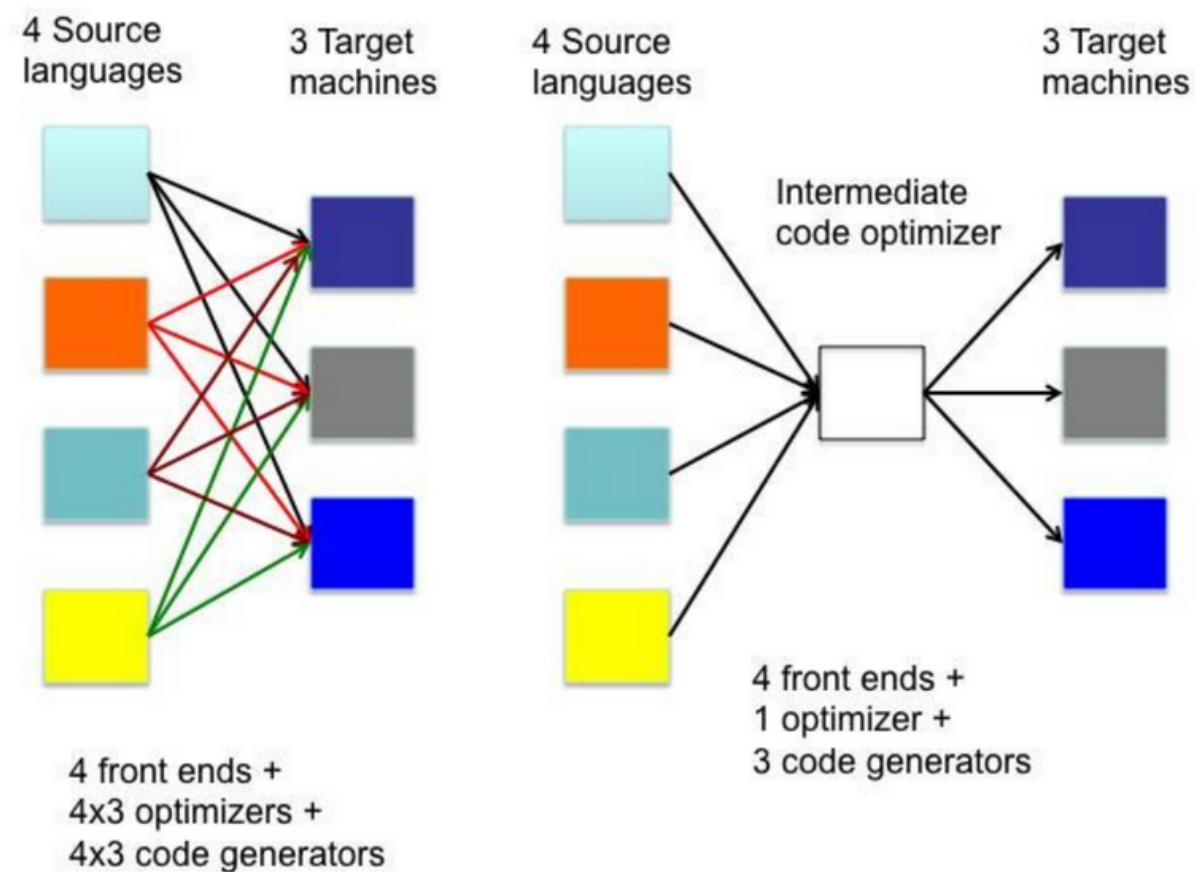
symbol name | Class | Type | Value | Line No. | Nesting Count | Count of Params |
-----|-----|-----|-----|-----|-----|-----|
a | Identifier | int |  | 5 | 99999 | -1 |
b | Identifier | int |  | 5 | 99999 | -1 |
i | Identifier | int | 1 | 5 | 99999 | -1 |
for | Keyword |  |  | 7 | 9999 | -1 |
main | Function | int |  | 3 | 9999 | -1 |
nume | Identifier | int | 3.45 | 6 | 99999 | -1 |
int | Keyword |  |  | 3 | 9999 | -1 |

