## Department of CSE, RV College of Engineering – Bengaluru FAFL – 18CS52

## Video based seminar topics

| Sl<br>No | USN        | Name             | Topic  |
|----------|------------|------------------|--|
| 1        | 1RV20CS065 | KARTHIK S HALLAD | Show that the CFG with given productions is ambiguous and find an equivalent unambiguous grammar.  a. $S \rightarrow a \mid Sa \mid bSS \mid SSb \mid SbS$ b. $S \rightarrow SS \mid a \mid b$ c. $S \rightarrow ABA$ $A \rightarrow aA \mid \epsilon$ $B \rightarrow bB \mid \epsilon$ d. $S \rightarrow aSb \mid aaSb \mid \epsilon$ e. $S \rightarrow aSb \mid abS \mid \epsilon$ |
| 2        | 1RV20CS066 | KASHISH NAYAN    | For $\Sigma = \{a, b\}$ , construct DFA for the language of all strings containing both aba and bab as substrings.   |
| 3        | 1RV20CS067 | KAUSHIK B A      | For each of the languages given, use the pumping lemma to show that it cannot be accepted by an FA.  a. $L = \{a^nba^{2n} \mid n \ge 0\}$ b. $L = \{a^ib^j \mid a^k \mid k > i + j\}$ c. $L = \{a^ib^j \mid j = i \text{ or } j = 2i\}$ d. $L = \{a^ib^j \mid j \text{ is a multiple of } i\}$   |
| 4        | 1RV20CS068 | KEERTHI P        | Consider the two regular expressions $r = a^* + b^* \qquad s = ab^* + ba^* + b^*a + (a^*b)^*$ a. Find a string corresponding to r but not to s. b. Find a string corresponding to s but not to r.  |

|   |            |                                 | For each of the FAs pictured in the below figure, give a simple verbal description of the language it accepts.  |
|---|------------|---------------------------------|---|
| 5 | 1RV20CS069 | KESANAPALLI<br>LAKSHMI PRIYANKA |   |
|   |            |                                 | $\begin{array}{c c} & & & \\ & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & &$ |

|   |            |                            | For each of the FAs pictured in the below figure, give a simple verbal description of the language it accepts.  |
|---|------------|----------------------------|---|
| 6 | 1RV20CS070 | KOKKALLA VAMSHI<br>KRISHNA | $\begin{array}{cccccccccccccccccccccccccccccccccccc$  |
| 7 | 1RV20CS072 | KRITADNYA KALING           | For each of the languages given, use the pumping lemma to show that it cannot be accepted by an FA.  a. $L = \{x \in \{a, b\} * \mid n_a(x) < 2n_b(x)\}$ b. $L = \{x \in \{a, b\} * \mid \text{no prefix of } x \text{ has more b's than a's}\}$ c. $L = \{a^{n^3} \mid n \ge 1\}$ d. $L = \{ww \mid w \in \{a, b\} * \}$                   |
| 8 | 1RV20CS073 | KRUTHIKA P                 | For each of the case below, write what language (a subset of $\{a, b\}^*$ ) is generated by the context-free grammar with the indicated productions.  a. $S \rightarrow aS \mid bS \mid \epsilon$ b. $S \rightarrow SS \mid bS \mid a$ c. $S \rightarrow SaS \mid b$ d. $S \rightarrow SaS \mid b \mid \epsilon$                            |
| 9 | 1RV20CS074 | KUNAL SATISH<br>MAHAJAN    | Find context-free grammars generating each of the languages below.<br>a. $L = \{w \in \{a, b\} * : n_a(w) \neq n_b(w)\}$<br>b. $L = \{w \in \{a, b\} * : n_a(v) \geq n_b(v), \text{ where } v \text{ is any prefix of } w\}$<br>c. $L = \{w \in \{a, b\} * : n_a(w) = 2n_b(w) + 1\}$<br>d. $L = \{w \in \{a, b\} * : n_a(w) = n_b(w) + 2\}$ |

