Data Analysis and Algorithm

Practical 2

Write a Program to implement Merge sort

&

Find the run time of the algorithm

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maite a Program to implement Menge sont & find the granny time of the orlgorithm 7 For algorithms such as selection sourt or insertion sourt when we increase the infert size the time taken also increase by a lig margin.

Thus we look at algorithm like menge sont which has I nontine of aloga in all cases ile with large then insention sort. Meage so at ves the Divide & Congressively a large appray into smaller section particularly 2 equal sections & southing those section independently & secursively & at last combining.

Thus, the algorithm has 3 parts. Divide, Conquen & combine & you get two annays annEppergrange, Ptgs

Trombing - Menge the souled gulanting where on steps 1-4 me momente & of Rentime of merge sout algorithm neage sont we look at the same -7 We consider the Pivide stop to have a constant time because in the stop all we do is calculate the midpoint of amay. to take O(n) time as it will act on all the element once. Thus divide & combine step together take O(2) on on on for some constant Twe consider the nuntine of mergosort of n elements as som of Planents array added with on it divides -7 Similarly the suntime of mengo sout of (N2) elements will be sun of twice the suntime of mengeront of (n/4) elements

Thuy we see, n/2 n/22 · cn/2 = (n b. Caln = Ca 0/4 1/4 1/4 1/4 MARIE SELECTER RES ESPIL $\frac{1}{1} \frac{1}{1} \frac{1}$ Constant. C110117 = 1777 = 1000 ive total time of menge sont will be sum of mengeing time of each level. Whene Lis given as

L = log_n + 1

Cg:- n=8

L = log_2 8+1

L = 4

i.e. for n=8 we have 4 levels. thus we suitskitely legant 1 in do on we get (ch (logant 1)

For Dig O notation we disand constants a low order ten (to)

thus get O(n (ogan))

Algorithm for metge sout The Meage sout is implement with the use of specularity functions.

I amm = input

BEGG = 0.0

ENCO = 220 Learns. END - ann. length.

menge Sont Cann, BEG, END) menge sont (ann, DEG, END) mid = (BEG + END)/2 Menge Sont (ann, BEG, MID)

Menge Sont (ann, MID+1, END)

Menge (Dann, BEG, MID, END) 3) Menge Cann, BEG, MIO, END)

