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# CS2106

# Introduction to OS

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## Lecture 1 Introduction

# Overview

## ■ Operating Systems basic concepts:

- ❑ What is OS?
- ❑ Brief History
  - Motivation for OS
- ❑ Overview of Modern OSes

## ■ Operating System Structures

- ❑ OS components
- ❑ Types of kernel

## ■ Virtual Machines

# What is OS?

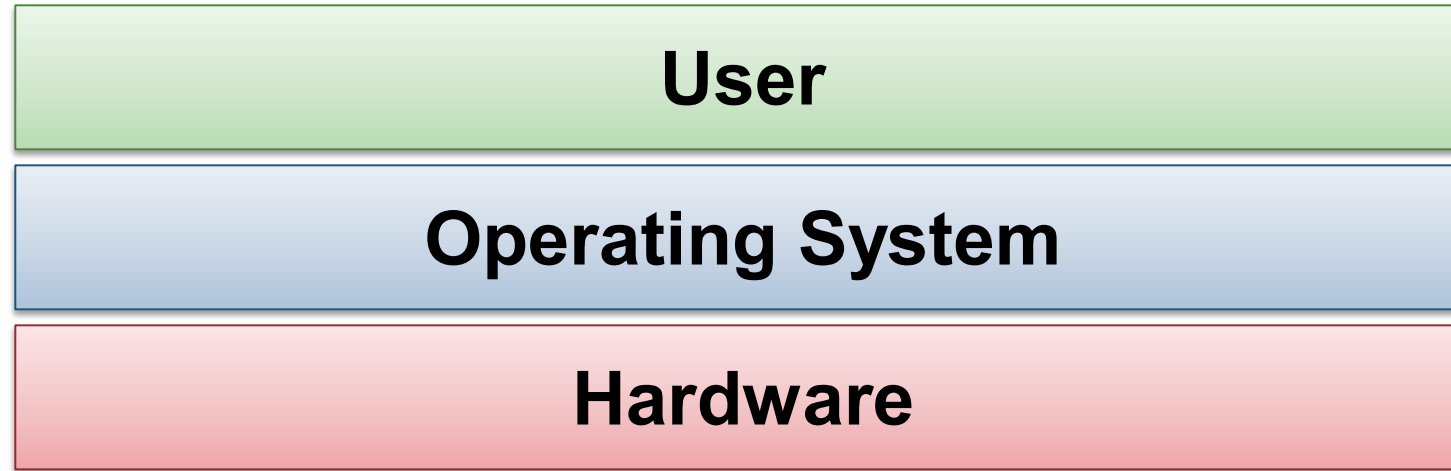
- Incorrect/Incomplete definition:

- ❑ It is the desktop when you boot up your PC
- ❑ The "thing" that stores your games
- ❑ Windows! (or Mac! ) (or Linux!)

- One simple definition:

- ❑ A **program** that acts as an **intermediary** between a **computer user** and the **computer hardware**
- ❑ **Wikipedia:** An operating system (OS) is system software that manages computer hardware, software resources, and provides common services for computer programs.

# Illustration: What is an OS?



- A simplified view:
  - Will be refined as we move along
- The most general version:
  - Hardware ( not only computer! )
  - User ( can be application programs or actual person! )

# Example of Common OS

- On Computer:

- ❑ Windows 10/8/XP
- ❑ Mac OS X
- ❑ Linux distros: Ubuntu, Redhat, Debian, Fedora, CentOS
- ❑ Solaris, FreeBSD

- On Smartphone:

- ❑ iOS, Android, Windows Mobile

- Other hardware with OS:

- ❑ Game console: PS4, Xbox, Nintendo Switch, ...
- ❑ Home appliance: Blu-ray/DVD Player, Mio Box, ...

To invent the future, you must understand the past

# BRIEF HISTORY OF OS

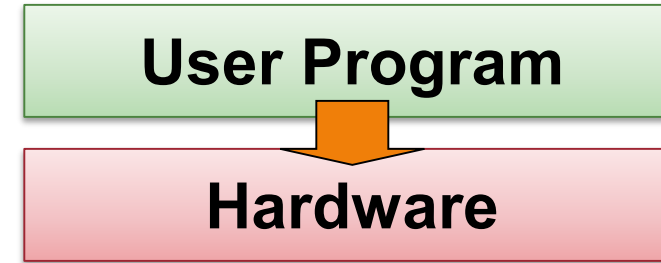
# Brief History of OS

- Essentially, OS evolves with:
  - ❑ Computer hardware
  - ❑ User application and usage pattern
- The "first" computers:
  - ❑ Electronic Numerical Integrator And Computer (**ENIAC**)
    - 1945
    - Program controlled by cables and switches
  - ❑ Harvard Mark I:
    - 1944
    - Program controlled by punched paper tape

# OS for the first computers

- OS Type:

- ❑ **NO OS**



- Programs directly interact with hardware

- ❑ Reprogram by changing **physical** configuration of hardware

- **Advantage:**

- ❑ Minimal overhead

- **Disadvantage:**

- ❑ Not portable

- ❑ Inefficient use of computer!



