程序的机器级表示: 基础知识

Machine-Level Programming: Basics



- □ Intel处理器体系结构的历史
 History of Intel processors and architectures
- C语言, 汇编语言和机器语言 C, assembly and machine code



History of Intel processors and architectures

Intel x86系列处理器 Intel x86 Processors

- 在笔记本、桌面和服务器市场占有统治低位 Totally dominate laptop/desktop/server market
- 一个不断进化的设计过程 Evolutionary design
 - 向后兼容至8086处理器(诞生于1978年) Backwards compatible up until 8086, introduced in 1978
 - 随着时间增加了许多新的特性
 Added more features as time goes on

- 复杂指令集计算机 (CISC) Complex instruction set computer (CISC)
 - 指令多、格式复杂 Many different instructions with many different formats
 - 在Linux程序中只使用其中一个子集 But, only small subset encountered with Linux programs
 - 理论上CISC的性能很难与精简指令集计算机 (RISC) 相匹敌 Hard to match performance of Reduced Instruction Set Computers (RISC)
 - 但是Intel采用了CISC But, Intel has done just that!
 - 在低功耗情况下,速度会有影响 In terms of speed. Less so for low power.

History of Intel processors and architectures

Intel x86系列的进化: 里程碑 Intel x86 Evolution: Milestones

<i>Name</i> ■ 8086	<i>Date</i> 1978	<i>Transistors</i> 29K	<i>MHz</i> 5-10
■ 第一个16位Intel处理器,IBM PC计算机, DOS操作系统			
First 16-bit Intel processor. Basis for IBM PC & DOS			
■ 1MB寻址空间			
1MB addre	•	^ ·/	40.00
386	1985	275K	16-33
■ 第一个32位Intel处理器,采用IA32体系结构			
First 32 bit Intel processor , referred to as IA32			
■ 增加了扁平寻址,可以运行Unix操作系统			
Added "flat addressing", capable of running Unix			
Pentium 4E	2004	125M	2800-3800
■ 第一个64位Intel x86处理器,采用x86-64体系架构			
First 64-bit Intel x86 processor, referred to as x86-64			
Core 2	2006	291M	1060-3500
第一个双核Intel处理器			
First multi-core Intel processor			
Core i7	2008	731M	1700-3900
Core i9	2017	7G	3300-4500

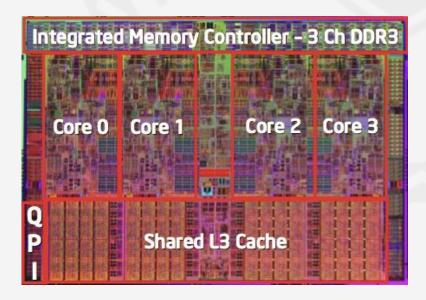
History of Intel processors and architectures

Intel x86系列处理器 Intel x86 Processors

新特性的加入

Added Features

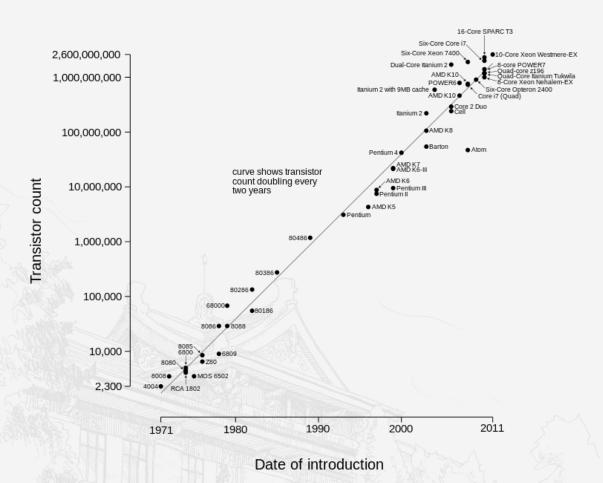
- 多媒体操作的指令支持
 Instructions to support multimedia operations
- 提供了更加有效率的条件操作指令
 Instructions to enable more efficient conditional operations
- 机器字长从32位变为64位 Transition from 32 bits to 64 bits
- 多核
 More cores





History of Intel processors and architectures

摩尔定律 Moore's Law



单位面积上可以容纳的晶体管数量 几乎每两年增加一倍

The number of transistors in a dense integrated circuit doubles approximately every two years.