|    |            |              | Move Number   | State                                  | Input  | Stack Symbol  | Move(s)  |
|----|------------|--------------|---|--|--|---|--|
|    |            |              | 1   | $q_0$                                  | а  | $Z_0$   | $(q_0, xZ_0), (q_1, aZ_0)$   |
|    |            |              | 2   | $q_0$                                  | b  | $Z_0$   | $(q_0, xZ_0), (q_1, bZ_0)$   |
|    |            |              | 3   | $q_0$                                  | a  | x   | $(q_0, xx), (q_1, ax)$   |
|    |            |              | 4   | $q_0$                                  | b  | x   | $(q_0, xx)(q_1, bx)$   |
|    |            | KUSHAJ KUMAR | 5   | $q_1$                                  | a  | a   | $(q_1, a)$   |
|    |            |              | 6   | $q_1$                                  | b  | b   | $(q_1, b)$   |
|    |            |              | 7   | $q_1$                                  | a  | b   | $(q_1,b),(q_2,\Lambda)$  |
| 10 | 1RV20CS075 |              | 8   | $q_1$                                  | b  | a   | $(q_1, a), (q_2, \Lambda)$   |
|    |            |              | 9   | $q_2$                                  | a  | x   | $(q_2,\Lambda)$  |
|    |            |              | 10  | $q_2$                                  | b  | x   | $(q_2,\Lambda)$  |
|    |            |              | 11  | $q_2$                                  | Λ  | $Z_0$   | $(q_3, Z_0)$   |
|    |            |              | (all  | other co                               | mbination  | ns)   | none   |
|    |            |              | symbol read. From<br>onto the stack the<br>option of ignoring | $q_0$ it all symbol the input to reach | so has the it has just symboth the according to the according to the according to the according the according the according to the according the according to t | ne choice of ente<br>st read. In state q<br>I that is read and<br>epting state it m | e stack for each inpuring $q_1$ , by pushing $q_1$ there is always the leaving the stack ust eventually be abl |
|    | 1RV20CS076 | KUSHI P      |   |  |  |   | in which every a (if there   |

|    |               |                 |  |   |   | given for a PDA ach case the langua                            |                                   |
|----|---------------|-----------------|--|---|---|--|-----------------------------------|
|    |               |                 | Move Number  | State   | Input   | Stack Symbol   | Move(s)                           |
|    |               |                 | 1  | $q_0$   | а   | $Z_0$  | $(q_1, a Z_0)$                    |
|    |               |                 | 2  | $q_0$   | b   | $Z_0$  | $(q_1, bZ_0)$                     |
|    |               |                 | 3  | $q_1$   | a   | a  | $(q_1, a), (q_2, a)$              |
|    |               |                 | 4  | $q_1$   | b   | a  | $(q_1, a)$                        |
|    |               |                 | 5  | $q_1$   | a   | b  | $(q_1, b)$                        |
|    |               |                 | 6  | $q_1$   | b   | b  | $(q_1, b), (q_2, b)$              |
| 12 | 1RV20CS077    | M M ARJUN       |  | (all other co   | ombinations)                                  |  | none                              |
| 12 | 110 20 250 77 | W W A A C I V   | Move Number  | State   | Input   | Stack Symbol   | Move(s)                           |
|    |               |                 | 1  | $q_0$   | a   | $Z_0$  | $(q_0, XZ_0)$                     |
|    |               |                 | 2  | $q_0$   | b   | $Z_0$  | $(q_0, XZ_0)$                     |
|    |               | 3               | $q_0$  | a   | X   | $(q_0, XX)$  |                                   |
|    |               |                 | 4  | $q_0$   | b   | X  | $(q_0, XX)$                       |
|    |               |                 | 5  | $q_0$   | c   | X  | $(q_1, X)$                        |
|    |               |                 | 6  | $q_0$   | c   | $Z_0$  | $(q_1, Z_0)$                      |
|    |               |                 | 7  | $q_1$   | a   | X  | $(q_1, \Lambda)$                  |
|    |               |                 | 8  | $q_1$   | b   | X  | $(q_1, \Lambda)$                  |
|    |               |                 | 9  | $q_1$   | Λ   | $Z_0$  | $(q_2, Z_0)$                      |
|    |               |                 |  | (all other  | combination                                   | s)   | none                              |
| 13 | 1RV20CS078    | M S RAGHAVENDRA | For the grammar G<br>a. $S \rightarrow aSa$ , $S \rightarrow b$<br>b. $S \rightarrow abB$ , $A \rightarrow$<br>c. $S \rightarrow AB \lambda$ , $A -$<br>d. $S \rightarrow aaB$ , $A \rightarrow$<br>that w is not in | oSb, $S \rightarrow \lambda$<br>aaBb, $B \rightarrow \alpha$<br>aB, $B \rightarrow \beta$<br>bBb $ \lambda$ , $B - \beta$ | , show a den  bbAa, A − Sb, draw de  Aa, draw | rivation tree for w = $\lambda$ , draw derivation tree for the | n tree for w = ab<br>e w = aabbbb |