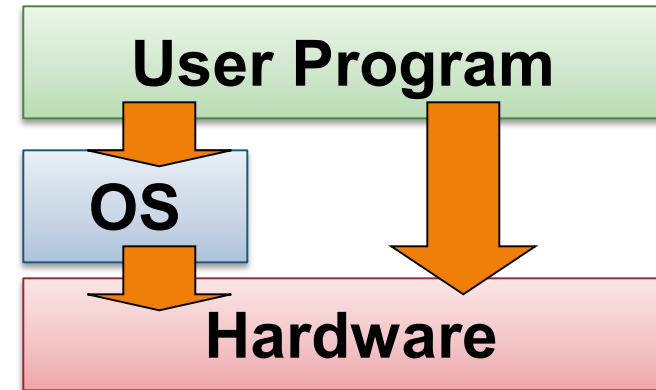
# Mainframes: The "Big Iron"

- Commonly used by large corporations in 60s, 70s
- Common features:
  - ❑ No interactive interface
  - ❑ Accept programs in the form of:
    - Paper tape, magnetic tape, punch card
  - ❑ Support batch processing only
  - ❑ Very costly
    - Usually "rented" instead of owned
- Example:
  - ❑ IBM 360
    - Cost 5 billion US dollars in 1964 to develop
    - Cost 130k US dollar in 1965 to buy

# OS for Mainframes

- OS Type:

- ❑ Batch OS



- Batch OS:

- ❑ Execute user program (a.k.a **job**) one at a time
    - Load job from media, execute, collect result

- User Job:

- ❑ Still interact with hardware directly
  - ❑ With additional information for the OS
    - Resource required
    - Job specification

# OS for Mainframes: Improvements

- Simple batch processing is inefficient:
  - ❑ CPU idle when perform I/O
- One possible Improvements:
  - ❑ Multiprogramming:
    - loads multiple jobs and runs other jobs when I/O needs to be done
    - Overlaps computation with I/O
- Another development of OS during this period (70s):
  - ❑ Time-Sharing OS

# Time-Sharing OS

## ■ Features:

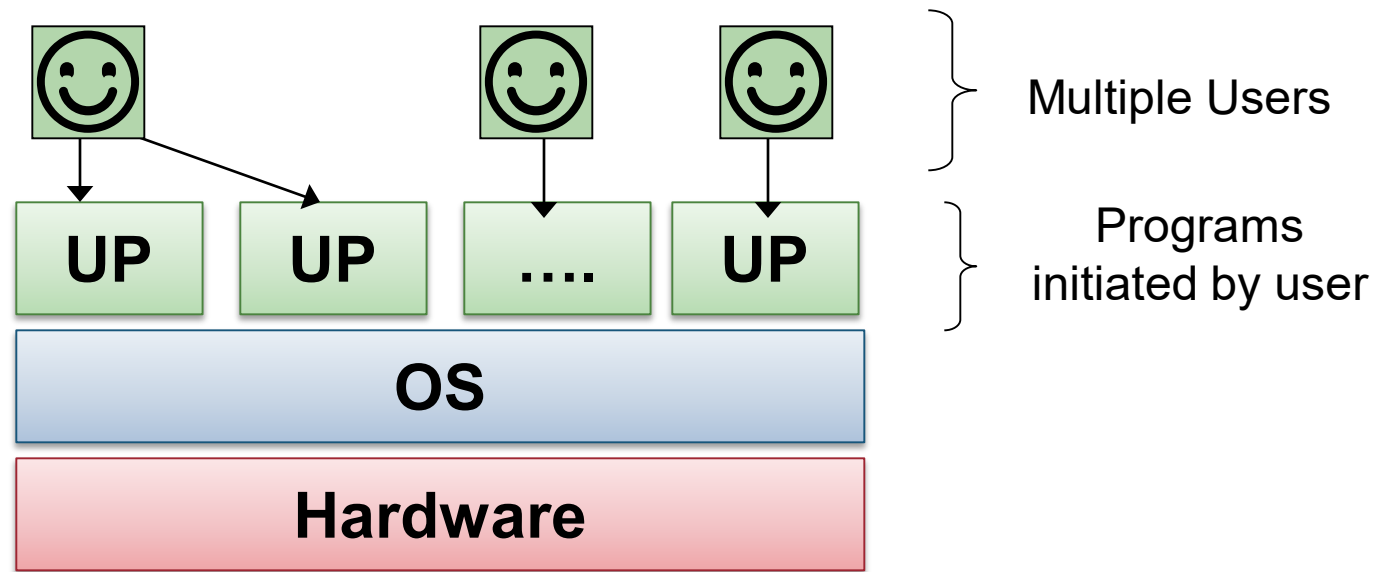
- ❑ Allow multiple users to interact with machine using terminals (***teletypes***)
- ❑ User job scheduling
  - Illusion of Concurrency
- ❑ Memory management

