

A Survey on Static Program Analysis

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Abstract

Static program analysis is an important and effective way to improve the quality of software systems. We conducted a lightweight survey by reviewing 16 papers related to static analysis published in the International Conference on Software Engineering (ICSE) between 2010 to 2015. Through this survey, we would like to get basic knowledges of the development of static analysis techniques in both academia and industry.

1 Introduction

Static program analysis, also called static analysis, is the analysis of computer software that is performed without actually executing programs. In most cases the analysis is performed on some version of the source code, and in the other cases, some form of the object code. Static analysis techniques range from the most mundane to the more complex, semantics-based analysis techniques.

In this survey, we review the papers related to static analysis published in ICSE between 2010 to 2015. The purpose of the survey is to investigate the research directions and the current achievements in the area of static analysis research. Hence, we constructed a Static Analysis publication repository [1] for this survey: we first searched two major online repositories, IEEE XPlore and ACM Portal, collecting 43 papers which have "static analysis", "static program analysis", "symbolic analysis", "program analysis", and "static + analysis" keywords in their title and abstract, and are published in the years between 2010 and

2015. Then, we filter out those unrelated to the static analysis techniques. In the end, we have 16 papers to review. Our discussion in the following sections is based on the analysis of these 16 papers, which is divided into three smaller categories according to the different applications of static analysis techniques: software testing, software maintenance and development, and software security.

2 The Application of Static Analysis

2.1 Using Static Analysis in Software Testing

In this section, we review 7 papers that made different innovations in the application of static analysis techniques in testing software.

In [2], Pradel *et al.* presented an automatic protocol conformance checker, which combines a dynamic miner of multi-object protocols[3][4] and a static checker of API constraints[5]. Compared with the existing approaches for API protocol testing, the new checker is built on two independent analyses. Such a strategy can guarantee the best results returned at different stages. In order to bridge these two analyzers, they designed a translation mechanism to connect two different formalisms used by the two analyzers respectively. Their approach finds bugs without any a priori known specifications and without human refinement or selection of the mined protocols. However, the cost of the automation is the introduction of false positives. To deal with the issue and improve the precision of results, they designed three pruning techniques: pruning by *Number of Confirmation*, by *Protocol Error Rate*, and by *Method Error Rate*[2].

Zhang presented a technique to combine results from a dynamic analysis and a static analysis to improve random test generation[6]. The combination can be summarized as three steps: dynamic inference, static analysis, and guided random test generation. The hybrid approach first takes a sample program as input, infers an call sequence model[7] by dynamic analysis, and enhances it with *direct state transition dependence constraint* and *abstract object profile constraint*. The enhanced call sequence model is used to capture the possible legal method

call sequences and argument values. Since the model may miss covering some interesting methods or certain method call invocation orders, a static method dependence analysis is designed to enrich the model, based on the hypothesis: two methods are related if the fields they read or write overlap. Finally, both the dynamically-inferred model and the statically-identified dependence information guide a random test generator to create legal and behaviorally-diverse tests. Compared with Randoop[8] (a pure random approach), Palulu[7] (a dynamic-random approach), RecGen[9] (a static-random approach), the author’s hybrid approach achieved higher line coverage on average.

Similarly, Ge *et al.* also proposed a hybrid defect-detection approach[10]. Their approach uses a novel way to combine both static verification and dynamic test generation so as to overcome the false positives and the path space explosion issues. Based on several existing tools, they constructed a tool chain called *DyTa* to implement the approach. According to their design, in the static phase, *DyTa*, first, applies the static checker of Code Contracts tool[11] to the program, and identifies a set of potential defects and collects their locations and descriptions. In order to make the program under analysis amenable for Dynamic Symbolic Execution (DSE)[12] -based test generation tool *Pex*[13], then *DyTa* performs two types of instrumentations to the program under analysis. In the dynamic phase, *DyTa* applies *Pex* to primarily explore those paths related to the potential defects detected by the static checker. Besides, different from other approaches, such as Check ‘n’ Crash [14], *DyTa* considers both locations of potential defects and error conditions so that the potential defect could be covered in the first place.

