

Lecture 2

Relational model & keys

Learning Objectives

- Understand relational model fundamentals
- Master different types of keys
- Learn ER modeling techniques
- Apply normalization (1NF-3NF, BCNF)

What is the Relational Model?

A data model based on relations (tables) where data is organized in rows and columns

Key Components:

- Relation (Table): Collection of related data entries
- Tuple (Row): Single data record
- Attribute (Column): Data field with specific meaning
- Domain: Set of valid values for an attribute

Example: University Database

- Student Table

StudentID	Name	Email	Major	GPA
12345	Alice Johnson	alice@edu.com	Computer Science	3.8
12346	Bob Smith	bob@edu.com	Mathematics	3.6
12347	Carol Davis	carol@edu.com	Physics	3.9

Relational Model Properties

1. Unique Relation Names

Each table must have a distinct name in the database

2. Atomic Attribute Values

Each cell contains single, indivisible values

✗ Phone: "123-456-7890, 987-654-3210»

✓ Separate PhoneNumber table

3. No Duplicate Tuples

Each row must be unique

4. Order Independence

Row order doesn't matter: (Alice, Bob, Carol) = (Bob, Carol, Alice)

Column order doesn't matter in logical sense

Example of Violations

BAD DESIGN:

Student(ID, Name, Contacts)

(1, "John", "john@email.com, 555-1234")

GOOD DESIGN:

Student(ID, Name)

Contact(StudentID, ContactType, ContactValue)

Understanding Keys - Superkey

Superkey: Any combination of attributes that uniquely identifies each tuple

Example:

Student Relation

Student(StudentID, SSN, Email, Name, Major, GPA)

Superkeys:

- {StudentID} ← Unique student identifier
- {SSN} ← Social Security Number
- {Email} ← University email is unique
- {StudentID, Name} ← More than needed but still unique
- {SSN, Major} ← More than needed but still unique
- {StudentID, SSN, Email} ← Definitely more than needed!

Key Point: If you can uniquely identify every student using these attributes, it's a superkey!

Understanding Keys - Candidate Key

Candidate Key: Minimal superkey (no proper subset is also a superkey)

From Previous Example:

Student(StudentID, SSN, Email, Name, Major, GPA)

Candidate Keys (minimal superkeys):

- {StudentID} ✓ Minimal - removing it breaks uniqueness
- {SSN} ✓ Minimal - removing it breaks uniqueness
- {Email} ✓ Minimal - removing it breaks uniqueness

Not Candidate Keys:

- {StudentID, Name} ✗ Not minimal
- Name is unnecessary - {SSN, Major} ✗ Not minimal
- Major is unnecessary

Understanding Keys - Primary & Foreign

**Primary Key: Selected
candidate key for the relation**

- Only ONE per relation
- Cannot be NULL
- Should be stable (rarely changes)

Foreign Key:

**Attribute that references primary
key of another relation**

Understanding Keys - Primary & Foreign

Example:

Student (StudentID, Name, Email, Major) - **Primary key**

Course (CourseID, Title, Credits, Department) - **Primary key**

Enrollment (StudentID, CourseID, Semester, Grade) - **Foreign keys**

StudentID references Student(StudentID)

CourseID references Course(CourseID)

Key Constraints in Action

Real University Database Example:

-- Primary Keys ensure uniqueness
Student(**StudentID**, Name, Email,
Major)

PK: StudentID

Course(**CourseID**, Title, Credits,
Department)

PK: CourseID

Professor(**ProfID**, Name, Department,
Salary)

PK: ProfID

Enrollment(StudentID, CourseID,
Semester, Grade)

Teaching(ProfID, CourseID, Semester,
Classroom)

Key Constraints in Action

Real University Database Example:

-- Primary Keys ensure uniqueness

Student(**StudentID**, Name, Email,
Major)

PK: StudentID

Course(**CourseID**, Title, Credits,
Department)

PK: CourseID

Professor(**ProfID**, Name, Department,
Salary)

