程序的机器级表示: 控制

Machine-Level Programming: Control

本章内容 Topic

- □ 条件码 Condition codes
- 条件分支 Conditional branches
- 循环 Loops
- switch语句
 Switch Statements



处理器状态(x86-64, 部分的) Processor State (x86-64, Partial)

- 当前程序的执行信息 Information about currently executing program
 - 临时数据 Temporary data (%rax, ...)
 - 运行时栈的位置(栈顶) Location of runtime stack (%rsp)
 - 当前代码控制点的位置(即将要执行的指令地址) Location of current code control point (%rip, ...)
 - 最近一次指令执行的状态 Status of recent tests (CF, ZF, SF, OF)

	1108.00010	
	%rax	%r8
	%rbx	%r9
	%rcx	%r10
	%rdx	%r11
	%rsi	%r12
	%rdi	%r13
	%rsp	%r14
当前栈顶	%rbp	%r15
Current stack top	%rip	指令指针(程序计数器) Instruction pointer
	CF ZF SF	OF 条件码 Condition codes

Registers



条件码(隐式设置) Condition Codes (Implicit Setting) Not set by lead instruction

i leaq 指令不修改条件码

- ■简单的位寄存器 Single bit registers
 - Carry Flag (for unsigned) 进位标志 (无符号数)
 - Sign Flag (for signed) 符号标志 (有符号数)
 - Zero Flag 零标志
 - Overflow Flag (for signed) 溢出标志 (有符号数)

- ■通过算术运算可以隐式设置条件码(可以把它看做是运算的副作用) Implicitly set (think of it as side effect) by arithmetic operations
 - 例如: addq Src, Dest \leftrightarrow t = a+b Example: addq Src, Dest \leftrightarrow t = a+b
 - CF 被置位,如果运算时出现了超出最高位的进位(无符号数运算溢出) **CF** set if carry out from most significant bit (unsigned overflow)
 - **ZF** 被置位. 如果 t == 0 **ZF** set if t == 0
 - SF 被置位,如果 t<0 (看做是有符号数)</p> **SF** set if t < 0 (as signed)
 - OF 被置位,如果有符号数运算出现了溢出 **OF** set if two's-complement (signed) overflow (a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)



条件码(显式设置:比较指令) Condition Codes (Explicit Setting: Compare)

- ■通过比较指令可以显式设置条件码 Explicit Setting by Compare Instruction
 - cmpq Src2, Src1
 - cmpq b,a 这条指令和a-b的作用类似,但不需要将结果写入目标寄存器 cmpq b,a like computing a-b without setting destination
 - CF 被置位,如果运算时出现了超出最高位的借位(用于无符号数比较) CF set if carry out from most significant bit (used for unsigned comparisons)
 - **ZF** 被置位,如果 a == b **ZF** set if a == b
 - **SF** 被置位,如果 (a-b) < 0 (看做是有符号数) **SF** set if (a-b) < 0 (as signed)
 - OF 被置位,如果有符号数运算出现了溢出 OF set if two's-complement (signed) overflow (a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)



条件码(显式设置:测试指令) Condition Codes (Explicit Setting: Test)

- ■通过测试指令也可以显式设置条件码 Explicit Setting by Test Instruction
 - testq Src2, Src1
 - testq b,a 这条指令和a&b的作用类似,但不需要将结果写入目标寄存器 testq b,a like computing a&b without setting destination
 - 根据 Src1&Src2 的结果设置条件码 Sets condition codes based on value of Src1 & Src2
 - 用于对一个操作数的某几个位进行掩码检测 Useful to have one of the operands be a mask
 - **ZF** 被置位,当 a&b == 0 **ZF** set if a&b == 0
 - SF 被置位,如果 (a&b) < 0 SF set if (a&b) < 0



