

Data Structures

Course code: IT623

HASH TABLE



LANGUAGE COMPLEXITY ORGANIZING DATA
INFORMATION ORGANIZING DATA
CIENT OPERATIONS ABSTRACT STRUCTURE

COMPUTER APPLICATION PROCEDURE
IMPLEMENTATION
DATABASE PERFORM
AMOUNTS

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Lectures 3

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3. Modulan function & Anithmetic
* For any integer a and any positive integer n, the value a mod n is the remainder (on residue)
  of the quotient alm.
                         nead as: "a modulo n"
* Morre exactly K (mod M) 1sthe unique integers of such that
  t When ( is positive, simple divide ( by M) to obtain remainden ( ). Thus,
   7. In add my 25 (mod 7) = 411; 25 (mod 5) = 0, 35 (mod 11)=2, 3 (mod 8) = 3
    IF (a mod n) = (b mod n), we write a = b (mod n) and say that a is equivalent to b, modulo w.
                          => conduration >
    The mathematical Congruence relation is defined as follows:

a=b (mod m) if and only if M divides h-a.
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- * In other words, a=b (mod n) if a and b have the same nemainder when divided by n.
- * Equivalently, a= 6 (mod w), if and only if n is a divisor of b-a.
- a = b (mod n) it a is not equivalent to b, modulo N.) one god you or is considered

4. Integen and Absolute value Functions: 0, 32 (moq 11)= 3 3 (moq 2) = 3

Let & be any neal number. The integen value of x, written INT(x), convents or into ow integen by deleting (truncating) the fractional port of the number.

INT(x) = Loc | when x is positive

INT(x) =
$$[x]$$
 when x is negative.

Solution value gives a positive integer.

+ similarly, obsolute value gives a positive integer.

5. Exponentials:

For all neal 0>0, m, and w, we have following identities:

$$a^{-1} = 1/a$$

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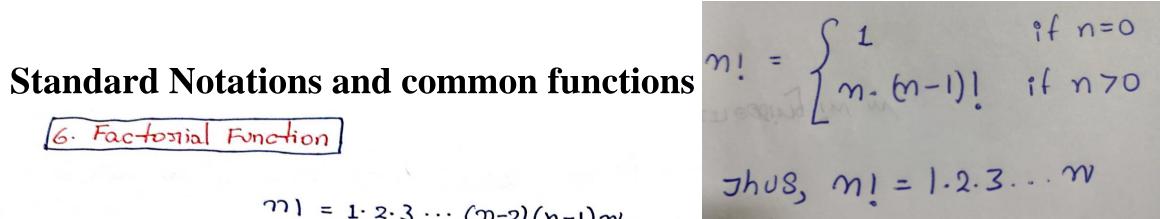
$$a^{-1} = 1/a$$

$$a^{-1} = a^{-1}$$

For all m and a 7,1, the function and is monotonically inconcessing in w.

6. Factorial Function

$$\gamma_1 = 1 \cdot 2 \cdot 3 \cdots (\gamma_{n-2})(\gamma_{n-1}) w$$
.



$$f^{(i)}(n) = \begin{cases} n & \text{if } i = 0 \\ f(f(i-1)(n)) & \text{if } i \neq 0 \end{cases}$$

Normal recurrence relation, not factorial

For example, if f(n) = 2n, then 2 w.

$$F_0 = 0$$
,
 $F_1 = 1$,
 $F_i = F_{i-1} + F_{i-2}$ for $i = 72$

8. Polynomials Given a nonnegative integer (a), a polynomial Pm n of degree d, is a function p(n) of the form $\phi(m) = \sum_{i=0}^{d} a_i m^i$; sometimes m is α . where the constants a, a, and are the coefficients of the polynomial and a for * A polynomial is osymptotically positive if and only if and only if and only if and only if * For an asymptotically positive polynomial b(n) of degree d, we have b(n) = O(nd) + k and not * For any real constant a 7,0, the function na is monotonically increasing * For any neal constant a < 0, the function not is monotonically decreasing. * A function f(n) is polynomially bounded if f(n) = O(nk) for some constant k.