History of Intel processors and architectures

Core i7 Broadwell 2015

桌面版

服务器版

Desktop Model

4核

4 cores

集成图形单元(显卡) Integrated graphics

3.3-3.8 GHz

65W

Server Model

8核

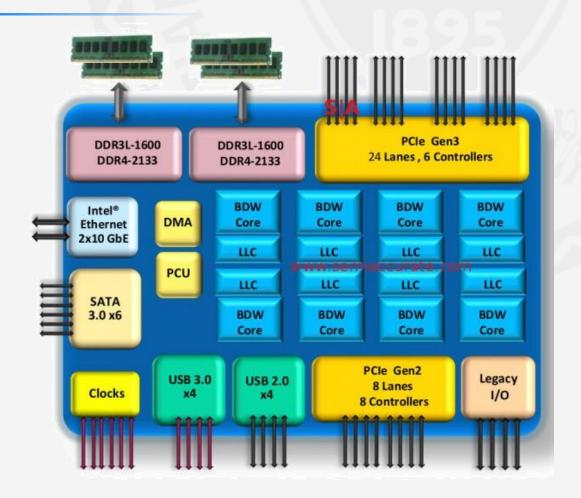
8 cores

■ 集成 I/O

Integrated I/O

2-2.6 GHz

45W



History of Intel processors and architectures

X86兼容处理器: AMD x86 Clones: Advanced Micro Devices (AMD)

- 一 历史上 Historically
 - AMD仅仅跟随着Intel
 AMD has followed just behind Intel
 - 性能稍差,价格更便宜 A little bit slower, a lot cheaper
- 随后 Then
 - 从DEC和其他业绩下降的公司招聘顶级电路设计师 Recruited top circuit designers from Digital Equipment Corp. and other downward trending companies
 - 是出了Opteron架构:成为Pentium 4的有力竞争对手 Built Opteron: tough competitor to Pentium 4
 - I 引入了x86-64架构, 自主的扩展到64位体系结构 Developed x86-64, their own extension to 64 bits

- 近年来 Recent Years
 - Intel重新迎头赶上,重新占据了半导体技术的世界主导地位 Intel got its act together, leading the world in semiconductor technology
 - AMD稍有落后,依赖外部的半导体代工厂 AMD has fallen behind, which relies on external semiconductor manufacturer

History of Intel processors and architectures

Intel 64位体系结构发展的历史 Intel's 64-Bit History

- 2001年: Intel试图从IA32彻底转变为IA64 2001: Intel Attempted Radical Shift from IA32 to IA64
 - 完全不同的体系结构(安腾)
 Totally different architecture (Itanium)
 - 以legacy(传统)模式执行IA32的指令 Executes IA32 code only as legacy
 - 性能令人失望 Performance disappointing
- 2003年: AMD提出了体系结构进化的解决方案 2003: AMD Stepped in with Evolutionary Solution
 - x86-64位体系结构(现称为AMD64) x86-64 (now called "AMD64")
- Intel觉得有义务专注于IA64
 Intel Felt Obligated to Focus on IA64
 - 难以承认技术路线的错误以及AMD方案更优 Hard to admit mistake or that AMD is better

- 2004年: Intel提出了EM64T体系结构实现对IA32的64 位扩展
 - 2004: Intel Announces EM64T extension to IA32
 - 扩展实现了64位内存寻址技术 Extended Memory 64-bit Technology
 - 几乎与x86-64相同Almost identical to x86-64!
- 2019年: 英特尔宣布放弃IA64架构 2019: Intel announced abandonment of IA64
- 除低端x86处理器外,其他处理器均支持x86-64 All but low-end x86 processors support x86-64
 - 但是目前许多程序仍然在32位模式下运行 But, lots of code still runs in 32-bit mode



- Intel处理器体系结构的历史
 History of Intel processors and architectures
- □ C语言, 汇编语言和机器语言 C, assembly and machine code

C, assembly and machine code

定义 Definitions

体系结构: (指令集体系结构, ISA) 编写汇编代码时需要理解的处理器设计部分。

Architecture: (also ISA: instruction set architecture) The parts of a processor design that one needs to understand to write assembly code.

- M如:指令集规范、寄存器组织 Examples: instruction set specification, registers.
- 微体系结构: 体系结构的具体实现
 Microarchitecture: Implementation of the architecture.
 - 例如: 高速缓存大小、核心频率 Examples: cache sizes and core frequency.