读取条件码 Reading Condition Codes

- SetX指令
 SetX Instructions
 - 根据条件码表达式将目标寄存器的最后 一个字节修改为0或1 Set low-order byte of destination to 0 or 1 based on combinations of condition codes
 - 一不会影响目标寄存器最高7个字节的值 Does not alter remaining 7 bytes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setns	~SF	Nonnegative
setg	~(SF^OF) &~ZF	Greater (Signed)
setge	~(SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF) ZF	Less or Equal (Signed)
seta	~CF&~ZF	Above (unsigned)
setb	CF	Below (unsigned)



x86-64 各寄存器中最后一个字节的名称 Referencing low-order byte of x86-64 Register

%rax	%al	%r8	%r8b
%rbx	%bl	%r9	%r9b
%rcx	%cl	%r10	%r10b
%rdx	%dl	%r11	%r11b
%rsi	%sil	%r12	%r12b
%rdi	%dil	%r13	%r13b
%rsp	%spl	%r14	%r14b
%rbp	%bpl	%r15	%r15b



读取条件码 Reading Condition Codes

```
int gt (long x, long y)
{
   return x > y;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
cmpq %rsi, %rdi # Compare x:y
setg %al # Set when >
movzbl %al, %eax # Zero rest of %rax
ret
```

在x86-64指令集中,32位操作指令 会将目标寄存器的高32位清0

In x86-64 ISA, 32-bit instructions also set upper 32 bits to 0

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跳转 Jumping

- ■jX指令 jX Instruction
 - 根据条件码跳转到代码的其他 位置执行 Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional (无条件)
je	ZF	Equal / Zero
jne	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF)&~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
j1	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF&~ZF	Above (unsigned)
jb	CF	Below (unsigned)



Conditional branches

跳转指令的编码 Jump Instruction Encoding

```
%rdi, %rax
4004d0:
               48 89 f8
                                              movq
4004d3:
               eb 03
                                                     4004d8<loop+0x8>
                                              jmp
4004d5:
               48 d1 f8
                                                     %rax
                                              sarq
               48 85 c0
4004d8:
                                              testq
                                                     %rax, %rax
                                                     4004d5<loop+0x5>
4004db:
               7f f8
                                              jg
4004dd:
               f3 c3
                                              repz retq
```



条件分支示例(早期模式) Conditional Branch Example (Old Style)

```
■生成汇编代码
Generation
> gcc -Og -S -fno-if-conversion control.c
```

```
long absdiff (long x, long y)
{
  long result;
  if (x > y)
    result = x-y;
  else
    result = y-x;
  return result;
}
  control.c
```

```
absdiff:
          %rsi, %rdi # x:y
  cmpq
  jle
          .L4
          %rdi, %rax
  movq
          %rsi, %rax
  subq
  ret
.L4:
        # x <= y
          %rsi, %rax
  movq
          %rdi, %rax
  subq
  ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

使用goto语句等价表示 Expressing with goto Code

- C语言允许使用goto语句 C allows goto statement
- 跳转至标签所在位置的语句继续执行 Jump to position designated by label

```
long absdiff (long x, long y)
{
  long result;
  if (x > y)
    result = x-y;
  else
    result = y-x;
  return result;
}
```

```
long absdiff_j (long x, long y)
    long result;
    int ntest = x <= y;
    if (ntest)
        goto Else;
    result = x-y;
    goto Done;
 Else:
    result = y-x;
Done:
    return result;
```



Conditional branches

条件表达式的翻译(使用分支) General Conditional Expression Translation (Using Branches)

c代码

C Code

val = Test ? Then_Expr : Else_Expr;

val = x>y ? x-y : y-x;

- ■为Then和Else表达式创建独立的代码块 separate code regions for then & else expressions
- ■根据条件选择合适的一个代码块并执行 Execute appropriate one

goto 语句版本

goto Version