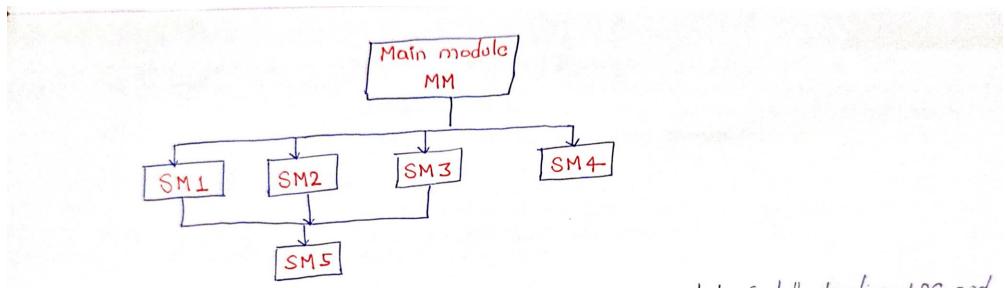
9. Logarithms lg n = log2n (binary algorithm) $ln w = loge n \quad (notwood logarithm)$ $lg^{\kappa} n = (lg n)^{\kappa} \quad (e \times ponentiation)$ For any 18 18 La = 18 (18 M) (Combosition) monotonically decreasing. For any 1 son constant 2,7,0, the Junction no is monotonically increasing An important notational convention we shall adopt is that logarithm functions will apply one to the next term in the formula, so that Ign + K will mean (Ign) + K and not 19(n+K) X. IF we bold b>1 constant, then for noo, the function logs of is Strictly in creasing. ... of one the coefficients of the low man for the

Fon all neal a70, 670, c70, and 27

$$* \quad a = b^{\log_b a},$$

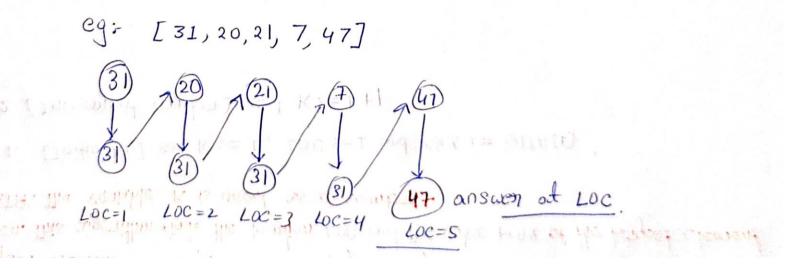
- * An algorithm, intuitively speaking, is a finite step-by-step list of well-defined instructions for solving a particular problem.
- * Simplifed one: " An Algorithm is a sequence of instructions for solving a problem"
- * Algorithms are implemented using programs.
- * An efficient program is organized into modules * Main Module

 * Sub module
- * Main module: Generial description of the algorithm. (MM)
- * Sub module: Detailed and specific information. (SM)



Example - An annay DATA of numerical values is in memory. We want to find the location LOC and

-the value MAX of the langest element of DATA.



Algorithm: (Langest Element in Annay) A non-empty annay DATA with N numerical volves is given. This algorithm finds the location LOC and the value MAX of the langest element of DATA. The vaniable K is used as a counterful.

Step 1. [Initialize.] Set K:= 1, LOC:=1 and MAX:= DATA[1]

