If kn is prime, then there is notwing more to prove. If not, then

It is composite: I integers b, c & 1 < b < C < k+1 and k+1 = b C. BY

HYPOTHESIS, b = p, p, ...p; AND C = q, q, ...qm ARE PRODUCTS OF PRIMES.

THEN k+1 = b C = p, p, ...p; q, q, ...qm is a product of primes.

:, every a ∈ N\ {i} is either prime or a product of primes.

Now we show untqueness, by compadition, Assume that a > 1

Is the product of primes in two different ways:

$$a = P_1 P_2 \cdots P_m = 9_1 9_2 \cdots 9_n$$

SINCE P, |a, EUCLID'S LEMMA SAYS P, DIVIDES ONE OF THE Q; WITHOUT LOSS OF GENGRATHY, LET P, |q, SINCE q, IS PRIME, ITS DIVISORS ARE I AND q, HENCE, P, =q, AND

 $\frac{q}{p_1} = p_2 p_3 \cdots p_m = p_2 p_3 \cdots p_n$

BY THE SAME LOGIC, P2 MUST DEVIDE ONE OF THE REMAINTING 9; WLOG P2/92.

 $\frac{a}{P_1 P_2} = P_3 P_4 \cdots P_m = P_3 P_4 \cdots P_n$

CONTINUENCE LIKE THIS, WE FIND THAT $M \leq n$ AND $p_i = q_i$ $\forall i \in \{1, 2, ..., m\}$. THE SAME ALGUMENT WITH THE p PRIMES AND q PAIMES REVERSED OFFVES US THAT $n \leq m$ AND $q_i = p_i$ $\forall i \in \{1, 2, ..., n\}$. THEREFORE, m = n AND WE HAVE THAT THE TWO FACTORIZATIONS ARE THE SAME.

EX: FIND THE PRIME FACTORIZATION.

a) 974 b) 1300 c) 2722 l) 50,193

GREATEST COMMON DIVISOR (GCO)

DEF: LET a, b E IL WITH AT LEAST ONE OF a, b NONZERO. THE

GREATEST COMMON DIVISOR OF a AND b, DENOTED BY gcd(a,b),

IS THE NUMBER CEIN SUCH THAT

a) c IS A COMMON DIVISOR OF a AND b: c/a AND c/b;

b) IF & IS A COMMON OIVISOR OF a AND b, THEN & & C.

EX: gcd (18,12) =6, SINCE 6/18 AND 6/12, AND THERE IS NO BIGGER INTEGER THAT DIVIDES THEM BOTH.

NOTE THAT gcd.(18,12) = gcd(-18,12) = gcd(18,-12) = gcd(-18,-12).

PRIME FACTORIZATIONS CAN BE USED TO FIND GCOS. IF

$$a = \rho_i^{\alpha_i} \dots \rho_{\mu}^{\alpha_{\mu}}$$
 AND $b = \rho_i^{\beta_i} \dots \rho_{\mu}^{\beta_{\mu}}$ (Some α_i , β_i can be zero), Then $gcd(a_ib) = \rho_i^{\delta_i} \dots \rho_{\mu}^{\delta_{\mu}}$, where $\delta_i = min\{\lambda_i, \beta_i\}$.

EX: GIVEN THAT 3720 = 22.5.7.23 AND 1155 = 3.5.7.11, WE HAVE

gcd(3720,1155) = 5.7 = [35]

EX: FIND gcd (35,100,6975)

LEAST COMMON MULTIPLES (LCM)

- DEF: LET a, b & IL WITH AT LEAST ONE OF a, 6 NONZEDO, THE LEAST COMMON MULTIPLE OF a AND b, DENOTED BY | cm (a, b), IS THE MIMBER CEN SUCH THAT
 - 1) CIS A COMMONMULTIPLE OF a AMO b; i.e., a/c AND b/c;
 - 2) IF I IS A COMMON MULTIPLE OF a AND b, THEN C & d.