| 14 | 1RV20CS079 | MADHVESH ACHARYA<br>M  | For each of the case below, write what language (a subset of $\{a,b\}^*$ ) is generated by the context-free grammar with the indicated productions.  a. $S \rightarrow T T$ $T \rightarrow aT \mid T \mid a \mid b$ b. $S \rightarrow aSa \mid bSb \mid aAb \mid bAa$ $A \rightarrow aAa \mid bAb \mid a \mid b \mid \epsilon$ c. $S \rightarrow aT \mid bT \mid \epsilon$ $T \rightarrow aS \mid bS$ d. $S \rightarrow aT \mid bT$ $T \rightarrow aS \mid bS \mid \epsilon$ |
|----|------------|------------------------|--|
| 15 | 1RV20CS080 | MALAVIKA<br>HARIPRASAD | Give a context-free grammar for generating all properly nested parentheses.  A properly nested parenthesis structures are the ones involving two kinds of parentheses, say () and []: example ([]), ([[]]) [()], but not ([)] or ((]].   |
| 16 | 1RV20CS081 | MANOJ M                | Find context-free grammars for the following languages:  a. $L = a^n b^n$ , n is even.  b. $L = a^n b^n$ , n is odd.   |
| 17 | 1RV20CS082 | MANOJKUMAR<br>BELLATTI | In each case below, find a context-free grammar with no $\epsilon$ -productions that generates the same language, except possibly for $\epsilon$ , as the given CFG.  f. $S \rightarrow AB \mid \epsilon$ $A \rightarrow aASb \mid a$ $B \rightarrow bS$ g. $S \rightarrow AB \mid ABC$ $A \rightarrow BA \mid BC \mid \epsilon \mid a$ $B \rightarrow AC \mid CB \mid \epsilon \mid b$ $C \rightarrow BC \mid AB \mid A \mid c$   |
| 18 | 1RV20CS083 | MAYA S RAO             | <ul> <li>Find a regular expression corresponding to each of the following subsets of {a, b}*.</li> <li>1. The language of all strings in which the number of a's is even and the number of b's is odd.</li> <li>2. The language of all strings in which both the number of a's and the number of b's are odd.</li> </ul>   |

|    |            |                    | For the NFAs below, find a regular expression corresponding to the language it accepts.  |
|----|------------|--------------------|--|
| 19 | 1RV20CS084 | MAYUR ANKLEKAR     |  |
| 20 | 1RV20CS085 | MAYUR S CHITTARAGI | Find context-free grammars generating each of the languages below.<br>a. $L = \{a^nb^mc^k, k = n + m\}$<br>b. $L = \{a^nb^mc^k, k =  n - m \}$   |
| 21 | 1RV20CS086 | MD ZEAUL HAQUE     | In each case below, find a context-free grammar with no $\epsilon$ -productions that generates the same language, except possibly for $\epsilon$ , as the given CFG.  a. $S \rightarrow AB \mid \epsilon$ $A \rightarrow aASb \mid a$ $B \rightarrow bS$ b. $S \rightarrow AB \mid ABC$ $A \rightarrow BA \mid BC \mid \epsilon \mid a$ $B \rightarrow AC \mid CB \mid \epsilon \mid b$ $C \rightarrow BC \mid AB \mid A \mid c$ |
| 22 | 1RV20CS087 | MEETH J DAVDA      | Suppose $M = (Q, \Sigma, q_0, A, \delta)$ is an FA, q is an element of Q, and x and y are strings in $\Sigma^*$ . Using structural induction on y, prove the formula $\delta^*(q, xy) = \delta^*(\delta^*(q, x), y) $  |