step 2. [Increment counter] set K:= K+1

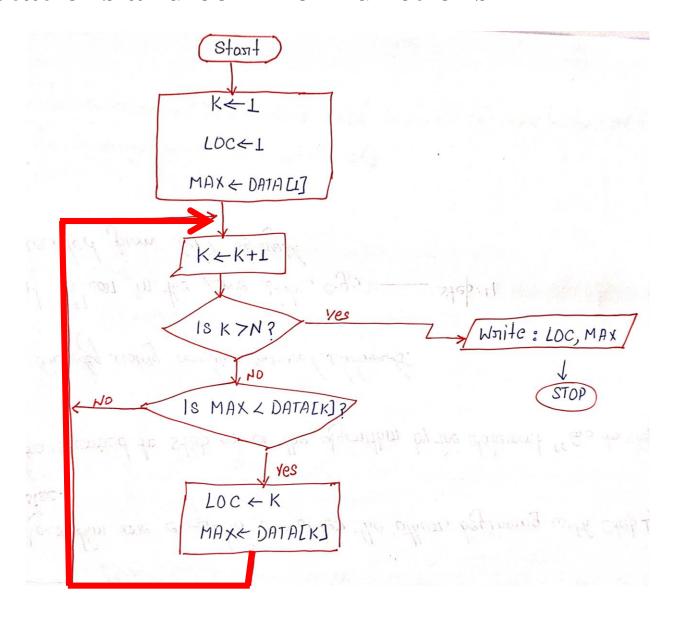
Step 3. [Test countern.] IF K>N, then:

Write: LOC, MAX, and Exit

Step 4. [compane and update.] IF MAX < DATA[K], then:

Set LOC: = K and MAX : = DATA[K].

Step 5. [Repeat Loop.] Go to step 2.



Steps, Control, Exit
The steps of the algorithm one executed one after the other, beginning with step 1, unless indicated otherwise.
Cg. Steb 5.
on so susting certain control statements.
* IF several statement appear in the same step, c.g., step:4
* IF several statement appear in the same step, c.g., step:4 They were executed from left to right.
* Exit completion

Comments Each step may contain or comment in brackets which indicates the main purpose of the step.
Usually appear at the begining on the end of the step.

Vaniable Names)

- * Voniables names will use capital letters, as MAX and DATA single letter names of voniables used as counters on subscript will also captalized in algority.
- * Lover cose con be used if a designate many a come somes a present prop

Assignment statement

went the pascal Opposition Read vanishles names (mant)

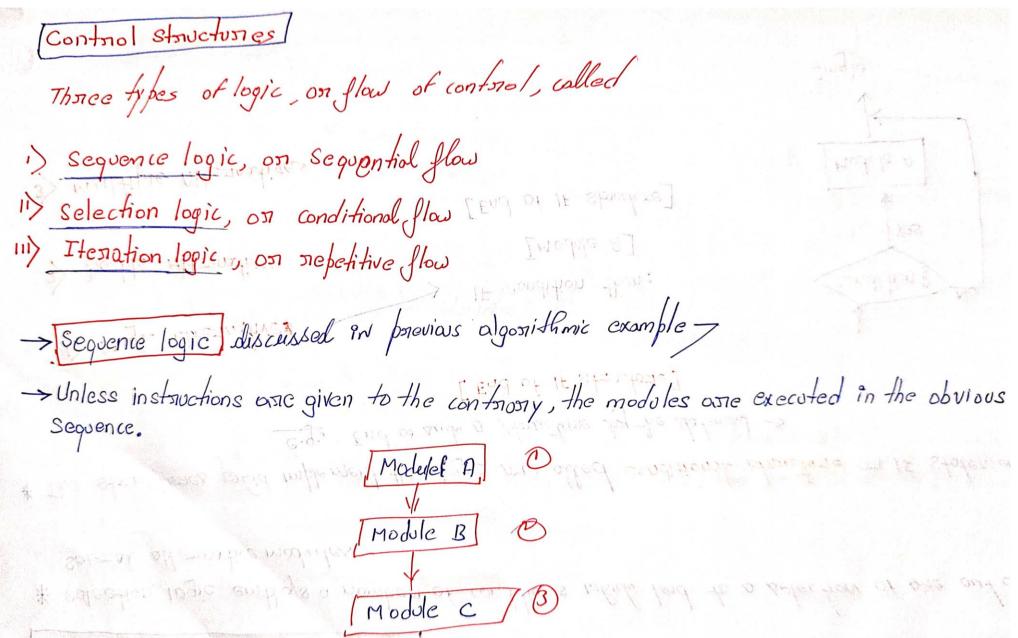
Input and output Input: Read Variables names (scanf) Output: Write Messages and/on variable names

Procedures

-> Used for an independent algorithmic module which solves a porticular problem.