PRINTING CONSTANT TABLE

constant name | constant type
-----|-----
3.45 | Floating Constant
10 | Number Constant
0 | Number Constant
1 | Number Constant

nithin@nithin1729s:~/Codes/Projects/C Compiler Phases/3_Semantic_Analyzer$ |
```

INTERMEDIATE CODE GENERATOR



- Can generate different intermediate representations like Abstract Syntax Trees (AST), Quadruples, or Direct Abstract Graphs (DAG).
- Handles complex language constructs like loops and conditionals.
- Prepares the code for optimization by simplifying and restructuring it.
- Optimizes control flow structures like loops and conditional statements.
- Generates temporaries for intermediate results.
- Handles function calls and parameter passing mechanisms.
- Converts expressions into a more manageable form for optimization

Input

```
// test case to check loop statements

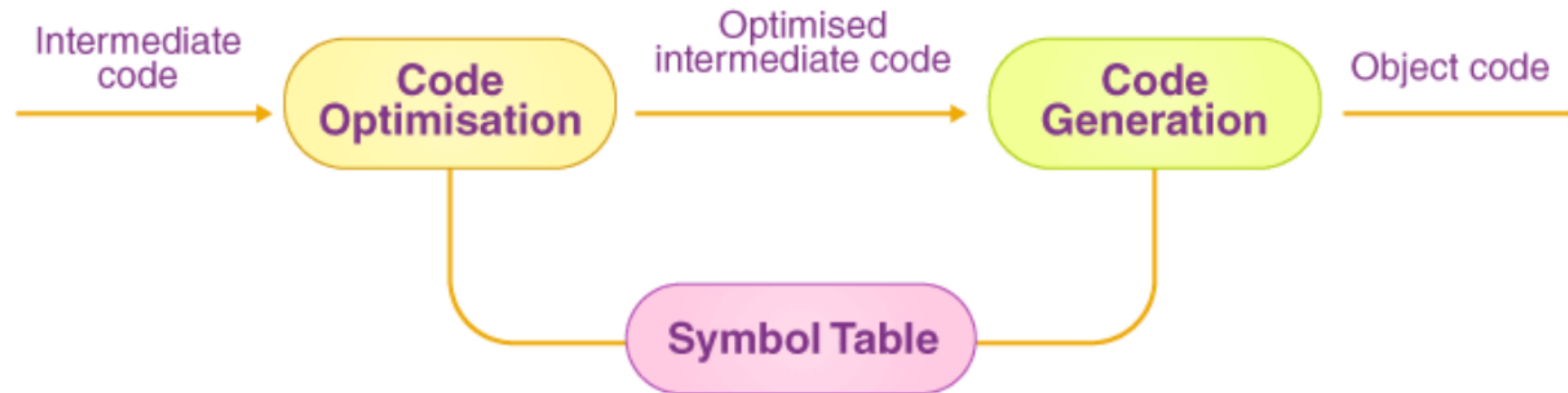
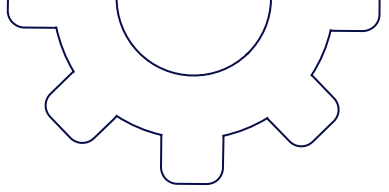
int main(){

    int i, a, b;
    int nume=3.45;
    for(i = 0; i < 10; i++){
        a=i;
    }
    i=1;
}
```

Output

```
nume = 3.45
i = 0
L0:
t0 = i < 10
ifFalse t0 goto L1
a = t0
t1 = i + 1
i = t1
goto L0
L1:
i = t1
```

CODE OPTIMIZER



- Exploits data locality for memory access optimization.
- Applies loop transformations such as loop unrolling and loop fusion.
- Utilizes profile-guided optimization for performance improvements.
- Considers instruction scheduling to minimize pipeline stalls.
- Incorporates inline expansion to reduce function call overhead.
- Implements loop vectorization for exploiting SIMD (Single Instruction, Multiple Data) instructions.
- Applies interprocedural optimizations across multiple translation units.
- Considers speculative execution and branch prediction strategies.

Input