| 23 | 1RV20CS088 | MEGHANSH MUNDRA                                      | For the FAs pictured below, use the minimization algorithm (Equivalence) to find a minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)   |
|----|------------|--|--|
| 24 | 1RV20CS089 | MINAL R D  | <ul> <li>Draw a FA accepting the language generated by the CFG having the given productions.</li> <li>a. S → aA   bC</li> <li>A → aS   bB</li> <li>B → aC   bA</li> <li>C →aB   bS   ε</li> <li>b. S → bS   aA   ε</li> <li>A → aA   bB   b</li> <li>B → bS</li> <li>c. S → abA   bB   aba</li> <li>A → b   aB   bA</li> <li>B → aB   aA</li> </ul>  |
| 25 | 1RV20CS090 | MOHAMMED KHALID<br>MOHAMMED<br>MINHAJUDDIN<br>ANSARI | Let r and s be arbitrary regular expressions over the alphabet $\Sigma$ . In each case below, find a simpler equivalent regular expression.  a. $r(r^* r + r^*) + r^*$ b. $(r + \varepsilon)^*$ c. $(r + s)^* rs(r + s)^* + s^*r^*$  |
| 26 | 1RV20CS091 | NAMAN N KARANTH                                      | Each of the following grammars, though not regular, generates a regular language. In each case, find a regular grammar generating the language:  a. $S \rightarrow AAS \mid ab \mid aab$ $A \rightarrow ab \mid ba \mid \epsilon$ b. $S \rightarrow AB$ $A \rightarrow aAa \mid bAb \mid a \mid b$ $B \rightarrow aB \mid bB \mid \epsilon$ c. $S \rightarrow AA \mid B$ $A \rightarrow AAA \mid Ab \mid bA \mid a$ $B \rightarrow bB \mid \epsilon$ |

| 27 | 1RV20CS092 | NAVANIKA J REDDY    | For the FAs pictured below, use the minimization algorithm (Equivalence) to find a minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)   |
|----|------------|---------------------|--|
| 28 | 1RV20CS093 | NAVEEN B TELI       | In each case, given the context-free grammar G, find a CFG G with no $\epsilon$ -productions and no unit productions that generates the language $L(G) = \{ \epsilon \}$ .  h. $S \to ABA$ $A \to aA \mid \epsilon$ $B \to bB \mid \epsilon$ i. $S \to aSa \mid bSb \mid \epsilon$ $A \to aBb \mid bBa$ $B \to aB \mid bB \mid \epsilon$ j. $S \to A \mid B \mid C$ $A \to aAa \mid B$ $B \to bB \mid bb$ $C \to aCaa \mid D$ $D \to baD \mid abD \mid aa$ |
| 29 | 1RV20CS094 | NEHA N              | For the following regular expressions, draw an NFA accepting the corresponding language, so that there is a recognizable correspondence between the regular expression and the transition diagram.  e. (a + b)(ab)*(abb)* f. (a + b)*(abba* + (ab)*ba) g. (a*bb)* + bb*a*  |
| 30 | 1RV20CS095 | NEHASHRI POOJAR S V | Construct NPDAs that accept the following languages:  a. $L = \{w : n_a(w) = n_b(w) + 1\}$ b. $L = \{w : n_a(w) = 2n_b(w)\}$   |