Although symbolic analysis has great potential for bug finding, such as better coverage than dynamic testing and higher precision than traditional static analysis, determining symbolic values for program variables especially related to loops and library calls is challenging, as the computation and data related to loops can have statically unknown bounds, and the library sources are typically not available at compile time. Le discussed the issues in details, and also proposed a novel concept of segmented symbolic analysis and its implementation, *Helium*, to deal with the issues in [15]. The core idea is to introduce dynamic

analysis to supply information that a pure static symbolic analyzer is slow or unable to produce. When a library call or a loop is encountered, the symbolic analysis suspends the propagation of the current symbolic condition and continues analyzing other paths. Meanwhile, *Helium* performs a structural and def-use analysis of the code and identifies the code segment for constructing a unit test. The dynamic analysis automatically synthesizes the unit test and generates the test inputs. After running the tests, a regression based inference is performed on test inputs and outputs to derive transfer functions for the unknown code segment. Notified with the newly discovered information, the symbolic analysis then updates the blocked symbolic conditions and resumes the analysis of this path. Different from other hybrid approaches, static and dynamic analyses are *concurrently* applied on different parts of the program to ensure that the capabilities of the analyses match the code characteristics, while the interaction between static and dynamic analyses are accomplished via a query based protocol.

Due to the difficulty of obtaining project-specific rules and manually writing checkers for them, commercial static program analysis tools are not fully put into use by software development organizations. To eliminate the obstacles and enhance the defect detection capability of static analysis, Sun *et al.* presented a hybrid approach in their research[16]. They introduced dependence-based pattern mining techniques[17] to capture frequent code patterns, which are in the form of generic System Dependence Graph (SDG)[18][19] subgraphs. Then, a rule extractor analyzes the SDG structure of a pattern and the abstract syntax trees (AST) of its statement instances, and generates the checking rules for the specific project. Finally, the mined code patterns are transformed into custom checkers that a commercial static analysis tool, such as Klocwork, can run against a code base to reveal defects.

[20] and [21] discuss the detection of concurrent-related errors, such as atomicity violations and data races. In [20], Zheng and Zhang proposed a context- and path-sensitive inter-procedural static analysis that detects atomicity violations in web applications regarding external resources in PHP code. The first novelty of the technique is that the authors developed a resource identity analysis, which

encodes the semantics of external operations into bit-vector logic constraints. These constraints, together with those generated from the regular PHP statements, are resolved by a SMT solver to determine if two given operations are accessing the same external resource. The second novelty is to define atomic regions by considering pair-wise atomicity of external operations, and by leveraging the resource identity analysis and program dependences. Marino *et al.* proposed an algorithm in their recent research[21] to extend an existing type-based static analysis so as to enhance the capability of deadlocks detection in programs written in a domain-specific language called *AJ*. The algorithm computes a partial order on atomic sets which is consistent with lock acquisition order. If such an order can be found, a program is deadlock-free. For programs that use recursive data structures, the approach is soundly extended to take into account a programmer-specified ordering between different instances of an atomic set.

2.2 Using Static Analysis in Software Development and Maintenance

Besides the wide application in the software testing phase of the entire software development life cycle, static analysis techniques also play significant role in other phases, such as development and maintenance phases.

Due to the dynamic nature of JavaScript, in IDEs for Javascript, code navigation and completion use heuristics that sometimes fail unexpectedly, while refactoring and analysis is all but unsupported. Different from statically typed programming languages, such as Java or C#, JavaScript is dynamically typed and uses prototype-based inheritance, therefore neither class hierarchy analysis nor its more refined variants that keep track of class instantiations [22] [23] are directly applicable. In [24], Feldthaus *et al.* argued that how to construct call graphs efficiently is the key to the problem. The existing flow analyses, which are either not fast enough for interactive use [25] [26] or only suitable for small framework-based applications [27], do not fit into the practical uses. Therefore, they proposed a fast field-based [28] [29] flow analysis. The approach

does not distinguish two functions that are assigned to properties of the same name. Besides, it only tracks function object and does not reason about any non-functional values, and meanwhile, it ignores dynamic property access (i.e., property reads and writes using JavaScript’s bracket syntax.) Like any flow analysis, this approach also faces a chicken-and-egg problem: to propagate (abstract) argument and return values between caller and callee, a call graph is required, yet a call graph is the construction target. To tackle this problem, the authors attempted two variants of the flow-based analysis: the standard optimistic analysis [30] (OA) and the pessimistic analysis (PA). OA starts out with an empty call graph, which is gradually extended as new flows are discovered until a fix point is reached, while PA does not reason about interprocedural flow at all and simply give up on call sites whose call target may depend on such flow, except in cases where the callee can be determined purely locally. According to their evaluations, PA may be preferable in many cases.