$$Ex: a)/cm(12,4) = 12$$

b) 1cm (18,15) = 90

WE CAN USE PRIME FACTORIZATION TO CALCULATE LCM. IF $a = \rho_{i}^{\alpha_{i}} \cdots \rho_{k}^{\alpha_{k}} \text{ and } b = \rho_{i}^{\beta_{i}} \cdots \rho_{k}^{\beta_{k}} \text{ (some α_{i}, β_{i} can $\beta \in zero), then}$ $|\operatorname{cm}(a,b) = \rho_{i}^{\beta_{i}} \cdots \rho_{k}^{\beta_{k}}, \text{ where } \delta_{i} = \max \{\alpha_{i}, \beta_{i}\}.$

Ex: GIVEN THAT 3220 = $2^2 \cdot 5 \cdot 7 \cdot 23$ AMD 1155 = $3 \cdot 5 \cdot 7 \cdot 11$, WE HAVE $|cm(3220, 1155) = 2^2 \cdot 3 \cdot 5 \cdot 7 \cdot 11 \cdot 23 = 106, 260$

EXEMETSE: FIND /cm (35100,6975). FIND gcd (268944,198466).

THE EUCLIDEM ALGORITHM

THE GUCLTDEAN ALGORITHM IS A PROCESS FOR PINDING GCD, IT WOLKS BECAUSE OF THE QUOTIENT-REMAINTER THEOREM AND THE FOLLOWING TWO LEMMAS,

LEMMA 1: FORALL rEN, gcd (r, 0) = r.

PAOOF: EXERCISE.

LEMMA 2: LET a, bet, b $\neq 0$, q, $r \in \mathbb{N}$ $\exists a = bq + r$. THEN gcd(a,b) = gcd(b,r).

PROOF: LET $D = \{d \in \mathbb{Z}: dla, dlb\}, \overline{D} = \{d \in \mathbb{Z}: dlb, dlr\}$. WE WILL SHOW THAT $D = \overline{O}$.

 $(\subseteq): LET + \in D$, THEN $\times |a|$ AND $\times |b|$, WE HAVE a = bq + r

=a-bq=r

SINCE $\frac{9}{x}$ AND $\frac{69}{x}$ AND WE HAVE $\int_{0}^{\infty} e^{7}L$, so x/Γ . HENCE, $x \in \overline{\mathbb{O}}$, AND WE HAVE $0 \subseteq \overline{\mathbb{O}}$.

(2): LET XED. THEN X/b AND X/P. WE HAVE

a=bq+r

 $\frac{a}{\lambda} = \frac{bg + r}{2}$

SINCE & AND & ARE INTEGERS, WE HAVE & EL, SO + 1a.

HENCE, LED, AND WE HAVE DED.

THEREFORE, D = D. SO EVERY COMMON DIVISOR OF a AM b FS ALSO A COMMON DIVISOR OF b AND P, AND VICE VERSA.

(igcd(a,b) = max d = max d = gcd(b,r),

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EUCLIDEAN ALGORITHM; 1) Let a < b 50. 2) CHECK-IF b=0, IF SO, LEMMA (SAYS gcd (a,b) = a. 3) IF 6 to, USE QUOTIENT-REMAINDER THEOREM TO FIND Q, I WITH O = r L b SUCH THAT a = bq +r. LEMMA 2 SAYS gcd(a,b) = gcd(b,r). 4) SET a=b, b=r AND GO TO STEP 2. THIS ALGORITHM WILL TERMINATE WITH 1=0, STICE EACH REMAINDER IS SMALLER THAN THE PREVIOUS ONE. Ex: FINO gcd (2772,2710) 2772 = 2310 · 1 + 462 2310 = 462.5 +0 inged (2772, 2310) = 462. EX: FIND gcd (-243,223) (REMEMBER god (a, b) = god (101, 161). 243 = 227 . 1+20 223=20.11+3 20 = 3.6+2

DEF; INTEGERS a, bARE CALLED COPPLINE (RELATIVELY PRIME, MUTUALLY PRIME) IF gcd(a,b) = 1.

EXERCISE: TRUE OR FALSE? FOR ALL XEIN THERE EXISTS YEN SUCH THAT gcd (x,y) = 1.

THM (BÉZOUT'S IDENTITY): LET a, b ETL\ [0], THEN d = gcd(a,b)
EXISTS, AND THERE EXIST M, N ETL SUCH THAT Ma + n b = d.

