

**MATH221: Mathematics for Computer Science**  
**Outline Solutions to Tutorial Sheet Week 3**  
Autumn 2017

1. (i) Note 1: “ $\vee$ ” is only False when both parts are False.  
Note 2: “ $\implies$ ” is only False when the first part is True and the second is False.

$$\begin{array}{ccccccc}
 (P & \implies & Q) & \vee & (P & \implies & \sim Q) \\
 & & & & \text{F} & & \\
 & \text{F} & & & & \text{F} & \text{(Note 1)} \\
 \text{T} & & \text{F} & & \text{T} & & \text{F} \quad \text{(Note 2)} \\
 & & & & & & \text{T} \quad \text{(Negation)}
 \end{array}$$

Therefore, we have a contradiction for  $Q$ . Thus the statement is a tautology.

- (ii) Note 1: “ $\vee$ ” is only False when both parts are False.  
Note 2: “ $\implies$ ” is only False when the first part is True and the second is False.

$$\begin{array}{ccccccc}
 \sim (P & \implies & Q) & \vee & (Q & \implies & P) \\
 & & & & \text{F} & & \\
 \text{F} & & & & & \text{F} & \text{(Note 1)} \\
 & \text{T} & & & & & \text{(Negation)} \\
 & & \text{T} & & \text{T} & & \text{F} \quad \text{(Note 2)} \\
 \text{F} & & \text{T} & & & & \text{(Transferring values across)}
 \end{array}$$

There is no contradiction. Therefore, the statement is NOT a tautology.

- (iii) Note 1: “ $\vee$ ” is only False when both parts are False.  
Note 2: “ $\implies$ ” is only False when the first part is True and the second is False.  
Note 3: “ $\wedge$ ” is only True when both parts are True.

$$\begin{array}{ccccccc}
 (P & \wedge & Q) & \implies & (\sim R & \vee & (P \implies Q)) \\
 & & & & \text{F} & & \\
 & \text{T} & & & \text{F} & & \text{(Note 2)} \\
 \text{T} & & \text{T} & & \text{F} & & \text{(Notes 3 \& 1)} \\
 & & & & \text{T} & & \text{F} \quad \text{(Note 2)}
 \end{array}$$

We have a contradiction for  $Q$ . Thus the statement is a tautology.

2. (i)  $((P \wedge Q) \implies R) \equiv (\sim (P \wedge Q) \vee R)$  [Thm 1.4.2, 6]  
 $\equiv ((\sim P \vee \sim Q) \vee R)$  [DeMorgan's Laws]  
 $\equiv (\sim P \vee \sim Q \vee R)$  [Thm 1.4.2, 2]

- (ii)  $(P \implies (P \vee Q)) \equiv (\sim P \vee (P \vee Q))$  [Thm 1.4.2, 6]  
 $\equiv (\sim P \vee P \vee Q)$  [Thm 1.4.2, 2]

3. (i) Firstly,  $(P \implies Q) \equiv (\sim P \vee Q)$  [Thm 1.4.2, 6]  
Thus,  $\sim (P \implies Q) \equiv \sim (\sim P \vee Q)$   
 $\equiv (\sim \sim P \wedge \sim Q)$  [DeMorgan's Laws, Thm 1.4.2, 5]  
 $\equiv (P \wedge \sim Q)$  [Thm 1.4.2, 4]

- (ii)  $((P \wedge \sim Q) \implies R) \equiv (\sim (P \wedge \sim Q) \vee R)$  [Thm 1.4.2, 6]  
 $\equiv ((\sim P \vee \sim \sim Q) \vee R)$  [DeMorgan's Laws, Thm 1.4.2, 5]  
 $\equiv ((\sim P \vee Q) \vee R)$  [Thm 1.4.2, 4]  
 $\equiv (\sim P \vee (Q \vee R))$  [Thm 1.4.2, 2]  
 $\equiv (P \implies (Q \vee R))$  [Thm 1.4.2, 6]

4. (i)  $P \implies (Q \vee P) \equiv (\sim P \vee (Q \vee P))$  [Thm 1.4.2, 6]  
 $\equiv (\sim P \vee Q \vee P)$  [Thm 1.4.2, 2]  
 $\equiv (\sim P \vee P \vee Q)$  [Thm 1.4.2, 1]

However,  $\sim P \vee P$  is a tautology, therefore, by the Conclusion, this statement is also a tautology, that is,  $P \implies (Q \vee P)$  is a tautology.

- (ii)  $(P \wedge Q) \implies (\sim R \vee (P \implies Q))$   
 $\equiv (\sim (P \wedge Q) \vee (\sim R \vee (P \implies Q)))$  [Thm 1.4.2, 6]  
 $\equiv ((\sim P \vee \sim Q) \vee \sim R \vee (\sim P \vee Q))$  [Thm 1.4.2, 5, 2 and 6]  
 $\equiv (\sim P \vee \sim Q \vee \sim R \vee \sim P \vee Q)$  [Thm 1.4.2, 2]  
 $\equiv (Q \vee \sim Q \vee \sim P \vee R)$  [Thm 1.4.2, 1]

As in part (i), we have a tautology as  $Q \vee \sim Q$  is a tautology.