In [31], Balachandran proposed a tool, Review Bot, to integrate the static analysis into the code review phase of software development. The tool attempts to solve two issues: the automations of both checking coding standard violations and comment defect patterns and generating reviewer recommendations. To answer the first question, Review Bot is built as an extension to Review Board, an open-source code review tool. Review Bot uses three static analysis tools, which are Checkstyle, PMD, and FindBugs, to automate the checks for coding standard violations and common defect patterns, and publish code reviews using the output from these tools. To solve the second problem, the authors designed a reviewer recommendation algorithm, which is based on line change history of source code. The line change history of a line in a filediff contained in a review request is the list of review requests which affected that line in the past.

[32], [33], and [34] discuss the use of static analysis in detecting potential software system configuration errors. Usually, a configuration error is silent, sometimes non-crashing, which does not exhibit a crashing point, dump a stack trace, output an error message, or indicate suspicious program variables that may have incorrect values. Lacking such information makes many existing techniques such

as dynamic slicing [35], dynamic information flow tracking [36], and failure trace analysis inapplicable.

In [32] and [33], Zhang and Ernst discussed a case of configuration errors, in which the bug reporter had already minimized the bug report: if any part of the configuration or input is removed, SUT either crashes or no longer exhibits this error. Although the root cause of the bug is that the user failed to set one configuration option, no previous configuration error diagnosis technique [37] [36] [38] [39] [40] [41] [42] can be applied directly in this case. For that, they designed a tool called *ConfDiagnoser*, which is used by system administrators and end-users when they encounter an error that they do not know how to fix. The process of linking the undesired behavior to specific root cause can be summarized as three steps: (1) *ConfDiagnoser* uses thin slicing [20] to statically identify the predicates each configuration option affects in the source code; (2) *ConfDiagnoser* selectively instruments the program-to-diagnose so that it records the run-time behaviors of affected predicates in an execution profile; (3) *ConfDiagnoser* identifies the predicates whose dynamic behaviors deviate the most between correct and undesired executions, and then identifies its affecting configuration options as the likely root causes. Finally, it outputs a ranked list of suspicious configuration options and explanations. Compared to previous approaches [36] [38] [39] [40] [43] [35], *ConfDiagnoser* is fully automated, can diagnose both non-crashing and crashing errors, and does not require OS-level support.

Rabkin and Katz also discussed the issues happened in the software system configuration management in their research [34]. They observed that the key-value style of configuration is widely adopted in software systems, as it is convenient for developers or users to add new options incrementally. However, such a convenience also introduces some kinds of errors. For example, user-written configuration files can assign values for options that a program never reads; documentation falls out of step with an evolving program, and etc.. Therefore, Rabkin and Katz proposed an approach based on static analysis. They classifies the program configuration options into 17 types. Then, [34]

Kim *et al.* discussed semantic clone detections in their research [44]. According to their investigation, most clone detectors [45] [46] [47] [48] [49] are based on textual similarity, which are not effective to detect semantic clones that are functionally similar but syntactically different, while existing approaches to detect semantic clones have limitations. For example, program dependence graphs-based approaches can be affected by syntactic changes so as to miss some semantic clones; the clone detectability of random testing-based approaches may depend on the limited test coverage. Hence, Kim *et al.* introduced semantic-based static analysis into the clone detection. They built a semantic-based static analyzer on top of commercialized analyzer SPARROW, which can summarize each procedure after analyzing the procedure based on the abstract interpretation framework [50], and these procedural summaries have been carefully tuned to capture all memory-related behaviors in real-world C programs [51]. In order to support path-sensitive analysis, they further extended SPARROW to be path-sensitive like [52] by adding guards and guarded values to the abstract domain. By comparing programs' abstract memory states, they achieve the goal of the semantic clone detection.

Static program slicing is an attractive option for performing routine change impact analysis of newly committed changesets. However, static program slicing suffers from performance issues when analyzing large systems. Besides, program slicing can have accuracy issues as well [53]. In their experiment [54], Acharya and Robinson identified and outlined the unique problems encountered in using static program slicing directly for change impact analysis of large industrial software: (1) the build time and the total slice time are excessive with HIGH setting being the most precise. (2) the safety with HIGH setting might come at the expense of the impact set being over conservative. (3) it's not clear how much of the impact set's accuracy is sacrificed for faster builds with the LOW setting. (4) The impact set can be very large so that it cannot be easily explored by the developers for routinely inspecting the impact of the committed changesets. The terms of HIGH and LOW denote two sets of CodeSurfer, which is a commercial static analyzer. In order to reduce the time and improve the accuracy, Acharya and Robinson proposed a framework based on static program slicing, *Imp*. The framework is implemented on top of CodeSurfer. The core ideas behind *Imp*

are (1) perform HIGH setting analysis infrequently; (2) obtain clues in the program for possible over conservativeness by computing the common global data structures among large impact sets and the high impact functions; (3) perform LOW setting analysis frequently; (4) use the clues obtained from the infrequent HIGH setting builds to guide the LOW setting builds. According to their evaluations, *Imp* is capable of perform change impact analysis on systems with over a million lines of code.