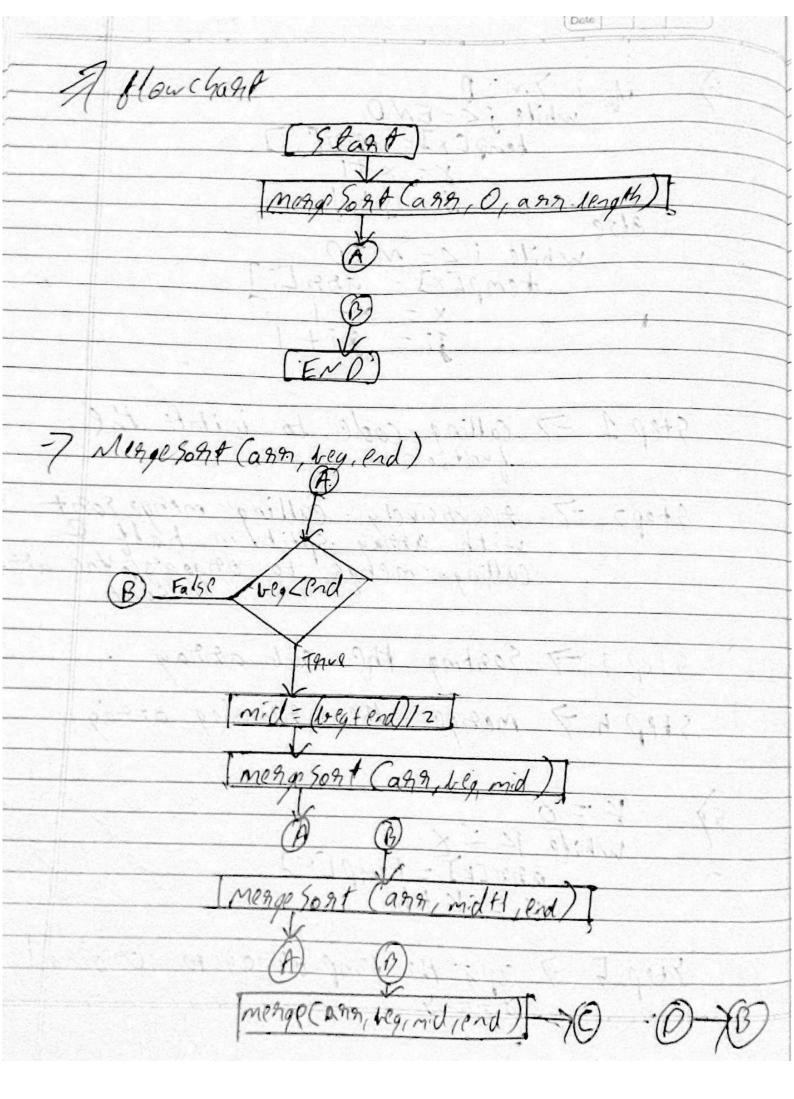
i = BEG

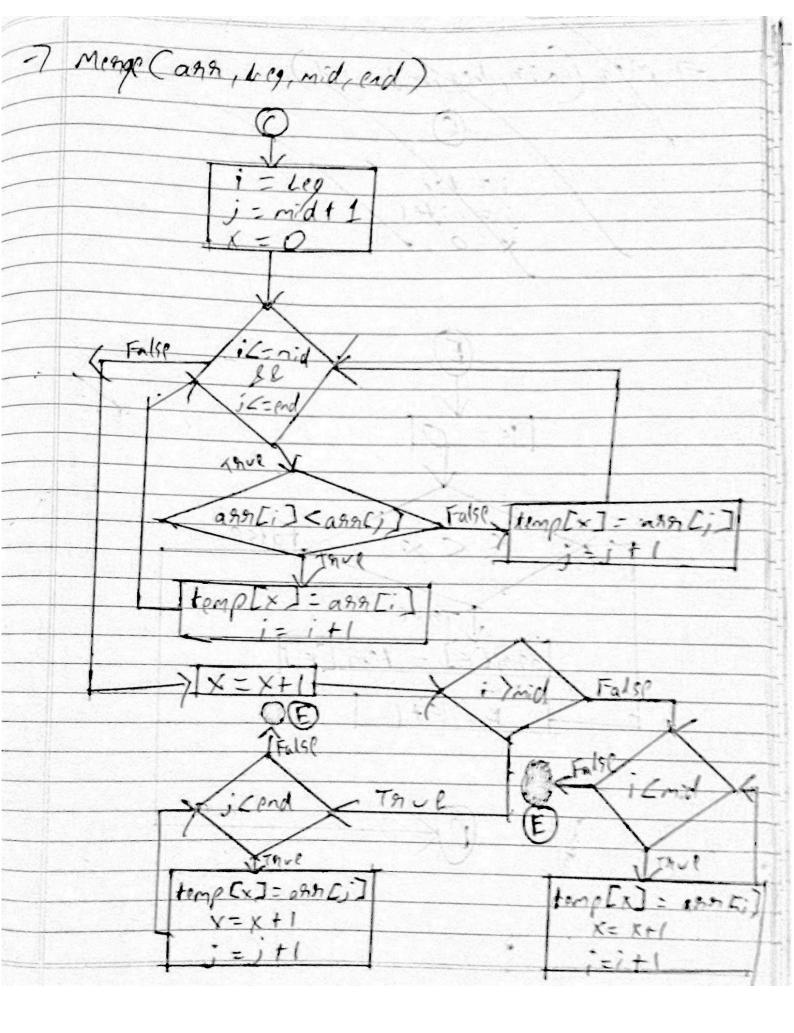
i = MID+1

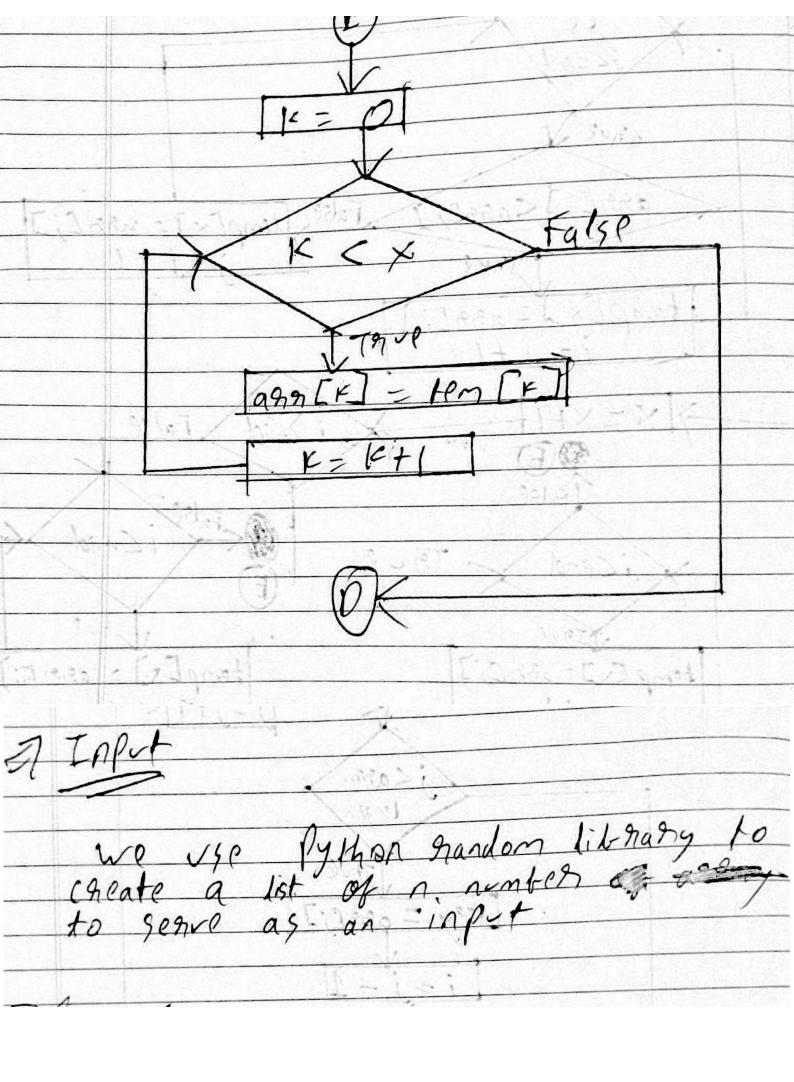
x = 0 While (i Z=MID 28 j Z=ENO) if ann[i] < ann [i]

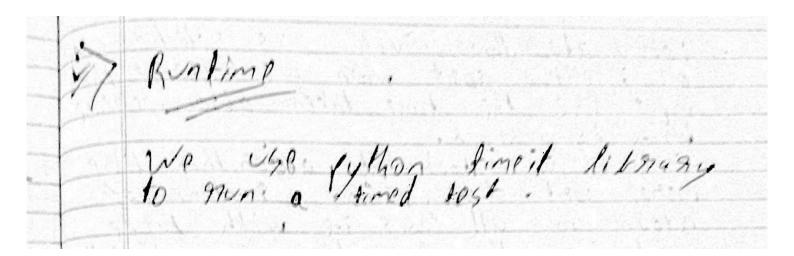
bemp[x] = ann[i] Plse temp[X] = ann [j] x = x+1

if i >MID while j <= ENO temp[x] = ann[;] X=x+1 j = j + j2150 while i K = MID temp[x] = ann[i] X = X+1 1 = 1 + 1 of calling code to initate the Step 2 7 specksively calling metage sont with annay split in half 2 culling metage to organize the assey Step 3 -7 Souting the sub atmay Steph 7 menge the hemaning annay K=0 While KXX ann(K) = temp(K) K= K+1 7+1p 5 -7 copy the temp atmay to original