## ■ Famous Examples:

- ❑ CTSS developed at MIT 1960s
- ❑ Multics (1970s)
  - Considered as the *parent of Unix*
- ❑ Pushed the state of art in virtual memory, security

## ■ Similar to Unix servers today but more primitive

# Time-sharing OS: Illustration



- OS manages the sharing of:
  - ❑ CPU time, memory and storage
- **Virtualization** of hardware:
  - ❑ Each program executes **as if** it has all the resources to itself

# Minicomputer and Unix

- Minicomputer follows the mainframe:
  - A "mini" version of mainframe:
    - Smaller and cheaper
  - Example:
    - Digital Equipment Corp (DEC) PDP-11
- Famous OS:
  - Unix
    - Developed by AT&T employees, including Ken Thompson, Dennis Ritchie, Douglas McIlroy, and Joe Ossanna
    - Ken Thompson and Dennis Ritchie
      - Invented the C programming language as well!!

# Personal Computer

- Apple II PC (1977):
  - ❑ First successfully produced mass home computer
  - ❑ Designed by Steve Wozniak (alone!)
- IBM PC (1981):
  - ❑ The first generic PC
  - ❑ PC becoming a collection of **commodity** hardware components
  - ❑ Leads to dominance of Microsoft OSes on PCs: MSDOS (1981) then Windows (1985)

# OS on Personal Computer

- Machine (can be) dedicated to user, not timeshared between multiple users
  - Give rise to **personal OS**
- Several Models:
  - Windows model:
    - Single user at a time but possibly more than 1 user can access
    - Dedicated machine
  - **Unix model:**
    - One user at the workstation but other users can access remotely
    - General time sharing model



Abstraction  
Resource Allocator  
Control Program

Why do we need OS?

# MOTIVATIONS OF OS

# Motivation for OS: Abstraction

- Large variation in hardware configurations
- Example (Hard disk):
  - Different capacity (500mb, 320gb, 1.5tb etc)
  - Different capabilities:
    - Rotation per minutes (RPM)
    - Access (read/write) speed
    - etc.
- However, hardware in the same category has **well defined and common functionality**
  - Example (Hard disk): Store and retrieve information

# Motivation for OS: Abstraction

- Operating System serves as an **abstraction**:
  - Hide the different low level details
  - Present the common high level functionality to user
- The user can then perform essential tasks **through** operating system
  - no need to concern with low level details
- Provides:
  - Efficiency, programmability and **portability**

Or if the program follows POSIX, then it can be used between different OS

# Motivation for OS: **Resource Allocator**

- Program execution requires multiple resources:
  - CPU, memory, I/O devices etc.
- For better utilization of resources, multiple programs should be allowed to execute simultaneously
- OS is a **resource allocator**
  - Manages all resources
    - CPU, Memory, Input/Output devices
  - Arbitrate potentially conflicting requests
    - for efficient and fair resource use

# Motivation for OS: **Control Program**

- Program can misuse the computer:
  - Accidentally: due to coding bugs
  - Maliciously: virus, malware etc.
- Multiple users can share the computer:
  - Tricky to ensure **separate user space**
- OS is a **control program**
  - Controls execution of programs
    - Prevent errors and improper use of the computer
    - Provides security, isolation and protection

# Motivation for OS: **Summary**

- Manage resources and coordination
  - process synchronization, resource sharing
- Simplify programming
  - abstraction of hardware, convenient services
- Enforce usage policies
- Security and protection
- User Program Portability:
  - Across different hardware
- Efficiency
  - Sophisticated implementations
  - Optimized for particular usage and hardware