2.3 Using Static Analysis in Software Security

Security of software systems touches on a vast and complex array of issues, making it difficult and expensive to implement a comprehensive security solution. In this section, we discuss three novel approaches to detect vulnerabilities.

Remote code execution (RCE) attacks are one of the most prominent security threats for web applications. However, the existing detection techniques failed or have a limited capability to detect and confirm RCE attacks. Most of the existing static analysis techniques [55] [56] [57] [58] [59] [20] [60] cannot cohesively reason about the string and non-string parts of an application and many lack path sensitivity, whereas RCE attacks require satisfying intriguing path conditions, involving both strings and non-strings. Even though some researchers model both strings and non-strings in dynamic symbolic execution of web apps [61], these new techniques simply focus on modeling the executed path, whereas RCE attack detection requires modeling all program paths. To solve the issue, Zheng and Zhang proposed a path- and context-sensitive inter-procedural static analysis in their research [62], which can detect RCE vulnerabilities in PHP code. Their system makes use of LLVM, the PHP compiler (*phc*) [63], the STP solver and the HAMPI string solver [55]. At the abstraction phase, the analysis first creates two abstractions of each PHP script: one for the string related behavior and the other for the non-string related behavior. The non-string abstraction includes additional taint semantics to reason about the input correlation for each variable. At the encoding and solving phase, two abstractions are encoded separately. An algorithm is designed for querying the STP solver and

the HAMPI string solver [55] iteratively and alternatively to derive a consistent path-sensitive solution for both sets of constraints. For a potentially vulnerable file write, the system queries the string constraints if the file name ends with the PHP extension, and queries the non-string constraints to determine if the written content is tainted and the file write is reachable. The solution has to consistently satisfy all these queries.

In the previous work [64] [65], the authors mined static code patterns that implement input validation and input sanitization methods to build vulnerability predictors based on supervised learning. However, the earlier work suffers from two major drawbacks: (1) the predictive capability of proposed static attributes is dependent on the precision of the classification of the input validation and sanitization code patterns; (2) the effectiveness of supervised learning-based approach is dependent on the availability of sufficient training data labeled with manually checked security vulnerabilities. In order to address the limitations, the authors proposed attributes [66], based on hybrid static and dynamic code analysis, which characterize input validation and sanitization code patterns for predicting SQL injection and XSS vulnerabilities. static analysis in their research is used to classify the nodes on the data dependence graph (DDG) of a sink that represents a program statement that interacts with database or web client. Then, their approach captures the classifications in a set of attributes on which vulnerability predictors are to be built.

Moller and Schwarz addressed their observation on web applications and argued that vulnerability involving client-state manipulation is strongly correlated to information flow from hidden fields or other kinds of client state to operations involving the shared application state on the server [67]. Based on the observation, they prototyped a static analysis tool, called *WARlord*, for mainstream web application frameworks to detect such vulnerabilities. The analysis has three major components: the first component infers the dataflow between the individual servlets and pages that constitute the application, in order to identify the client-state parameters; the second component finds out which objects represent shared application state by analyzing the program codes; the third component performs an information flow analysis to identify the possible flow

of user controllable input from client-state parameters to shared application state objects. Compared to the manual inspection in their evaluation, this approach has great precision, and it can analyze between 10 and 200 pages per minutes.

3 Conclusion and Future Work

In this paper, we reviewed 16 conference papers related to static program analysis based on three application categories: software testing, software maintenance and development, and software security. In the software testing community, the general research trend is to combine other assistant methods with pure static analysis, because pure static analysis inevitably produces false positives. Besides, the static analysis techniques are also adopted by researchers to detect resource contention problems in concurrent computing. In the software maintenance and development community, researchers prototyped different toolchain or extended existing automation tools to maximize the automation in coding/reviewing, managing configuration options, and monitoring and analyzing code changes. In the software security community, static analysis is used to obtain the overall path information of program under test.

Due to the time limitation, our survey simply covers 16 papers from a single international conference in only recent 5 years. Therefore, the help of the survey is very limited for a future research. However, it is a great start. We plan to extend our current survey to cover more research works in a wider timeline so that the survey can help pinpoint those valuable research points.

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