

This handout is meant to serve as a representative for the prerequisite knowledge you might need in this module. Not all of it will be needed. You are **not** expected to recollect all these background details immediately and you will **not** be tested on these topics assessments directly. Use these questions to fill any gaps in your systems knowledge. A good “systems security” student should know these kinds of things.

If you're curious about the answers, use the LumiNUS class forum to discuss it with others.

1. A process on a computer has allocated only the virtual memory region 0x6000000 - 0x8000000. What happens when it accesses 0x5000000 at runtime?

- a. The processor will raise a General Protection Fault (GPF) and jump to the kernel (ring-0)
- b. The processor will raise a General Protection Fault (GPF) and jump to process A (ring-3)
- c. The processor will halt.
- d. The processor will kill the process A.

2. The OS wishes to switch context between process A (currently running) to process B (next to run). An efficient OS should perform which of the following steps :

- a. Copy all memory pages of A from RAM to disk, and load all memory pages of B from disk.
- b. Copy the page tables of B from disk to RAM, and write all page tables of A to disk
- c. Only switch the page-table-base-register (PDBR) to point to B's page table instead of A's
- d. No changes to page-table related registers / data is necessary.

3. How does an OS share a memory page between two processes? *IPC memset*

4. Here is the output of the `/proc/pid/maps` of a process. Can you identify the code segment, data segment of the `/bin/cat` binary executable and the stack segment?

Code -

Data -

Stack -

```

00400000-0040b000 r-xp 00000000 08:01 655384 /bin/cat
0060a000-0060b000 r--p 0000a000 08:01 655384 /bin/cat
0060b000-0060c000 rw-p 0000b000 08:01 655384 /bin/cat
009a4000-009c5000 rw-p 00000000 00:00 0 [heap]
7fdbcb3b000-7fdbcb421d000 r--p 00000000 08:01 663756 /usr/lib/locale/locale-archive
7fdbcb421d000-7fdbcb43d9000 r-xp 00000000 08:01 1446080 /lib/x86_64-linux-gnu/libc-2.19.so
7fdbcb43d9000-7fdbcb45d8000 ---p 001bc000 08:01 1446080 /lib/x86_64-linux-gnu/libc-2.19.so
7fdbcb45d8000-7fdbcb45dc000 r--p 001bb000 08:01 1446080 /lib/x86_64-linux-gnu/libc-2.19.so
7fdbcb45dc000-7fdbcb45de000 rw-p 001bf000 08:01 1446080 /lib/x86_64-linux-gnu/libc-2.19.so
7fdbcb45de000-7fdbcb45e3000 rw-p 00000000 00:00 0
7fdbcb45e3000-7fdbcb4606000 r-xp 00000000 08:01 1446056 /lib/x86_64-linux-gnu/ld-2.19.so
7fdbcb47ec000-7fdbcb47ef000 rw-p 00000000 00:00 0
7fdbcb4803000-7fdbcb4805000 rw-p 00000000 00:00 0
7fdbcb4805000-7fdbcb4806000 r--p 00022000 08:01 1446056 /lib/x86_64-linux-gnu/ld-2.19.so
7fdbcb4806000-7fdbcb4807000 rw-p 00023000 08:01 1446056 /lib/x86_64-linux-gnu/ld-2.19.so
7fdbcb4807000-7fdbcb4808000 rw-p 00000000 00:00 0
7fff3b99d000-7fff3b9be000 rw-p 00000000 00:00 0 [stack]
7fff3b9fe000-7fff3ba00000 r-xp 00000000 00:00 0 [vdso]
ffffffff600000-ffffffff601000 r-xp 00000000 00:00 0 [vsyscall]

```

5. What order are the following tools invoked in to prepare an executable? What is the input-output of each component? Which of these, if any, may be present in the process when the executable runs?

- A. Compiler, Linker, Assembler, Loader
- B. Compiler, Linker, Loader, Assembler
- C. Compiler, Assembler, Linker, Loader
- D. Compiler, Loader, Linker, Assembler

6. Consider the following six C program snippets. Which of these will fault on line 3 in function `foo`?

```

void foo (x) {
    int a, *b, c[1] ;
    b = (int*) malloc (sizeof (int));
    scanf ("%d", a);
}

```

```

void foo (x) {
    int a, *b, c[1] ;
    b = (int*) malloc (sizeof (int));
    scanf ("%d", b);
}

```

```

void foo (x) {
    int a, *b, c[1] ;
    b = (int*) malloc (sizeof (int));
    scanf ("%d", c);
}

```

6.

```

void foo (x) {
    int a, *b, c[1] ;
    b = (int*) malloc (sizeof (int));
    scanf ("%d", &a);
}

```

```

void foo (x) {
    int a, *b, c[1] ;
    b = (int*) malloc (sizeof (int));
    scanf ("%d", &b);
}

```

```

void foo (x) {
    int a, *b, c[1] ;
    b = (int*) malloc (sizeof (int));
    scanf ("%d", &c);
}

```

7. The 'call 0x5000' instruction on an x86 CPU has the following semantics :

```
RetVal := EIP+4; // EIP is the instruction pointer register
ESP := ESP - 4; // ESP is the stack pointer register
[ESP] := RetVal; // [REG] denotes access to value pointed by REG
EIP := 0x5000;
```

What do you think is the right semantics of the `ret` instruction?

✓	<pre>RetVal := [ESP]; ESP := ESP + 4; EIP:= RetVal;</pre>	<pre>ESP := ESP + 4; RetVal := [ESP]; EIP:= RetVal;</pre>
	<pre>RetVal := [ESP]; EIP:= RetVal;</pre>	<pre>RetVal := [ESP]; ESP := ESP - 4; EIP:= RetVal;</pre>

Hint: Intel x86: <http://www.cs.virginia.edu/~evans/cs216/guides/x86.html>

8. We want to dynamically allocate a 2D integer array (matrix) in C. Let the base pointer be `int** x`;
Which of the following is the safe way of doing this --- A, B, or C?

A. `x = malloc (100 * 50 * sizeof(int));`
`... x[i][j] ...`

B. `x = malloc(100 * sizeof(int));`
`for (int i = 0; i < 100; i++) {`
`x[i] = malloc(50 * sizeof(int));`
`}`
`... x[i][j] ...`

C. `x = malloc(100 * sizeof(int*));`
`for (int i = 0; i < 100; i++) {`
`x[i] = malloc(50 * sizeof(int));`
`}`
`... x[i][j] ...` ✓

(Hint: Code taken from [here](#))

9. Given a secret key k and a message m , which of the following is a secure one-time encryption technique? Explain why (ideally, write down a proof using the conditional probabilities $\Pr[m|E(m)]$).

- A. $E(m) := m \text{ bitwise-and } k$
- B. $E(m) := m \text{ bitwise-or } k$
- C.** $E(m) := m \text{ bitwise-xor } k$
- D. None of the above.

10. You are asked to debug the following program.

```
void main()
{
    int x = 10, y = 20;
    f(x, y);
}

void f(int x, int y)
{
    g(x, y);
}

void g(int x, int y)
{
    qsort(x, y);
}

void qsort(int x, int y)
{
    if (x > y)
        printf("%d, %d", y,x);
    else
        printf("%d, %d", x,y);
}
```

Write down the GDB commands to perform the following steps:

- a. Put a breakpoint on the entry to function g.
- b. Print the address of function g
- c. Print the call-stack when you hit the breakpoint set at step a.
- d. Print the address of variable x
- e. Print the contents of memory 0x5000
- f. Print the return address at breakpoint set at step a.
- g. Print the assembly instructions of function g.