The families of modern OS

# OVERVIEW OF MODERN OS

# Modern OS: Overview

PC



Mobile



iOS

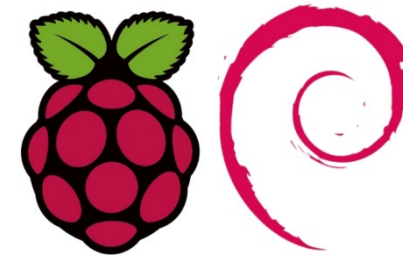


ANDROID

Real-Time

freeRTOS

Embedded



Raspbian



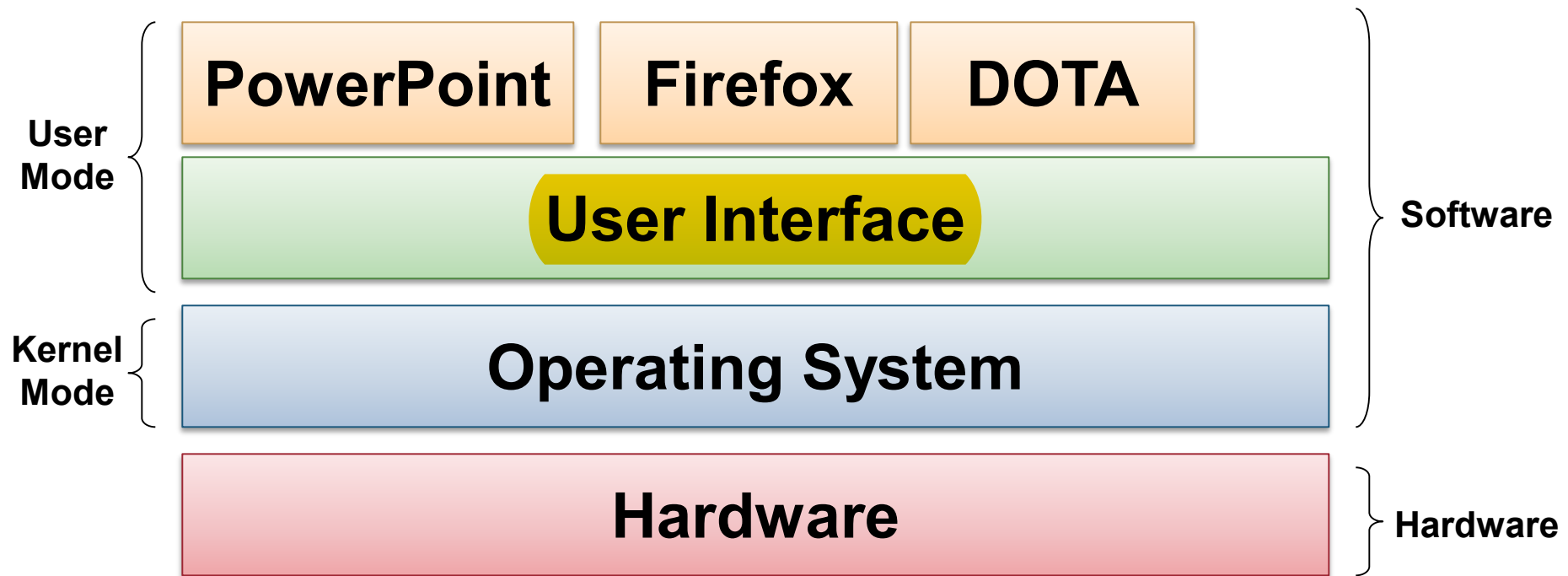
Common Architecture for OS

# OS STRUCTURE

# Operating System Structures

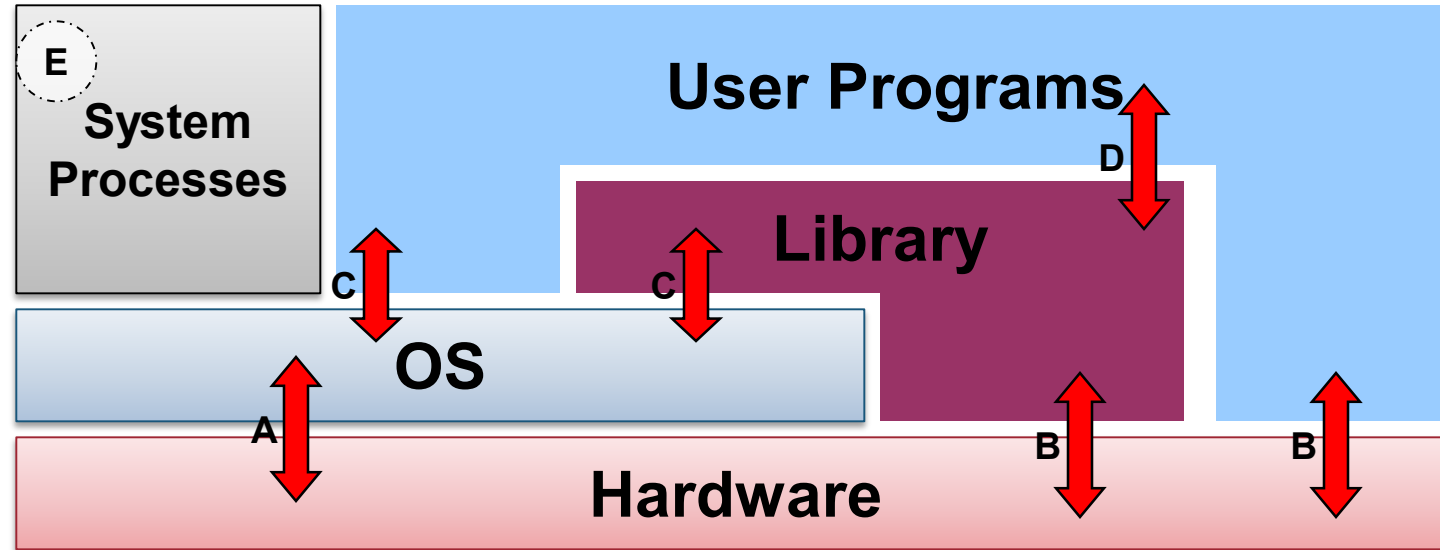
- We have identified the major capabilities of an OS
  - i.e. the *specification* of the OS
- Let us now consider:
  - The best way to provides these capabilities
  - i.e. the *implementations* of the OS
- Operating system structure:
  - *Organization* of the various components
  - Important factors:
    - Flexibility
    - Robustness
    - Maintainability
    - Performance