```
ntest = !Test;
if (ntest)
        goto Else;
val = Then_Expr;
goto Done;
Else:
   val = Else_Expr;
Done:
   . . .
```

条件分支Conditional branches

使用条件数据移动指令 Using Conditional Moves

- 条件数据移动指令 Conditional Move Instructions
 - Instruction supports: if (Test) Dest ← Src
 - 1995年后的x86处理器开始支持 Supported in post-1995 x86 processors
 - GCC在编译时会尝试使用这个指令翻译条件分支 GCC tries to use them
 - 仅当保证逻辑安全的时候使用 But, only when known to be safe
- 为什么使用条件数据移动指令? Why?
 - 分支会破坏流水线的指令流,影像处理器性能
 Branches are very disruptive to instruction flow through pipelines
 - 条件数据移动指令不需要改变控制流 Conditional moves do not require control transfer

C 代码 C Code

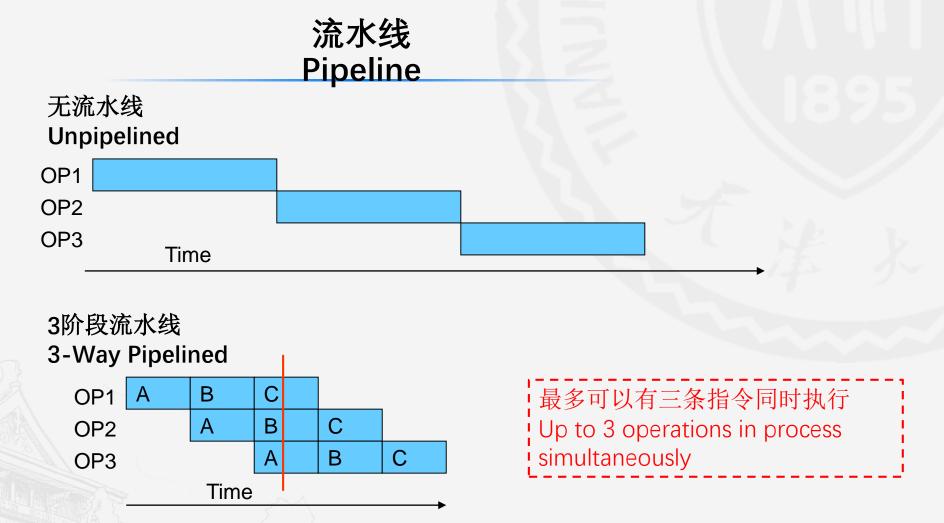
```
val = Test
? Then_Expr
: Else_Expr;
```

goto 语句版本 goto Version

```
result = Then_Expr;
eval = Else_Expr;
nt = !Test;
if (nt) result = eval;
return result;
```

条件分支

Conditional branches





举例:条件数据移动指令 Conditional Move Example

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
long absdiff (long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```



不能使用条件数据移动指令的情况 Bad Cases for Conditional Move

■大量的计算 Expensive Computations

- val = Test(x) ? Hard1(x) : Hard2(x);
- 条件数据移动指令会将所有的结果提前计算出来 Both values get computed
- 只有计算都非常简单的时候,使用条件数据移动指令才会有意义 Only makes sense when computations are very simple
- ■存在风险的计算 Risky Computations

- val = p ? *p : 0
- 可能导致程序出错
 May have undesirable effects
- 有副作用的计算 Computations with side effects

$$val = x > 0 ? x*=7 : x+=3;$$

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Do-While循环示例 "Do-While" Loop Example

C Code

```
long pcount_do (unsigned long x)
{
  long result = 0;
  do {
    result += x & 0x1;
    x >>= 1;
  } while (x);
  return result;
}
```

goto Version

```
long pcount_goto (unsigned long x)
{
  long result = 0;
  loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto loop;
    return result;
}
```

计算x编码中"1" 的个数 Count number of 1's in argument x 使用条件分支决定继续 或退出循环 Use conditional branch to either continue looping or to exit loop



Do-While循环编译后的结果 "Do-While" Loop Compilation

Register	Use(s)
%rdi	Argument x
%rax	result

goto Version

```
long pcount_goto (unsigned long x)
{
  long result = 0;
loop:
  result += x & 0x1;
  x >>= 1;
  if(x) goto loop;
  return result;
}
```