Program

```
import random
    def merge(arr, 1, m, r):
3.
        n1 = m - 1 + 1
4.
        n2 = r - m
5.
6.
        # create temp arrays
7.
        L = [0] * (n1)
8.
        R = [0] * (n2)
9.
10.
        # Copy data to temp arrays L[] and R[]
11.
        for i in range(0, n1):
            L[i] = arr[l + i]
12.
13.
        for j in range(0, n2):
14.
15.
            R[j] = arr[m + 1 + j]
16.
17.
        # Merge the temp arrays back into arr[1..r]
                  # Initial index of first subarray
18.
        i = 0
                  # Initial index of second subarray
19.
        j = 0
20.
                  # Initial index of merged subarray
21.
22.
        while i < n1 and j < n2:
            if L[i] <= R[j]:</pre>
23.
24.
                arr[k] = L[i]
25.
                i += 1
26.
            else:
                arr[k] = R[j]
27.
28.
                j += 1
29.
            k += 1
30.
31.
        # Copy the remaining elements of L[], if there
        # are any
32.
33.
        while i < n1:
            arr[k] = L[i]
34.
35.
            i += 1
            k += 1
36.
37.
        # Copy the remaining elements of R[], if there
38.
39.
        # are any
40.
        while j < n2:
41.
            arr[k] = R[j]
42.
            j += 1
43.
            k += 1
45. # l is for left index and r is right index of the
46. # sub-array of arr to be sorted
47. def mergeSort(arr, 1, r):
48.
        if 1 < r:
49.
            # Same as (1+r)//2, but avoids overflow for
50.
            # large 1 and h
51.
            m = 1+(r-1)//2
52.
            # Sort first and second halves
```

```
mergeSort(arr, 1, m)
53.
54.
            mergeSort(arr, m+1, r)
            merge(arr, 1, m, r)
55.
56.
            print(arr)
57.
         return(arr)
59. #driver code
60. if __name__ == "__main__":
61.
        arr=[]
62.
        for i in range (1,10):
63.
            n = random.randint(0,1000)
64.
            arr.append(n)
65.
        mergeSort(arr, 0, len(arr)-1)
66.
              print(mergeSort)
```

```
====== RESTART: C:\Users\ADMIN\Desktop\
      771, 932, 266, 393, 59, 781, 553,
[527,
      771, 932, 266, 393, 59, 781, 553,
      771, 932, 266, 393, 59,
                              781, 553,
[527,
      393, 527, 771, 932,
                          59, 781, 553,
[266,
      393, 527, 771, 932, 59, 781, 553,
[266, 393, 527, 771, 932, 59,
                              781, 553,
[266, 393, 527, 771, 932, 59, 553, 781,
[59, 266, 393, 527, 553, 771, 781, 932, 949]
[59, 266, 393, 527, 553, 771, 781, 932, 949]
>>>
```

Timing code

```
1. SETUP_CODE = '''
2. from __main__ import mergeSort,merge
3. import random
4.
5.
6. TEST_CODE = '''
7. arr=[]
8. input_size = 10
9. for i in range (1,input_size):
        n = random.randint(0,1000)
10.
       arr.append(n)
11.
12. mergeSort(arr, 0, len(arr)-1)'''
14. times = timeit.timeit(setup = SETUP_CODE,
15.
                           stmt = TEST_CODE,
16.
                           number = 1)
17.
18. print(times)
```

Output time with array of size 10

Output time with array of size 100000

```
0.8353626
```

Onclusion

We observe that merge soft take

very small amount of time when the apput

size in incressed as compared to other lineary

lalgorithms.