5. (i) The proposition is true:  
 If  $x$  is a positive integer, then  $x^2 \leq 3 \implies x \leq \sqrt{3}$ .  
 Now  $\sqrt{3} \approx 1.7$  and so  $x = 1$ .
- (ii) The proposition is false. You should have tried proving it using DeMorgan's Laws and failed.  
 Now, find values of  $x$  and  $y$  that make the statement false.  
 Let  $x = 0$  and  $y = 1$ .  $\sim (x > 1) \vee \sim (y \leq 0)$  is True.  
 $(x \leq 1) \wedge (y > 0)$  is also True.  
 Thus,  $\sim ((x \leq 1) \wedge (y > 0))$  is False and the proposition is False.

6. (i) (a)  $(\sim (x > 1) \implies \sim (y \leq 0)) \equiv (\sim \sim (x > 1) \vee \sim (y \leq 0))$  [Thm 1.4.2, 6]  
 $\equiv ((x > 1) \vee (y > 0))$  [Thm 1.4.2, 4]
- (b)  $((y \leq 0) \implies (x > 1)) \equiv (\sim (y \leq 0) \vee (x > 1))$  [Thm 1.4.2, 6]  
 $\equiv ((y > 0) \vee (x > 1))$

The statements are the same!

- (ii)  $\sim (\sim (P \vee Q) \wedge \sim Q) \equiv (\sim \sim (P \vee Q) \vee \sim \sim Q)$  [DeMorgan's Laws]  
 $\equiv ((P \vee Q) \vee Q)$  [Thm 1.4.2, 4]  
 $\equiv (P \vee Q)$  [Thm 1.4.2, 2]

7. (i) Every real number that is not zero is either positive or negative.  
 The statement is true.
- (ii) The square root of every natural number is also a natural number.  
 The statement is false (consider  $n = 2$ ).
- (iii) Every student in MATH221 can correctly solve at least one assigned problem.  
 Lecturers are yet to work out if this is true or false!

8. (i)  $\forall x \in \mathbb{R}, \forall y \in \mathbb{R}, (xy = 0 \implies (x = 0 \wedge y = 0))$   
 The statement is false (consider  $x = 1$  and  $y = 0$ ).
- (ii)  $\forall x \in \mathbb{R}, \exists y \in \mathbb{Z}, x \leq y$   
 The statement is true.
- (iii)  $\exists$  student  $S$  in MATH221,  $\forall$  lecturer's jokes  $j$ ,  $S$  hasn't laughed at  $j$   
 Caz hopes this is false.

9. Let  $H$  be the set of all people (human beings).

- (i)  $P \equiv \exists p \in H, \forall q \in H, p \text{ loves } q.$   
 $\sim P \equiv \sim (\exists p \in H, \forall q \in H, p \text{ loves } q).$   
 $\equiv \forall p \in H, \sim (\forall q \in H, p \text{ loves } q).$   
 $\equiv \forall p \in H, \exists q \in H, p \text{ doesn't love } q.$

In a nice world,  $P$  is true!

- (iii)  $P \equiv \exists p \in H, \exists q \in H, p \text{ loves } q.$   
 $\sim P \equiv \sim (\exists p \in H, \exists q \in H, p \text{ loves } q).$   
 $\equiv \forall p \in H, \sim (\exists q \in H, p \text{ loves } q).$   
 $\equiv \forall p \in H, \forall q \in H, p \text{ doesn't love } q.$

$P$  is definitely true!

- (v)  $P \equiv \forall x \in \mathbb{Q}, x \in \mathbb{Z}$   
 $\sim P \equiv \sim (\forall x \in \mathbb{Q}, x \in \mathbb{Z})$   
 $\equiv \exists x \in \mathbb{Q}, x \notin \mathbb{Z}$   
 $\sim P$  is true.

- (vii)  $P \equiv \exists n \in \mathbb{N}, n \text{ is not prime.}$   
 $\sim P \equiv \sim (\exists n \in \mathbb{N}, n \text{ is not prime})$   
 $\equiv \forall n \in \mathbb{N}, n \text{ is prime.}$

$P$  is true.

10. (i)  $\forall x \in \mathbb{R}, (x > 1 \implies x > 0)$

This statement is **true**.

Clearly,  $0 < 1 < x$ , so  $x > 0$ .

(iii)  $\exists x \in \mathbb{R}, (x > 1 \implies x^2 > x)$

This statement is **true**.

Let  $x = 2$ .

Then  $x > 1$  and  $x^2 = 4 > 2 = x$

(v)  $\forall x \in \mathbb{R}, \forall y \in \mathbb{R}, x^2 + y^2 = 9$

This statement is **false**.

Let  $x = 1$  and  $y = 1$ ,

then  $x^2 + y^2 = 2 \neq 9$ .

(vii)  $\exists x \in \mathbb{R}, \forall y \in \mathbb{R}, x^2 + y^2 \geq 0$

This statement is **true**.

Let  $x = 0$ . For each  $y \in \mathbb{R}, y^2 \geq 0$ ,

and we have  $x^2 + y^2 = y^2 \geq 0$ .