| 31 | 1RV20CS096 | NIKHIL BENNUR              | For the FAs pictured below, use the minimization algorithm (Table filling) to find a minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)  |
|----|------------|----------------------------|---|
| 32 | 1RV20CS097 | NIKHIL TAVANAPPA<br>BELAVI | For the following regular expressions, draw an NFA accepting the corresponding language, so that there is a recognizable correspondence between the regular expression and the transition diagram.  h. (b + bba)*a i. (a + b)*(abb + ababa)(a + b)* |
| 33 | 1RV20CS098 | NIMISHA DEY                | For the NFAs below, find a regular expression corresponding to the language it accepts.   |

| 34 | 1RV20CS099 | NISHAL H N                 | <ul> <li>Find a CFG generating the given language.</li> <li>a. The set of odd-length strings in {a, b}* with middle symbol a.</li> <li>b. The set of even-length strings in {a, b}* with the two middle symbols equal.</li> <li>c. The set of odd-length strings in {a, b}* whose first, middle, and last symbols are all the same.</li> </ul> |
|----|------------|----------------------------|--|
| 35 | 1RV20CS100 | NISHITH S SHETTY           | Find a regular expression corresponding to each of the following subsets of {a, b}*.  a. The language of all strings containing exactly two a's.  b. The language of all strings containing at least two a's.  |
| 36 | 1RV20CS101 | NITHISH S                  | Construct NPDAs that accept the following languages:<br>a. $L = \{w : n_a(w) + n_b(w) = n_c(w)\}$<br>b. $L = \{w : 2n_a(w) \le n_b(w) \le 3n_a(w)\}$<br>c. $L = \{w : n_a(w) < n_b(w)\}$   |
| 37 | 1RV20CS102 | NITIN SINGH                | Find context-free grammars for the following language:<br>$L = a^n b^n$ , n is a multiple of three.  |
| 38 | 1RV20CS103 | P BHUVANESHWAR             | Find a regular expression corresponding to each of the following subsets of {a, b}*.  a. The language of all strings that do not end with ab.  b. The language of all strings that begin or end with aa or bb.   |
| 39 | 1RV20CS104 | PARI RAHEJA                | For $\Sigma = \{a, b\}$ , construct DFA for the language of all strings containing no more than one occurrence of the string aa. (The string aaa contains two occurrences of aa.)  |
| 40 | 1RV20CS105 | PAVAN R                    | Draw a transition diagram for an FA that accepts the string <b>abaa</b> and no other strings.  |
| 41 | 1RV20CS106 | PEDDISETTY VARAD<br>NITHIN | Find a regular expression corresponding to each of the following subsets of {a, b}*.  a. The language of all strings not containing the substring aa.  b. The language of all strings in which the number of a's is even.  |

| 42 | 1RV20CS107 | PETA SIVA DEEKSHITH<br>REDDY | In each case, given the context-free grammar G, find a CFG G with no $\epsilon$ -productions and no unit productions that generates the language $L(G) - \{ \epsilon \}$ .  a. $S \to ABA$ $A \to aA \mid \epsilon$ $B \to bB \mid \epsilon$ b. $S \to aSa \mid bSb \mid \epsilon$ $A \to aBb \mid bBa$ $B \to aB \mid bB \mid \epsilon$ c. $S \to A \mid B \mid C$ $A \to aAa \mid B$ $B \to bB \mid bb$ $C \to aCaa \mid D$ $D \to baD \mid abD \mid aa$ |
|----|------------|------------------------------|--|
| 43 | 1RV20CS108 | PRADHAAN R<br>KEDLAYA        | Find a regular grammar generating the language L(M), where M is the FA shown below:  |