PK: ProfID

-- Composite Primary Key

Enrollment(**StudentID**, **CourseID**, **Semester**,
Grade)

PK: (StudentID, CourseID, Semester)

FK: StudentID → Student(StudentID)

FK: CourseID → Course(CourseID)

Teaching(**ProfID**, **CourseID**, **Semester**,
Classroom)

PK: (ProfID, CourseID, Semester)

FK: ProfID → Professor(ProfID)

FK: CourseID → Course(CourseID)

ER Model Introduction

Entity-Relationship Model: Conceptual design tool for database structure

Main Components:

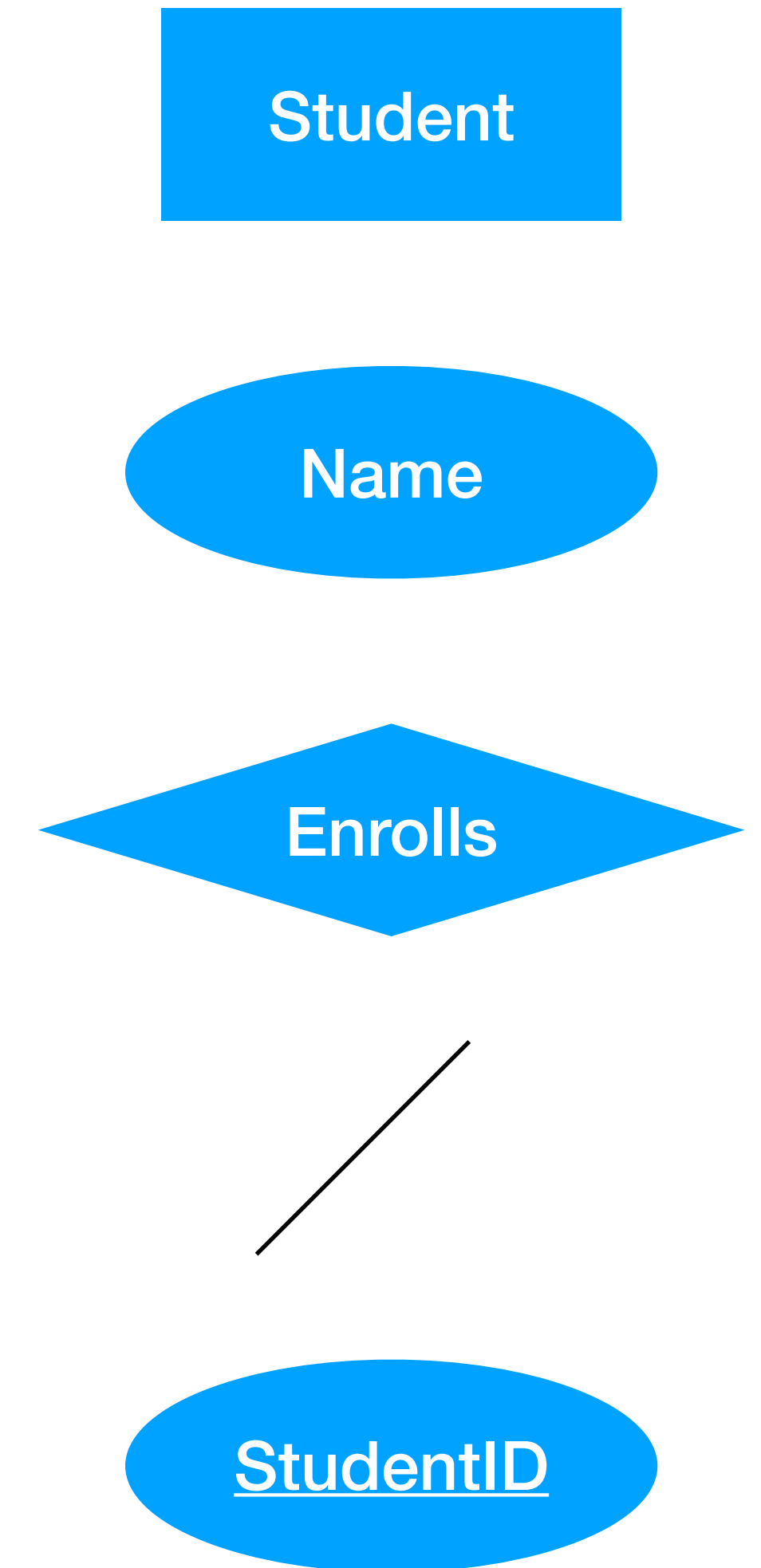
1. Entities - Things in the real world
2. Attributes - Properties of entities
3. Relationships - Associations between entities