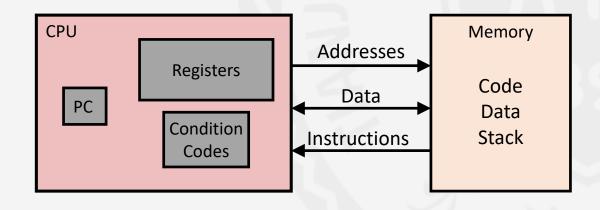
- 代码格式 Code Forms
 - 机器语言: 处理器可以直接执行的字节级的程序 Machine Code: The byte-level programs that a processor executes
 - **汇编语言:** 文本形式的机器语言 **Assembly Code:** A text representation of machine code
- 举例:常见的指令集体系结构 Example ISAs:
 - Intel: x86, IA32, Itanium, x86-64
 - ARM: 几乎所有的移动电话中都使用 ARM: Used in almost all mobile phones



汇编语言/机器语言视角下的计算机 Assembly/Machine Code View

C, assembly and machine code

程序员可见的状态 Programmer-Visible State



- ■PC:程序计数器
 - PC: Program counter
 - 存储下一条要执行指令的地址 Address of next instruction
 - 在x86-64中的名称为 RIP Called "RIP" (x86-64)
- 寄存器文件 Register file
 - 程序的数据会频繁地使用它来存储 Heavily used program data

■ 条件码

Condition codes

■ 存储最近一次算术逻辑运算的状态 信息

Store status information about most recent arithmetic or logical operation

■ 用于条件分支
Used for conditional branching

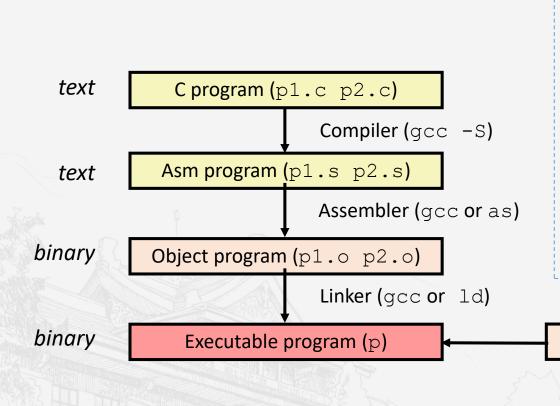
存储器

Memory

- 基于字节寻址的阵列
 Byte addressable array
- 存储程序和用户数据
 Code and user data
- F储栈数据(以实现过程的支持)
 Stack to support procedures

C, assembly and machine code

将C语言代码转换为机器语言 Turning C into Object Code



- 代码文件: p1.c p2.c Code in files p1.c p2.c
- 编译命令: gcc –Og p1.c p2.c -o p
 Compile with command: gcc –Og p1.c p2.c -o p
 - 使用基本的编译优化选项 –Og (最新版本GCC支持)
 Use basic optimizations –Og [New to recent versions of GCC]
 - 将编译结果写入文件 p
 Put resulting binary in file p

Static libraries (.a)

C, assembly and machine code

将C语言代码编译为汇编代码 Compiling Into Assembly

C代码

C Code (sum.c)

```
long plus(long x, long y);

void sumstore(long x, long y, long *dest)
{
    long t = plus(x, y);
    *dest = t;
}
```

X86-64汇编

Generated x86-64 Assembly

```
sumstore:
   pushq %rbx
   movq %rdx, %rbx
   call plus
   movq %rax, (%rbx)
   popq %rbx
   ret
```

- 使用下面的命令生成 Obtain with command gcc –Og –S sum.c
- 生成文件: sum.s Produces file sum.s

警告:由于编译选项的不同和gcc版本的不同可能会得到不同的编译结果

Warning: Will get very different results on other machines due to different versions of gcc and different compiler settings.

C, assembly and machine code

汇编语言的特征:数据类型 Assembly Characteristics: Data Types

- **整数:** 1、2、4或8字节的 "Integer" data of 1, 2, 4, or 8 bytes
 - 数据的值 Data values
 - 地址 (无类型的指针) Addresses (untyped pointers)
- **浮点数:** 4、8或10字节 Floating point data of 4, 8, or 10 bytes
- **代码:** 指令的字节序列编码 Code: Byte sequences encoding series of instructions
- **没有聚合类型**,例如:数组或结构体 No aggregate types such as arrays or structures
 - 这些在汇编语言中都都表现为在内存中连续分配的字节 Just contiguously allocated bytes in memory

C, assembly and machine code

汇编语言的特征: 操作 Assembly Characteristics: Data Types

- 对寄存器或存储器数据执行算术/逻辑运算
 Perform arithmetic function on register or memory data
- 在寄存器和存储器间传输数据
 Transfer data between memory and register
 - 将数据从存储器加载至寄存器
 Load data from memory into register
 - 将寄存器的数据存储至存储器 Store register data into memory
- 转移控制
 Transfer control
 - 无条件跳转 至/从 过程
 Unconditional jumps to/from procedures
 - 条件分支 Conditional branches