COPOLLARY! IF a AND b ARE RELATIVELY PRIME, THEN THERE EXIST M, N & ZL SUCH THAT Ma+nb=1.

HOW DO WE FIND M, N? WE USE THE EUCLEDEAN ALGORITHM IN LEVERSE.

Ex: FIND M, N & 7L SUCH THAT gcd (330, 156) = 330m + 1561.

O 330 = 156.2 + 18

156 = 18.8 +12

3 18 = 12.1+6

912=6.2+0 => gcd (330,156)=6.

NOW STARTING WITH THE SECOND-LAST LINE, ISOLATE THE GED AND USE EACH PREVIOUS LINE TO SUBSTITUTE FOR THE OTHER FACTORS.

3 18=12.1+6=76=18-12

(1) 12 = 156 - 18.8 => 6 = 18 - (156 - 18.8) = -156 + 18.9

(1) 18 = 330 -156.2 => 6 = -156 + (330-156.2).9 = 330.98 - 156.19

i. m = 9, n = -19.

Ex: FIM
$$M, N \in \mathbb{Z}$$
 Such THAT $243m + 213m = 1$.
 $243 = 223 \cdot 1 + 20$
 $223 = 20 \cdot 11 + 3$
 $20 = 3 \cdot 6 + 2$
 $3 = 2 \cdot 1 + 1$
 $\Rightarrow 1 = 3 - 2$

$$= 7 - 3 - 2$$

$$= 3 - (20 - 3 \cdot 6) = -20 + 3 \cdot 7$$

$$= -20 + (223 - 20 \cdot 11) \cdot 7 = 223 \cdot 7 - 20 \cdot 76$$

$$= 223 \cdot 7 - (243 - 223) \cdot 78 = -243 \cdot 78 + 223 \cdot 85$$

$$\therefore m = -77, n = 85.$$

THE PIGEONHOLE PRINCIPLE! LET K, NEW, K<N. IF

N PIGEONS FLY INTO K PIGEONHOLES, THEN SOME PIGEONHOLE

CONTAINS AT LEAST TWO PIGEONS.

PROOF: SUPPOSE THAT EACH PIGEONHOLE CONTAINS AT MOST

ONE PIGEON. THEN THE TOTAL NUMBER OF PIGEONS IS

AT MOST & 1 = K < 11, A CONTRABILITION. THEREFORE,

THERE EXISTS A PIGEONHOLE THAT CONTAINS MORE THAN

ONE PIGEON.

EX AMPLES :

- a) YOU HAVE A DRAWER FULL OF SOCKS, OF 3 DIFFERENT COLOURS. HOW MANY SOCKS MUST YOU PICK AT RAMOM TO BE SUPE YOU HAVE A MATCHENG PATR?
- A: 4. THE FIRST 3 COULD POSSIBLY BE ALL 3 DIFFERENT COLOUPS, BUT
 THE FOURTH WILL MATCH ONE OF THOSE (ORELSE THERE'S A PREVIOUS PAIR)

- b) IN A ROOM OF 367 PEOPLE (ALLOWING FOR LEAP YEAR), AT LEAST 2 OF THEM SHAPE A BIRTHDAY.
- C) HUMANS HAVE A MAXIMUM OF ABOUT 500,000 HATPLS, IS IT

 GUARAMEED THAT 2 RESTDENTS OF WOLLONGONG HAVE EXACTLY

 THE SAME NUMBER OF HATPLS? HOW ABOUT 2 RESTDENTS OF SYDNEY?

SOME FORMAL EQUIVALENT STATEMENTS TO THE PIGGONHOLE PRINCIPLE:

- 1) LET A BE ASET OF NELEMENTS. IF A IS PARTITIONED INTO K PATRWISE DISJOINT SUBSETS, WHERE K<N, THEN AT LEAST ONE SUBSET CONTAINS MORE THAN ONE ELEMENT.
- 2) A FUNCTION FROM ONE PINITE SET TO A SMALLGRISGT CANNOT BE ONE-TO-ONE. THERE MUST BE AT LEAST TWO ELEMENTS THAT MAP TO THE SAME POINT.
- EX: TNA GROUP OF 700 PEOPLE, MUST THERE BE TWO WHOSE FIRST NAMES HAVE THE SAME FIRST AND LAST LETTERS?
 - A: AT MOST 26 PEOPLE CAN HAVE DIFFERENT FIRST LETTERS, AND AT MOST 26 PEOPLE CAN HAVE DIFFERENT LAST LETTERS. SO AT MOST 26.26 = 676 PEOPLE CAN HAVE DIFFERENT EITHER FIRST OR LAST LETTERS. SO YES, A GROUP OF 700 HAS 2 THAT SHAPETHE SAME FIRST AMBLAST LETTERS.
- PROBLEMS OF THES SORT ENVOLVE FIGURING OUT HOW TO FORM THE PIGGONHOLES PROPERLY (HOW TO PARTITION THE SET).

- Ex: 5 DIFFERENT NUMBERS APLE SELECTED FROM THE SET

 S= {1,2,3,4,5,6,7,8}. SHOW THAT 2 OF THE SELECTED

 NUMBERS SUM TO 9.
- A: PARTITION THE SET INTO PATRS THAT SUM TO 9 (THE PIGGONHOLES):

 S = {1,8} U{2,7} U{3,6} U{4,5}. THERE ARE 4 SUBSETS, SO

 IT'S POSSIBLE TO SELECT 4 NUMBERS FROM S SUCH THAT ONLY

 ONE OF THEM BELONGS TO EACH SUBSET. CHOOSING 5 NUMBERS,

 BY THE PIGGONHOLE PATRICIPLE, RESULTS FN 2 NUMBERS CHOSEN

 BELONGING TO THE SAME SUBSET.
- EX: A FESTAMRANT SERVES 3 OTFFERENT SALADS, 6 OTFFERENT
 MATINS AND 4 DIFFERENT DESSERTS. HOW MANY PEOPLE MUST EAT
 THERE TO ENSURE THAT AT LEAST 2 OF THEM HAVE THE SAME MEAL
 - A: THERE ARE 3.6.4 DIFFERENT MEALS (THIS IS COMBINATORICS, MORE ON THIS LATER), SO THERE MUST BE 73 PEOPLE.
- GENERALIZED PIGEONHOLE PRINCIPLE: IF N PIGEONS FLY INTO K PIGEONHOLES, AND N>KM FOR SOME MEN, THEN SOME PIGEONHOLE CONTAINS AT LEAST M+1 PIGEONS.
- EX: SHOW THAT IN A GROUP OF 85 PEOPLE, THE FIRST NAME OF AT LEAST 4 OF THEM MUST START WITH THE SAME LETTER.
 - A: SS PIGEONS, 26 PIGEON HOLES, SS = 26.3+7. SO SS > 26.3 => M = 3, AND SOME PIGEON HOLE CONTAINS AT LEAST M+1=4 PIGEONS.
 - : AT LEAST 4 PEOPLE'S MAMES START WITH THE SAME LETTER.

EXERCISE: FIND THE MINIMUM NUMBER OF STUDENTS IN A CLASS TO BESURE 3 OF THEM ARE BORN IN THE SAME MONTH.

EX: WE WANT TO ASSIBN 70 STUDENTS TO 11 CLASSES SO THAT NO CLASS HAS MORE THAN 15 PEOPLE. SHOW THAT THERE MUST BE AT LEAST 3 CLASSES WITH 5 OR MORE PEOPLE.

A: ASSUME ONLY & CLASSES HAVE 5 OF MORE PEOPLE, AND SHOW A COMPADICITION.

THE BEST CASE IS THAT THOSE 2-CLASSES ARE FULL (LEAVING THE FEWEST POSSIBLE PEOPLE FOR THE OTHER CLASSES), IS STUDENTS EACH.
THEN 40 PEOPLE REMAIN, FOR 9 CLASSES. SINCE

40=9.4+4, M=4 AND THE PRINCIPLE SAYS AT LEASTONE OF THE 9 CLASSES HAS 5 PEOPLE OR MORE INIT.

MODULAR ARITHMETIC

DEF: LET NEW, a ETL. WE DEFINE a mod n TOBE THE REMAINDER WHEN a IS DIVIDED 134 N.