```
// test case to check loop statements

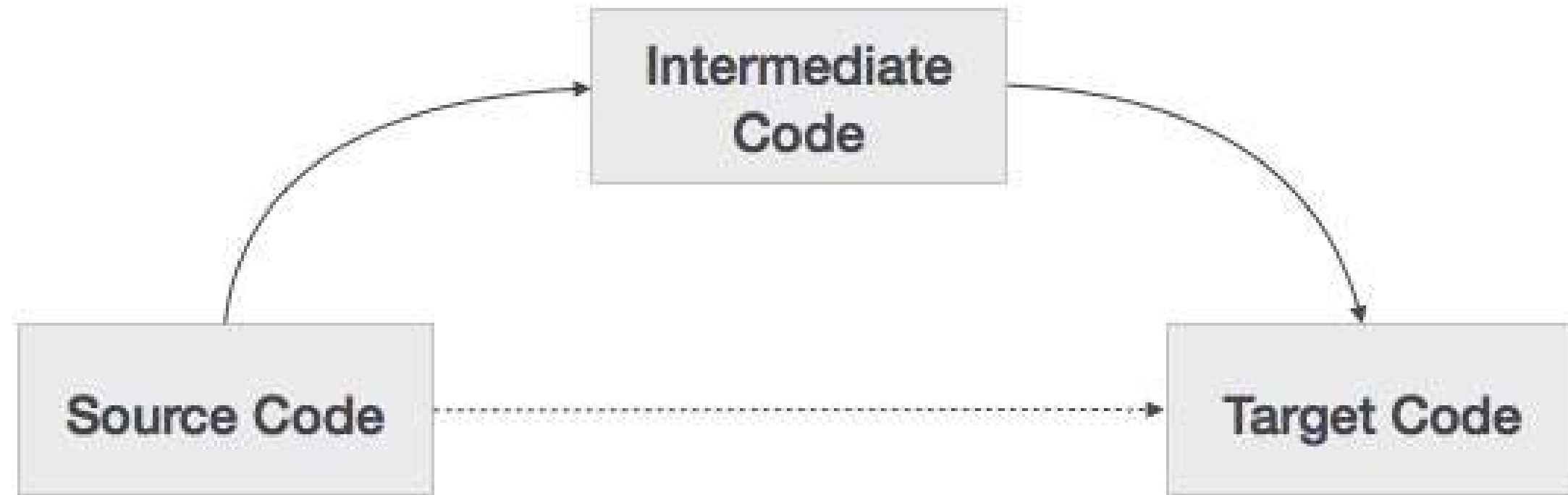
int main(){

    int i, a, b;
    int nume=3.45;
    for(i = 0; i < 10; i++){
        a=i;
    }
    i=1;
}
```

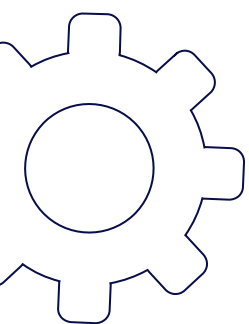
Output

```
1 nume = 3.45
2 i = 0
3 L0:
4 t0 = True
5 ifFalse t0 goto L1
6 a = t0
7 t1 = 1
8 i = t1
9 goto L0
10 L1:
11 i = t1
```

TARGET CODE GENERATOR



- Resolves addressing modes and memory layout considerations.
- Manages register allocation and spill code generation.
- Optimizes instruction selection and scheduling for the target architecture.
- Integrates platform-specific instruction sets and features.
- Handles platform-specific binary formats such as ELF or COFF.
- Manages symbol resolution and relocation for linking.
- Supports generation of position-independent code (PIC) for shared libraries.
- Integrates runtime support for exception handling and dynamic memory management.



Input