```
movl
         $0, %eax
                    # result = 0
.L2:
                    # loop:
         %rdi, %rdx
  movq
        1, \%edx # t = x & 0x1
  andl
  addq
         %rdx, %rax
                   # result += t
        %rdi
  shrq
                    # x >>= 1
  jne
         .L2
                    # if (x) goto loop
  rep; ret
```



Do-While循环通用的翻译方式 General "Do-While" Translation

C Code

```
do

Body

while (Test);
```

Goto Version

```
loop:
   Body
if (Test)
   goto loop
```

```
Body: {
          Statement<sub>1</sub>;
          Statement<sub>2</sub>;
          ...
          Statement<sub>n</sub>;
}
```



While循环通用的翻译方式 #1 General "While" Translation #1

- ■"跳转到中间"翻译方法
 "Jump-to-middle" translation
- 使用 -Og 编译优化选项生成代码 Used with -Og

While version

while (*Test*) *Body*



Goto Version

```
goto test;
loop:
   Body
test:
   if (Test)
     goto loop;
done:
```



举例: While循环 #1 While Loop Example #1

Jump to Middle

与 do-while 循环相比,循环开始 前先跳转至循环条件检测的位置 Compared with do-while version of function, initial goto starts loop at test

C Code

```
long pcount_while (unsigned long x)
{
  long result = 0;
  while (x) {
    result += x & 0x1;
    x >>= 1;
  }
  return result;
}
```

```
long pcount_goto_jtm (unsigned long x)
   long result = 0;
   goto test;
 loop:
   result += x \& 0x1;
  x >>= 1;
test:
   if(x) goto loop;
   return result;
```



While循环通用的翻译方式 #2 General "While" Translation #2

While version

```
while (Test)
Body
```





Do-While Version

```
if (!Test)
    goto done;
    do
        Body
        while(Test);
done:
```



Goto Version

```
if (!Test)
    goto done;
loop:
    Body
    if (Test)
       goto loop;
done:
```



举例: While循环 #2 While Loop Example #2 Do-While Version

与 do-while 循环相比,循环开始 前先检测循环条件,再进入循环 Compared with do-while version of function, initial conditional

guards entrance to loop

C Code

```
long pcount_while (unsigned long x)
{
  long result = 0;
  while (x) {
    result += x & 0x1;
    x >>= 1;
  }
  return result;
}
```

```
long pcount_goto_dw (unsigned long x)
  long result = 0;
  if (!x) goto done;
 loop:
  result += x \& 0x1;
 x >>= 1;
  if(x) goto loop;
done:
  return result;
```



Loops

一般形式

General Form

For循环一般形式 "For" Loop Form

```
for (Init; Test; Update )

Body
```

```
#define WSIZE 8*sizeof(int)
long pcount_for (unsigned long x)
{
    size_t i;
    long result = 0;
    for (i = 0; i < WSIZE; i++) {
        unsigned bit = (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

```
初始化
        i = 0
Init
检测
        i < WSIZE
Test
更新
        i++
Update
循环体
          unsigned bit =
Body
             (x >> i) & 0x1;
           result += bit;
```



"For" Loop → While Loop → Goto

For Version

```
for (Init; Test; Update )
Body
```



While Version

```
Init;
while (Test ) {
    Body
    Update;
}
```



```
Init;
if (!Test)
  goto done;
do
  Body
  Update
  while(Test);
done:
```

```
Init;
  if (!Test)
    goto done;
loop:
    Body
    Update
  if (Test)
    goto loop;
done:
```





For与While之间的转换 For-While Conversion

```
Init
   Test
          i < WSIZE
Update
          i++
  Body
            unsigned bit = (x >> i) \& 0x1;
            result += bit;
```

```
long pcount_for_while (unsigned long x)
  size_t i;
  long result = 0;
  i = 0;
 while (i < WSIZE)
    unsigned bit = (x >> i) \& 0x1;
    result += bit;
    i++;
  return result;
```



Loops

举例: For循环编译后的结果 "For" Loop Conversion Example

C Code

```
#define WSIZE 8*sizeof(int)
int pcount_for(unsigned x) {
   int i;
   int result = 0;
   for (i = 0; i < WSIZE; i++) {
     unsigned mask = 1 << i;
     result += (x & mask) != 0;
   }
   return result;
}</pre>
```