EX:

$$10 \mod 4 = 2$$
 (STINCE $\frac{10}{4} = 2 \text{ wITH REMAINDER } 2$)
 $18 \mod 3 = 0$
 $-8 \mod 6 = 4$
 $10 \mod 1 = 0$

DEF: a, b ARE CONGRUENT MODULON, WRITTEN Q = b (moda), IF n/(a-b). EQUIVALENTLY, a = b (modn) IFF a modn = b modn. EX: TRUE OR FALSE? a) 154 = 56 (mod 11) b) 7 = -9 (modf) a) 11 (154-56) \$ 11 | 98 FALSE. b) 8/(7-(-a)) => 8/16 TRUE. EX: FINO X SUCH THAT 12= x (mod 5) WE NEED 5 (12-x), SOX IS ANY OF [..., -8, -3, 2, 7, 12,...} EXERCISE! IF M = 0 (mod 2), WHAT CAN YOU SAY ABOUT M? THM (CONGRUENCE ARITHMETEC): LGT MAN, a, b, G, d & Z. IF a = c (modn) AND b = d (modn), THEN (1)(a+b)=(c+d)(modn);(2) (a-b) = (c-d) (modn); (3) ab = cd (modn); (4) a = c (modn) + m ∈ N. PROF! 1 (a-c) = a-c=1p, pell. n1(b-d)=)|b-d=ng,g∈7L. (1) a+b=(np+c)+(nq+d)=n(p+q)+c+d

(a+b)(modn) = [n(p+q)+c+d](modn) = (c+d)(modn)

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- (2) a-b = (np+c)-(nq+d) = n(p-q) + (c-d)(a-b)(mod n) = [n(p-q) + (c-d)](mod n) = (c-d)(mod n)
- (3) $ab = (np+c)(nq+d) = n^2pq + npd + nqc + cd = n(npq+pd+qc) + cd$ ab(modn) = [n(npq+pd+qc) + cd] (modn) = cd(modn)
- (4) INDUCTION.
 - a) m = 1! a' = c' (modn)
- EX: a) GIVEN THAT 2064=1715+349, FIMD 2064 (mod 17).
 - b) GIVEN THAT 713064 = 803. SH, FIND 713064 (mod 8).
 - c) FIND × SUCHTHAT 3 =× (mod 5).
- a) $1715 = 15 \pmod{17}$ AND $349 = 9 \pmod{17}$ $\Rightarrow (1715 + 349) = (15 + 9) \pmod{17} = 7 \pmod{17}$.
- b) 803 = 3 (mod 8) AND 88 = 0 (mod 8) => 803.88 = 3.0 (mod 8) = 0 (mod 8).
- c) $3^9 = 3^4 \cdot 3^4 \cdot 3 = 81 \cdot 81 \cdot 3$. $81 = 1 \pmod{5}$ $\Rightarrow 81 \cdot 81 \cdot 3 \pmod{5} = 1 \cdot 1 \cdot 3 \pmod{5} = 3 \pmod{5}$.

EXERCITE: FIND THE REMARKABLER WHEN 7 IS DINIBED BY 16.

THM (CANCELLATION LAW): LET NETL, a,b, CETL, IF gcd(a,n) = 1

PROOF: ab = ac (modn)

(ab-ac) = 0 (modn)

 $a(b-c) \equiv O(m o d n)$

= a(b-c) = KN FOR SOME KETL. BUT a AND N ARE

COPPINE BY (B), SO THE FACTOR OF N ON THE LEFT-HAND STOE

IS CONTAINED FN (b-c). HENCE,

 $(b-c) \equiv o \pmod{n}$ $b \equiv c \pmod{n}$

NOTE: (A) IS ESSENTIAL. COUNTEREXAMPLE:

60 = 90 (mod 15)

10.6 = 10.9 (mod 15), BUT

6 \$ 9 (mod 15).

SINCE 10 AND 15 ARE NOT COPRIME, THE CANCELLATION LAW DOES NOT APPLY.

EX: GIVEN 10904=32 (mod9), FIND THE SMALLEST XEIN SUCH THAT X = 1363 (mod 9).

