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### Problem Solving - 2

- ① Draw a space-time diagram for a six-segment Pipeline showing the time it takes to process 10 tasks.

soln

Space time diagram:

★ The behavior of Pipeline can be defined through space-time diagram. It shows the segment utilization as a function of time.

★ The horizontal axis has time in clock cycles and vertical axis has segment number.

Number of segment pipeline  $k = 6$

Number of task  $n = 10$ .



Number of clock cycles requires to complete  $n$  tasks and  $k$  segment Pipeline is as shown

$$\Rightarrow k+n-1$$

$$\Rightarrow 6+10-1$$

$$\Rightarrow 15$$

$$\boxed{\text{Total clock cycles} = 15}$$

Space-time diagram for the 6 segment Pipeline

segments	TASK														
	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
segment 1	T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	T <sub>6</sub>	T <sub>7</sub>	T <sub>8</sub>	T <sub>9</sub>	T <sub>10</sub>					
segment 2		T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	T <sub>6</sub>	T <sub>7</sub>	T <sub>8</sub>	T <sub>9</sub>	T <sub>10</sub>				
3			T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	T <sub>6</sub>	T <sub>7</sub>	T <sub>8</sub>	T <sub>9</sub>	T <sub>10</sub>			
4				T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	T <sub>6</sub>	T <sub>7</sub>	T <sub>8</sub>	T <sub>9</sub>	T <sub>10</sub>		
5					T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	T <sub>6</sub>	T <sub>7</sub>	T <sub>8</sub>	T <sub>9</sub>	T <sub>10</sub>	
6						T <sub>1</sub>	T <sub>2</sub>	T <sub>3</sub>	T <sub>4</sub>	T <sub>5</sub>	T <sub>6</sub>	T <sub>7</sub>	T <sub>8</sub>	T <sub>9</sub>	T <sub>10</sub>

This diagram shows 10 Tasks T<sub>1</sub> through T<sub>10</sub> execute in 6 segments.



② Determine the number of clock cycles that it takes in an eight segment Pipeline with 300 tasks.

Soln

→ Let there be 'n' tasks to be performed in the 'pipelined processor'.

→ The first instruction will be taking

'k' segment Pipeline cycles to exit out of the 'pipeline' but the other 'n-1' instructions will take only '1' cycle each. i.e., a total of 'n-1' cycles.

→ So, to perform 'n' instructions in a 'pipelined processor' time taken is

$k + (n-1)$  cycles.

→ So, in our case no. of clock cycles

$$\Rightarrow 8 + (300 - 1)$$

$T_p \Rightarrow 307$ clock cycles.
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Q) Non-pipe.

Solution:

The speed up ratio (S):

It is defined as the speedup of a Pipeline Processing with respect to the equivalent non-pipeline processing

Formula:-

$$S = \frac{nt_n}{(k+n-1)t_p}$$

Number of tasks  $n=200$

For Non pipeline:

Time taken by non-pipeline to process a task  $t_n = 60 \text{ ns}$ .

Total time taken by non-pipeline to Process 200 task  $= nt_n$ .

$$\Rightarrow 200 \times 60$$

$$t_n \Rightarrow 12000 \text{ ns.}$$



For Pipeline:

No. of segments pipeline,  $k=9$

Time period of 1 clock cycle  $t_p = 20\text{ns}$ .

Total time requires to complete "n" task in k segments pipeline with  $t_p$  clock cycle time.

$$\Rightarrow (k+n-1)t_p$$

$$\Rightarrow (9+200-1)20$$

$$\Rightarrow (208)20$$

$$\Rightarrow 4160 t_p$$

Speed up ratio:

When total time taken by the pipeline to process 200 tasks is divided by the total time required to complete 'n' tasks in k segment pipeline with  $t_p$  clock cycle time then speed up ratio is obtain

$$S = 200/4160$$

$$\therefore S = 2.884 \approx 3$$



(A)

Solution:-

Time for operands to be read from memory into registers.

$$\therefore R_1 \text{ and } R_2 = 30\text{ns.}$$

Time for the signal to propagate through multiplier = 35ns.

Time for the transfer data to  $R_3 = 5\text{ns.}$

Time to add 2 numbers and store to

$$\underline{R_5 = 15\text{ns}}$$

The operation performing in this case can be grouped in 3 segments which is shown as.