C, assembly and machine code

Code for sumstore

0x0400595:

0x53

0x48

0x89

0xd3

0xe8

0xf2

0xff

0xff

0xff

0x48

0x89

0x03

0x5b • 起始地址 0x0400595

• 14个字节

Total of 14 bytes

0xc3

Starts at address 0x0400595

• 不等长指令,1、3或5个字节

Each instruction 1, 3, or 5 bytes

目标码 Object Code

■ 汇编器 Assembler

- 将 .s 翻译为 .o Translates .s into .o
- 对每条指令进行二进制编码
 Binary encoding of each instruction
- 几乎是完整的可执行代码 Nearly-complete image of executable code
- 缺少了不同文件的链接信息
 Missing linkages between code in different files

链接器 Linker

- 实现了不同文件间的引用 Resolves references between files
- 与静态链接库进行了结合
 Combines with static run-time libraries
 - M如:代码中的 malloc、printf E.g., code for malloc, printf
- 某些库是需要动态链接的
 Some libraries are dynamically linked
 - 链接将出现在程序开始执行时 Linking occurs when program begins execution

C, assembly and machine code

举例: 机器指令 Machine Instruction Example

*dest = t;

movq %rax, (%rbx)

0x40059e: 48 89 03

C代码 C Code

> ■ 将 t 的值存储至 dest 指向的地址 Store value t where designated by dest

汇编

Assembly

- 将8字节的数据(在x86-64中称为四字)移动至存储器 Move 8-byte value (Quad words in x86-64 parlance) to memory
- 操作数 Operands

 t:
 Register
 %rax

 dest:
 Register
 %rbx

 *dest:
 Memory
 M[%rbx]

- 目标码(机器指令) Object Code
 - 3字节指令 3-byte instruction
 - 存储于地址 0x40059e Stored at address 0x40059e



C, assembly and machine code

反汇编目标码 Disassembling Object Code

反汇编

Disassembled

0000000000400595 <sumstore>: %rbx 400595: push 400596: 48 89 d3 %rdx,%rbx mov 400599: e8 f2 ff ff ff 400590 <plus> calla 48 89 03 %rax,(%rbx) 40059e: mov %rbx 4005a1: pop 4005a2: retq

- 反汇编器
 Disassembler **objdump -d sum**
 - 将 t 的值存储至 dest 指向的地址
 Store value t where designated by dest
 - 探索目标码的一个十分有用的工具 Useful tool for examining object code
 - 可以分析指令的编码序列 Analyzes bit pattern of series of instructions
 - 根据目标码重新生成汇编代码 Produces approximate rendition of assembly code
 - 可以对任何可执行程序文件和.o文件进行反汇编 Can be run on either a.out (complete executable) or .o file

C, assembly and machine code

目标码 Object

另一种反汇编方法 Alternate Disassembly

反汇编 Disassembled

0x0400595: 0x53 0x48 0x89 0xd3 0xe8 0xf2 0xff 0xff 0xff 0x48 0x89 0x03 0x5b 0xc3

```
Dump of assembler code for function sumstore:

0x000000000000400595 <+0>: push  %rbx

0x00000000000400596 <+1>: mov  %rdx,%rbx

0x00000000000400599 <+4>: callq  0x400590 <plus>
0x0000000000040059e <+9>: mov  %rax,(%rbx)

0x000000000004005a1 <+12>:pop  %rbx

0x0000000000004005a2 <+13>:retq
```

■ 使用gdb调试器 Within gdb Debugger

gdb sum

■ 反汇编过程(函数) Disassemble procedure

disassemble sumstore

反汇编从sumstore开始的的14个字节目标码 Examine the 14 bytes starting at sumstore

x/14xb sumstore



C, assembly and machine code

什么文件可以被反汇编? What Can be Disassembled?

```
% objdump -d WINWORD.EXE
```

WINWORD.EXE: file format pei-i386

No symbols in "WINWORD.EXE". Disassembly of section .text:

30001000 <.text>:

30001000: 55 push %ebp

30001001: 8b ec mov %esp,%ebp

30001003: 6a ff push \$0xffffffff

30001005: 68 90 10 00 30 push \$0x30001090

3000100a: 68 91 dc 4c 30 push \$0x304cdc91

- 任何可以解释为可执行 代码的文件 Anything that can be interpreted as executable code
- 反汇编程序分析字节并 重构为汇编代码 Disassembler examines bytes and reconstructs assembly source