# Illustration: High level view of OS



- Operating System is essentially a software
  - ❑ Runs in **kernel mode**: complete access to all hardware resources
- Other software executes in **user mode**
  - ❑ With limited (or controlled) access to hardware resources

# Illustration: Generic OS Components



- **A**: OS executing machine instructions
- **B**: normal machine instructions executed (program/library code)
- **C**: calling OS using **system call interface**
- **D**: user program calls library code
- **E**: system processes
  - Provide high level services, usually part of OS

# OS as a Program

- OS is also known as the **kernel**
  - ❑ Just another program with some special features
    - Deals with hardware issues
    - Provides system call interface
    - Special code for interrupt handlers, device drivers
- Kernel code has to be different than normal programs:
  - ❑ no use of system call in kernel code
  - ❑ can't use normal libraries
  - ❑ no normal I/O
- Consider this:
  - ❑ Normal programs use OS: what does OS use? 😊

# Implementing Operating System

## ■ Programming Language:

- ❑ Historically in assembly/machine code
- ❑ Now in **HLLs:**
  - Especially C/C++
- ❑ Heavily hardware architecture dependent

## ■ Common code organization:

1. Machine independent HLL
2. Machine dependent HLL
3. Machine dependent assembly code

## ■ Challenges:

- ❑ “No one else” to rely on for nice services
- ❑ Debugging is hard
- ❑ Complexity
- ❑ Enormous Codebase

# OS Structures

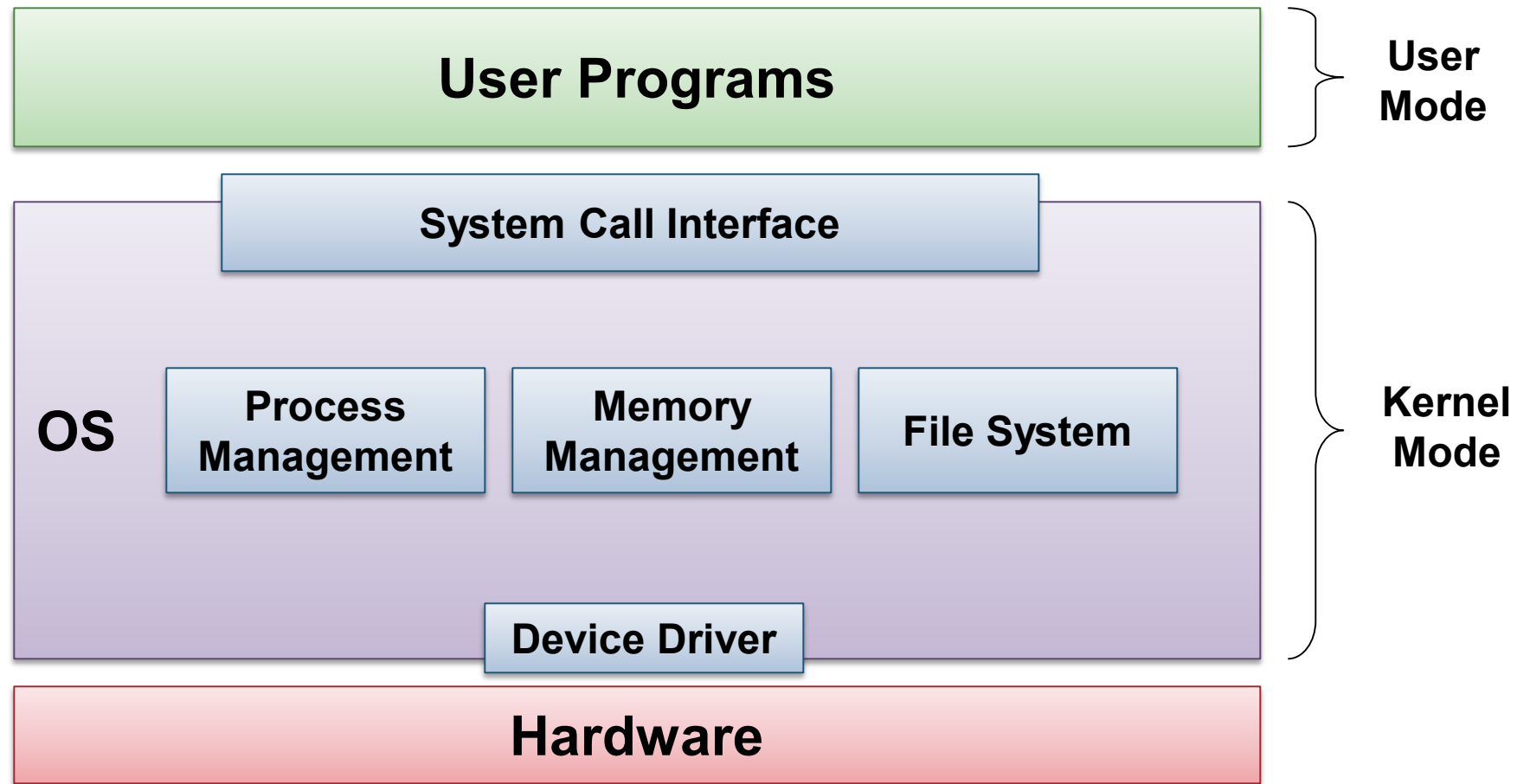
- Several ways to structure an OS:
  - ❑ **Monolithic**
  - ❑ **Microkernel**
  - ❑ Layered
  - ❑ Client-Server
  - ❑ Exokernel
  - ❑ Hybrid
  - ❑ etc.
- We will cover the first two in details:
  - ❑ They represent the whole range of possibilities
  - ❑ Most other approaches are variant or improvement

