MATHEMATICAL PROGRAMMING-II

PROJECT DOCUMENT:

SECOND ORDER CONE PROGRAMMING RELAXATION OF NON-CONVEX QUDRATIC OPTIMIZATION PROBLEMS

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ABSRTACT:

Nonconvex minimax separable quadratic optimization problems with multiple separable quadratic constraints and their second-order cone programming (SOCP) relaxations. Under suitable conditions, we establish exact SOCP relaxation for minimax nonconvex separable quadratic programs. We show that various important classes of specially structured minimax quadratic optimization problems admit exact SOCP relaxations under easily verifiable conditions. These classes include some minimax extended trust-region problems, minimax uniform quadratic optimization problems, max dispersion problems, and some robust quadratic optimization problems under bounded data uncertainty. The present work shows that nonconvex minimax separable quadratic problems with quadratic constraints, which contain a hidden closed and convex epigraphical set, exhibit exact SOCP relaxations.

INTRODUCTION

Nonconvex quadratic optimization problems involving multiple quadratic constraints are a class of important and computationally hard global optimization problems that arise in many practical applications. They have been extensively studied in the literature. Especially, in recent years, a great deal of attention has been focused on studying them using their semidefinite programming (SDP) relaxation problems and second-order cone programming (SOCP) relaxation problems [2, 3, 5, 10, 21, 23, 28, 27, 30]

$$\begin{split} (P) \inf_{x \in \mathbb{R}^n} \max_{1 \leq i \leq p} \left\{ \frac{1}{2} x^T \left(U \Sigma_i U^T \right) x + a_i^T x + \alpha_i \right\} \\ \text{s.t. } \frac{1}{2} x^T \left(U \Lambda_j U^T \right) x + b_j^T x + \beta_j \leq 0, j = 1, \dots, q, \end{split}$$

where U is an

orthogonal matrix; $\Sigma i, i=1,...,p$, and $\Lambda j, j=1,...,q$, are diagonal matrices with diagonal elements given by $\sigma 1$ i ,..., σn i and $\mu 1$ j ,..., μn j , respectively, that is, $\Sigma i=diag(\sigma 1$ i ,..., σn i) and $\Lambda j=diag(\mu 1$ j ,..., μn j); ai, i=1,...,p, and bj , j=1,...,q, are n-dimensional vectors; $\alpha i, i=1,...,p$, and βj , j=1,...,q, are real numbers.

PROGRAMMING LANGUAGES/TECHNIQUES USED: -

Python with machine learning/SOCP relaxations and epigraphical sets

LITURATURE SURVEY

AUTHOR	YEAR	SURVEY
Masakazu Kojima	2015	This paper proposes a SOCP (second-order-cone programming) relaxation method, which strengthens the lift-and-project LP (linear programming) relaxation method by adding non-convex quadratic valid inequalities for the positive semidefinite cone involved in the SDP relaxation
RujunJiang(Fudan University) Duan Li(City University of Hong Kong)	2019	Second order cone constrained convex relaxations for nonconvex quadratically constrained quadratic programming
FranzRendl	2016	LP (linear programming) relaxation method by adding non-convex quadratic valid inequalities for the positive semidefinite cone involved in the SDP relaxation

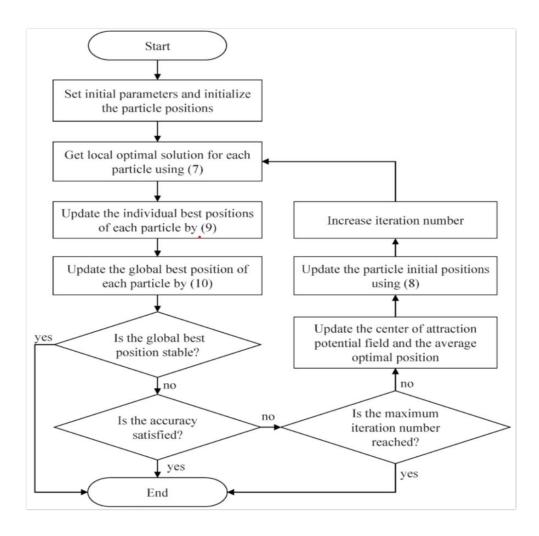
CONVEX OPTIMIZATION

 Convex optimization is a subfield of mathematical optimization that studies the problem of minimizing convex functions over convex sets. Many classes of convex optimization problems admit polynomial-time algorithms, whereas mathematical optimization is in general NP-hard.