```
// test case to check loop statements

int main(){

    int i, a, b;
    int nume=3.45;
    for(i = 0; i < 10; i++){
        a=i;
    }
    i=1;
}
```

Output

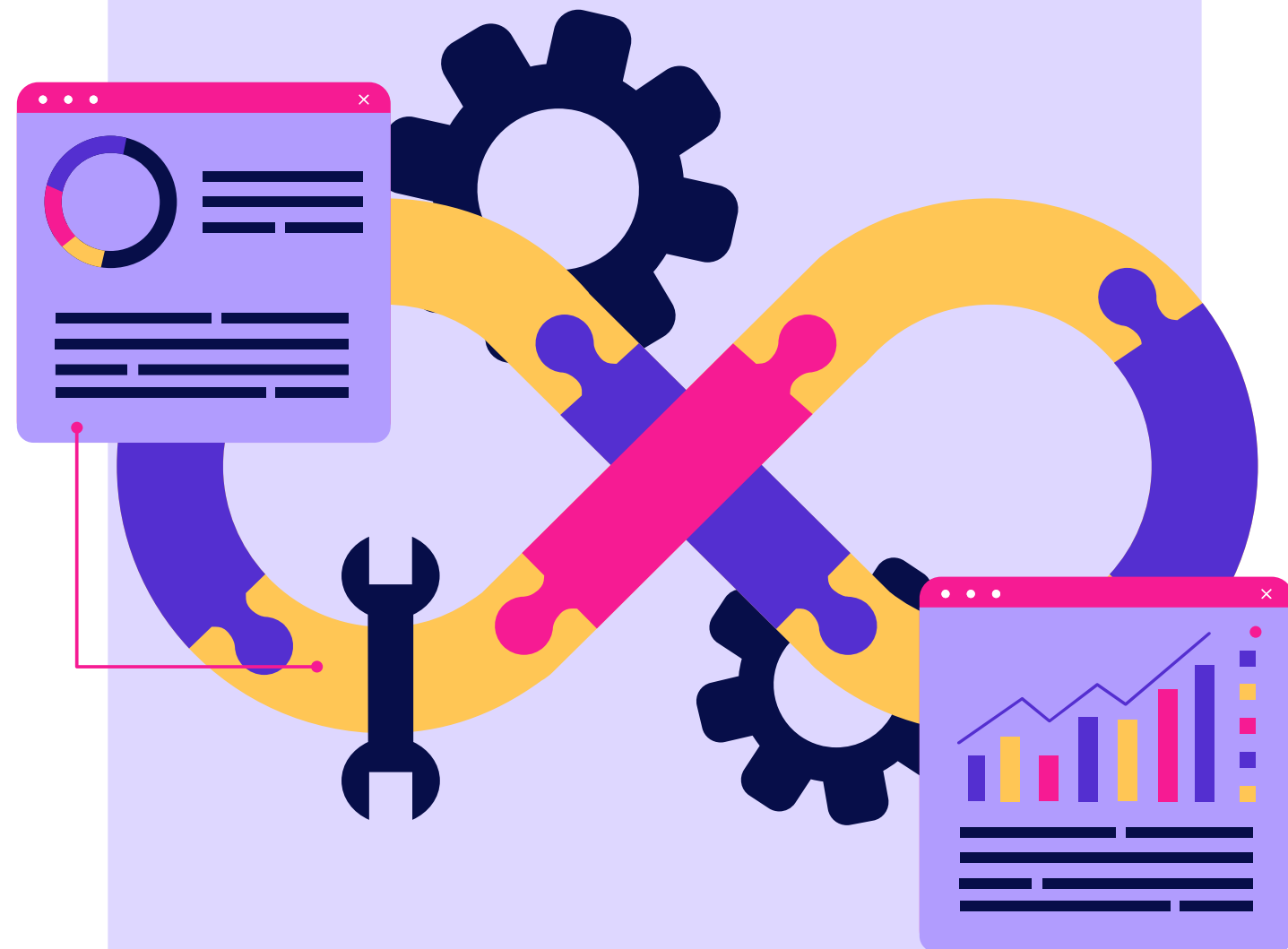
```
.text
L0:
MOV R0,=i
MOV R1,[R0]
CMP R1,#10
BGE L1
MOV R2,=i
MOV R3,[R2]
MOV R4,=t1
MOV R5,[R4]
ADD R5,#3,R1
STR R5, [R4]
MOV R6,=i
MOV R7,[R6]
MOV R8,#t1
STR R8, [R6]
B L0
L1:
MOV R9,=i
MOV R10,[R9]
MOV R11,#t1
STR R11, [R9]
SWI 0x011
.DATA
nume: .WORD 3.45
i: .WORD 0
a: .WORD t0
```

Conclusion

- With the lex and yacc tool one can create its own compiler, wherever one is required.
- It is basically procedural language compiler tools and to support object oriented one need to work on structure of C language to support object oriented which makes the compiler quite complex.
- To use lex and yacc on UNIX is easy as compared to other operating systems. Gcc is the basic compiler to generate the executable from lex and yacc compiled files.
- By studying these tools one can understand the basic structure of the compiler designed in C and go forward from it.

Future Scope

- Machine Learning Based Optimization
- Quantum Computing Compilation
- Heterogenous Computing Compilation
- High Level Synthesis
- Domain Specific Compilation



THANK YOU