Why Use ER Models?

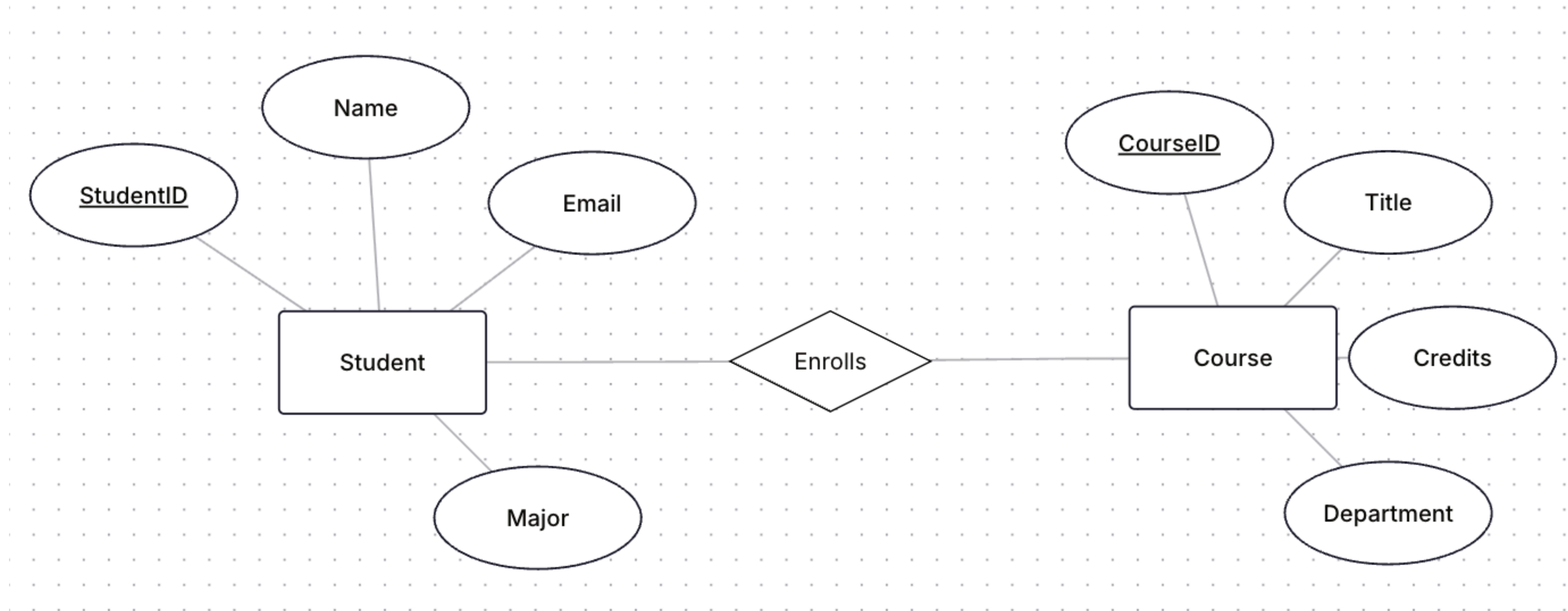
- Visual representation of database structure
- Communication tool between stakeholders
- Foundation for relational database design - Helps identify business rules and constraints

ER Notation

- Rectangles: Entities
- Ovals: Attributes (double oval = multi-valued, dashed = derived)
- Diamonds: Relationships
- Lines: Connect components
- Underlined: Primary key attributes