L在这个例子中,初始的循环条件上检测可以编译器被优化掉 Linitial test can be optimized away

Goto Version

```
int pcount for gt(unsigned x)
  int i;
  int result = 0;
  i = 0;
                     Init
 if (!(i < WSIZE))
                     ! Test
 loop:
                               Body
    unsigned mask = 1 << i;
    result += (x & mask) != 0;
  i++;
                     Update
  if (i < WSIZE)
    goto loop;
                     Test
done:
 return result;
```

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Switch Statements

Switch语句示例 Switch Statement Example

- ■多个case共用同一语句块 Multiple case labels
 - 5 & 6
- Case贯穿 Fall through cases
 - 2
- ■Case缺失 (case值不连续)
 Missing cases
 - 4

```
long switch_eg (long x, long y, long z)
    long w = 1;
    switch(x) {
    case 1:
        W = y*z;
        break;
    case 2:
        W = y/z;
        /* Fall Through */
    case 3:
        W += Z;
        break;
    case 5:
    case 6:
        W -= Z;
        break;
    default:
        W = 2;
    return w;
```

Switch Statements

跳转表的结构 Jump Table Structure

Switch的一般形式 Switch Form

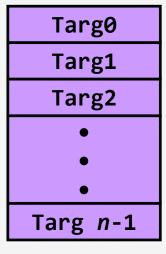
```
switch(x) {
  case val_0:
    Block 0
  case val_1:
    Block 1
    .....
  case val_n-1:
    Block n-1
}
```

翻译后(扩展C) Translation (Extended C)

```
goto *JTab[op];
```

跳转表 Jump Table

jtab:



跳转目标(语句块) Jump Targets

Code Block 2

Targ 0: Code Block 0

Targ 1: Code Block 1

Targ *n*-1: Code Block *n*-1

Targ 2:

Switch Statements

分析跳转表 1 Understanding Jump Table 1

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

setup:

switch_eg:

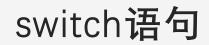
movq %rdx, %rcx

cmpq \$6, %rdi # x:6

ja .L8 # Use default

jmp *.L4(,%rdi,8) # goto *JTab[x]

默认值的范围是多少? What range of values takes default? 注意:w并没有在switch 开始前初始化 Note that w not initialized here



Switch Statements

间接跳转

Indirect jump

分析跳转表 2 Understanding Jump Table 2

```
long switch_eg(long x, long y, long z)
{
    long w = 1;
    switch(x) {
        . . .
    }
    return w;
}
```

Jump table

```
.section
          .rodata
  .align 8
.L4:
          .L8 \# x = 0
  .quad
          .L3 \# x = 1
  .quad
  .quad
          .L5 \# x = 2
          .L9 \# x = 3
  .quad
          .L8 \# x = 4
  .quad
  .quad
          .L7 \# x = 5
          .L7 \# x = 6
  .quad
```

Switch Statements

Setup部分汇编语句说明 Reading Condition Codes

Jump table

```
.section
         .rodata
 .align 8
.L4:
         .L8 \# x = 0
 .quad
         .L3 \# x = 1
 .quad
         .L5 \# x = 2
 .quad
 .quad
         .L9 \# x = 3
          .L8 \# x = 4
 .quad
          .L7 \# x = 5
 .quad
          .L7 \# x = 6
 .quad
```

- 跳转表结构 Jump Table Structure
 - 基地址是 .L4 Base address at .L4
 - 每个跳转目标需要8个字节(指向目标语句块的地址) Each target requires 8 bytes
- 直接跳转: jmp .L8
 Direct: jmp .L8
 - 直接跳转至.L8标签所指向地址的指令 Jump target is denoted by label .L8
- 间接跳转: jmp *.L4(,%rdi,8)
 Indirect: jmp *.L4(,%rdi,8)
 - 跳转表起始地址.L4
 Start of jump table: .L4
 - 缩放因子必须是8的整倍数(每个地址是8个字节) Must scale by factor of 8 (addresses are 8 bytes)
 - 从地址 .L4 + x*8 处获得跳转目标的位置 Fetch target from effective Address .L4 + x*8
 - **仅限于** $0 \le x \le 6$ 的情况 Only for $0 \le x \le 6$