A: NOTE THAT 10904 = 1363.8.

1363.8 = 4.8 (mod 9)

1363 = 4 (mol9) (SINCE 8 +9 ARE COPRIME)

:x=4.

(449, SO THERE CAN BE NO SMALLER CONGRUENCE 70.

CONGRUENCE CLASSES MODILO A

THE QUOTIENT-REMATINDER THEOREM GIVES US THE FOLLOWING.

FACT: LET NETL. EVERY INTEGER XETL IS CONGRUENT MODILO N TO EXACTLY ONE ELEMENT IN [0, 1, 2, ..., n.if.

THIS ALLOWS US TO GROUP INTEGERS ACCORDING TO THEIR DEALTHOURS AFTER DEVEDENGBY N.

DEF: LET DEIN. THE CONGRUENCE CLASS (RESTOUE) OF a & 7L MODULO N IS THE SET [a] = {x \in 72 : x \equinodn)}.

EX: WRITE THE CONGRUENCE CLASSES FOR N = 4. HOW MANY OF THEM AME THERE?

THM: LET NEN, THERE ARE EXACTLY NOTSTENCT CONGRUENCE CLASSES: [0], [1], ..., [n-1].

PROOF: FIRST, SHOW THAT NO TWO OFO, I,..., n-1 ARE CONGRUENT MODULO n. LET $0 \le a < b < a$, $a,b \in \mathbb{N}$. THEN $b-a \in \mathbb{N}$ AND b-a < a. THUS, $n \nmid (b-a)$, so $b \not\equiv a \pmod{a}$. THEREFORE, NO TWO OF 0,1,...,n-1 ARE CONGRUENT, AND WE HAVE THAT [0], [i],..., [n-i] ARE ALL DISTINCT RESIDUES.

NEXT, SHOW THAT EVERY \times ETL IS IN ONE OF THESE RESTRUES. THE QUOTIENT-REMAINDER THEOREM GIVES $\times = ng + r$, $o \le r < n$. So $r \in \{0,1,...,n-i\}$, and $x - r = ng = 7 \times = r \pmod{n}$. THEREFORE, EVERY \times ETL IS IN ONE OF [0], [1], ..., [n-i].

DEF: LGT NEW. THE COMPLETE SET OF RESTONES MODILION IS THE SET Zn = {[0], [1], ..., [n-1]}.

 $E \times : 7L_3 = \{ LOI, LII, L2I \}, SO IN 7L_3, WE HAVE$ [4] = [1]; [-1] = [2]; [30] = [0].

EX: IN Zn,

- a) [0] U[] U... U[n-1] =
- P) [0]U[i]U…U[u-i]=

OPERATIONS ON The

WE WANT TO DEFINE A DOITION AND MULTIPLICATION ON TLA. SINCE DIFFERENT NUMBERS CAN CIEVE THE SAME RESTRUCT, WE MUST BE CAMEFUL WITH THE DEFINITIONS.

THM: LET NEW. THE OPERATION +:

IS WELL-DEFINED ADDITION ON The, i.e. IF [a] = [c] AND [b] = [d],
THEN [a+b] = [c+d]. SIMILARLY, THE OPERATION .:

IS WELL-DEPINED MULTIPLIEATION ON The, i.e. IF [a] = [c] AND [b] = [d], THEN [ab] = [cd].

PROOF !

$$[a] = [c] \Rightarrow a = C \pmod{n} \Rightarrow \exists k, \in 7L \Rightarrow a = c + k, n.$$

$$a+b=(c+k,n)+(d+k+n)=c+d+n(k,+k+n)$$

$$ab = (C+k,n)(d+k+n) = cd+n(k+2)(d+n+k+k+2)$$

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EX: WAITE ADDITION AND MULTIPLIFEATION TABLES FOR 763.

+	[o]	E1]	口打
[0]			
[2]			

	0	[0]	[I]	[2]
	[0]			
ľ	[门			
	[2]			

PROPERTIES OF 7Ln

- 1) + AND · ARE CLOSED (BINARY) OPERATORS.
- 2) + AND . ARE COMMUTATIVE.
- 3) + AMD . APE ASSOCIATIVE.
- 4) · IS OISTRIBUTEVE OVER +.
- 5) IDENTITIES ARE [O] UMBERT, [I] UNDER ..
- 6) THE ADDITIVE INVERSE OF [x] IS [n-x].
- 7) MULTIPLICATIVE TOVERSES EXEST ONLY FOR X FIGCO (x, n) = 1.