ER Notation



ER Model - Entities

Entity: Real-world object with independent existence

Types of Entities:

1. **Strong Entity:** Independent existence

Examples: Student, Course, Professor, Department

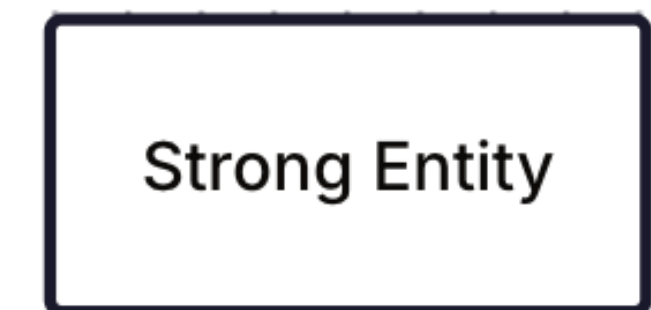
2. **Weak Entity:** Depends on another entity for identification

Examples:

- Dependent (depends on Employee)
- Room (depends on Building)
- OrderItem (depends on Order)

ER Notation:

Strong Entity: Rectangle



Weak Entity: Double Rectangle



ER Model - Attributes

Entity: Real-world object with independent existence

Attribute Types:

1. **Simple Attributes:** Atomic, indivisible Examples:
Name, Age, Salary, GPA
2. **Composite Attributes:** Can be divided into subparts
Address → Street, City, State, ZipCode
FullName → FirstName, MiddleName, LastName
3. **Multi-valued Attributes:** Multiple values allowed
PhoneNumbers: {555-1234, 555-5678, 555-9012}
Languages: {English, Spanish, French}
4. **Derived Attributes:** Calculated from other attributes
Age (derived from BirthDate and current date)
TotalSalary (derived from BaseSalary + Bonus)

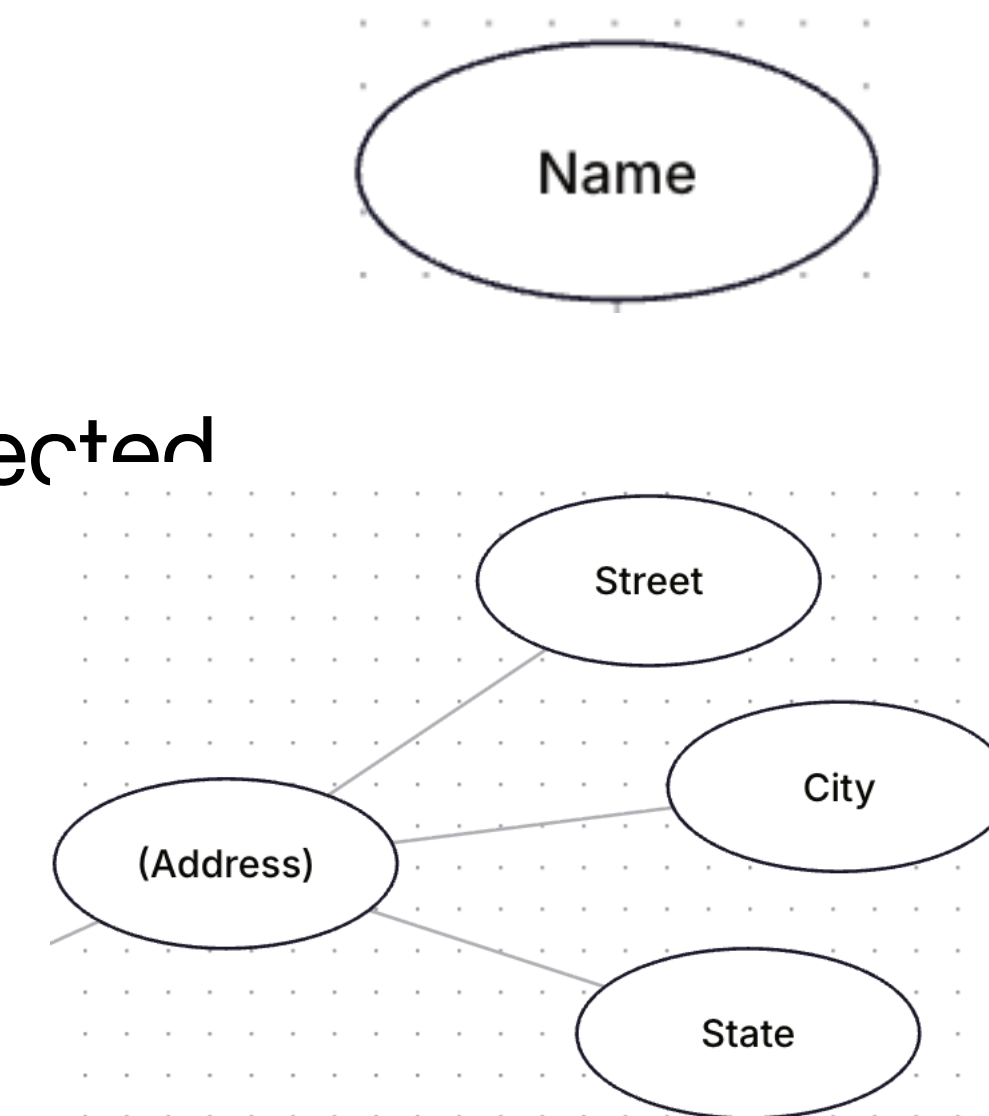
ER Notation:

Simple: Oval

Composite: Oval with connected sub-ovals

Multi-valued: Double Oval

Derived: Dashed Oval



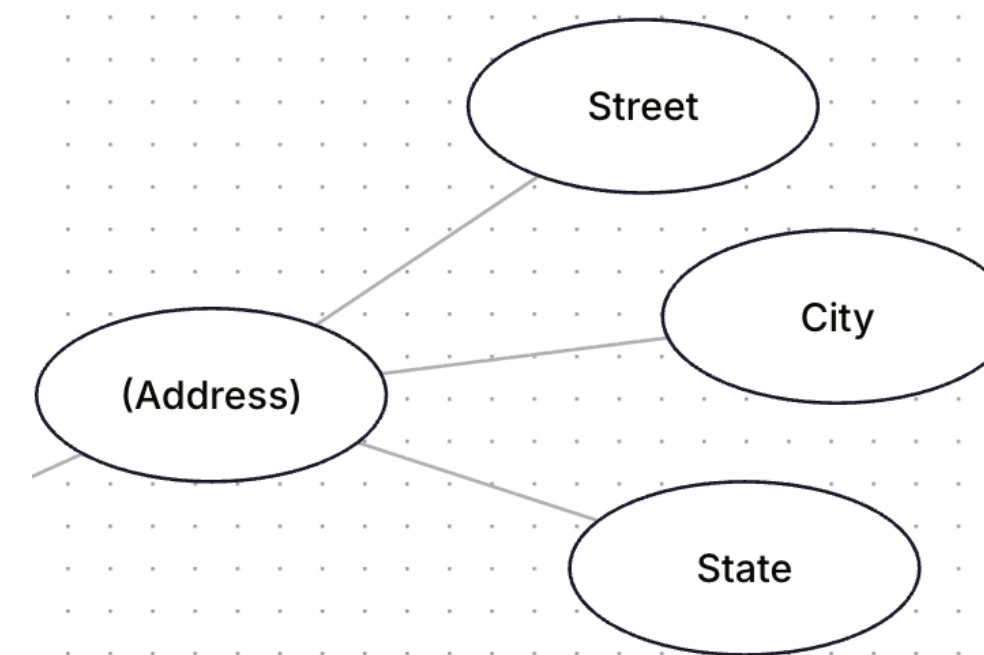
ER Model - Attributes

ER Notation:

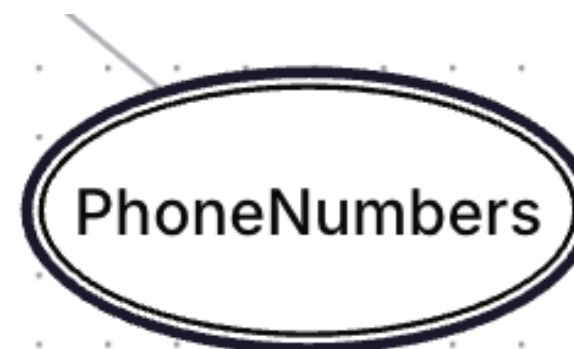
Simple: Oval



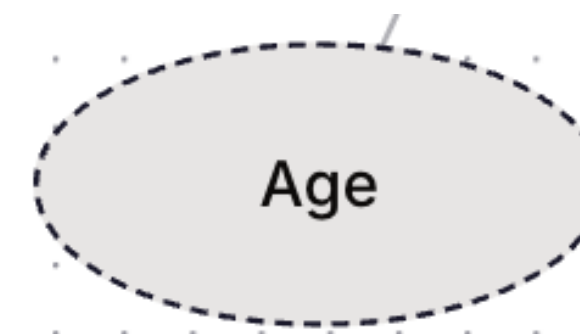
Composite: Oval with connected sub-ovals



Multi-valued: Double Oval



Derived: Dashed Oval



ER Model - Relationships

Relationship: Association between entities

Types by Degree:

- Binary: Between 2 entities (most common)
- Ternary: Between 3 entities (Student TAKES Course FROM Professor)
- n-ary: Between n entities

Cardinality Constraints:

- 1:1 (One-to-One)
- 1:N (One-to-Many)
- M:N (Many-to-Many)

Student — ENROLLS — Course
(0,N) (0,N)

↑ ↑
min,max participation

Professor — TEACHES — Course
(1,N) (0,1)

Department — HAS — Professor
(1,1) (5,50)

Normalization

Normalization - Why Do We Need It?

Problems with Unnormalized Tables:

Example: StudentCourseInfo Table

StudentID	Name	Major	CourseID	Title	Credits	Grade
1001	John	Computer Science	CS101	Intro	3	A
1001	John	Computer Science	CS102	00P	4	B
1002	Mary	Mathematics	CS101	Intro	3	A
1002	Mary	Mathematics	MATH201	Calculus	4	A

Problems:

- 1. **Redundancy:** John's info repeated for each course
- 2. **Update Anomaly:** Change John's major → update multiple rows
- 3. **Insert Anomaly:** Can't add new student without enrollment
- 4. **Delete Anomaly:** Delete last enrollment → lose student info

First Normal Form (1NF)

Rule: All attributes must contain atomic (indivisible) values

Violation Examples:

✗ Multi-valued Attributes:

Student (ID, Name, **PhoneNumbers**, **Skills**)

(1, "John", "555-1234, 555-5678",
"Java, Python, SQL")

✗ Composite Attributes:

Employee (ID, Name, **Address**, Department)

(101, "Alice", "123 Main St, Boston, MA
02101", "CS")

1NF Solutions:

✓ Separate Multi-valued:

Student (ID, Name)

StudentPhone (StudentID, PhoneNumber)

StudentSkill (StudentID, Skill)

✓ Break Composite:

Employee (ID, Name, **Street, City, State, Zip**, Department)

Second Normal Form (2NF)

Rule: Must be in 1NF + No partial dependencies

Partial Dependency: Non-key attribute depends on part of composite primary key

Violation Example:

✗ Enrollment (StudentID, CourseID,
StudentName, CourseName, Grade)

Primary Key: (StudentID, CourseID)

Partial Dependencies:

- StudentName depends only on StudentID
- CourseName depends only on CourseID
- Grade depends on full key ✓

2NF Solution:

✓ Student (StudentID, StudentName)

Course (CourseID, CourseName)

Enrollment (StudentID, CourseID, Grade)

Key Point: If primary key is single attribute, you automatically satisfy 2NF!