|    |            |              | For the NFAs below, find a regular expression corresponding to the language it accepts.   |
|----|------------|--------------|---|
| 44 | 1RV20CS109 | PRADHAN A N  |   |
| 45 | 1RV20CS110 | PRAGAM JAIN  | Consider the two regular expressions $r = a^* + b^* \qquad s = ab^* + ba^* + b^*a + (a^*b)^*$ a. Find a string corresponding to both r and s. b. Find a string in {a, b}* corresponding to neither r nor s.   |
| 46 | 1RV20CS111 | PRAGATHI B C | Find a regular expression corresponding to each of the following subsets of {a, b}*.  1. The language of all strings containing no more than one occurrence of the string aa.  (The string aaa should be viewed as containing two occurrences of aa.)  2. The language of all strings in which every a is followed immediately by bb. |
| 47 | 1RV20CS112 | PRAJWAL C R  | For the NFAs below, find a regular expression corresponding to the language it accepts. $ \begin{array}{cccccccccccccccccccccccccccccccccc$   |
| 48 | 1RV20CS113 | PRAJWAL P    | Find a regular expression corresponding to each of the following subsets of {a, b}*.  a. The language of all strings containing both bb and aba as substrings.  b. The language of all strings not containing the substring aaa.  |

| 49 | 1RV20CS114 | PRAJWAL T S                | Find context-free grammars generating each of the languages below.<br>c. $L = \{a^n b^m c^k : n = m \text{ or } m \le k,  n \ge 0,  m \ge 0,  k \ge 0 \}$<br>d. $L = \{a^n b^m c^k : n = m \text{ or } m \ne k,  n \ge 0,  m \ge 0,  k \ge 0 \}$<br>For the FAs pictured below, use the minimization algorithm (Table filling) to find a |
|----|------------|----------------------------|--|
| 50 | 1RV20CS115 | PRANAMYA MADY              | minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)  |
| 51 | 1RV20CS116 | PRANSHU PRAKHAR<br>SINGH   | For $\Sigma = \{a, b\}$ , construct DFA for the language of all strings containing both bb and aba as substrings.  |
| 52 | 1RV20CS117 | PRASAD PATIL               | Find context-free grammars generating each of the languages below.<br>a. $L = \{a^nb^mc^k, k = n + 2m\}$<br>b. $L = \{a^nb^mc^k, k \neq n + m\}$   |
| 53 | 1RV20CS118 | PRASANNA SURESH<br>NAIK    | For the FAs pictured below, use the minimization algorithm (Table filling) to find a minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)   |
| 54 | 1RV20CS119 | PRASHANTH REDDY<br>GUNDALA | Find a regular expression corresponding to each of the following subsets of {a, b}*.  a. The language of all strings not containing the substring bba.  b. The language of all strings containing both bab and aba as substrings.  |

| 55 | 1RV20CS120 | PRATEEK PANDA                  | Show that the CFG with given productions is ambiguous and find an equivalent unambiguous grammar.  a. $S \rightarrow a \mid Sa \mid bSS \mid SSb \mid SbS$ b. $S \rightarrow SS \mid a \mid b$ c. $S \rightarrow ABA$ $A \rightarrow aA \mid \epsilon$ $B \rightarrow bB \mid \epsilon$   |
|----|------------|--------------------------------|---|
| 56 | 1RV20CS121 | PRATHEEK M                     | Describe the language generated in each case by the CFG with productions a. $S \rightarrow ST \mid \epsilon$ $T \rightarrow aS \mid bT \mid b$ b. $S \rightarrow aaS \mid bbS \mid Saa \mid Sbb \mid abSab \mid abSba \mid baSab \mid baSab \mid \epsilon$ c. $S \rightarrow aSB \mid bSA \mid \epsilon$ $A \rightarrow a$ $B \rightarrow b$ d. $S \rightarrow aaSbb \mid SS \mid \epsilon$ |
| 57 | 1RV20CS122 | PRATHIKSHA K R                 | For $\Sigma = \{a, b\}$ , construct DFA for language of all strings in which both the number of a's and the number of b's are even.   |
| 58 | 1RV20CS123 | PRATIKSHA<br>NARASIMHA NAYAK G | Show that the CFG with given productions is ambiguous and find an equivalent unambiguous grammar. $ a.  S \to aSb \mid aaSb \mid \epsilon                                  $  |
| 59 | 1RV20CS124 | PRATYUSH KISHORE               | For the FAs pictured below, use the minimization algorithm (Equivalence) to find a minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)  |