# Monolithic OS

- Kernel is:
  - ❑ One **BIG** special program
    - Various services and components are integral part
  - ❑ Good SE principles are still possible with:
    - modularization
    - separation of interfaces and implementation
- This is the traditional approach taken by:
  - ❑ Most Unix variants, DOS, Windows 9x
- **Advantages:**
  - ❑ Well understood
  - ❑ Good performance
- **Disadvantages:**
  - ❑ Highly coupled components
  - ❑ Usually devolved into very complicated internal structure



# Monolithic Kernel Illustration



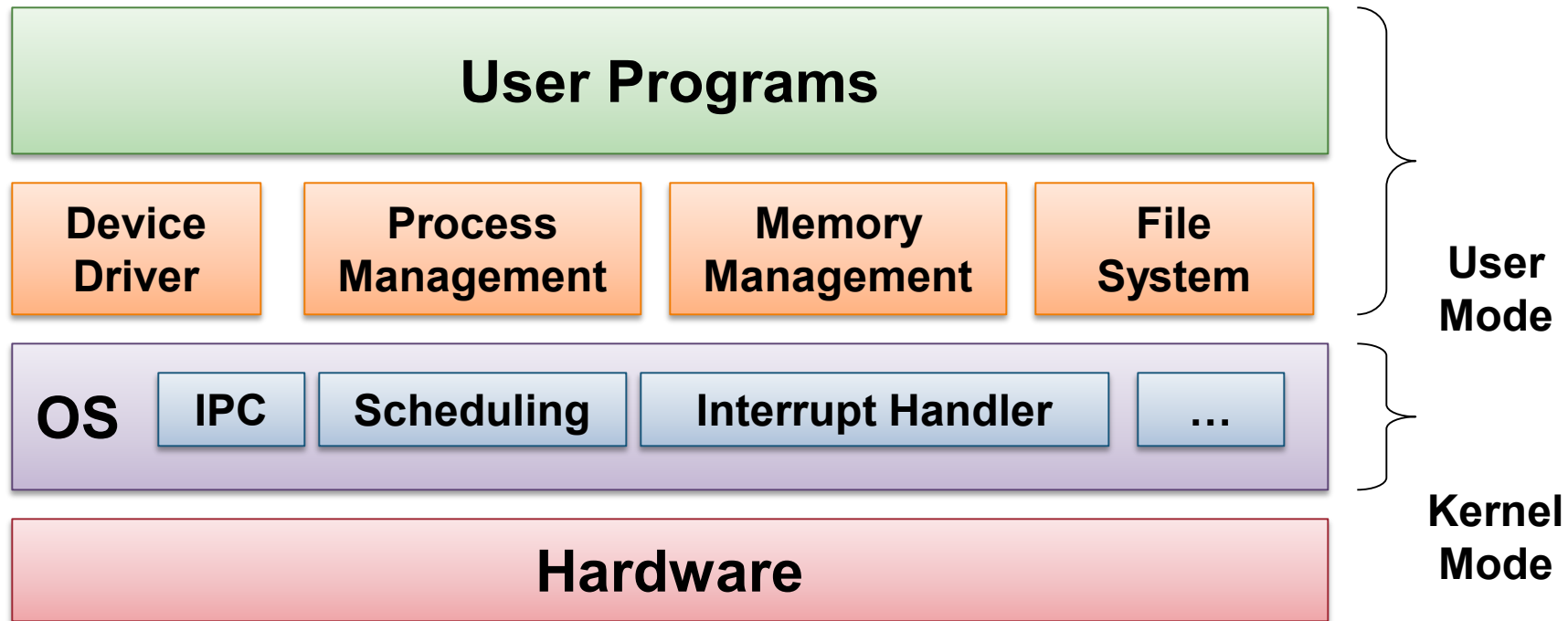
Generic Architecture of Monolithic OS Components