Third Normal Form (3NF)

Rule: Must be in 2NF + No transitive dependencies

Transitive Dependency: $A \rightarrow B \rightarrow C$ (where A is key, C is non-prime)

Violation Example:

✗ `Employee(EmpID, Name, Department, DeptHead, DeptBudget)`

Dependencies:

`EmpID → Department` (direct dependency ✓)

`Department → DeptHead` (transitive dependency ✗)

`Department → DeptBudget` (transitive dependency ✗)

Chain: `EmpID → Department → DeptHead`

3NF Solution:

✓ `Employee(EmpID, Name, Department)`
`Department(Department, DeptHead, DeptBudget)`

Another Example:

✗ `Order(OrderID, CustomerID, CustomerName, CustomerCity)`

`OrderID → CustomerID → CustomerName`

`OrderID → CustomerID → CustomerCity`

✓ `Order(OrderID, CustomerID)`
`Customer(CustomerID, CustomerName, CustomerCity)`

Boyce-Codd Normal Form (BCNF)

Rule: For every functional dependency $A \rightarrow B$, A must be a superkey

When 3NF \neq BCNF: Overlapping candidate keys

`StudentAdvisor(StudentID, Major, Advisor)`

Business Rules:

- Each student has one major
- Each advisor works in only one major -
- Students can have multiple advisors in their major

Functional Dependencies:

- $(\text{StudentID}, \text{Advisor}) \rightarrow \text{Major}$
- $(\text{StudentID}, \text{Major}) \rightarrow \text{Advisor}$
- $\text{Advisor} \rightarrow \text{Major} \leftarrow$ This violates BCNF!

Candidate Keys: $\{\text{StudentID}, \text{Advisor}\}$, $\{\text{StudentID}, \text{Major}\}$

BCNF Solution:

`StudentMajor(StudentID, Major)`

`AdvisorMajor(Advisor, Major)`

`StudentAdvisor(StudentID, Advisor)`

Normalization Practice - Library System

Given Unnormalized Table:

LibraryRecord(BookID, Title, AuthorName,
AuthorCountry, PublisherName, PublisherCity,
StudentID, StudentName, StudentMajor,
CheckoutDate, DueDate, FineAmount)

Sample Data:

(B1, "Database Systems", "Elmasri", "USA",
"Pearson", "Boston", S1, "John", "CS",
"2024-01-15", "2024-02-15", 5.00)

Step 1: Identify Functional Dependencies

BookID → Title, AuthorName,
AuthorCountry, PublisherName,
PublisherCity

StudentID → StudentName, StudentMajor

(StudentID, BookID, CheckoutDate) →
DueDate, FineAmount

AuthorName → AuthorCountry (assuming
each author from one country)

PublisherName → PublisherCity

Normalization Practice - Solution

2NF Decomposition:

BookInfo(BookID, Title, AuthorName, AuthorCountry, PublisherName, PublisherCity)

StudentInfo(StudentID, StudentName, StudentMajor)

CheckoutRecord(StudentID, BookID, CheckoutDate, DueDate, FineAmount)

3NF Decomposition:

Book(BookID, Title, AuthorName, PublisherName)

Author(AuthorName, AuthorCountry)

Publisher(PublisherName, PublisherCity)

Student(StudentID, StudentName, StudentMajor)

Checkout(StudentID, BookID, CheckoutDate, DueDate, FineAmount)

Summary & Key Takeaways

Relational Model:

- Tables (relations) with rows (tuples) and columns (attributes)
- Various types of keys ensure data integrity

ER Modeling:

- Visual design tool for database structure
- Entities, attributes, and relationships
- Foundation for relational database design

Normalization Benefits:

- 1NF: Atomic values
- 2NF: No partial dependencies
- 3NF: No transitive dependencies
- BCNF: All dependencies from superkeys

Key Skills Developed:

- ✓ Identify different types of keys
- ✓ Create ER diagrams from requirements
- ✓ Map ER diagrams to relational schemas
- ✓ Apply normalization techniques
- ✓ Recognize and fix data anomalies