

Question - 1

Problem: Minimum Lead Blocks to Contain Radiation

You are given an integer array `radiation` where `radiation[i]` represents the **initial radiation level** at index `i`. Radiation attenuates (decreases) by **1 unit per step** as it spreads to neighboring indices.

You are also given an integer `m`. You may place **lead blocks** between two consecutive indices in the array. A lead block **prevents radiation from spreading further** across that boundary.

Your task is to determine the **minimum number of lead blocks** required such that, after radiation attenuation and blocking, the effective radiation level at **every index** is strictly less than `m`.

Example

Input

Radiation = [3, 1, 2, 1, 3]

M = 4

Output

2

Explanation

Step 1: Compute total radiation without blocks

Radiation at each index (summing contributions from every source):

- Index 0 → 3
- Index 1 → 4
- Index 2 → 4
- Index 3 → 4
- Index 4 → 3

So total = [3, 4, 4, 4, 3]

Step 2: Compare with the requirement (strictly less than $m = 4$)

- Index 0 $\rightarrow 3 < 4$ ✓
- Index 1 $\rightarrow 4$ ✗ (not strictly less)
- Index 2 $\rightarrow 4$ ✗
- Index 3 $\rightarrow 4$ ✗
- Index 4 $\rightarrow 3 < 4$ ✓

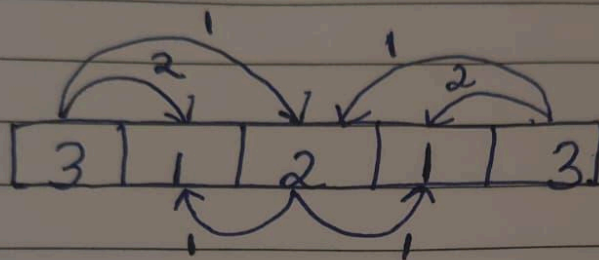
Indexes 1, 2, 3 are invalid.

Step 3: Place lead blocks to reduce radiation

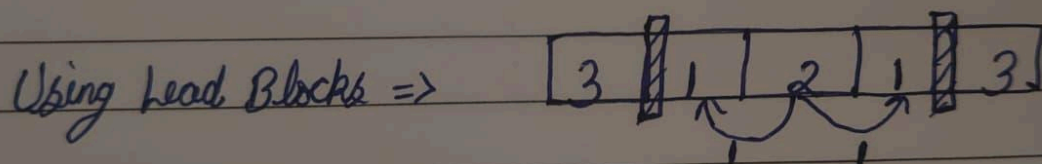
- Place a block between index 0 and 1:
New radiation = [3, 2, 2, 2, 3]
- Check again with $m = 4$:
 - Index 0 $\rightarrow 3 < 4$ ✓
 - Index 1 $\rightarrow 2 < 4$ ✓
 - Index 2 $\rightarrow 2 < 4$ ✓
 - Index 3 $\rightarrow 2 < 4$ ✓
 - Index 4 $\rightarrow 3 < 4$ ✓

Now all are strictly less than 4.

IMAGE GIVEN BELOW FOR REFERENCE



Radiation \Rightarrow 3 4 4 4 3



Radiation \Rightarrow 3 2 2 2 3

Question - 2

Problem: Maximum Sum of Differences of Two Subarrays

You are given an integer array `nums` and an integer `m`.

Your task is to select **two non-overlapping subarrays of size `m`** from `nums`. Define the "difference sum" as the sum of absolute differences of corresponding elements between the two subarrays.

Formally, if $A = \text{nums}[i \dots i+m-1]$ and $B = \text{nums}[j \dots j+m-1]$ (with $i+m-1 < j$ or $j+m-1 < i$ so they don't overlap), then the score is:

$$\text{score}(A, B) = \sum_{k=0}^{m-1} |A[k] - B[k]|$$

Return the **maximum possible score**.

Example 1:

Input: `nums = [1, 5, 2, 9, 6, 3]`, `m = 2`

Output: `10`

Explanation:

Choose subarray $A = [2, 9]$ and subarray $B = [6, 3]$.

The sum of absolute differences is:

$$|2-6| + |9-3| = 4+6 = 10$$

NOTE => BRUTE FORCE IS TOO INEFFICIENT AS COMPLEXITY = $O((N^2)*M)$.

TRY USING SLIDING WINDOW OR DYNAMIC PROGRAMMING

Question - 3

Based on Greedy Knapsack