THM: IF a AND n ARE COPRIME, THEN THERE EXTSTS $b \in \mathbb{Z}$ Such that $ab \equiv l \pmod{n}$, we call b the multiplicative to verse of a modulo n. b to unique modulo n. We write $b = q^{-l} \pmod{n}$.

PLOOF: CONSTOCA THE SET [0,a,la,3a,..., (n-1)a]. IF WE CAN SHOW

THAT THESE ARE ALL DISTINCT MODULO N, THEN EXACTLY ONE OF THEM

IS EQUAL to 1 (modn).

SUPPOSE THE CONTRARY: I c, d & N v sos, c, d < n > ca = da (modi), c fd

THEN (c-d) a = 0 (modin), so IKE/L > (c-d) a = Kn. But a AMD N

ARE COPPLIME, SO N (c-d). THIS IS A CONTRADICTION, SINCE

C AND d ARE DISTINCT NONNEGATIVE IMEGERS LESS THAN N. II

EX: FIM A MUTIPUT CATIVE INVERSE OF 43 MODULO 60.

A:WENEED X EN SUCH THAT 43x = 1 (mod 60), NOTICE THAT

43x MUST HAVE LAST DIGIT 1, SINCE 60 IS A MULTIPLE OF 10.

SO ANY X SUCH THAT 43x ENDS IN 1 HAS LAST DIGIT 7. THE

POSTABILITIES ARE 7, 17, 27, 37, 47, 57.

43.7=301=1 (mod 60).

: 43 = 7 (mod 60).

Ex: FIND 3 (mod 40).

A: BY THE SAME REASONANG AS ABOVE, THE POSSIBILITIES ARE 7, 17, 27, 37.

3.7=21=21 (mod 40)

3.17=51=11 (nod40)

3.27=81 = 1 (mod 40)

:. 3 -1 = 27 (mod 40).

CRYPTOGRAPHY IS THE STUDY OF METHODS FOR SENDING SECRET MESSAGES. THERE ARE MANY TECHNIQUES FOR ENCRYPTION AND DECRYPTION, ONE OF WHICH IS PUBLIC-KEY CRYPTOGRAPHY. THE METHOD USES BIG PRIME NUMBERS AND MODULAR ARITHMETIC. RSA IS ONE SUCH PUBLIC-KEY METHOD.

RSA

- 1) CHOOSE 2 LARGEPRIMES P, 9.
- 2) CHOOSE EELL THAT IS COPPINE WITH (P-1)(q-1).
- 3) CHOOSE d & 12 3 ed = 1 (mod (p-1)(q-1)).
- 4) THE PUBLIC IVEY IS (e, pg). THIS IS AVAILABLE TO EVERYONE FOR ENCRYPTION.
- 5) THE PRIVATE KEY IS (d, pg). THIS IS AVAILABLE ONLY TO THOSE WHO THE SEMBER WANTS TO BE ABLE TO DECRYPT.

ENCRYPTION STEP: LET THE MESSAGE TO BE ENCRYPTED BE METL,

0 < M < pq (A computer uses BINARY cole FOR EVERYTHING, SO

ENCRYPTING INTEGERS IS SUFFICIENT). THE ENCRYPTED MESSAGE

IS C = M (modpq).

DECRYPTION STEP: M IS RECOVERED BY $M = C \pmod{pq}.$

WE WILL NOT SEE THE PROOF. P AND Q ARE CHOSEN TO BE SEVERAL HUNDRED DIGITS LONG EACH, MAKING IT IMPOSSIBLE FOR A COMPUTER TO FIND THE FACTORS (P-1)(q-1) IN REASONABLE TIME, WE WILL SEE SOME EXAMPLES WITH SMALL PRIMES.

EX: LET A=1, B=2,..., Z=26, PUBLICKEY (3,55). ENCRYPT AND DECRYPT THE MESSAGE "HEY."