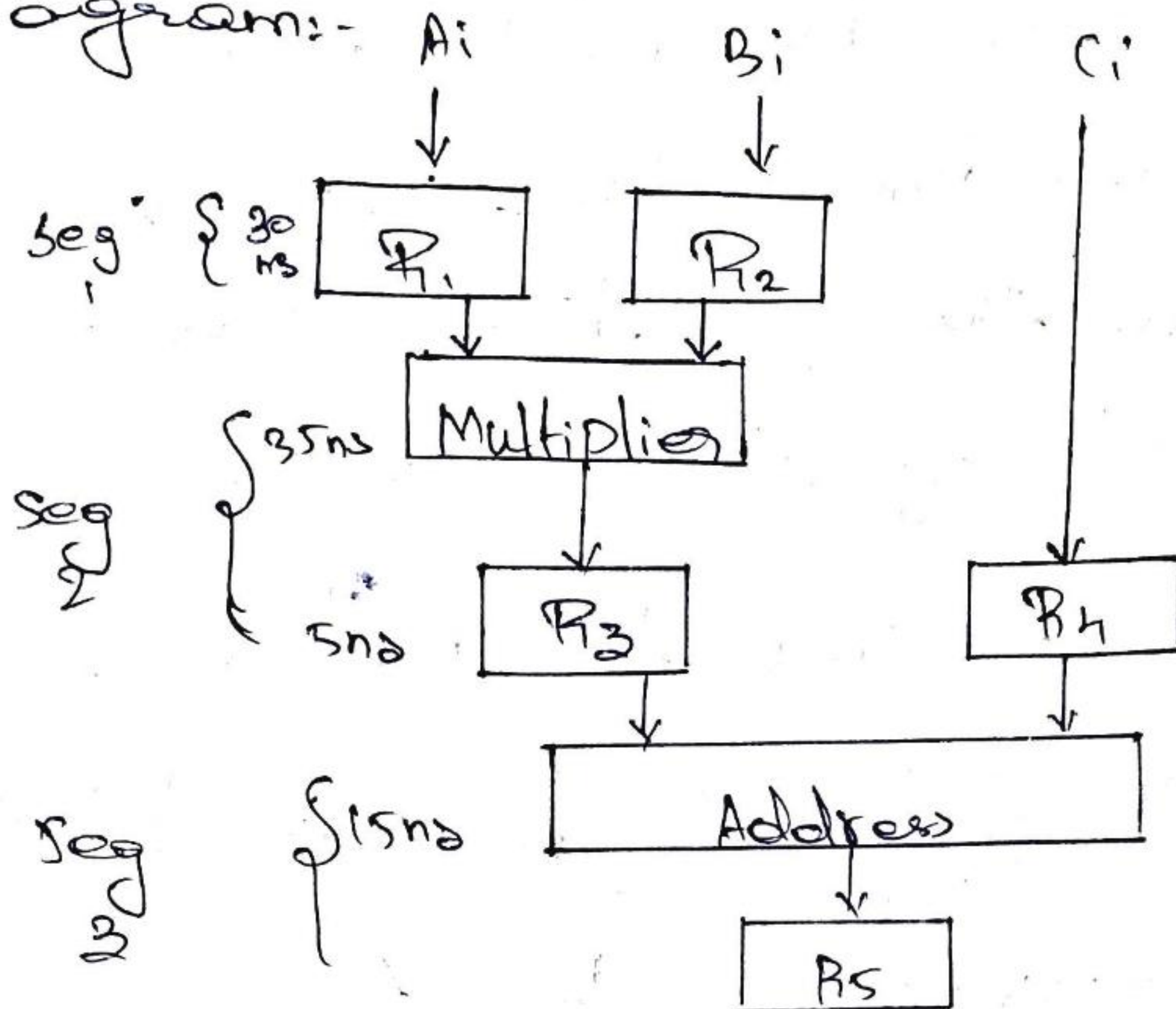
Segment 1:  $R_1 \leftarrow A_i \quad R_2 \leftarrow B_i$

Segment 2:  $R_3 \leftarrow R_1 * R_2 \quad R_4 \leftarrow C_i$

Segment 3:  $R_5 \leftarrow R_3 + R_4$



Diagram:-



(A)

The minimum clock cycle time that can be used is max propagation time of the 3 segments.

segment 1: 30 ns

segment 2: 35 + 5 = 40 ns

segment 3: 15 ns

Hence the maximum propagation time is segment 2 so clock cycle time is

$$t_P = 40 \text{ ns}$$



⑥ For non-Pipeline  $P_3$  and  $P_4$  is removed.

In this case time taken to multiply and add the operands... will be summation of 30, 35 and 15.

Time taken to add the operands without use of the Pipeline is as follows.

$$\Rightarrow 30 + 35 + 15$$

$$\Rightarrow \underline{80}$$

∴ Hence time taken to add the operands without the use of the Pipeline is

$$T_n = 80$$

⑦

Speed up ratio:

$$S = \frac{nt_n}{(K + (n-1))t_p}$$

No. of segment Pipeline,  $K = 3$

No. of task  $n = 15$

$$S = 15 \times 80 / (3 + (15-1))40$$

$$S = \frac{1200}{680}$$



$\therefore$  The time taken for the speed up of the Pipeline for 15 tasks is.

$$S = 1.76 \text{ ns}$$

For tasks,  $n = 150$

$$S = \frac{150 \times 80}{(3 + 150 - 1) \times 40} \Rightarrow \frac{12000}{6080} = 1.97 \text{ ns}$$

$\therefore$  Pipeline for 150 task is

$$\boxed{S = 1.97 \text{ ns} \approx 2}$$

D)

Maximum speed up ratio:-

$$\begin{aligned} S_{\max} &= t_n / t_p \\ &= 80 / 40 \end{aligned}$$

$$\boxed{\therefore S_{\max} = 2}$$

$\therefore$  Hence, the maximum speed up ratio is 2.

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