NON-CONVEX OPTIMIZATION

• A non-convex optimization problem is any problem where the objective or any of the constraints are non-convex,. Such a problem may have multiple feasible regions and multiple locally optimal points within each region

FLOW CHART



CODE:

```
In []: import pandas as pd
import numpy as np

import matplotlib.pyplot as plt
%matplotlib inline

import scipy
from scipy.sparse import rand
import scipy.sparse as sparse

import cvxpy as cp

from sklearn.model_selection import ParameterGrid

import copy
from time import time
import random

import warnings
warnings.filterwarnings('ignore')
```

SDP

SOCP

Main part

```
In [4]: def mySDP_id(M0, n, m, density):
    X = cp.Variable((n+1, n+1), PSD=True)
              problem = cp.Problem(cp.Minimize(cp.trace(M0@X.T)), constraints)
              start = time()
              problem.solve(solver = 'MOSEK')
              end = time()
Time = end - start
              print('n =', n, '\t m =', m, '\t density =', density)
print('Time:', Time)
print(problem.value)
              val = problem.value
return n, m, density, Time, val
\label{eq:constraints} $$ = [cp.trace(Mp[i]@X.T) <= 0 for i in range(0, m)]$$ constraints <math>+= [X[0][0] == 1]$$ constraints <math>+= [G@cp.diag(X)[1:] <= h]$$
              for j in range(1, n + 1):
                  for k in range(0, j):
    if Mp[0][k, j] < -0.0000001 or Mp[0][k, j] > 0.0000001:
        constraints += [cp.norm(cp.vstack([x[k][k] - X[j][j], 2*X[k][j]])) <= X[k][k] + X[j][j]]</pre>
   constraints = [cp.trace(Mp[i]@X.T) \leftarrow 0 \text{ for } i \text{ in } range(0, m)]
                 constraints += [X[0][0] == 1]
constraints += [G@cp.diag(X)[1:] <= h]
                 problem = cp.Problem(cp.Minimize(cp.trace(M0@X.T)), \ constraints)
```

```
start = time()
problem.solve(solver = 'MOSEK')
end = time()
Time = end - start
print('n =', n, '\t m =', m, '\t density =', density)
print('Time:', Time)
print(problem.value)
val = problem.value
return n, m, density, Time, val
```

```
In [12]: N, M, Density = [50, 100], [50], [0.1]
                 fResults = []
for n in N:
                        G = np.random.uniform(0, 100, size=(n+3, n))
h = G @ np.random.uniform(0, 100, size=(n,))
                        for density in Density:
    M0 = get_matrix(n, density)
for m in M:
    Results = []
    Mp = []
```

```
n = 50 m = 50
                         density = 0.1
Time: 1.807164192199707
-28567.588126761653
n = 50 \quad m = 50
                         density = 0.1
Time: 0.3102538585662842
-28567.585759113666
                          density = 0.1
                m = 50
n = 100
Time: 50.20961308479309
-193564.1780348802
                m = 50 density = 0.1
70571899
n = 100
Time: 6.866625070571899
-193564.17182150309
     n m density Time_SDP value_SDP Time_SOCP
                                                      value_SOCP
0 \quad 50 \quad 50 \quad \quad 0.1 \quad 1.807164 \quad -28567.588127 \quad \quad 0.310254 \quad -28567.585759
1 100 50 0.1 50.209613 -193564.178035 6.866625 -193564.171822
```

CONCLUSION:

we will established exact SOCP relaxations for nonconvex minimax separable quadratic optimization problems with multiple separable quadratic constraints under an epigraphical condition. We exploited hidden convexity in the form of a convex epigraphical set to achieve our results. We have also provided various classes of minimax problems for which our results hold under easily verifiable conditions

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