A: $pq = 55 \Rightarrow p = 5$, q = 11. C = 3, which is copains with (5-1)(11-1) = 40. The unencrypted MESSAGE IS 8525.

 $8^{3} = 64.8 \equiv 9.8 \pmod{55} = 72 \pmod{55} \equiv 17 \pmod{55}$ $5^{3} = 125 \equiv 15 \pmod{55}$

253=125.125 = 15.15 (mod 55)=225 (mod 55) = 5 (mod 55).

THE ENCRYPTED MESSAGE IS 17 15 5.

FROM A MENTOUS EXAMPLE, 3 (MOL40) = 27.

17²⁷ = 289¹³, 17 = 14¹³, 17 (mod 55) = 196⁶, 14.17 (mod 55) = 196⁶, 238 (mod 55) = 31⁶, 18 (mod 55) = 961³, 18 (mod 55) = 26³, 18 (mod 55)

=676.468 (mod 55) =16.28 (mod 55) = 448 (mod 55) =8 (mod 55)

SO THE DECAYPTED 17 IS 8, THE ORIGINAL "H".

SIMILARLY, 15 DECRYPTED IS 5, AND 5 DECRYPTED IS 25.

EXERCISE: DECRYPT THE MESSAGE 41 83 36 THAT WAS ENCRYPTED WITH PUBLIC KEY (5,91).

SET THEORY

A "SET" IS A LOOSELY-DEPINED COLLECTION OF ITEMS

CALLED "ELEMENTS". SETS ARE COMPLETELY DETERMINED BY THERR

GLEMENTS, i.e. TWO SETS WITH EXACTLY THE SAME ELEMENTS ARE

THE SAME SET. THE ORDER IN WHICH ELEMENTS ARE LISTED IS

IR RELEVANT, AND ELEMENTS MAY BE LISTED MORE THAN ONCE

WITHOUT CHANOTRO THE SET.

THE COLLECTION OF ALL PEOPLE IN THIS ROOM IS A SET.

THE COLLECTION OF YOUR FAVOURITE SONGS IS A SET.

THE COLLECTION OF ALL REAL MUMBERS IR IS A SET.

SETS COME FROM A UNIVERSE OF ELEMENTS U. FOR EXAMPLE, THE SET OF EVENNUMBERS COMES FROM THE UNIVERSE TL. SETS CAN BE CONTAINED IN OTHER SETS, AND CAN BE FINITE OR INFINITE.

SOME IMPORTANT SETS OF NUMBERS ARE

$$Q = \left\{\frac{a}{b} : a, b \in TL, b \neq 0\right\} \qquad (RATIONAL)$$

A SET CAN BE DEFINED BY A PROPERTY OF ELEMENTS OF A BIGGER SET, GIVEN A SET SET TBY

T = {+ ES: p(x)}, ALL THE GLEMENTS OF STHAT

EX: THE SET {+ER: -2<+ \leq 5} IS THE SET OF ALL REAL NUMBERS

BETWEEN -2 AND 5, NOT INCLUDING -2. THIS SET IS AN INTERVAL,

WHICH CAN BE DENOTED AS (-2, 5].

THE EMPTY SET IS THE SET WITH NO ELEMENTS, DENOTED BY Ø. IT

CAN BE REPRESENTED IN DIFFERENT WAYS:

[+EN:+ ++}; {+E12:3<+<2}.

A SET IS FINITE IF IN EN SUCH THAT THERE IS A ONE-TO-ONE CORPERDONDENCE WITH THE SET [1,2,..., n], FOR A SET S OF THIS SIZE, WE WRITE |S| = N AND SAY THAT S HAS CARDINALITY N. NOTE: |p| = 0.

A SET THAT IS NOT PINITE IS SAFO TO BE INFINITE.

SUBSETS

SATISFY P.

DEF: LET A AMD B BE SETS. WE SAY A IS A SUBSET OF B, WRITTEN

A SB, IFF EVERY EVERENT OF A IS ALSO AN ELEMENT OF B.

SYMBOLICALLY,

 $A \subseteq B \iff \forall x, x \in A \Rightarrow x \in B.$