11. (i) (a)  $\sim (\forall \varepsilon > 0, \exists x \neq 0, |x| < \varepsilon) \equiv \exists \varepsilon > 0, \sim (\exists x \neq 0, |x| < \varepsilon)$

$\equiv \exists \varepsilon > 0, \forall x \neq 0, |x| \geq \varepsilon$

(b) The negation of the statement is false.

(c) For any  $\varepsilon > 0$ , we can take  $x = \frac{\varepsilon}{2}$  and we have  $x \neq 0$  but  $|x| < \varepsilon$ .

(ii) (a)  $\sim (\exists y \in \mathbb{R}, \forall x \in \mathbb{R}, y < x^2) \equiv \forall y \in \mathbb{R}, \sim (\forall x \in \mathbb{R}, y < x^2)$

$\equiv \forall y \in \mathbb{R}, \exists x \in \mathbb{R}, y \geq x^2$

(b) The negation of the statement is false.

(c) Let  $y = -1$ . We know  $x^2 > 0$  for all  $x \in \mathbb{R}$  i.e.  $x^2 > y$ .

(ii)  $P \equiv \forall p \in H, \forall q \in H, p \text{ loves } q.$

$\sim P \equiv \sim (\forall p \in H, \forall q \in H, p \text{ loves } q).$

$\equiv \exists p \in H, \sim (\forall q \in H, p \text{ loves } q).$

$\equiv \exists p \in H, \exists q \in H, p \text{ doesn't love } q.$

In a perfect world,  $P$  is true!

(iv)  $P \equiv \forall p \in H, \exists q \in H, p \text{ loves } q.$

$\sim P \equiv \sim (\forall p \in H, \exists q \in H, p \text{ loves } q).$

$\equiv \exists p \in H, \sim (\exists q \in H, p \text{ loves } q).$

$\equiv \exists p \in H, \forall q \in H, p \text{ doesn't love } q.$

In our world,  $P$  is probably true!

(vi)  $P \equiv \sim (\forall n \in \mathbb{N}, \exists p \in \mathbb{N}, n = 2p)$

$\equiv \exists n \in \mathbb{N}, \sim (\exists p \in \mathbb{N}, n = 2p)$

$\equiv \exists n \in \mathbb{N}, \forall p \in \mathbb{N}, n \neq 2p$

$\sim P \equiv \sim \sim (\forall n \in \mathbb{N}, \exists p \in \mathbb{N}, n = 2p)$

$\equiv \forall n \in \mathbb{N}, \exists p \in \mathbb{N}, n = 2p$

$P$  is true.

(viii)  $P \equiv \forall \text{ triangle } T, T \text{ is a right triangle.}$

$\sim P \equiv \sim (\forall \text{ triangle } T, T \text{ is a right triangle})$

$\equiv \exists \text{ triangle } T, T \text{ is not a right triangle.}$

$\sim P$  is true.

(ii)  $\forall x \in \mathbb{R}, (x > 1 \implies x > 2)$

This statement is **false**.

Let  $x = 1.5$ . Then  $x > 1$  but  $x < 2$ .

(iv)  $\exists x \in \mathbb{R}, \left( x > 1 \implies \frac{x}{x^2 + 1} < \frac{1}{3} \right)$

This statement is **true**.

Let  $x = 3$ . Then  $x > 1$  and

$\frac{x}{x^2 + 1} = \frac{3}{10} < \frac{1}{3}.$

(vi)  $\forall x \in \mathbb{R}, \exists y \in \mathbb{R}, x^2 < y + 1$

This statement is **true**.

For  $x \in \mathbb{R}$ , let  $y = x^2$ .

Then clearly  $x^2 < y + 1$ .

(viii)  $\exists x \in \mathbb{R}, \exists y \in \mathbb{R}, (x < y \implies x^2 < y^2)$

This statement is **true**.

Let  $x = 0$  and  $y = 1$ . Then  $x < y$  and

$x^2 = 0 < 1 = y^2$ .

$$\begin{aligned}
\text{(iii) (a)} \quad & \sim \left( \forall y \in \mathbb{R}, \forall x \in \mathbb{R}, \left( x < y \implies x < \frac{x+y}{2} < y \right) \right) \\
& \equiv \exists y \in \mathbb{R}, \sim \left( \forall x \in \mathbb{R}, \left( x < y \implies x < \frac{x+y}{2} < y \right) \right) \\
& \equiv \exists y \in \mathbb{R}, \exists x \in \mathbb{R}, \sim \left( x < y \implies x < \frac{x+y}{2} < y \right) \\
& \equiv \exists y \in \mathbb{R}, \exists x \in \mathbb{R}, \left( x < y \wedge \left( \frac{x+y}{2} \leq x \vee \frac{x+y}{2} \geq y \right) \right) \\
& \equiv \exists y \in \mathbb{R}, \exists x \in \mathbb{R}, (x < y \wedge (y \leq x \vee x \geq y))
\end{aligned}$$

(b) The negation of the statement is false.

(c) Clearly,  $x < y \wedge (y \leq x \vee x \geq y)$  is equivalent to  $x < y \wedge x = y$ , which is impossible.