分析跳转表 3 Understanding Jump Table 3

```
Jump table
                                          switch(x) {
                                                    // .L3
                                          case 1:
                                              W = y*z;
 .section .rodata
                                              break;
   .align 8
                                          case 2:
                                                    // .L5
 .L4:
                                              W = y/z;
           .L8 \# x = 0
   .quad
                                              /* Fall Through */
          .L3 \# x = 1
   .quad
                                          case 3:
                                                  // .L9
          .L5 \# x = 2
   .quad
                                              W += Z;
          .L9 \# x = 3
                                              break;
   .quad
                                          .case 5:
          .L8 \# x = 4
   .quad
                                          case 6: // .L7
   .quad
           .L7 \# x = 5
                                              W -= Z;
           .L7 \# x = 6
   .quad
                                              break;
                                          default:
                                                    // .L8
                                              W = 2;
```



代码块(x==1) Code Blocks (x == 1)

```
.L3:

movq %rsi, %rax # y

imulq %rdx, %rax # y*z

ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value



处理落入另一个case的情况 Handling Fall-Through

```
long w = 1;
switch(x) {
                                  case 2:
                                      w = y/z;
case 2:
                                      goto merge;
    w = y/z;
    /* Fall Through */
case 3:
    W += Z;
    break;
                                              case 3:
                                                       W = 1;
                                              merge:
                                                       W += Z;
```



代码块(x==2, x==3) Code Blocks (x==2, x==3)

```
long w = 1;
switch(x) {
 case 2:
     W = y/z;
     /* Fall Through */
 case 3:
     W += Z;
      break;
```

```
.L5:
                 # Case 2
       %rsi, %rax
  movq
  cqto
  idivq %rcx # y/z
         .L6 # goto merge
  jmp
.L9:
                # Case 3
       $1, %eax # w = 1
  movl
.L6:
                 # merge:
  addq %rcx, %rax # w += z
  ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value



代码块(x==5, x==6, 缺省) Code Blocks (x==5, x==6, default)

```
.L7:  # Case 5,6
  movl $1, %eax # w = 1
  subq %rdx, %rax # w -= z
  ret
.L8:  # Default:
  movl $2, %eax # 2
  ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value



没有从0处开始的情况 Case does not Start From 0

```
void switch_eg_2 (long x) {
    switch(x) {
    case 10000:
    case 10002:
    case 10003:
    case 10005:
    case 10006:
    default:
```

```
switch_eg_2:
    leaq -$10000(%rdi), %rsi # %rsi=%rdi-10000
    cmpq $6, %rsi # x:6
    ja .L8 # Use default
    jmp *.L4(,%rsi,8) # goto *JTab[x]
```



稀疏的switch语句 Sparse Switch Statements

```
if () {
  if () {
  } else {
} else {
  if () {
  } else {
```

- ■将翻译为二分查找的语句 O(log n) may translate into binary search
- 而不是退化为 if-elseif-elseif-else O(n) not if-elseif-elseif-else O(n)



总结 Summarizing

- C语言的控制方法 C Control
 - if-then-else
 - do-while
 - while, for
 - switch
- 汇编语言的控制方法 Assembler Control
 - 条件跳转 Conditional jump
 - ■条件数据移动 Conditional move
 - 间接跳转(利用跳转表实现) Indirect jump (via jump tables)
 - 编译器通过自动生成代码序列实现更复杂的控制 Compiler generates code sequence to implement more complex control

■标准技术 Standard Techniques

middle form

- 将循环转换为do-while 或 jump-to-middle 形式 Loops converted to do-while or jump-to-
- 大规模的switch语句可以使用跳转表实现 Large switch statements use jump tables
- 稀疏的switch语句可以使用决策树(二分 查找)实现 Sparse switch statements may use decision trees