| 60 | 1RV20CS125 | PRIKSHIT          | Each of the following grammars, though not regular, generates a regular language. In each case, find a regular grammar generating the language: d. $S \rightarrow SSS \mid a \mid ab$ e. $S \rightarrow AabB$ A $\rightarrow aA \mid bA \mid \epsilon$ B $\rightarrow Bab \mid Bb \mid ab \mid b$ |
|----|------------|-------------------|---|
| 61 | 1RV20CS126 | PRITVISH R        | For a string $x \in \{a, b\}$ * with $ x  = n$ , how many states are required for an FA accepting x and no other strings? For each of these states, describe the strings that cause the FA to be in that state.   |
| 62 | 1RV20CS128 | PYDI VENKAT       | For the FAs pictured below, use the minimization algorithm (Equivalence) to find a minimum-state FA recognizing the same language. (It's possible that the given FA may already be minimal.)  |
| 63 | 1RV20CS192 | AKANSHA A PAI     | Construct NPDAs that accept the following languages:<br>d. $L = \{a^nb^mc^{n+m} : n \ge 0, m \ge 0\}$<br>e. $L = \{a^nb^{n+m}c^m : n \ge 0, m \ge 1\}$<br>f. $L = \{a^3b^nc^n : n \ge 0\}$  |
| 64 | 1RV20CS193 | NANDINI MOONKA    | Find context-free grammars generating each of the languages below.<br>c. $\{a^nb^m:n\leq m+3\}$<br>d. $\{a^nb^m:n=m-1\}$<br>e. $\{a^nb^m:n\neq 2m\}$<br>f. $\{a^nb^m:2n\leq m\leq 3n\}$   |
| 65 | 1RV20CS194 | RAHUL ANBALAGAN   | Find context-free grammars generating each of the languages below.<br>a. $\{a^ib^j \mid i \leq j\}$<br>b. $\{a^ib^j \mid i \leq j\}$<br>c. $\{a^ib^j \mid j = 2i\}$<br>d. $\{a^ib^j \mid i \leq j \leq 2i\}$  |
| 66 | 1RV20CS195 | SHRIKAR SWAROOP R | Construct NPDAs that accept the following languages:<br>g. $L = \{a^nb^{3n} : n \ge 0 \}$<br>h. $L = \{wcw^R : w \in \{a, b\}*\}$<br>i. $L = \{a^nb^m : n \le m \le 3n\}$   |

| 67 | 1RV20CS196 | DEEPTHA GIRIDHAR | Find context-free grammars generating each of the languages below.<br>a. $\{a^ib^j\mid j\leq 2i\}$<br>b. $\{a^ib^j\mid j<2i\}$<br>c. $\{a^ib^jc^k\mid i\neq j+k\}$<br>d. $\{a^nb^n\mid n \text{ is not a multiple of }3\}$                             |
|----|------------|------------------|--|
| 68 | 1RV20CS197 | SHARAN THOMAS    | Describe the language generated in each case by the CFG with productions $ \begin{array}{ccccccccccccccccccccccccccccccccccc$  |
| 69 | 1RV20CS198 | SHREYASA JOSHI   | For a string $x \in \{a, b\}$ * with $ x  = n$ , how many states are required for an FA accepting the language of all strings in $\{a, b\}$ * that begin with x? For each of these states, describe the strings that cause the FA to be in that state. |
| 70 | 1RV20CS199 | HARINI K S       | Construct NPDAs that accept the following regular languages: $ \begin{array}{l} j.  L_1 = L(aaa^*bab) \\ k.  L_2 = L(aab^*aba^*) \\ l.  L_1 \ U \ L_2 \\ m.  L_1 - L_2 \\ n.  L_1 \cap L_2 \end{array} $   |
| 71 | 1RV18CS198 | NAMAN SOOD       | Find a PDA that accepts the language $L = \{a^nb^{2n} : n \ge 0 \}$<br>Show the sequence of instantaneous descriptions for the acceptance of aabbbb  |