# Microkernel OS

- Kernel is:
  - ❑ Very small and clean
  - ❑ Only provides basic and essential facilities:
    - Inter-Process Communication (IPC)
    - Address space management
    - Thread management
    - etc.
- Higher level OS services:
  - ❑ Built on top of the basic facilities
  - ❑ Run as server process **outside** of the kernel
  - ❑ Use IPC to communicate
- Advantages:
  - ❑ Kernel is generally more robust and more extendible
  - ❑ Better isolation and protection between kernel and high level services
- Disadvantages:
  - ❑ Lower Performance

such as file system

# Microkernel Components



Generic Architecture of Microkernel OS Components

# Other Operating System Structure

## ■ Layered Systems:

- ❑ Generalization of monolithic system
- ❑ Organize the components into hierarchy of layers
  - Upper layers make use of the lower layers
  - Lowest layer is the hardware
  - Highest layer is the user interface

## ■ Client-Server Model

- ❑ Variation of microkernel
- ❑ Two classes of processes:
  - Client process request service from server process
  - Server Process built on top of the microkernel
  - Client and Server process can be on separate machine!

Ways of running OSes

# VIRTUAL MACHINES

# Motivation: Why Virtual Machines

- OS assumes total control of the hardware:
  - ❑ What if we want to run several OSes on the same hardware at the same time?
  - ❑ Cloud computing (IaaS)
- OS is hard to debug / monitor:
  - ❑ How do we observe the working of the OS?
  - ❑ How do we test a potentially destructive implementation?

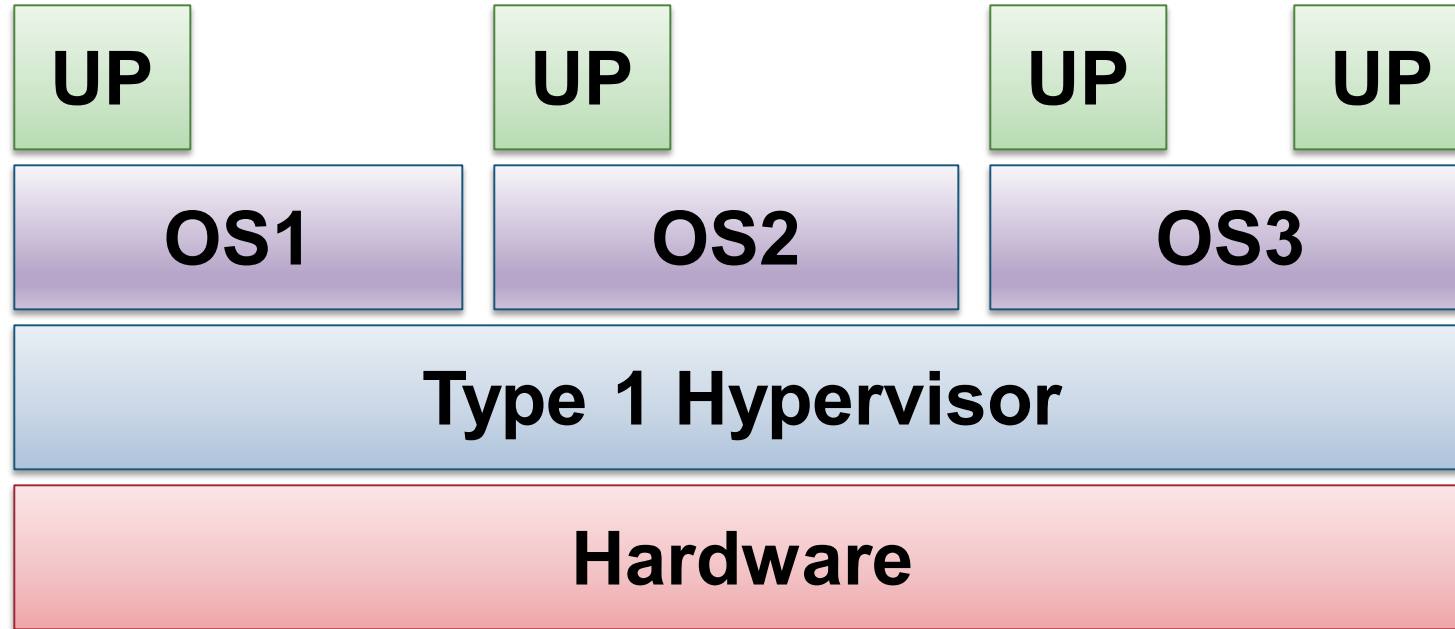
# Definition: Virtual Machine

- Virtual Machine:

- ❑ A software emulation of hardware
- ❑ **Virtualization** of underlying hardware
  - Illusion of complete hardware to level above: memory, CPU, hard disk etc...
- ❑ Normal (primitive) operating system can then run on top of the virtual machine

- Created and managed by **Hypervisor**
  - ❑ Also known as **Virtual Machine Monitor (VMM)**
  - ❑ Two classes of implementations shown next

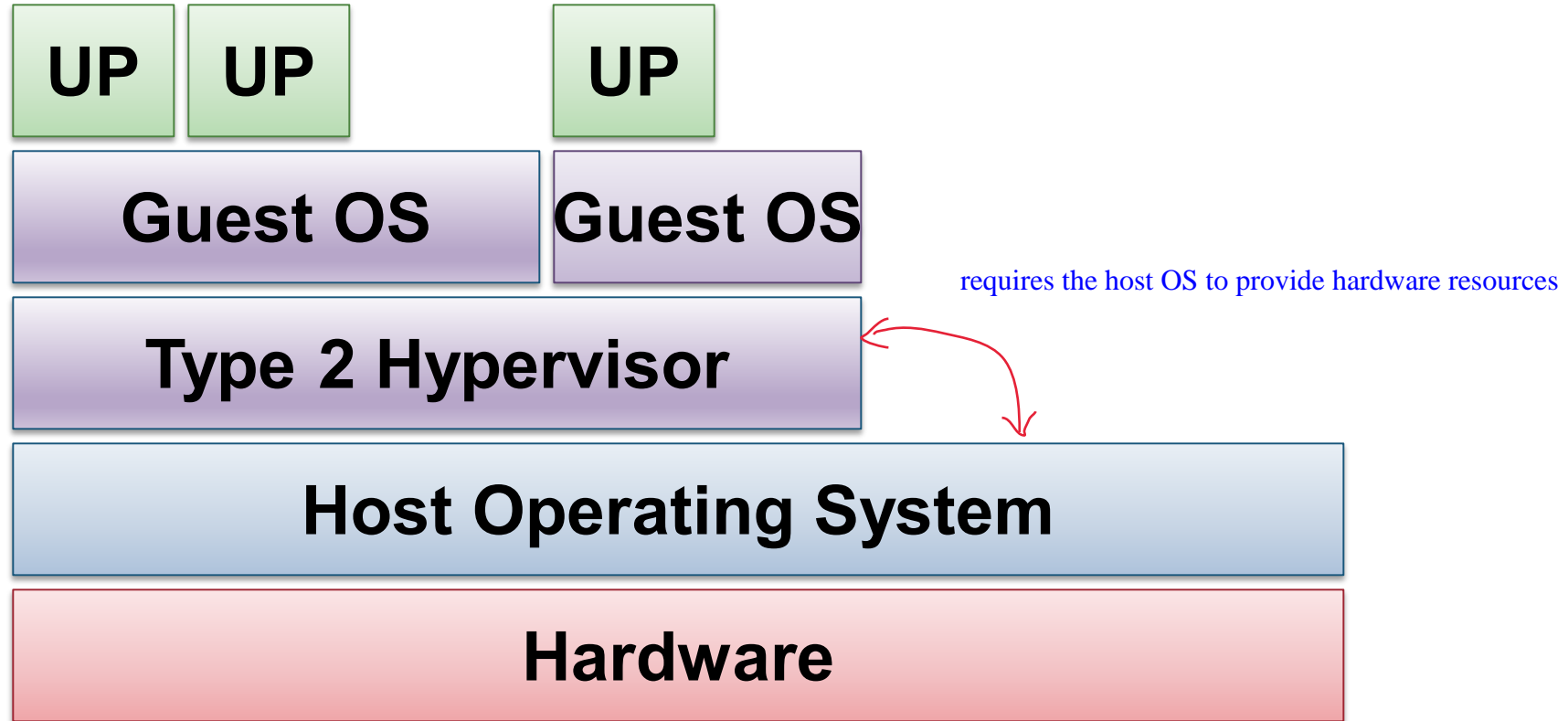
# Type 1 Hypervisor



- Type 1 hypervisor:
  - ❑ Provides individual *virtual machines* to guest OSes
  - ❑ eg. IBM VM/370



# Type 2 Hypervisor



- Type 2 hypervisor OS
  - ❑ Runs in host OS
  - ❑ Guest OS runs inside Virtual Machine
  - ❑ e.g. VMware

# Summary

- Definition of Operating System
- Roles of Operating System
- Common Operating System families
- Operating System structure

# Reference

- Modern Operating System (4<sup>th</sup> Edition)
  - By Andrew S.Tanenbaum
  - Published by Pearson
- Operating System Concepts (8<sup>th</sup> Edition)
  - By Abraham Silberschatz, Peter Baer Galvin & Greg Gagne
  - Published by McGraw Hill