Procedures / Modules, / Algorithm -> Interchangoble

Specific



Times times of imple on first of contract course

Selection Logic (Conditional Flow)

- * selection logic employs a number of conditions which lead to a selection of one out of Several alternative modules.
- * The structures which implement this logic one collect conditional structure on IF statement. C.g. End of such a structure by the statement ->

[End of IF structure]

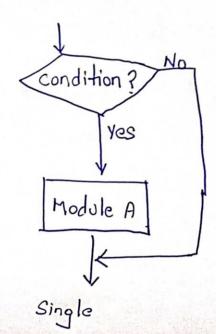
Single Alternatives of house complete

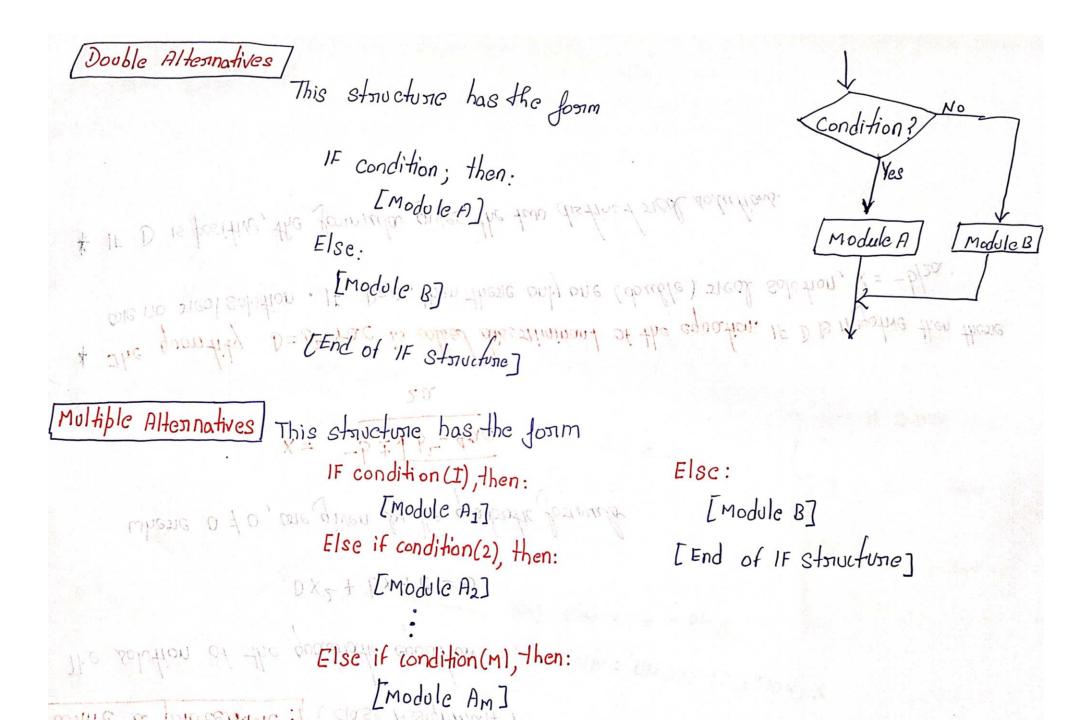
Double, Alternatives [module A]

IF condition, then:

[End of It stanctime]

3> multiple Alternatives





To be done by the student once...

Standard Notations and common functions

Where
$$a \neq 0$$
, are given by the quadratic formula

The solution of the quadratic equation(1) then:

where $a \neq 0$, are given by the quadratic formula

It countition(1) then:

Else:

To be done by the stude

The solution of the quadratic equation(1) then:

[End of the stude of the stude of the stude of the state of

The guarantity $D=b^2$ -fac is colled disconiminant of the equation. If D is negative then there are no neal solution. If D=0, then there only one (double) neal solution, x=-b/2a.

* IF D is positive, the formula gives the two distinct ned solutions.

Ir condition; then

Quality